

Algorithms and Data Structures Searching Abstract Data Structures, Symbol Tables

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Outline



Arrays

Symbol Tables

List-based Array-based



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Symbol Tables List-based Array-based



Arrays are by nature of fixed length. How can we make them expandable and still have direct memory access?

When adding a new element to a full array we could:

- 1. Copy the array to an array one bigger
- 2. Copy the array to an array m elements bigger
- 3. Copy the array to an array of double size



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- 1.
- 2.
- 3.



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- 2. O(n) all elements are copied each $m^{\rm th}$ time $O(\frac{n}{m}) = O(n)$
- 3. O(1) how can that be?



Money in the bank

Balance: 4 we hope that is enough to pay for future expansions

array is full

7 9 13 2



Money in the bank

Balance: 4 - 0 = 4 (creating new array considered free here)

array is full

7

9

13

create new array











Money in the bank

Balance:
$$4 - 4 = 0$$
 (using 1 per copy)

array is full 7 9 13 2
4 copies
create new array 7 9 13 2



Money in the bank

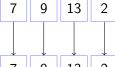
Balance:
$$0+3-1=2$$
 (charging 3 for an insert, using 1)



Money in the bank

Balance:
$$2+3-1=4$$
 (charging 3 for an insert, using 1)

array is full



4 copies

create new array



 $\mathsf{add}\ \mathsf{element}\ 4$

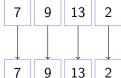


Money in the bank

Balance:
$$4+3-1=6$$
 (charging 3 for an insert, using 1)

array is full

4 copies



create new array

add element 4

add element 5

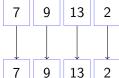


16

Money in the bank

Balance: 6+3-1=8 (charging 3 for an insert, using 1)

array is full



4 copies

create new array

add element 4

add element 5

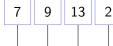
add element 8



Money in the bank

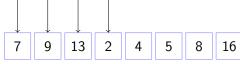
Balance: 8 enough to pay for 8 copies





4 copies

create new array



add element 4

add element 5

add element 8

$$O(3) = O(1)$$



Constant payload

Payload: 0

array is full

7 9 13



Constant payload

Payload: 0 (creating new array considered free here)

array is full 7 9 13 2

create new array



Constant payload

Payload:
$$1+1=2$$
 (1 for copying and 1 for inserting)

array is full 7 9 13 2
create new array 7 4

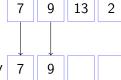
 $\mathsf{copy}\ 7\ \mathsf{insert}\ 4$



Constant payload

Payload:
$$1 + 1 = 2$$
 (1 for copying and 1 for inserting)

array is full



create new array

 $\operatorname{copy} 7 \operatorname{insert} 4$

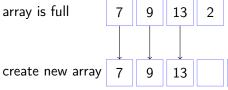
 $\mathsf{copy}\ 9\ \mathsf{insert}\ 5$



Constant payload

Payload: 1 + 1 = 2 (1 for copying and 1 for inserting)

array is full



copy 7 insert 4

copy 9 insert 5

copy 13 insert 8

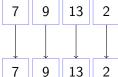


16

Constant payload

Payload: 1 + 1 = 2 (1 for copying and 1 for inserting)

array is full



create new array

 $\operatorname{copy} 7 \operatorname{insert} 4$

 $\mathsf{copy}\ 9 \ \mathsf{insert}\ 5$

copy 13 insert 8

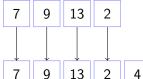
copy 2 insert 16



Constant payload

Payload: 8 in total for 4 insertions

array is full



create new array

 $\mathsf{copy}\ 7\ \mathsf{insert}\ 4$

 $\mathsf{copy}\ 9\ \mathsf{insert}\ 5$

copy 13 insert 8

 $\operatorname{copy}\ 2 \ \operatorname{insert}\ 16$

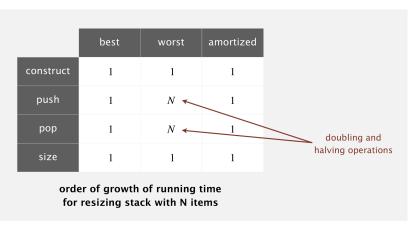
$$O(2) = O(1)$$

16

Resizing Arrays (according to Alg. 4th ed.)



Resizing Arrays - Cost



Q: discuss the differences from what we just had in slides?

Exercise 1 - Complexity



- □ What would the complexity (big-O) be if we:
 - Triple the array size instead of doubling it?
 - □ Only made the new array 50% bigger?
- □ Bearing in mind that most modern memory is paged¹, consider why doubling the array size is not such a bad idea?

 $^{^{1}}$ typically in 2^{n} sized pages

Exercise 2 - Implementation



1. Create a Java class FlexibleArray that uses the "Constant payload" algorithm.

```
public class FlexibleArray<T> {
    ...
    public T get(int index) { ... }
    public void set(int index, T element) { ... }
    public void add(T element) { ... }
    public int size() { ...}
}
```

Note that to create a new array of type T you must:

```
private T[] arrayOfT = (T[])new Object[1000];
```

- 2. Measure the time it takes to add 10.000, 100.000, and 1.000.000 elements.
- 3. Measure Javas build-in ArrayList with the same data.



Arrays

Symbol Tables List-based Array-based

Symbol tables



Examples of Symbol Tables

application	purpose of search	key	value				
dictionary	find definition	word	definition				
book index	find relevant pages	term	list of page numbers				
file share	find song to download	name of song	computer ID				
financial account	process transactions	account number	transaction details				
web search	find relevant web pages	keyword	list of page names				
compiler	find properties of variables	variable name	type and value				
routing table	route Internet packets	destination	best route				
DNS	find IP address	domain name	IP address				
reverse DNS	find domain name	IP address	domain name				
genomics	find markers	DNA string	known positions				
file system	find file on disk	filename	location on disk				

Symbol tables



Map in Java

- ☐ Lookup a value based on a key
- □ No null keys
- No null values
- □ No duplicate keys

```
public interface SymbolTable < K, V > {
  void put(K key, V value);
  V get(K key);
  int size();
  Iterable < K > keys();
  default void delete(K key) { put(key, null); }
  default boolean contains(K key) {
    return get(key) != null;
    }
  default boolean isEmpty() { return size() == 0; }
}
```

Ordered Symbol Tables

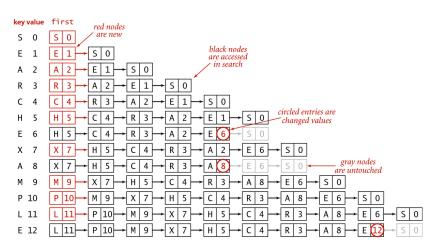


OrderedMap in Java

```
public interface
    OrderedSymbolTable <K extends Comparable <K>, V>
    extends SymbolTable < K, V > {
  K min():
  K max();
  K floor(K key);
  K ceiling(K key);
  int rank(K key);
  K select(int rank);
  void deleteMin():
  void deleteMax();
  int size(K low, K high);
  Iterable <K> keys(K low, K high);
  }
```

List-based Symbol Tables



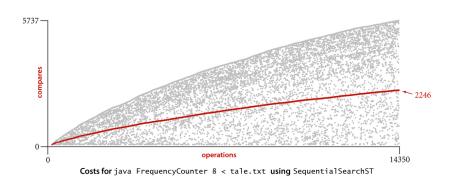


Trace of linked-list ST implementation for standard indexing client

List-based Symbol Tables



Costs



Q: go to the book (Ch.3.1, subsect. 'Performance Client') and figure out what the input argument '8' does here?

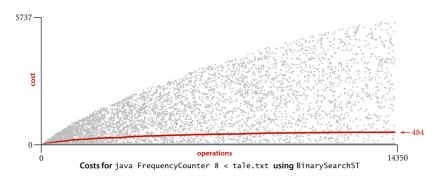


						بميا												1 - 53	ı			
						key	S[]										٧a	ls[]				
key	value	0	1	2	3	4	5	6	7	8	9	N	0	1	2	3	4	5	6	7	8	9
S	0	S										1	0									
Ε	1	Ε	S			0	ntrie	c in 1	rod			2	1	0					tries ved to			
Α	2	Α	Ε	S			vere i					3	2	1	0		/	, 1110	<i>чеи</i> и) ine	rign	
R	3	Α	Ε	R	S							4	2	1	3	0						
C	4	Α	C	Ε	R	S			en	tries	in gr	av 5	2	4	1	3	0					
Н	5	Α	\subset	Е	Н	R	S				t mo		2	4	1	5	3	0		:led e ange		s are
Ε	6	Α	\subset	Ε	Н	R	S					6	2	4	6	5	3	0	Cri	unge	u ru	ines
Χ	7	Α	\subset	Е	Н	R	S	Χ				7	2	4	6	5	3	0	7			
Α	8	Α	\subset	Е	Н	R	S	X				7	8	4	6	5	3	0	7			
М	9	Α	\subset	Е	Н	М	R	S	Χ			8	8	4	6	5	9	3	0	7		
Р	10	Α	\subset	Е	Н	M	Р	R	S	Χ		9	8	4	6	5	9	10	3	0	7	
L	11	Α	\subset	Ε	Н	L	М	Р	R	S	Χ	10	8	4	6	5	11	9	10	3	0	7
Ε	12	Α	\subset	Е	Н	L	$[\mathbb{M}$	Р	R	S	X	10	8	4	(12)	5	11	9	10	3	0	7
		Α	C	Ε	Н	L	Μ	Р	R	S	Χ		8	4	12	5	11	9	10	3	0	7

Trace of ordered-array ST implementation for standard indexing client



Costs



Q: What happened here? Look at that decrease in cost... Discuss in plenum.



Costs Comparison

Unordered sequential search vs. Ordered array binary search.

algorithm		ase cost inserts)	average-c (after N rand		efficiently support ordered		
(data structure)	search	insert	search hit	insert	operations?		
sequential search (unordered linked list)	N	N	N/2	N	no		
binary search (ordered array)	$\lg N$	2N	$\lg N$	N	yes		

Cost summary for basic symbol-table implementations

Q: What about insert (or put())? What is *its* complexity? And can we achieve better performance? Discuss in plenum, 'tales.txt' and '.



Pros & Cons in Searching

underlying data structure	implementation	pros	cons
linked list (sequential search)	SequentialSearchST	best for tiny STs	slow for large STs
ordered array (binary search)	BinarySearchST	optimal search and space, order-based ops	slow insert
	BST	easy to implement, order-based ops	no guarantees space for links
	RedBlackBST	optimal search and insert, order-based ops	space for links
	SeparateChainingHashST LinearProbingHashST	fast search/insert for common types of data	need hash for each type no order-based ops space for links/empty
	Pros and cons of symbo	l-table implementatio	ns

Q: how can we get better e.g. put() performance? Need more than linked lists and binary searches=> Datastructures.

Exercise 3 - Sequential Search



NB: These last exercises are optional, but will definitely help you to understand the linked lists, search complexity, and timing issues better.

- □ Add to the Algorithm 3.1 in the Book (seq. search in unordered list), the methods mentioned in the text, namely
 - □ size(),
 - □ keys(),

so essentially exercise 3.1.5, yet, without the eager delete().

- Exercise 3.1.6 of the Book.
- □ What's frequent most word of 10+ letters in 'Tale of two cities' (ex. 3.1.8 of the book)?