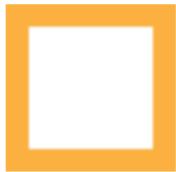


Algorithms and Data Structures

Graphs – Part III – Prim & Dijkstra

Jacob Trier Frederiksen & Anders Kalhauge

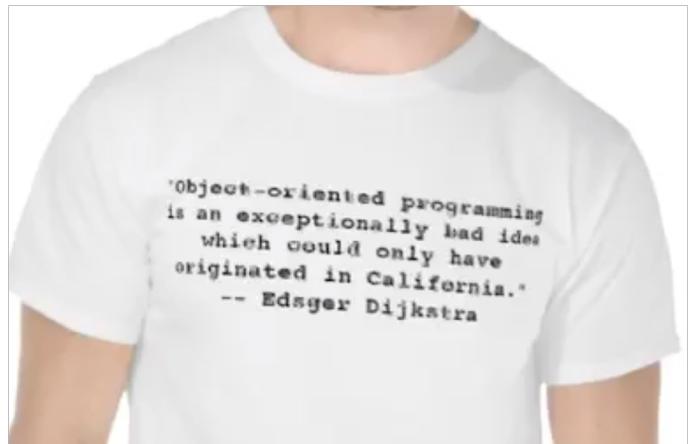


cphbusiness

Spring 2019

Shortest Paths

1. Path Relaxation
2. Dijkstra's Algorithm – today for simple digraphs only.



The Shortest Paths Problem (single-source)

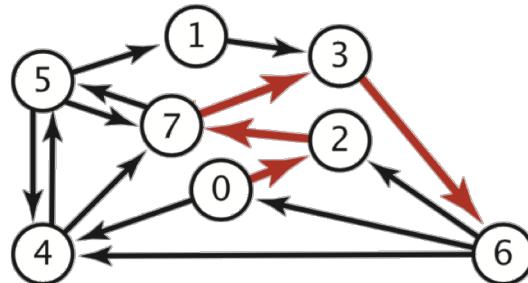
Given a digraph, find the shortest path between a vertex 's' (source) and one other vertex, or several (all) other vertices.

Assume simple digraphs only:

- *Single source*: from one vertex *s* to every other vertex.
- Weights positive
- No negative cycles

edge-weighted digraph

4->5	0.35
5->4	0.35
4->7	0.37
5->7	0.28
7->5	0.28
5->1	0.32
0->4	0.38
0->2	0.26
7->3	0.39
1->3	0.29
2->7	0.34
6->2	0.40
3->6	0.52
6->0	0.58
6->4	0.93

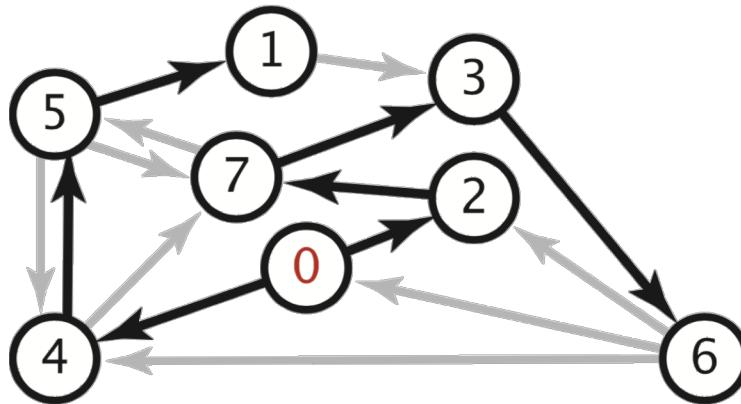


shortest path from 0 to 6

0->2	0.26
2->7	0.34
7->3	0.39
3->6	0.52

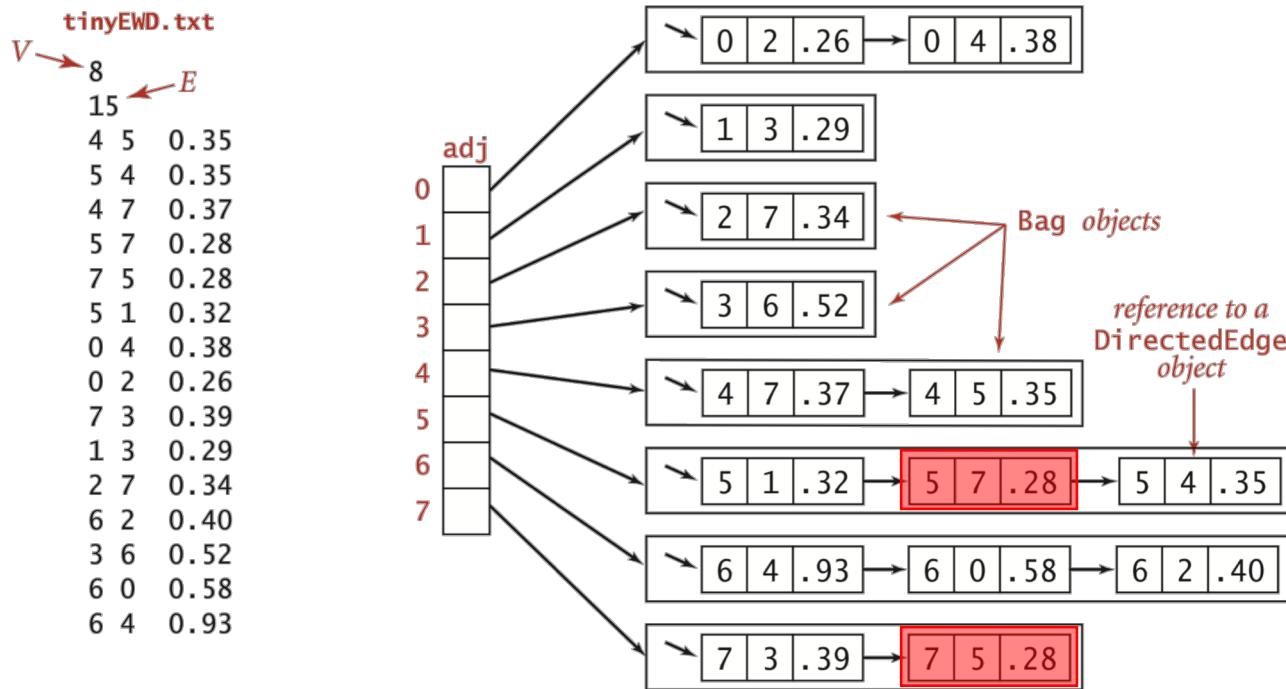
Shortest Path Tree

- Shortest Paths Tree, from computing all single-source paths from 's'.
- Represented by two arrays, `edgeTo[]` and `distTo[]`.



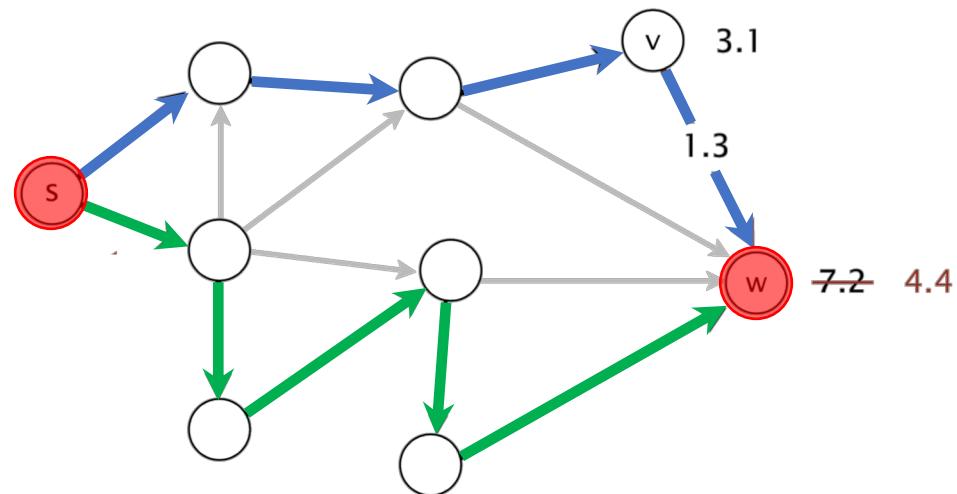
	<code>edgeTo[]</code>	<code>distTo[]</code>
0	null	0
1	5->1	0.32
2	0->2	0.26
3	7->3	0.37
4	0->4	0.38
5	4->5	0.35
6	3->6	0.52
7	2->7	0.34

What Data Structure for Dijkstra *et al.*?



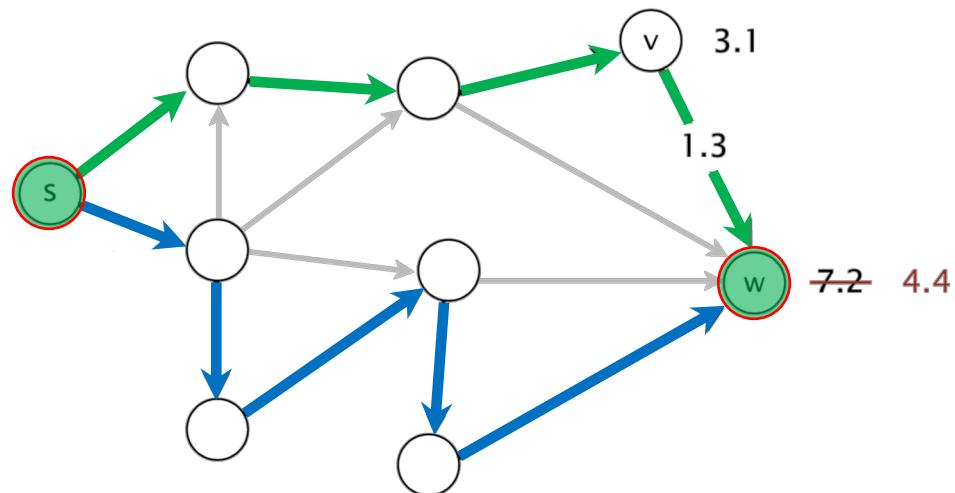
Edge Relaxation

If $e = v \rightarrow w$ gives shorter path to w through v , update both `distTo[w]` and `edgeTo[w]`.



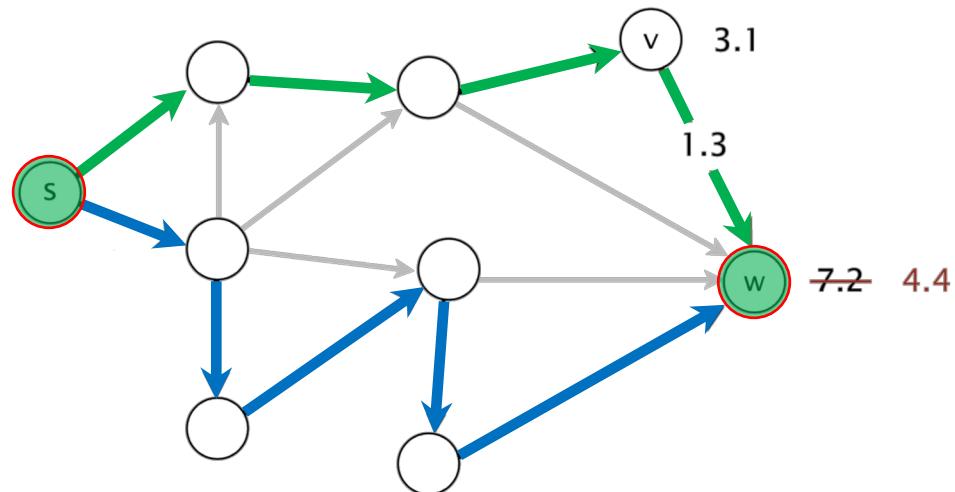
Edge Relaxation

If $e = v \rightarrow w$ gives shorter path to w through v , update both `distTo[w]` and `edgeTo[w]`.



A Generic SPT Algorithm

1. $\text{distTo}[s] = 0$
2. $\text{distTo}[v] = \infty$
3. While not all optimal:
Relax edges
4. Done.



Exercise 4

- $e = v \rightarrow w$ is an edge with weight 17.0.
- At some point, during the generic shortest paths algorithm, `distTo[v] = ∞` and `distTo[w] = 15.0`. What will `distTo[w]` be after calling `relax(e)`?

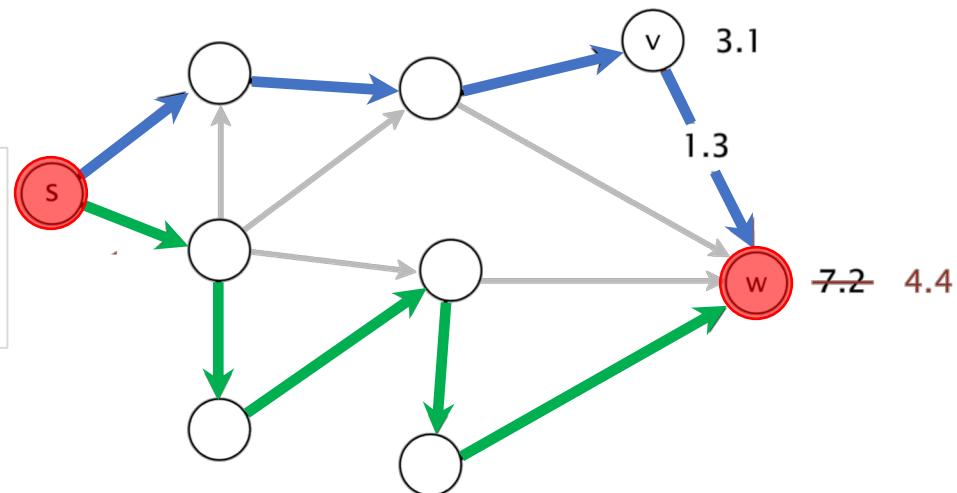
- The program will encounter a Java runtime exception
- 15.0
- 17.0
- ∞

Edge Relaxation HOWTO?

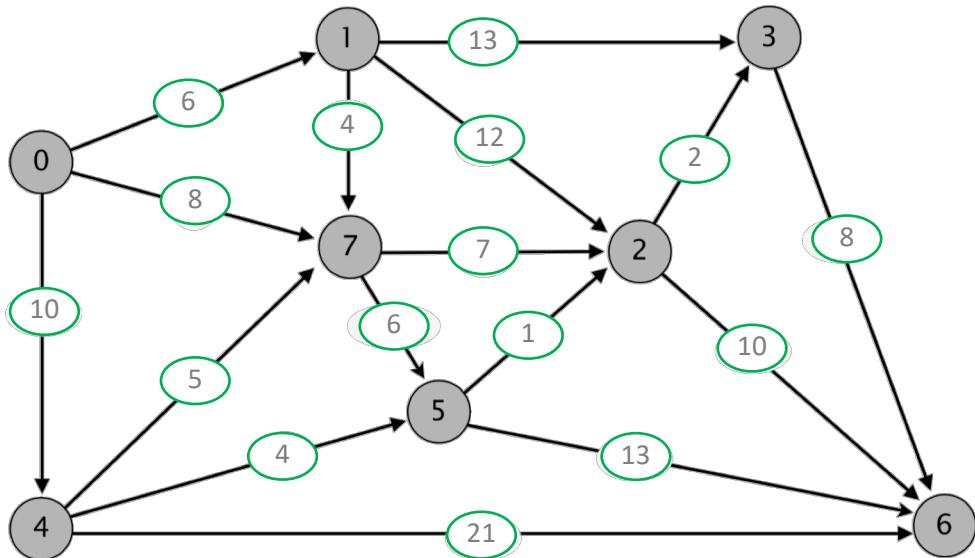
Three central algorithms exist:

1. Dijkstra's Algorithm
2. Topological Sorting
3. Bellman-Ford Algorithm

We might just yet
do something
about number 2....



Exercise 6: Using Dijkstra's Algorithm



v	distTo[]	edgeTo[]
0	0.0	-
1	6.0	0 → 1
2		
3		
4		
5		
6		
7		

Dijkstra's Algorithm

- Keeping the queued vertices/edges on a minPQ (heap).

PQ implementation	insert	delete-min	decrease-key	total
unordered array	1	V	1	V^2
binary heap	$\log V$	$\log V$	$\log V$	$E \log V$
d-way heap	$\log_d V$	$d \log_d V$	$\log_d V$	$E \log_{E/V} V$
Fibonacci heap	1^\dagger	$\log V^\dagger$	1^\dagger	$E + V \log V$

Exercise 7

For the five example graphs

- `tinyEWG.txt` ($V=8, E=16$)
- `mediumEWG.txt` ($V=250, E=$)
- `1000EWG.txt` ($V=1000, E=8433$)
- `10000EWG.txt` ($V=10000, E=61731$)
- `largeEWG.txt` ($V=1000000, E=7586063$)

compare running times, using the `DijkstraSP()` class and the `SP()` class.

Time your computations. You can use fx the Book's `StopWatch()` class for this purpose. Or a system time diff, or you pick...

Exercise 5

Q: What is the BigO running time of Dijkstra's algorithm, using a binary heap?

- V
- E
- V^2
- $E \ Log \ V$

Dijkstra's Algorithm; complexity

- Space: $O(V)$
- Time: $O(E \ Log \ V)$



Acknowledgement

- Some material has been taken from the Princeton course 'Algorithms 4th Edition', by Sedgewick & Wayne.