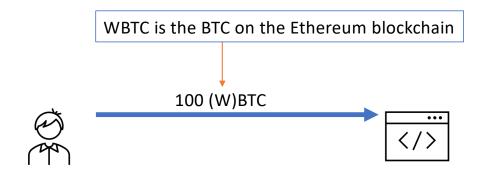
## Robust (Decentralized) Oracle Design

Leifu Zhang

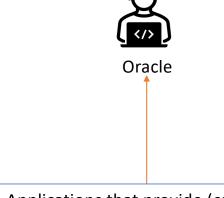
The Hong Kong University of Science and Technology (Guangzhou)

July 2024 @ CMID

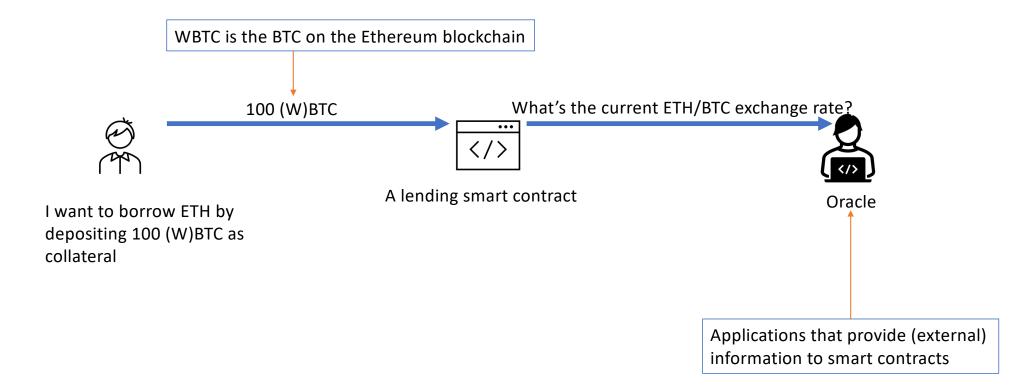


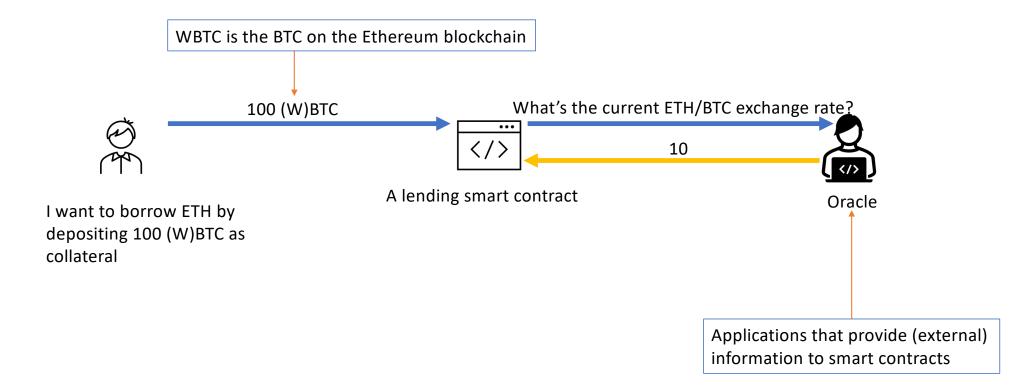
A lending smart contract

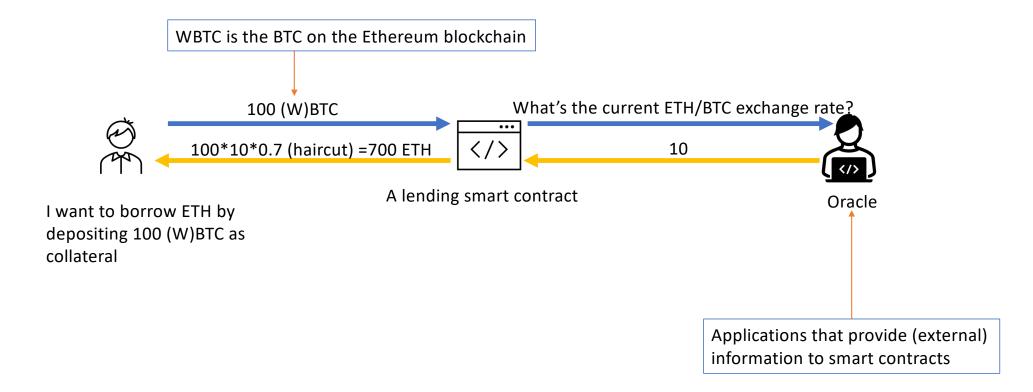
I want to borrow ETH by depositing 100 (W)BTC as collateral

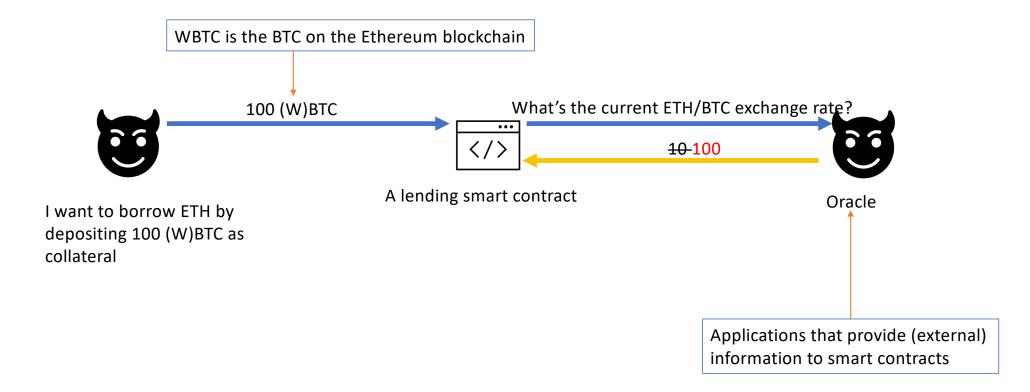


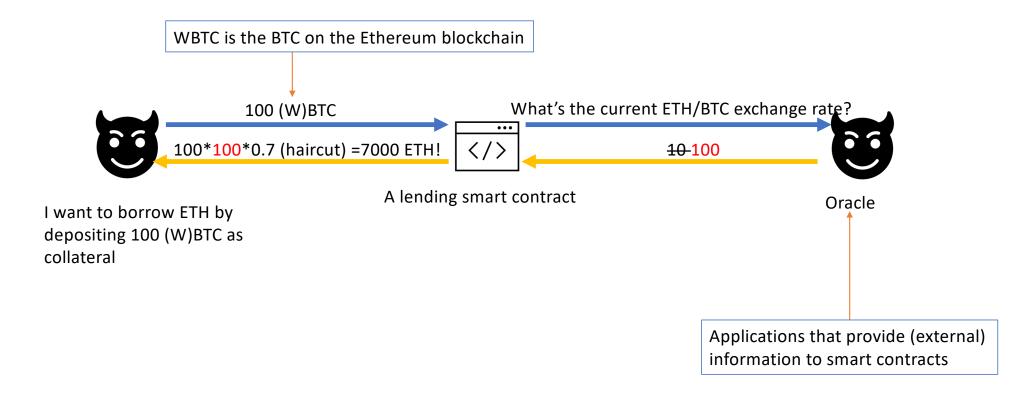
Applications that provide (external) information to smart contracts

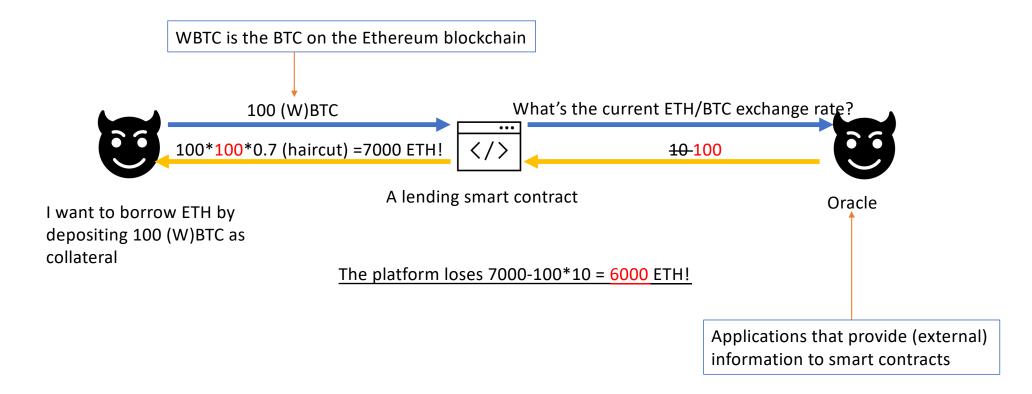


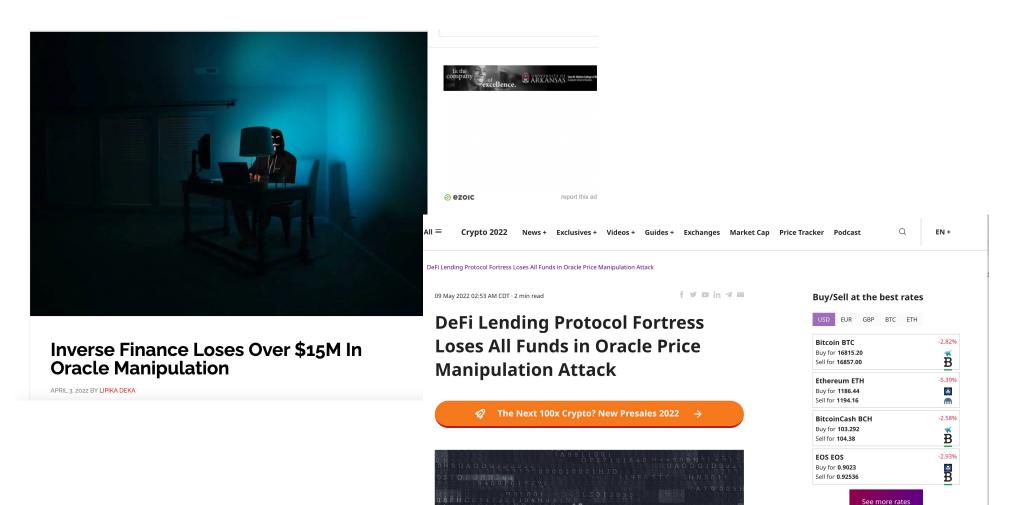












Sources: <a href="https://www.tronweekly.com/inverse-finance-loses-15m-oracle-manipulation/">https://www.tronweekly.com/inverse-finance-loses-15m-oracle-manipulation/</a>
<a href="https://cryptonews.com/news/defi-lending-protocol-fortress-loses-all-funds-oracle-price-manipulation-attack.htm">https://cryptonews.com/news/defi-lending-protocol-fortress-loses-all-funds-oracle-price-manipulation-attack.htm</a>

## The importance of oracles

- Oracles are the cornerstone of DeFi
  - Decentralized lending platforms
  - Prediction markets
  - Insurance contracts
  - NFT games
  - (Many) stablecoins

• ...



## The oracle problem

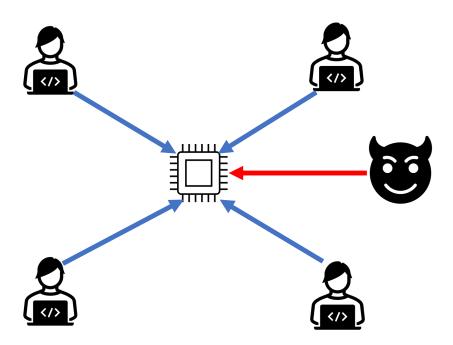
• How can we ensure the information provided by oracles is accurate?

# The oracle problem

- How can we ensure the information provided by oracles is accurate?
- Single source → single point of failure

### The oracle problem

- How can we ensure the information provided by oracles is accurate?
- Single source → single point of failure → decentralization!



Q: Can we find a robust compensation mechanism?

#### **Definition**

A compensation mechanism is robust if, under that mechanism, there is an equilibrium in which truthful reporting is the unique optimal response for strategic nodes regardless of the adversary's strategy.

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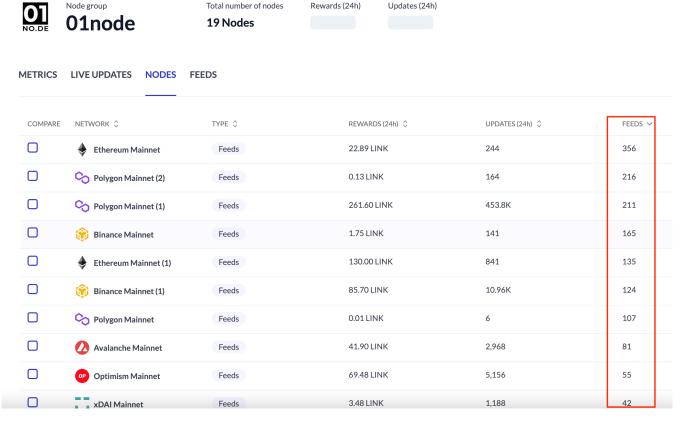
- A: Without identifying an honest node, generally no
- Takeaway: "A" limit of decentralization

• Q: What is the optimal way to aggregate information under the worst-case scenario?

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- Key observations:
  - 1. Obtaining consensus = unsupervised learning with contaminated data
  - 2. The popular aggregating method ignores the multi-dimensional structure of decentralized oracles---each node usually covers many cryptocurrencies

## The high-dimensional structure



Source: https://market.link/nodes/568cedcc-46f3-49e4-84c7-a9d7d5e23a0d/nodes

• Q: What is the optimal way to aggregate information under the worst-case scenario?

- A: A filtering algorithm can dramatically improve the consensus by utilizing this multi-dimensional structure
  - Adversarial nodes which look "normal" in every single dimension could be detected from a "global" view
  - Approaching the theoretical limit

#### Related literature

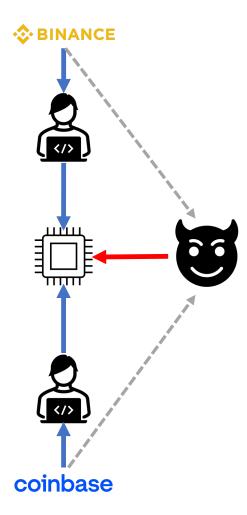
- · Oracle design
  - F. Zhang et al. (2016), F. Zhang et al. (2020), Breidenbach et al. (2021)
  - Contribution: 1) "A" limit of decentralization; 2) connecting machine learning to oracle design
- Information elicitation
  - McCarthy (1956), Savage (1971), Prelec (2004), Miller et al. (2005), P. Zhang and Chen (2014), Lambert (2019), Gao et al. (2019)
  - · Contribution: Getting an impossible result under the adversarial environments
- Manipulation in traditional capital markets
  - Gandhi et al. (2019), A. Zhang (2022)
  - Contribution: Shedding light on designing replacements for the London Inter-Bank Offered Rate (LIBOR)
- Byzantine fault tolerance
  - Lamport et al. (1982), Amoussou-Guenou et al. (2021), Halaburda et al. (2021)
- Machine learning
  - Lai et al. (2016), Diakonikolas et al. (2016, 2017, 2019), Charikar et al. (2017), Zhu et al. (2022)

#### Setting

- n (a large number of) nodes;  $\varepsilon n$  nodes are controlled by an adversary
- The rest nodes are risk-neutral and strategic: Maximizing the expected payoffs given by the designer
- Ground truth  $X \sim U(\mathbb{R}^d)$
- Each strategic node has a private signal

$$\mathbf{s}_i = \mathbf{X} + \mathbf{e}_i$$

- $\mathbb{E}[\mathbf{e}_i] = \mathbf{0}$  and  $\mathbf{e}_i$  has a bounded covariance matrix
- Key assumption: The adversary observes strategic nodes' private signals



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#### Timing:

- 1. The designer announces a compensation mechanism
- 2. Each node submits a report
- The designer pays each node and outputs a consensus

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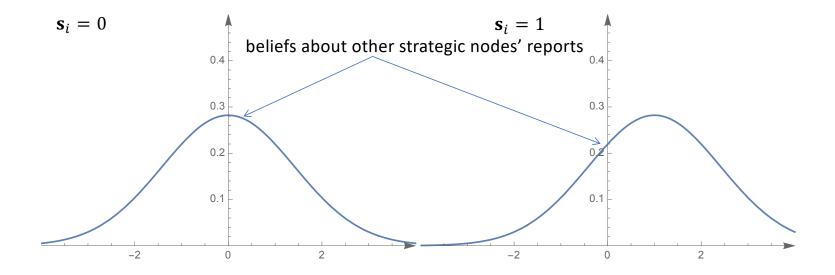
#### Goals:

- Find a robust compensation mechanism 1.
- Find a robust consensus  $\hat{\mathbf{X}}$  that is close to  $\mathbf{X}$

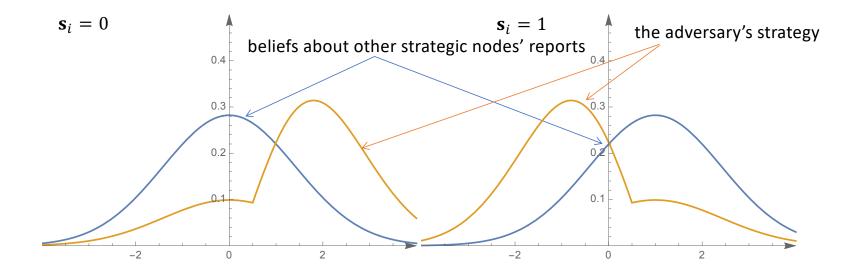
robust = good given the adversary's any strategy

Part 1: (No) Robust compensation mechanism

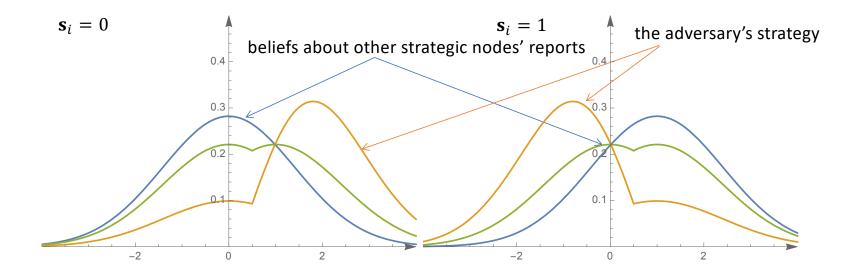
• Suppose  $\mathbf{s}_i = \mathbf{X} + \mathcal{N}(0,1)$  and consider node i's decision problem



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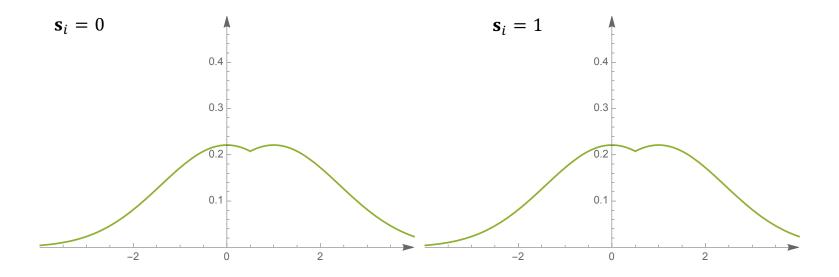


- Suppose  $\mathbf{s}_i = \mathbf{X} + \mathcal{N}(0,1)$  and consider node i's decision problem
- Greenline =  $(1 \varepsilon)$  \* Blueline +  $\varepsilon$  \* Orangeline
- Node *i*'s beliefs about other nodes' reports

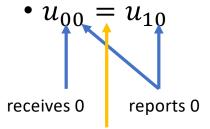


**Lemma** [implied by an observation in robust statistics]

Under a mild sufficient condition, the adversary has a reporting strategy such that even if node i may have different private information, node i's beliefs about other nodes' reports are unchanged.



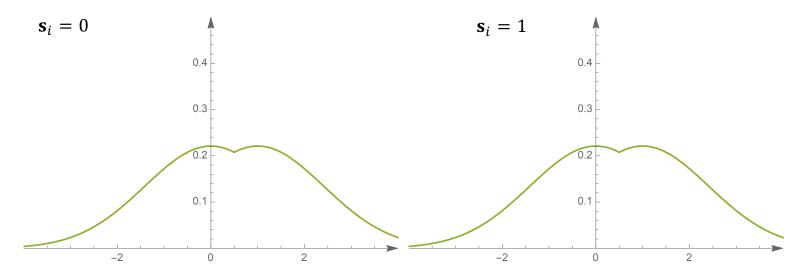


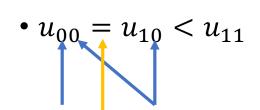


Private signal  $\mathbf{s}_i$ 

_	0	1
0	$u_{00}$	$u_{01}$
1	$u_{10}$	$u_{11}$

Report  $\mathbf{r}_i$ 





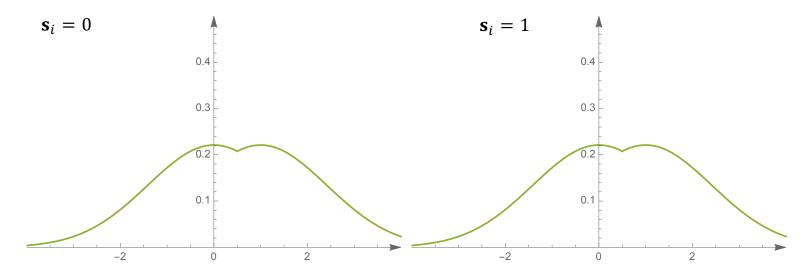
reports 0

receives 0

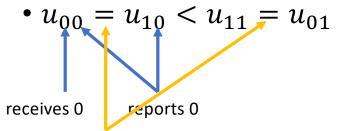
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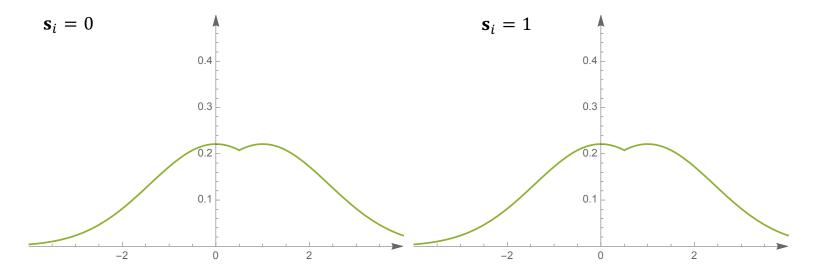




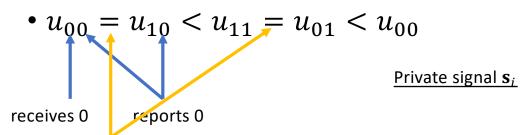
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Report  $\mathbf{r}_i$ 

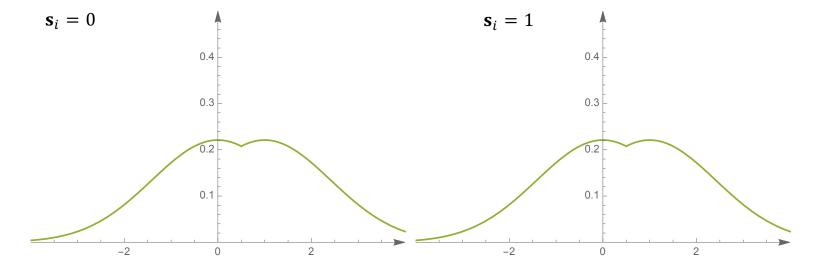






	0	1
0	$u_{00}$	$u_{01}$
1	$u_{10}$	$u_{11}$

Report  $\mathbf{r}_i$ 



- Let  $Q(\cdot; \mathbf{s})$  be a strategic node's posterior belief about another strategic node's private signal after observing  $\mathbf{s}$
- Let  $d_{\mathrm{TV}}$  denotes the total variation distance

$$d_{\mathrm{TV}}(P, P') \coloneqq \sup_{E \in \mathfrak{B}} [P(E) - P'(E)]$$

• Let  $\mathcal{D}$  be the dataset of all reports

#### **Theorem**

If there are are two different signal realizations, s and s', such that

$$d_{\mathrm{TV}}(Q(\cdot;\mathbf{s}),Q(\cdot;\mathbf{s}')) \leq \frac{\varepsilon}{1-\varepsilon},$$

then for any compensation mechanism  $\mathcal M$  as a function of  $\mathcal D$ ,  $\mathcal M$  cannot be robust.

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the private signal's precision

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then for any compensation mechanism  $\mathcal{M}$  as a function of  $\mathcal{D}$ ,  $\mathcal{M}$  cannot be robust.

the adversary's power

 Economic intuition: Has to reward truth-telling and/or punish misreporting; but no way to check whether node i misreports or not given the adversary's strategy

• Mathematical "intuition": Data contamination breaks the stochastic relevance condition [which is the necessary condition to have a strict truth-telling eqm (P. Zhang and Chen, 2014)]

Part 2: Robust consensus

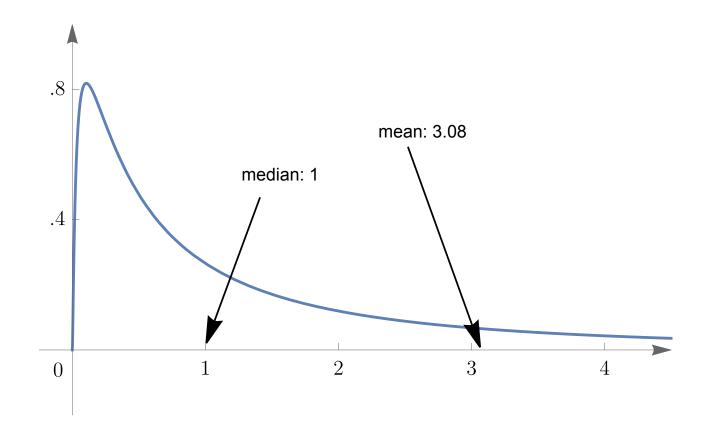
#### Robust consensus: Overview

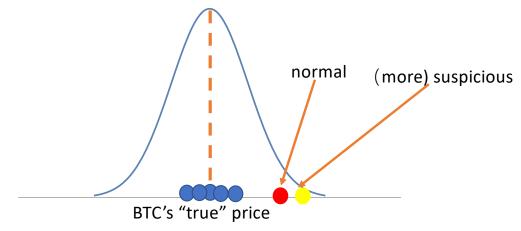
• The most popular consensus mechanism:

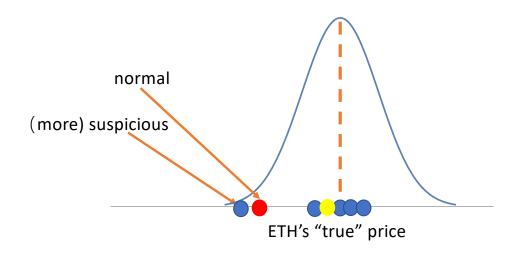
Taking the (coordinate-wise) median

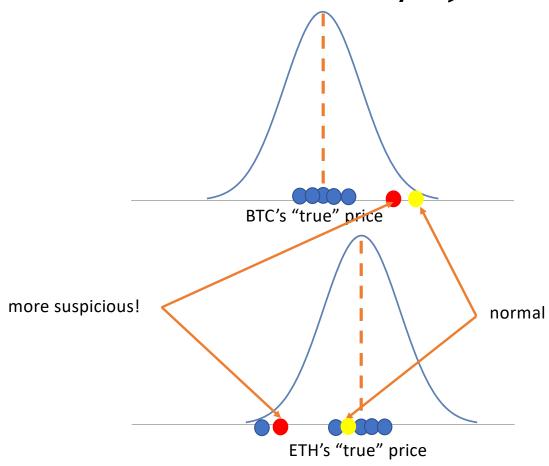
- Bad if the noise term is asymmetric even without an adversary!
- Not a bad estimator if symmetric; but is far from optimal under a highdimensional environment!
  - Even the best 1-d estimator can yield a  $L^2$ -norm error  $\geq C\sqrt{\varepsilon d}$  (Folklore)
- Recent machine learning algorithms---unsupervised learning with contaminated datasets--- could yield a consensus that nearly achieves the error's theoretical lower bound without assuming symmetry!

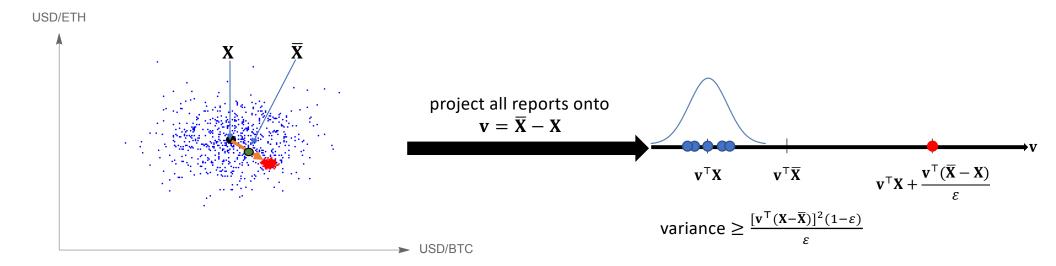
## The current method may fail



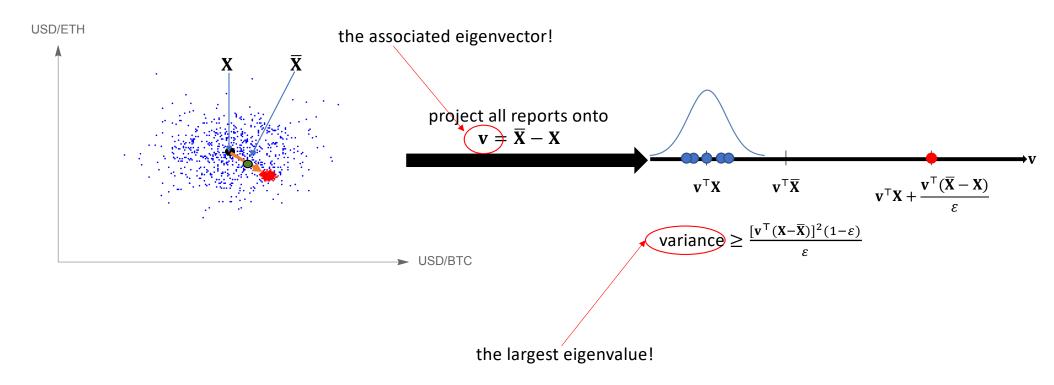








[The high-level idea (Diakonikolas et al., 2016, 2017; Diakonikolas and Kane, 2021): Using the covariance matrix!]



[The high-level idea (Diakonikolas et al., 2016, 2017; Diakonikolas and Kane, 2021): Using the covariance matrix!]

The filtering algorithm (Diakonikolas et al., 2016, 2017; Zhu et al., 2022)

- 1. Calculate the empirical covariance of the dataset  ${\mathcal D}$  and find the largest eigenvalue
- 2. If the largest eigenvalue is small, then return the empirical mean of  $\mathcal{D}$
- 3. Otherwise,
  - project  $\mathcal{D}$  onto the eigenvector that is associated with the largest eigenvalue;
  - Downweight each point according to the distance between its projection and the projection of the empirical mean, and obtain a new dataset  $\widetilde{\mathcal{D}}$ ;
  - replace  ${\mathcal D}$  with  $\widetilde{{\mathcal D}}$  and return to Step 1

theoretical lower bound

#### Theorem (Zhu et al., 2022)

The filtering algorithm will output a consensus  $\widehat{X}$  such that

$$\|\widehat{\mathbf{X}} - \mathbf{X}\|_{2} \le \sigma \sqrt{\varepsilon} \left( \frac{1}{\sqrt{1-\varepsilon}} + \frac{\sqrt{2}}{1-2\varepsilon} \right),$$

where  $\sigma^2$  is an upper bound on the  $L^2$ -norm of the noise term's covariance matrix.

theoretical lower bound

Theorem (Zhu et al., 2022)

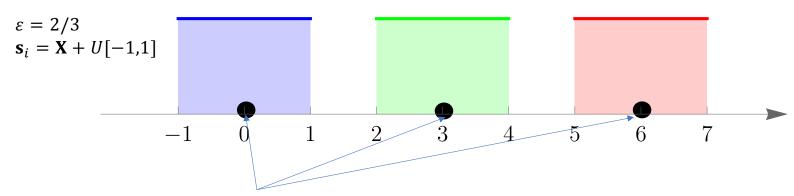
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Best 1-d estimator:  $\geq \sigma \sqrt{\varepsilon d}$ 

- Charikar et al. (2017)
  - There is no algorithm can return a unique consensus that is close to the ground truth
  - But we can return a list of candidates, in which at least one of them is "good"
  - A clever clustering algorithm



Which point is the ground truth?

#### Concluding remarks

• In general, no perfect decentralized solution to the oracle problem

Machine learning can improve the consensus substantially

All results also shed light on designing replacements for LIBOR

# Thank you! 🕥