

Approximate Matching

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Try the code together with your neighbour

02input/raw_to_input.R





String distance

Default (Optimal String Alignment distance)

Count number of character deletions, insersions, substitutions and transpositions (of adjacent characters)

```
library(stringdist)
stringdist("Ross Ihaka", "Robert Gentleman")
```

```
## [1] 12
```





Exact Matching with match

```
lookup <- c("Alice", "Bob", "Carol", "Danny")
raw <- c("Bob", "Carl", "Rob", "bob", "Dan", "Alice")
i <- match(raw, lookup)
data.frame(raw=raw, matched=lookup[i])</pre>
```

```
##
     raw matched
## 1
      Bob
              Bob
## 2
     Carl <NA>
## 3
      Rob
             <NA>
## 4
     bob
             <NA>
## 5
             <NA>
     Dan
## 6 Alice
            Alice
```





Approximate Matching with stringdist::amatch

```
library(stringdist)
j <- amatch(raw, lookup, maxDist=2)
data.frame(raw=raw, matched=lookup[i], amatched=lookup[j])</pre>
```

```
##
      raw matched amatched
## 1
      Bob
             Bob
                      Bob
## 2
     Carl <NA>
                    Carol
## 3
    Rob < NA >
                      Bob
## 4
     bob <NA>
                      Bob
## 5
             <NA>
      Dan
                    Danny
## 6 Alice
            Alice
                    Alice
```



 \rightarrow Match with closest match, and distance < 2.



Optimal string alignment?

```
stringdist("Robert Gentleman", "Gentleman, Robert")
## [1] 15
stringdist("Robert Gentleman", "Ross Ihaka")
## [1] 12
\rightarrow OSA wil give a false match (if we allow maxDist of 12)
```





Alternative: cosine distance

```
## [1] 0.9139337
```

Notes

- Based on counting co-occurrence of character *q*-grams (here: pairs).
- Always between 0 and 1





More on amatch

Example

```
amatch(raw, lookup, method="cosine", maxDist=0.5, q=3)
```





Assignment

Merge data from the companies dataset with data from backbone.csv.

- Using approximate matching on the "name" and "company" column.
- Think about and try different distance functions and maxDist
- Keep your best solution
- · Remove rows that cannot be matched
- Write to 02input/myinput.csv





More on string distances

Main idea

Define a sence of distance between two text strings

Distance

A function d(s, t) that takes two arguments and

- returns a nonnegative number,
- returns zero if and only if s = t
- is symmetric: d(s,t) = d(t,s)
- is the length of a shortest path between s and t: $d(s,t) \leq d(s,u) + d(u,t)$

Note

Some string distances violate one or more of the above assumptions.





Distance types

- Edit based
- q-gram based
- Heuristic





Edit-based string distances

Idea

- 1. Choose basic steps to alter a string
- 2. Find the smallest nr of steps that changes s into t
- 3. The distance equals the nr of steps needed.

Basic steps

- deletion: hihi o hii
- insertion: $hihi \rightarrow hihih$
- substitution: hihi ightarrow hiha
- transposition: $\mathsf{hihi} \to \mathsf{ihhi}$





Edit-based string distances

Distance	Allowed operation			
	substitution	deletion	insertion	transposition
Hamming	~	×	×	*
LCS	×	~	V	*
Levenshtein	✓	~	~	×
OSA	✓	~	~	✓ *
Damerau-	✓	~	~	✓
Levenshtein				

^{*}Substrings may be edited only once.





Example

Levenshtein distance

$$leia \xrightarrow[+1]{\text{sub}} lela \xrightarrow[+1]{\text{ins}} leela$$

Longest common subsequence distance

$$\textit{leia} \xrightarrow[+1]{\text{del}} \textit{lea} \xrightarrow[+1]{\text{ins}} \textit{leea} \xrightarrow[+1]{\text{ins}} \textit{leela}$$





q-gram based distances (I)

Algorithm

- Tabulate substrings of length q (= q-gram profile)
- Compute a distance between the profiles

Example

```
2-gram profile of banana
ba an na
1 1 2
```





q-gram based distances (II)

q-gram distance

Manhattan distance between q-gram profiles

$$\sum_{\text{qgram}} |n_{\text{qgram}}(s) - n_{\text{qgram}}(t)|$$

Cosine-distance

1 minus the cosine of the angle between the profiles

$$1 - \frac{\mathbf{n}(s) \cdot \mathbf{n}(t)}{\|\mathbf{n}(s)\| \|\mathbf{n}(t)\|}$$

Node

- Does not satisfy the 'identity' demand.
- Often one chooses q = 2 of q = 3





Jaro-Winkler distance

Jaro distance

$$d_j(s,t) = 1 - rac{1}{3}\left(rac{m}{|s|} + rac{m}{|t|} + rac{m+T}{m}
ight)$$

- m number of matching characters (within a window)
- T number of matches that need swapping
- |s|, |t| number of characters in s, t.

Jaro-Winkler distance

$$d_{jw}(s,t) = [1 - p\ell(s,t)]d_j(s,t)$$

- $\ell(s,t)$ length of longest equal prefix (up to 4 characters)
- p a number between 0 and 0.25 (usually 0.1)





Soundex

Algorithm

- Strings are appointed a code: same code means 'sounds the same'
- Equal codes: distance zero, otherwise 1

Example

- Farnsworth \rightarrow H652
- Fnarswort \rightarrow H562

Note

- Based on English pronunciation
- Many extensions exist (see the phonics R package)





Which one to use?

Considerations

- Fixed versus variable structure/length
- Why would strings differ? (typos, speech, deliberate changes)
- Performance



