

# Algorithms & Data Structures

## Graphs – Part I

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cphbusiness

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# What's your favorite Mid-term Quiz?

- What would you like to test yourself on for the mid-term quiz?

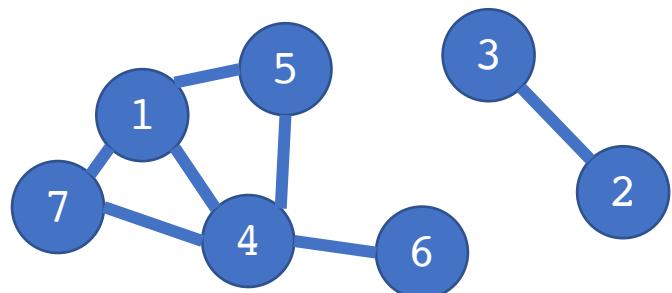
- ...
- ...
- ..
- ....

# Today's Quiz – the white board "peleton".

1. Can you name three generic queue topologies?
2. What is a Complete Binary Tree?
3. What is a max-oriented heap?
4. What are the basic steps in a heap sort (given random data)?
5. What is the complexity of heap sorting (worst case)?

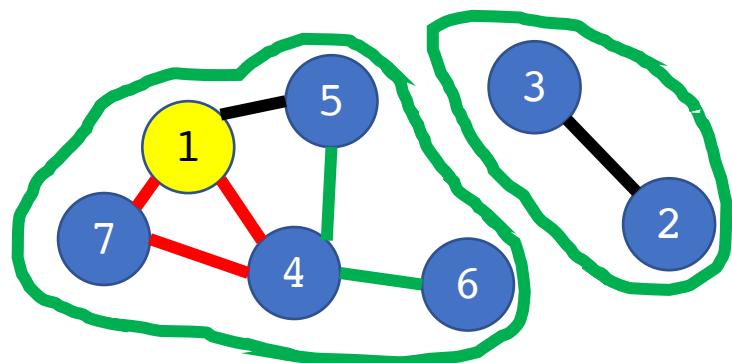
# Graphs – Very simple struct.s.

- Set of vertices connected by edges.
- Edges connected in any possible way
- $G = (V, E)$



# Graphs – Definitions

- **Path**: sequence of vertices connected by edges.
- **Cycle**: path whose first and last vertices are the same.
- **Connectivity**:
  - two vertices are connected if a path exists between them.
  - a graph is a set of connected components.
- **Degree**: a vertex with – say – three connections has degree 3.
- **Acyclic** (not shown): has no cycles.
- **Directed**(not shown): has directionality.



# Copenhagen S-Trains

**This is a graph.**

Q: what is the degree of Frederikssund?

Q: is Hundige connected with Egedal?

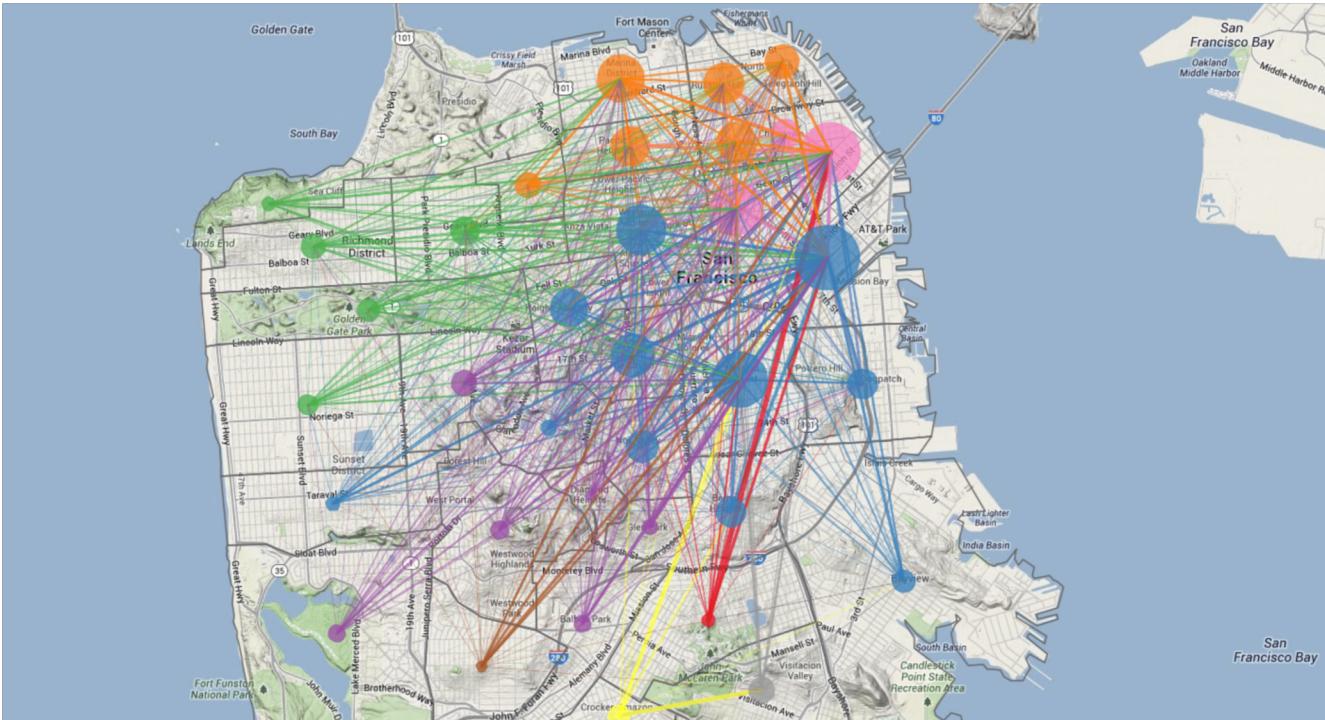
Q: are there any cycles on the S-Train net?



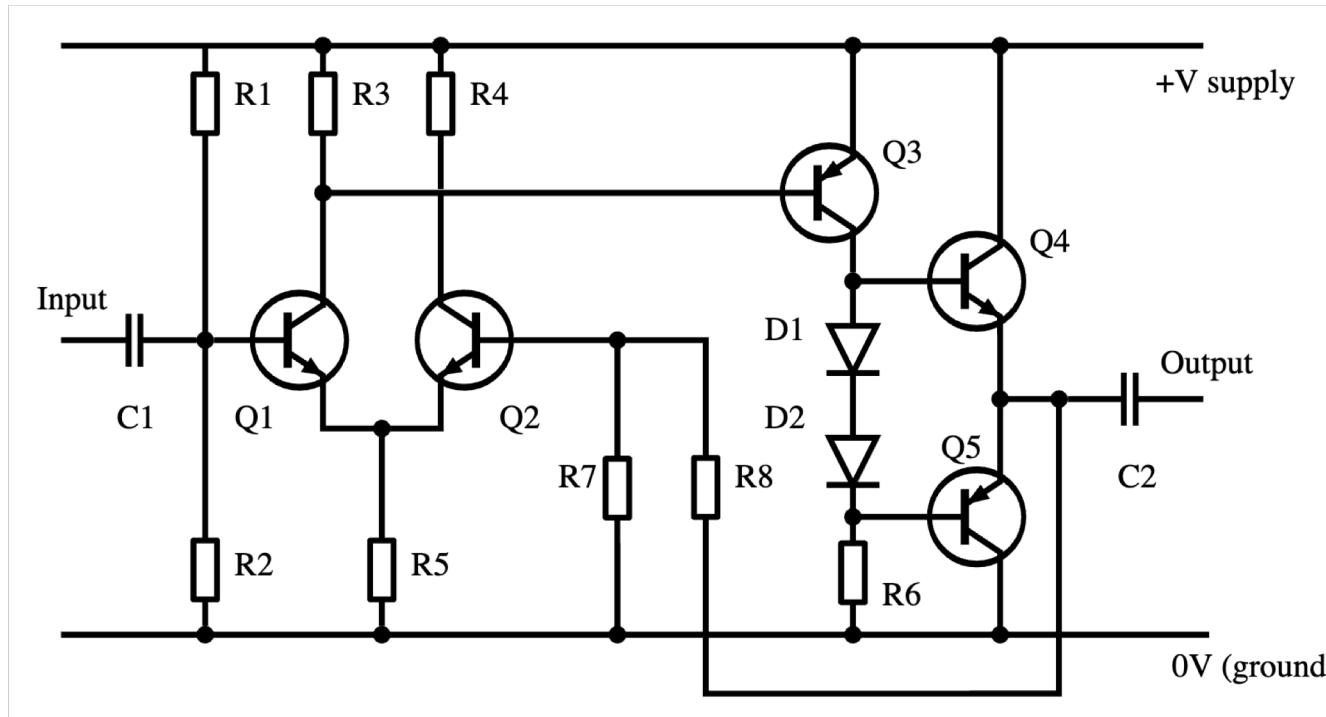
# Applications of graphs – Facebook



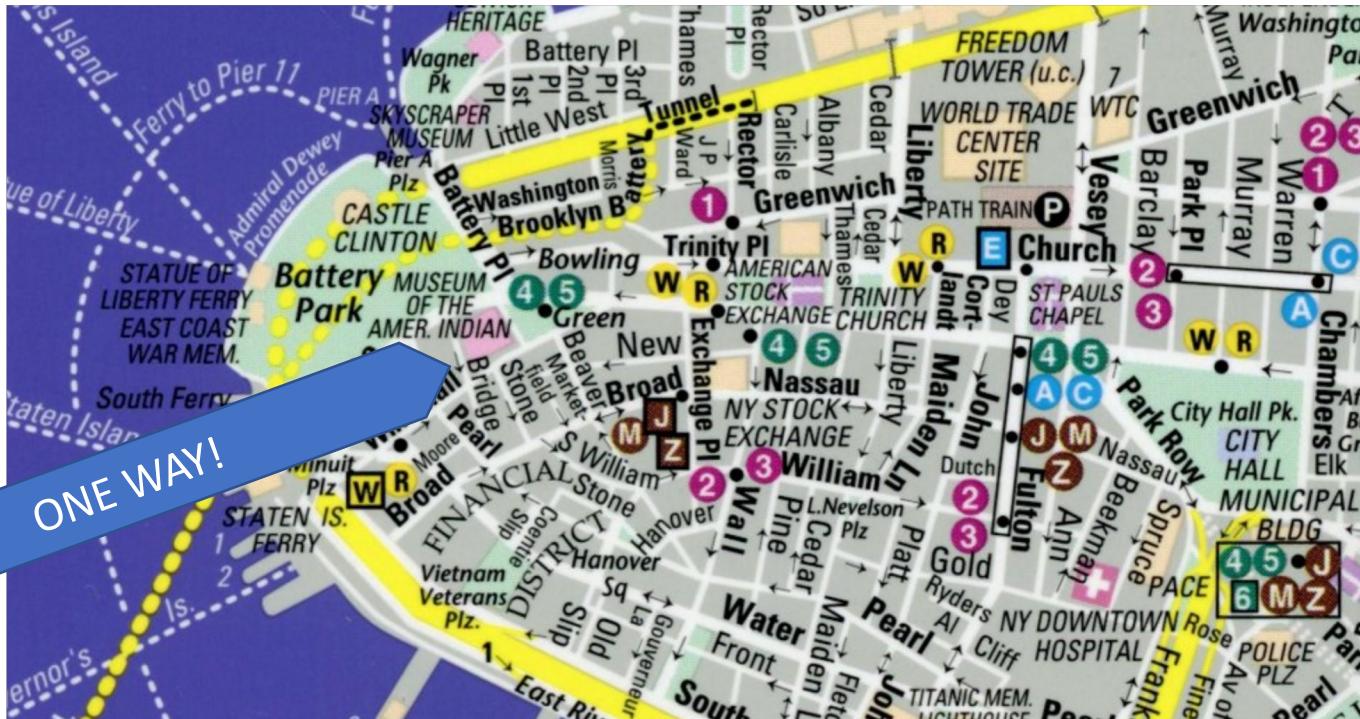
# Uber traffic



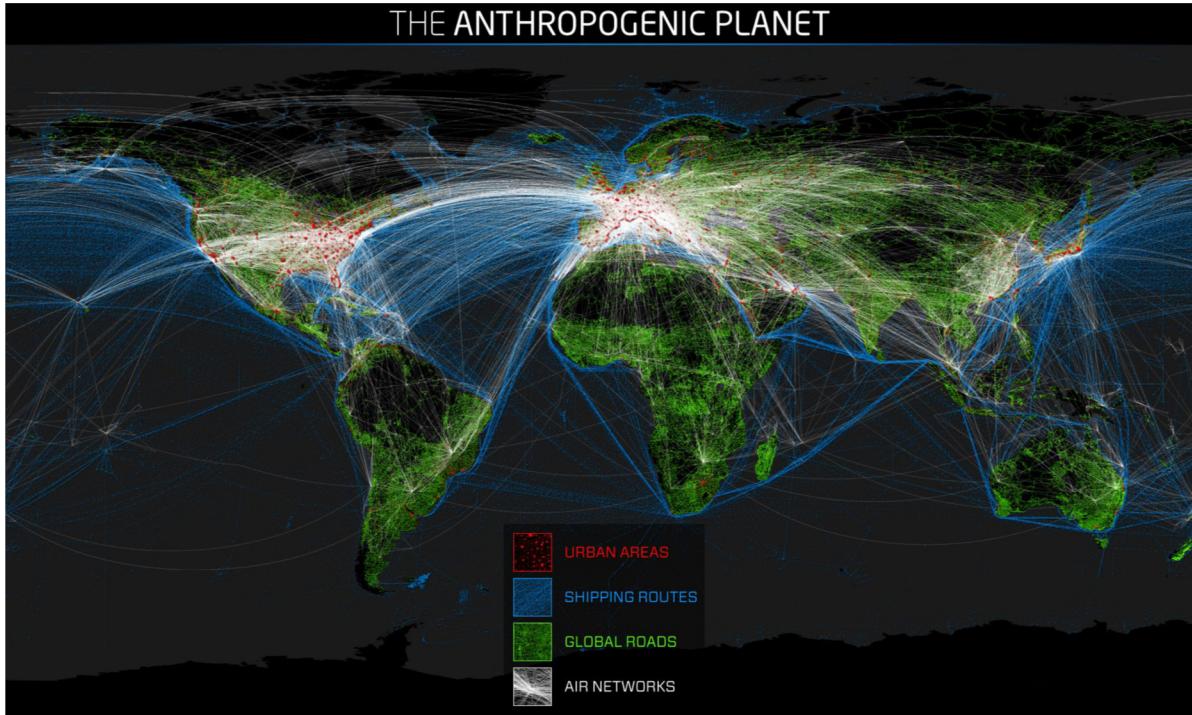
# Electronic circuits



# New York Street Map

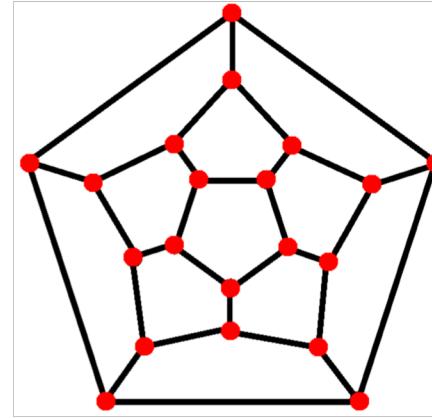
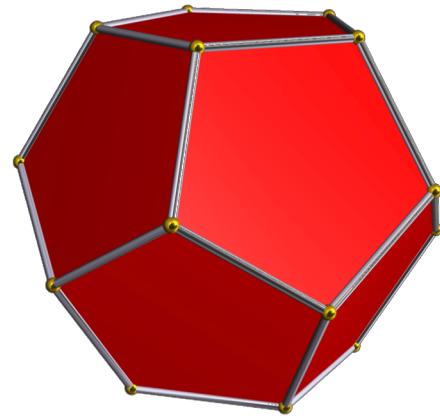


# The Whole Smack



# Applications of Graphs

- **Biology:** vertices could represent regions where certain species exist and the edges represent migration path or movement between the regions, i.e. **neural networks**.
- **Maps:** Various locations are represented as vertices and the roads are represented as edges. Graph theory is used to **minimize path distance** between nodes.
- **E-commerce:** recommender engines use graph theory to **recommend** items of similar type to (previous) user choices.



arXiv.org

<https://xxx.lanl.gov/> check it out, in particular Computer Science / Algs & data struct.s.

*There is still lots of research going on out there!*

# Typical graph theory questions

| problem                  | description   |
|--------------------------|---|
| <b>s-t path</b>          | <i>Is there a path between s and t ?</i>                            |
| <b>shortest s-t path</b> | <i>What is the shortest path between s and t ?</i>                  |
| <b>cycle</b>             | <i>Is there a cycle in the graph ?</i>                              |
| <b>Euler cycle</b>       | <i>Is there a cycle that uses each edge exactly once ?</i>          |
| <b>Hamilton cycle</b>    | <i>Is there a cycle that uses each vertex exactly once ?</i>        |
| <b>connectivity</b>      | <i>Is there a way to connect all of the vertices ?</i>              |
| <b>biconnectivity</b>    | <i>Is there a vertex whose removal disconnects the graph ?</i>      |
| <b>planarity</b>         | <i>Can the graph be drawn in the plane with no crossing edges ?</i> |
| <b>graph isomorphism</b> | <i>Do two adjacency lists represent the same graph ?</i>            |

# Graphs

- Undirected
- Directed
- Bi-partite
- Weighted and unweighted

## White board exercise

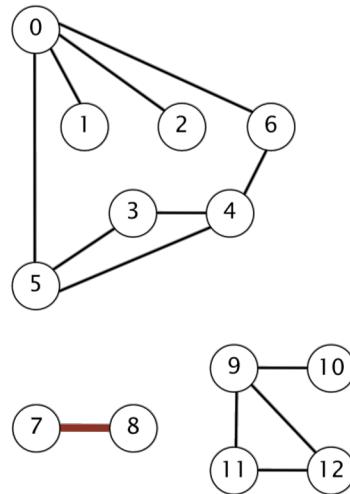
(4 people)

Draw:

1. **Undirected cyclic graph**, with 3 connected components.
2. **Directed graph** with 1 acyclic and 2 cyclic connected components.
3. A **bi-partite graph**, single connected component
4. A **weighted graph** example.

# Graphs – representations, data structures.

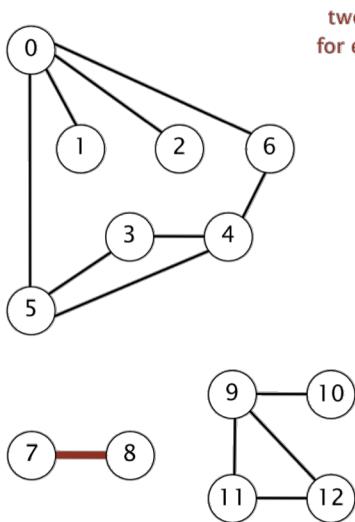
- Array of edges, E? Nope.



|    |    |
|----|----|
| 0  | 1  |
| 0  | 2  |
| 0  | 5  |
| 0  | 6  |
| 3  | 4  |
| 3  | 5  |
| 4  | 5  |
| 4  | 6  |
| 7  | 8  |
| 9  | 10 |
| 9  | 11 |
| 9  | 12 |
| 11 | 12 |

# Graphs – representations, data structures.

- Array of edges,  $E$ ? Nope.
- Adjacency matrix,  $V \times V$ ? Nope.

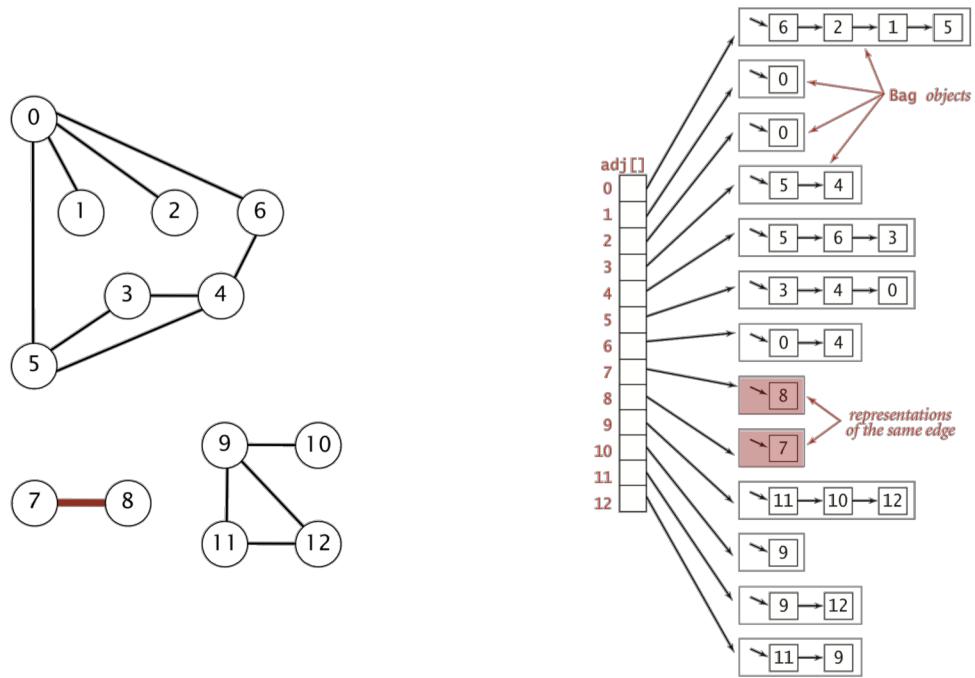
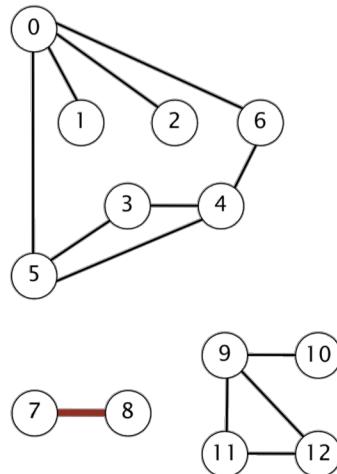


two entries  
for each edge

|    | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
|----|---|---|---|---|---|---|---|---|---|---|----|----|----|
| 0  | 0 | 1 | 1 | 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0  | 0  | 0  |
| 1  | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0  | 0  | 0  |
| 2  | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0  | 0  | 0  |
| 3  | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0  | 0  | 0  |
| 4  | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 0 | 0  | 0  | 0  |
| 5  | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0  | 0  | 0  |
| 6  | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0  | 0  | 0  |
| 7  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0  | 0  | 0  |
| 8  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0  | 0  | 0  |
| 9  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1  | 1  | 1  |
| 10 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1  | 0  | 0  |
| 11 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0  | 0  | 1  |
| 12 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0  | 1  | 0  |

# Graphs – representations, data structures.

- Array of edges,  $E$ ? Nope.
- Adjacency matrix,  $V \times V$ ? Nope.
- Array of adjacency lists? Yep.



# Graphs – representations, data structures.

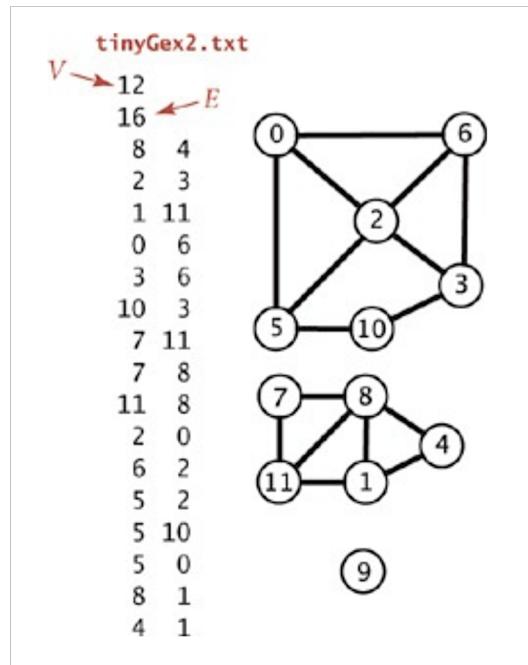
- Array of edges,  $E$ ? Nope.
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| representation   | space   | add edge | edge between v and w? | iterate over vertices adjacent to v? |
|------------------|---------|----------|-----------------------|--------------------------------------|
| list of edges    | $E$     | 1        | $E$                   | $E$                                  |
| adjacency matrix | $V^2$   | 1 *      | 1                     | $V$                                  |
| adjacency lists  | $E + V$ | 1        | $degree(v)$           | $degree(v)$                          |

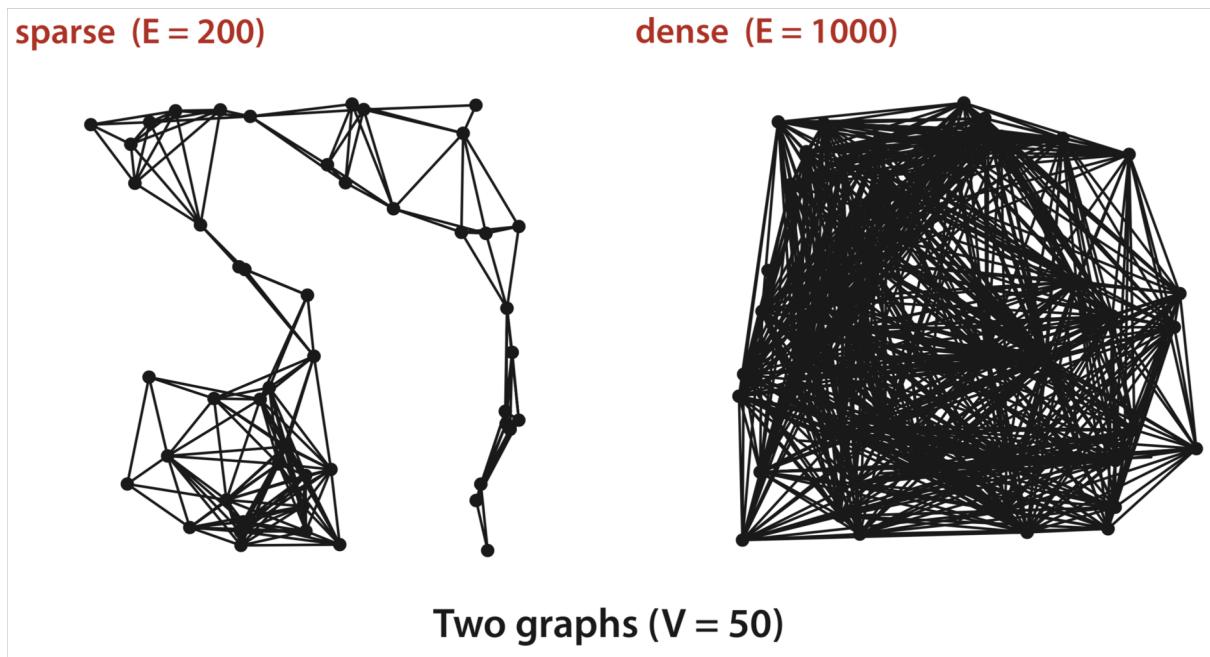


# Exercise 1

- Draw the **array of adjacency lists** diagram for the built by the Graph's input stream constructor for the graph, `tinyGex2.txt` to the right.
- Hint: you'll find all needed bits in Chapter 4.1, subsections '*Representation Alternatives*' and '*Adjacency-lists data structure*'.



# Graphs – dense vs. sparse

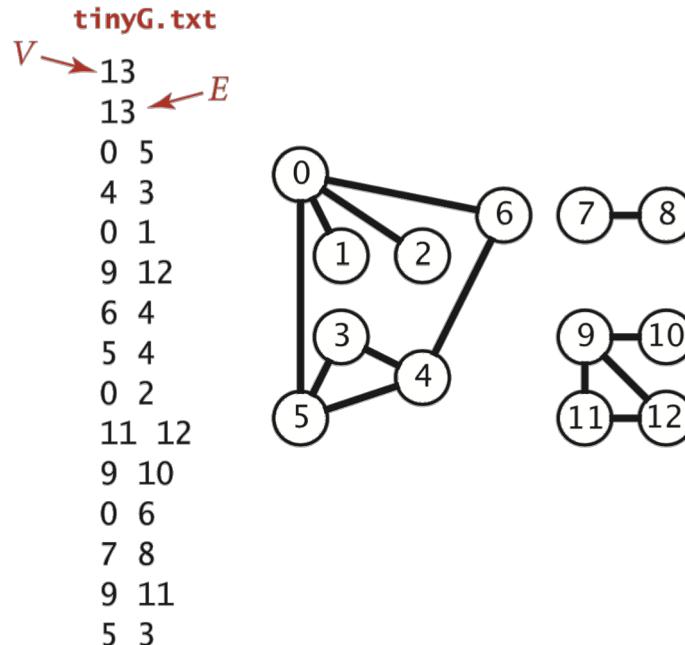


# Exercise 2

- Construct – rather synthesize – two undirected graphs from scratch for two cases
  1. a sparse graph,
  2. a dense graph,

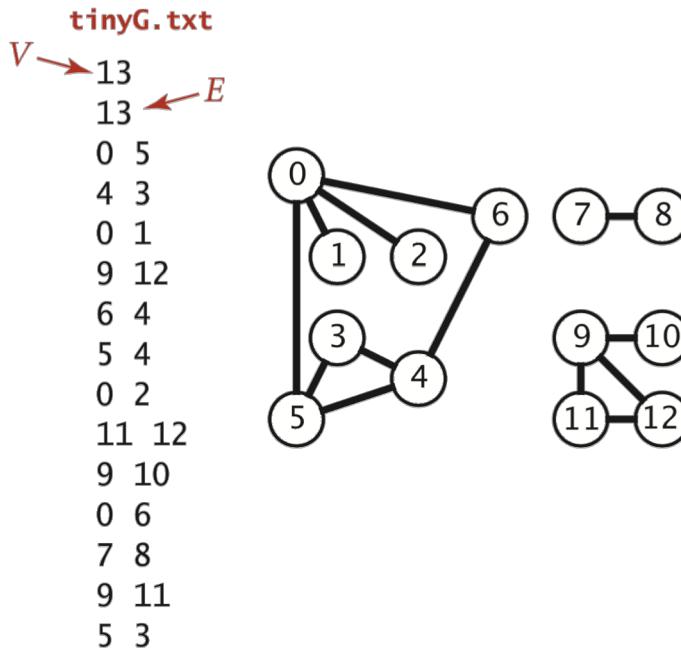
by any means that you choose. Java code would work, but so would Python, 'R' and many other scripting languages. Your choice.

- It should be saved to a text file, similar to the file 'tinyG.txt' from the Algorithms data library.
- What to consider when making a synthetic graph?
- How to parametrize?



# Exercise 3

- Produce the adjacency matrices for your two graphs.
- Save the matrix data (binary) to file and make an image representation of the two matrices respectively.
- For the two cases, do you have full connectivity, in a single connected component, or do you have more than one?
  - Can we even check that from an adjacency matrix? If so, howto?



# Depth-first Search, DFS

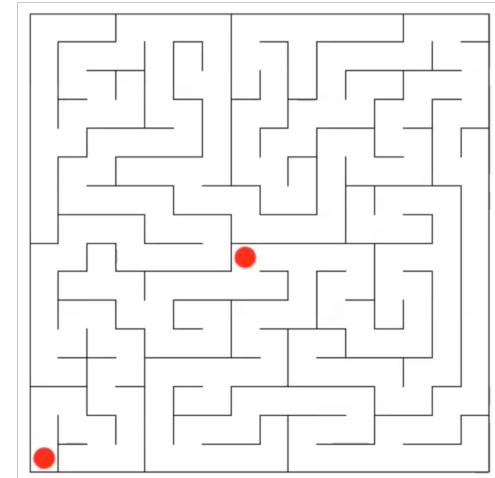
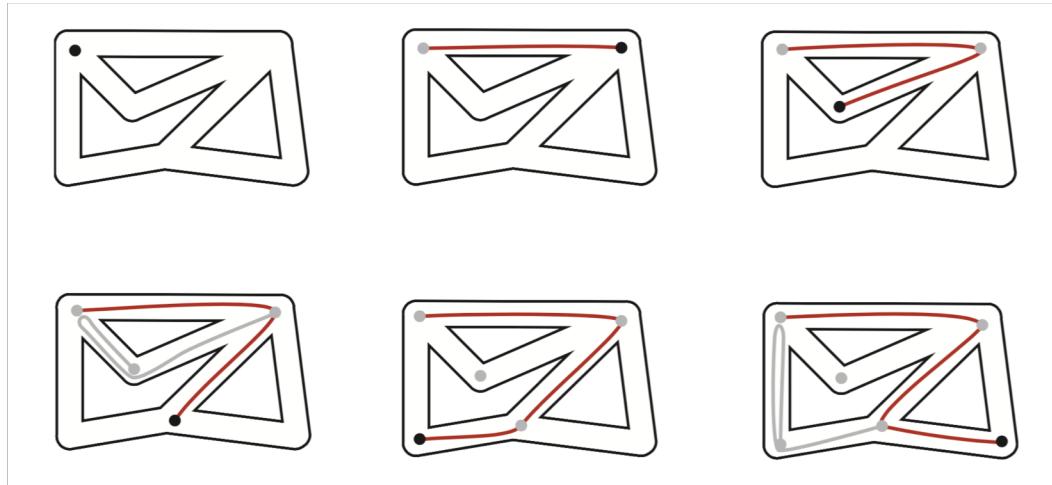
## **Applications:**

- Finding connected components.
- Finding strongly connected components.
- Topological sorting.
- Finding the bridges of a graph.
- Planarity testing.
- Solve mazes.
- Maze generation, use a randomized depth-first search.
- Finding biconnectivity in graphs.

# Tremaux Maze – depth-first search.

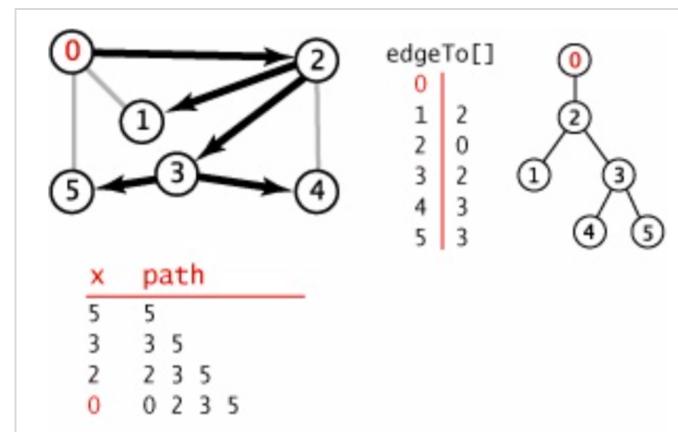
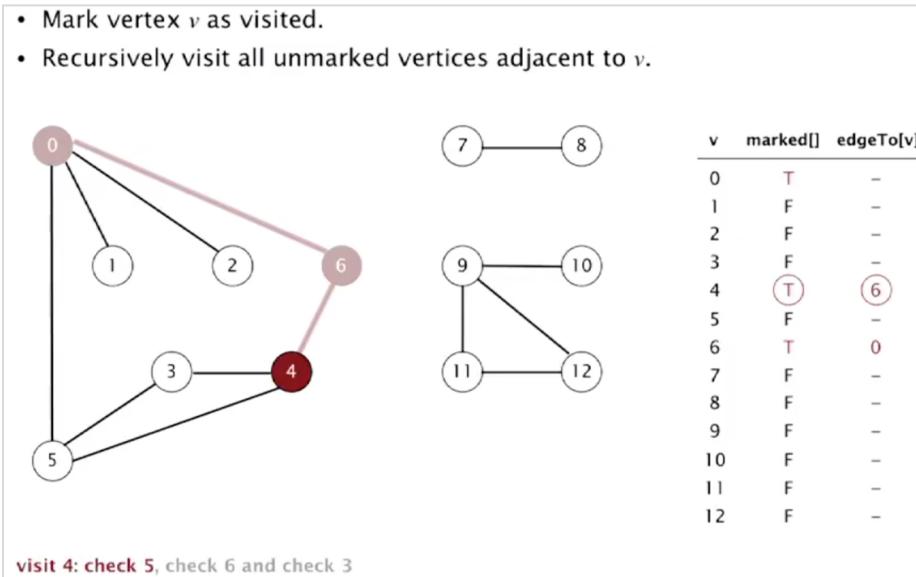
- Or the Theseus ... or just common sense.

How to find your way?



# Depth-first Search, DFS

- **pathTo[ v ]** : DFS uses a **stack** to store the **edgeTo[ v ]** values.

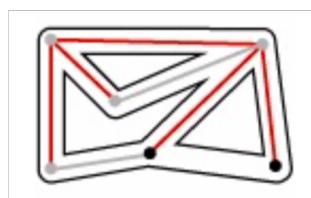
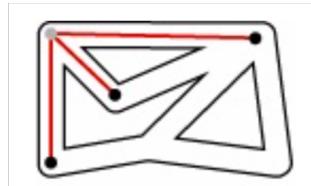


# Breadth-first Search, BFS

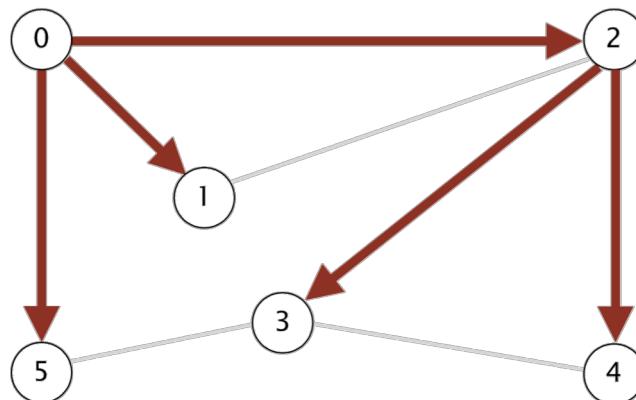
## Applications include:

- Finding the shortest path between two nodes  $u$  and  $v$ , with path length measured by number of edges (an advantage over DFS)
- Testing bipartiteness of a graph.
- (Reverse) Cuthill–McKee mesh numbering
- Ford–Fulkerson method for computing the maximum flow in a flow network
- Serialization/Deserialization of a binary tree vs serialization in sorted order, allows the tree to be re-constructed in an efficient manner.
- Copying garbage collection, Cheney's algorithm

# Breadth-first Search, BFS



Uses a **queue** to keep track of vertices visited but whose adjacent vertices not yet visited.



| v | edgeTo[] | distTo[] |
|---|----------|----------|
| 0 | -        | 0        |
| 1 | 0        | 1        |
| 2 | 0        | 1        |
| 3 | 2        | 2        |
| 4 | 2        | 2        |
| 5 | 0        | 1        |

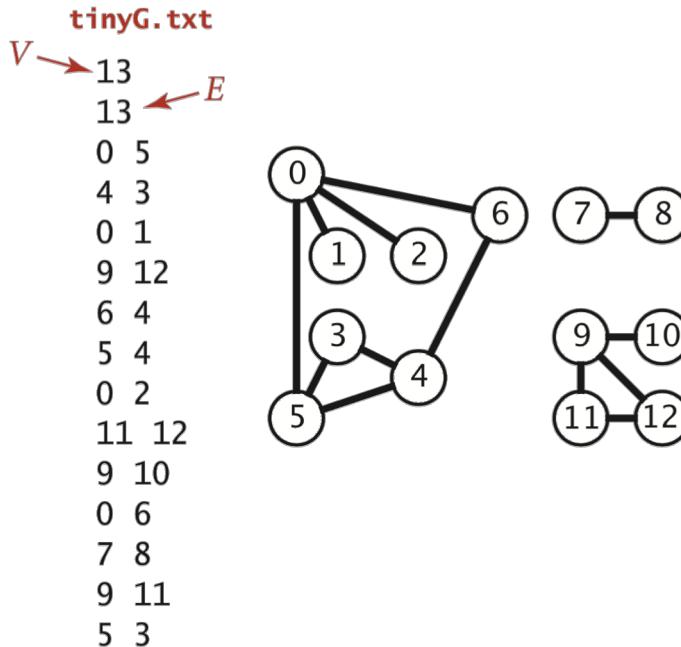
# Exercise 4

Out of time, use `mediumG.txt` from "Algorithms 4th Edition".  
<https://algs4.cs.princeton.edu/41graph/mediumG.txt>

1. Check whether your graph is connected.  
Hint: `TestSearch`
2. Find DFS paths for your synthetic graphs  
Hint: `DepthFirstPaths`
3. Find BFS paths for your synthetic graphs  
Hint: ...
4. Are your results for DFS and BFS different?
5. Find connected components for your graphs  
Hint: ...
6. Check if your graphs have cycles
7. Check whether your graphs are bi-partite.

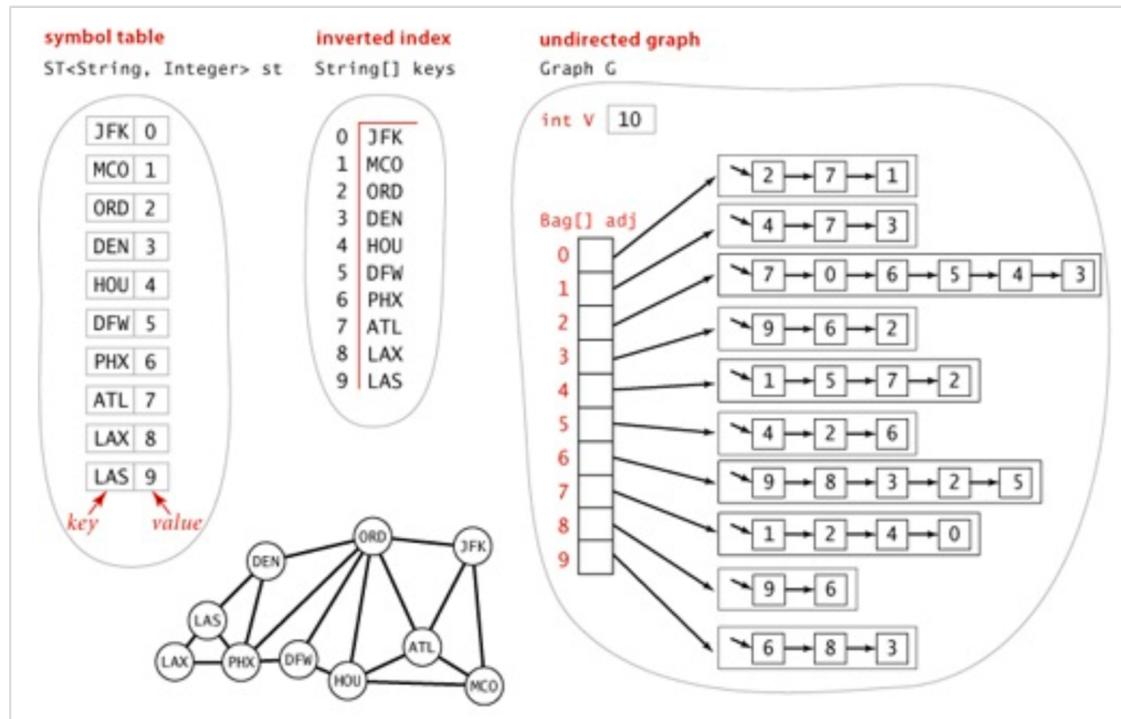
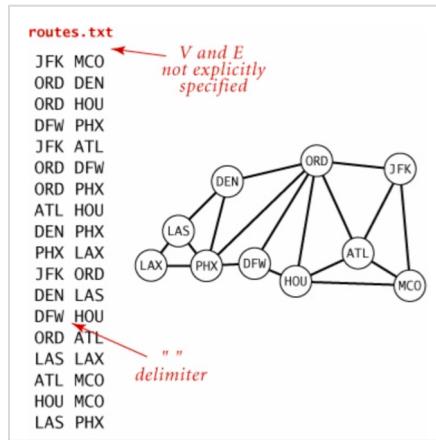
# Exercise 5

- For both of your synthesized graphs (sparse and dense) – which are in your previously generated text files (right?), find
  - by DFS whether 10 randomly chosen points are connected to an other chosen single point. Record running time for the 10 cases.
  - the shortest path between the same one chosen point and the randomly selected 10 other points, using BFS. Record running time for the 10 cases.



# A note on Symbol Graphs and symbol tables

- Not so different....



# Pop 1 (NB: directed graphs, week 12):

- How many different digraphs are there on  $V$  vertices?  
(allow for self-loops but do not allow parallel edges)

  $V$

  $2^V$

  $2^{(V^2)}$

  $2^{(2^V)}$

# Pop 2 (NB: directed graphs, week 12)

During a DFS in a digraph  $G$ ,  $\text{dfs}(v)$  is called after  $\text{dfs}(w)$  is called but before  $\text{dfs}(w)$  returns. What *must* be true about the graph  $G$ ?

- There exists a directed path  $v \rightarrow w$ .
- There exists a directed path  $w \rightarrow v$ .
- There does not exist a directed path  $v \rightarrow w$ .
- There exists a directed cycle containing  $v \rightarrow w$ .

# Pop 3 (NB: directed graphs, week 12)

A DFS is done on a DAG which contains edge  $v \rightarrow w$ .

Which of the following is impossible at the time of calling `dfs(v)`?

- `dfs(w)` has not yet been called.
- `dfs(w)` will be the next recursive call after `dfs(v)`.
- `dfs(w)` has already been called but not yet returned.
- `dfs(w)` has already been called and returned.



# Acknowledgement

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- Some material has been taken from the Princeton course 'Algorithms 4th Edition', by Sedgewick & Wayne.