

# **Approximate Matching**

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Tokyo | Institute of Statistical Mathemathics | 2020





# Try the code together with your neighbour

02input/raw\_to\_input.R





# **String distance**

### **Default (Optimal String Alignment distance)**

Count number of character deletions, insersions, substitutions and transpositions (of adjacent characters)

```
library(stringdist)
stringdist("Ross Ihaka", "Robert Gentleman")
```

```
## [1] 12
```





### Exact Matching with match

```
##
      raw matched
## 1
      Bob
               Bob
## 2
      Carl <NA>
      Rob
## 3
              <NA>
## 4
      bob
              < NA >
## 5
      Dan
              <NA>
  6 Alice
             Alice
##
```





# Approximate Matching with stringdist::amatch

```
library(stringdist)
j <- amatch(raw, lookup, maxDist=2)
data.frame(raw=raw, matched=lookup[i], amatched=lookup[j])</pre>
```

```
raw matched amatched
##
## 1
     Bob
            Bob
                    Bob
## 2
     Carl <NA> Carol
## 3
   Rob <NA>
                    Bob
## 4
    bob <NA>
                    Bob
## 5
     Dan <NA>
                  Danny
## 6 Alice
         Alice
                 Alice
```

 $\rightarrow$  Match with closest match, and distance  $\leq 2$ .





# **Optimal string alignment?**

```
stringdist("Robert Gentleman", "Gentleman, Robert")

## [1] 15

stringdist("Robert Gentleman", "Ross Ihaka")

## [1] 12

→ OSA wil give a false match (if we allow maxDist of 12)
```





### **Alternative: cosine distance**

#### **Notes**

## [1] 0.9139337

- Based on counting co-occurrence of character *q*-grams (here: pairs).
- Always between 0 and 1





#### More on amatch

### **Example**

```
amatch(raw, lookup, method="cosine", maxDist=0.5, q=3)
```





# **Assignment**

Merge data from the companies dataset with data from backbone.csv.

- Using approximate matching on the "name" and "company" column.
- Think about and try different distance functions and maxDist
- Keep your best solution
- Remove rows that cannot be matched
- Write to O2input/myinput.csv





# More on String distances





# More on string distances

#### Main idea

Define a sence of distance between two text strings

#### **Distance**

A function d(s, t) that takes two arguments and

- returns a nonnegative number,
- returns zero if and only if s = t
- is symmetric: d(s,t) = d(t,s)
- is the length of a shortest path between s and t:  $d(s,t) \leq d(s,u) + d(u,t)$

#### Note

Some string distances violate one or more of the above assumptions.



# **Distance types**

- Edit based
- q-gram based
- Heuristic





# **Edit-based string distances**

#### Idea

- 1. Choose basic steps to alter a string
- 2. Find the smallest nr of steps that changes s into t
- 3. The distance equals the nr of steps needed.

### **Basic steps**

- deletion:  $hihi \rightarrow hii$
- insertion: hihi  $\rightarrow$  hihih
- substitution: hihi  $\rightarrow$  hiha
- transposition: hihi  $\rightarrow$  ihhi





# **Edit-based string distances**

	Allowed operation			
Distance	substitution	deletion	insertion	transposition
Hamming	<b>~</b>	×	×	×
LCS	×	<b>V</b>	<b>V</b>	×
Levenshtein	<b>~</b>	<b>~</b>	<b>~</b>	×
OSA	V	<b>~</b>	<b>~</b>	<b>✓</b> *
Damerau-	<b>✓</b>	<b>~</b>	<b>~</b>	<b>✓</b>
Levenshtein				

<sup>\*</sup>Substrings may be edited only once.





# **E**xample

#### Levenshtein distance

$$leia \xrightarrow[+1]{\text{sub}} lela \xrightarrow[+1]{\text{ins}} leela$$

### Longest common subsequence distance

$$leia \xrightarrow[+1]{\text{del}} lea \xrightarrow[+1]{\text{ins}} leea \xrightarrow[+1]{\text{ins}} leela$$





# *q*-gram based distances (I)

### **Algorithm**

- Tabulate substrings of length q (= q-gram profile)
- Compute a distance between the profiles

### **Example**

```
2-gram profile of banana
ba an na
1 1 2
```





# q-gram based distances (II)

#### q-gram distance

Manhattan distance between q-gram profiles

$$\sum_{ ext{qgram}} |n_{ ext{qgram}}(s) - n_{ ext{qgram}}(t)|$$

#### Cosine-distance

1 minus the cosine of the angle between the profiles

$$1 - \frac{\mathbf{n}(s) \cdot \mathbf{n}(t)}{\|\mathbf{n}(s)\| \|\mathbf{n}(t)\|}$$

#### Note

- Does not satisfy the 'identity' demand.
- Often one chooses q = 2 of q = 3





### Jaro-Winkler distance

#### Jaro distance

$$d_j(s,t)=1-rac{1}{3}\left(rac{m}{|s|}+rac{m}{|t|}+rac{m+T}{m}
ight)$$

- *m* number of matching characters (within a window)
- T number of matches that need swapping
- |s|, |t| number of characters in s, t.

#### Jaro-Winkler distance

$$d_{jw}(s,t) = [1 - p\ell(s,t)]d_j(s,t)$$

- $\ell(s,t)$  length of longest equal prefix (up to 4 characters)
- p a number between 0 and 0.25 (usually 0.1)





### **Soundex**

### **Algorithm**

- Strings are appointed a code: same code means 'sounds the same'
- Equal codes: distance zero, otherwise 1

### **Example**

- Farnsworth  $\rightarrow$  H652
- Fnarswort  $\rightarrow$  H562

#### Note

- Based on English pronunciation
- Many extensions exist (see the phonics R package)





### Which one to use?

#### **Considerations**

- Fixed versus variable structure/length
- Why would strings differ? (typos, speech, deliberate changes)
- Performance



