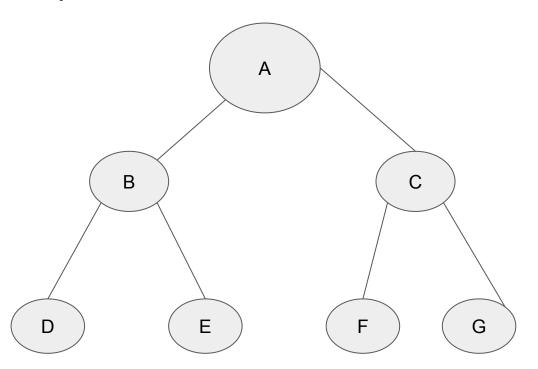
CWB-Module- 4

Tree and Graph Traversals

Traversals

Pre-order (Root, LST, RST)
Post-order (LST, RST, Root)
In-order (LST, Root, RST)

Example 1

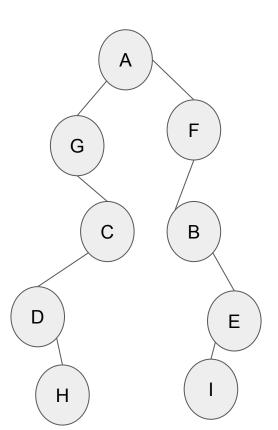


Pre-order - ABDECFG

Post-order - DEBFGCA

In-order - DBEAFCG

Example 2



Pre-order: AGCDHFBEI

Inorder: GDHCABIEF

Post-order: HDCGIEBFA

Binary Search Tree

A Binary Search Tree (BST) is a node-based binary tree data structure that maintains the following properties:

- 1. Left Subtree Property: All nodes in the left subtree of a node have values less than the node's value.
- 2. Right Subtree Property: All nodes in the right subtree of a node have values greater than the node's value.
- 3. No Duplicate Nodes: Typically, BSTs do not contain duplicate values (though some implementations may allow them in either the left or right subtree).
- 4. Recursive Structure: Both the left and right subtrees must also be binary search trees.

BST Example

```
8
3 10
 6
4 7 13
```

- All left children are smaller than their parent.
- All right children are larger than their parent.

Important Notes

1. Inorder traversal of BST will always produce ascending order.

```
Preorder(root)
{

print(root->data)

Preorder(root->lc)

Preorder(root->rc)
}
```

Postorder

```
Postorder(root)
postorder(root->lc)
postorder(root->rc)
print(root->data)
```

Inorder

```
Inorder(root)
Inordert(root->lc)
print(root->data)
Inorder(root->rc)
```

Queue Data Structure

A queue is a linear data structure that follows the First-In-First-Out (FIFO) principle, meaning the first element added is the first one to be removed. It is widely used in:

- Breadth-First Search (BFS/BFT)
- Level-order traversal in trees
- CPU scheduling
- Print spooling

Basic Queue Operations

Operation	Description	Time Complexity
Enqueue	Adds an element to the rear (end) of the queue	O(1)
Dequeue	Removes and returns the front element	0(1)
Peek/Front	Returns the front element without removing it	0(1)
IsEmpty	Checks if the queue is empty	0(1)
Size	Returns the number of elements in the queue	0(1)

Graph Traversal - BFT

```
BFT(v)
{ visited(v)=1
add(v,Q)
while(Q is not empty)
{ x=delete(Q)
print(x)
for all w adjacent to x
{ if (w is not visited)
visited(w)=1
add(w,Q) } } }
```

BFT Details

- 1. To implement BFT we are using queue data structure.
- 2. Time Complexity = O(V+E)
- 3. It is also called Level Order Traversal.

Applications of BFT

- 1. We can verify given graph is connected or not.
- 2. Using BFT, we can find connected components in the given graph.
- 3. Using BFT, we can check given graph contain cycle or not.
- Using BFT, we can find out shortest path from the given source to every vertex in the given unweighted graph.

DFT (Depth First Traversal)

```
DFT(V)
visited(v)=1
print(v)
for all w adjacent to v
if w is not visited
DFT(w)
```

DFT Details

- 1. To implement DFT, we are using stack.
- 2. Time Complexity = O(V+E)

Applications of DFT

- We can verify given graph contain cycle or not.
- 2. We can verify given graph is connected or not.
- 3. We can find out the number of connected components.
- 4. DFT cannot work finding single shortest path in the given unweighted graph.