



Computational Logic Exercises Module II – World models and logic systems



Preamble: World Models

A world model is a triple $\mathscr{W} = \langle L, D, I \rangle$ where:

(intensional world model)

 $L = \{a\}$ L is an assertional language, with descriptive assertions

D = {f} D is a domain of interpretation, a set of facts

I: $L \rightarrow D$ I is an interpretation function that maps assertions to facts

(extensional world model)

 $L = \langle E, \{C\}, \{P\} \rangle$ L is an assertional language with assertions about: $E = \{e\}$ is a set of (names

of) entities, {C} is a set of concepts, where a concept is a name of a class, {P}

and a set of (object/data) properties.

 $D = \langle E, \{C\}, \{P\} \rangle$ D is a domain of interpretation, a set of facts about the entities $E = \{e\}$, their

classes C⊆ E and properties P⊆ E× • • • × E

I: $L \rightarrow D$ I is an interpretation function that maps assertions to facts

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Preamble: Logic Systems

A **logic system** is a triple $\mathcal{I} = \langle L, D, I, \models \rangle$ or even $\mathcal{I} = \langle \mathcal{W}, \models \rangle$ where:

 $W = \langle L, D, I \rangle$ W is a world model

 $\models \subseteq D \times L$ \models is an entailment relation between subsets of the domain D (i.e. a specific

model M) and assertions in the language L (i.e. a specific theory)

A (logic) language L can be formally defined by a set of atomic formulas and a grammar, i.e. a set of connectives and formation rules to combine atomic formulas into complex formulas.

For instance:

<wff> ::= <atomic>

<wff> ::= <wff> □ <wff>

<wff> ::= <wff> □ <wff>



World models (I)

Define an extensional world model W constituted by a language L, a domain D and an interpretation function I to formalize the following models M1, M2 and M3.







ANSWER: A possible W is as follows

 $L = \langle E = \{i1, i2, i3\}, C = \{INS, R, G, B, SL\}, R = \emptyset \rangle$ $D = \langle E = \{a, b, c\}, C = \{insect, red, green, blue, sixlegs\}, R = \emptyset \rangle$ $I(i1) = 1; I(i2) = b; I(i3) = c; I(INS) = insect = \{a, b, c\}; I(R) = red = \{a\}; I(G) = green = \{b\}; I(B) = blue = \{c\}; I(SL) = sixlegs = \{a, b, c\}.$



World models (II)

Define a language L' ⊇ L and an interpretation function l' for the same problem that includes at least a complex formula.







ANSWER:

$$L' = L \cup \{INS \sqcap R\}$$

$$I' = (INS \sqcap R) = insect \cap red = \{a\}$$



Logic systems (I)

Define a logic system $L = \langle W, \models \rangle$, and in particular the entailment relation \models starting from the world model W defined before.







ANSWER: we need to provide it for the models M1, M2 and M3:

 $M1 \models INS(i1), M1 \models R(i1), M1 \models SL(i1)$

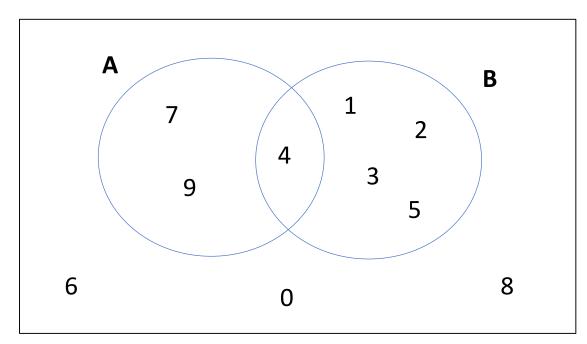
M2 = INS(i2), M2 = G(i2), M2 = SL(i2)

M3 = INS(i3), M3 = B(i3), M3 = SL(i3)



World Models (II)

Define an intentional world model W constituted by a language L, a domain D and an interpretation function I to formalize the following Venn Diagram V



ANSWER:

$$L = \{A, B, A \sqcap B, A \sqcup B\}$$

$$D = \langle E, \{C\}, \{P\} \rangle$$

$$E = \{0, 1, 2, 3, 4, 5, 6, 7, 8\}$$

$$I(A) = \{4, 7, 9\} \in \{C\};$$

$$I(B) = \{1, 2, 3, 4, 5\} \in \{C\};$$

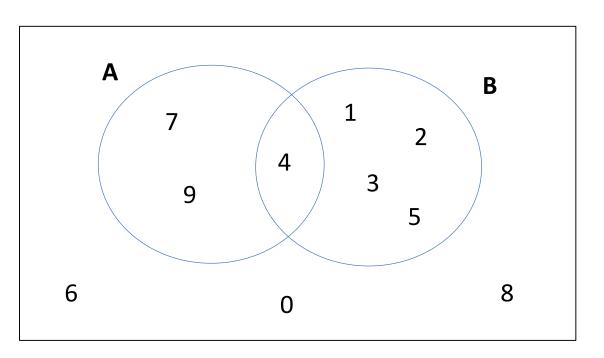
$$I(A \sqcap B) = I(A) \cap I(B) = \{4\} \in \{C\};$$

$$I(A \sqcup B) = I(A) \cup I(B) = \{1, 2, 3, 4, 5, 7, 9\} \in \{C\}^{7}$$



Logic systems (II)

Define a logic system $L = \langle W, \models \rangle$, and in particular the entailment relation \models starting from the world model W defined before.



ANSWER:

 $V \models A$

 $V \models B$

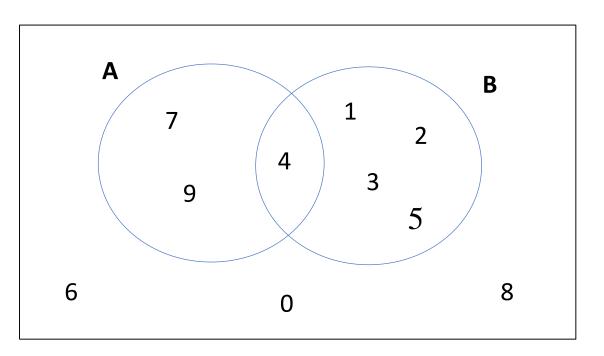
 $V \models A \sqcap B$

 $V \models A \sqcup B$



Logic systems (III)

Define a theory T such that $V \models T$.



ANSWER: In this case, we may take any $T \subseteq L$, for instance $\{A, B, A \sqcup B\}$



Preamble: Reasoning problems

Model checking. Given T and M, check whether $M \models T$

Satisfiability. Given T, check whether there exists M such that $M \models T$

Validity. Given T, check whether for all M, $M \models T$

Unsatisfiability. Given T, check whether there is no M such that $M \models T$

Logical consequence. Given T1, T2 and a set of reference models {M}, check whether

$$T1 \models \{M\} T2$$

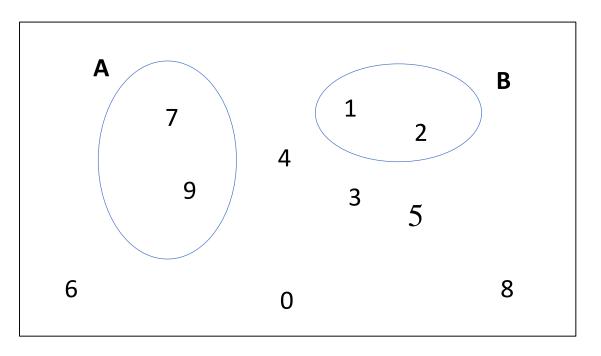
Logical equivalence. Given T1, T2 and a set of reference models {M}, check whether

$$T1 \models \{M\} T2 \text{ and } T2 \models \{M\} T1$$



Logic systems (IV): reasoning – model checking (a)

Given the theory $T = \{A, B, A \sqcup B\}$, and the model V below, check if $V \models T$

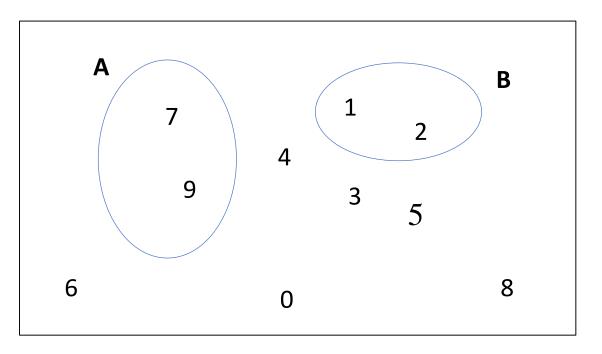


ANSWER: yes



Logic systems (IV): reasoning – model checking (b)

Given the theory $T = \{A, B, A \sqcap B\}$, and the model V below, check if $V \models T$



ANSWER: no



Logic systems (IV): reasoning – model checking (c)

Provide and example of Venn Diagram V' containing two sets A and B, and of a theory $T' \subseteq \{A, B, A \sqcap B, A \sqcup B\}$ such that $V' \not\models T'$.

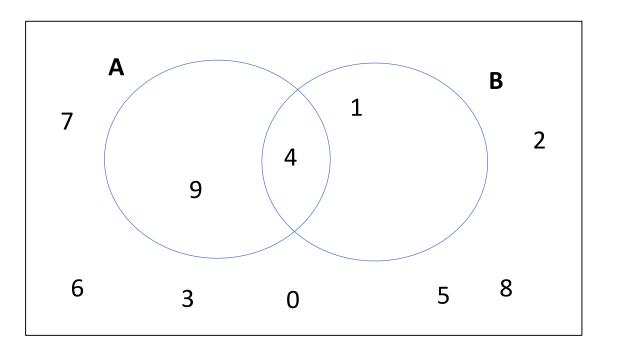
ANSWER: We may take the theory $T' = \{A \sqcap B\}$ and any V' in which A and B are disjoint.



Logic systems (IV): reasoning – satisfiability

Given the theory $T = \{A, B, A \sqcap B\}$, check whether there exists V such that $V \models T$

ANSWER

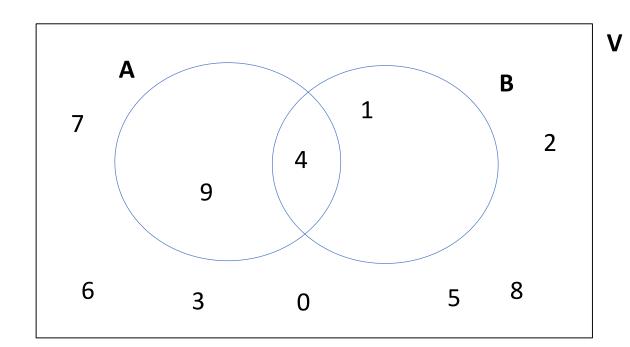




Logic systems (V)

Say which of the following statements are true given the formalization provided before.

- a) $V \models D$
- b) $V \models A$
- c) $V \models \emptyset$
- d) $I((A \sqcap B) \sqcap B) = I(A \sqcap B)$
- e) A ⊓ B is an atomic formula



ANSWER: b, d

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Homework

Answer to the following questions

- 1. What are the main characteristics of a representation language?
- 2. What is the difference between extensional and intentional representations?
- 3. What is the difference between an atomic formula and a complex formula?
- 4. What is an interpretation function?
- 5. What is entailment and what are its properties?
- 6. What are the desired properties of logic languages?
- 7. When it is the case that a theory is correct and complete?
- 8. Can you describe the main reasoning problems?