

Geometry driven collapses for simplifying Čech complexes

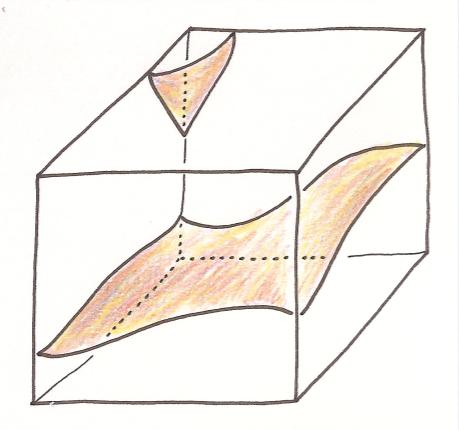
Dominique Attali (*) and André Lieutier ()**

(*) Gipsa-lab

(**) Dassault Systèmes

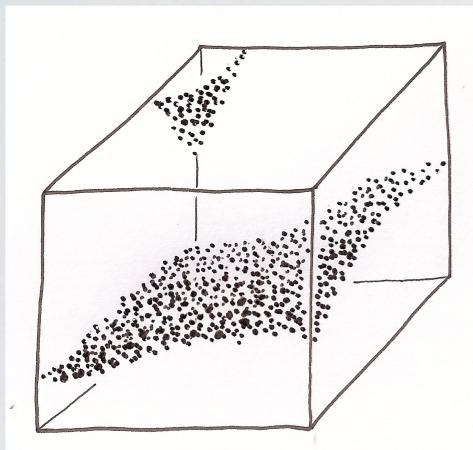
*Workshop on Applied and Computational Algebraic Topology
July 15-19, 2013 - Bremen, Germany*

Shape



Approximation

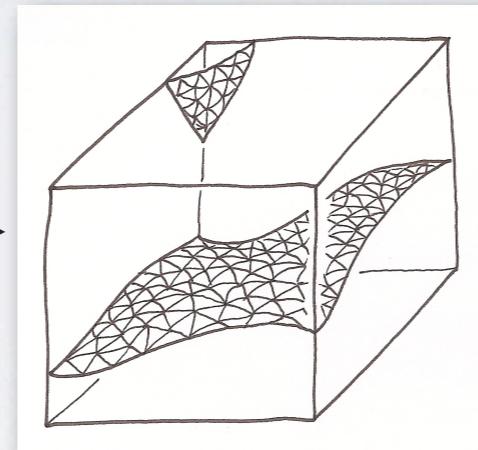
n points



Input

RECONSTRUCTION

Simplicial complex

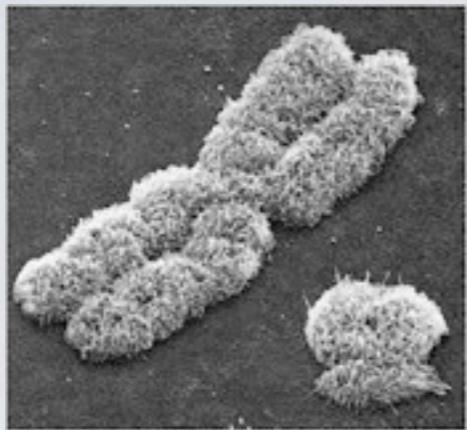


Output

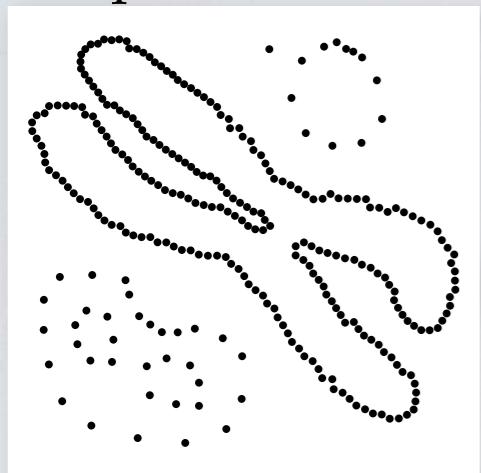
PROCESSING

Betti numbers
Volume
Medial axis
Signatures
...

in 2D



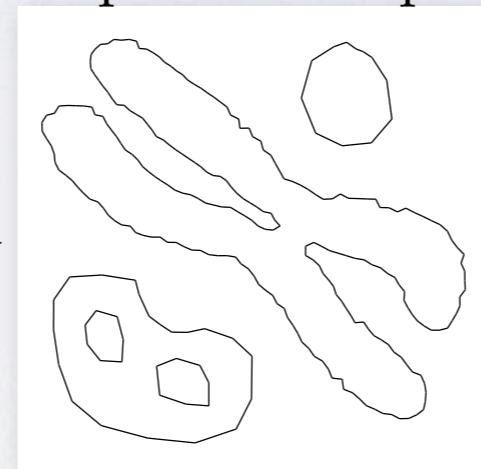
n points in \mathbb{R}^2



Input

RECONSTRUCTION

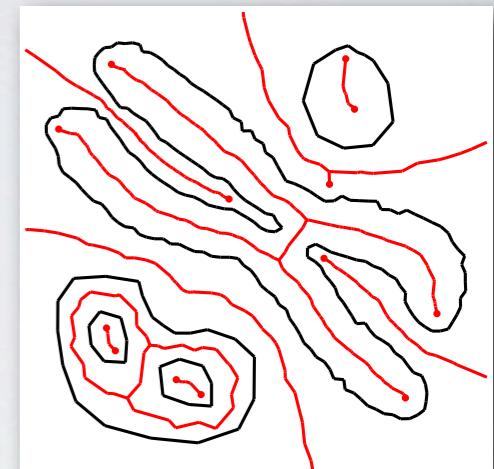
Simplicial complex



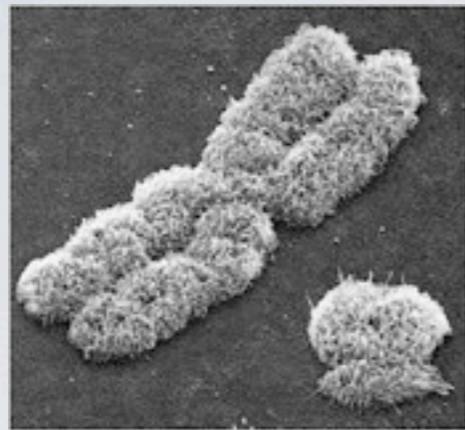
Output

PROCESSING

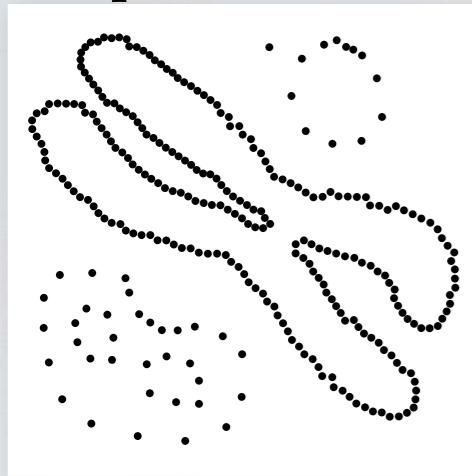
Medial axis



in 2D

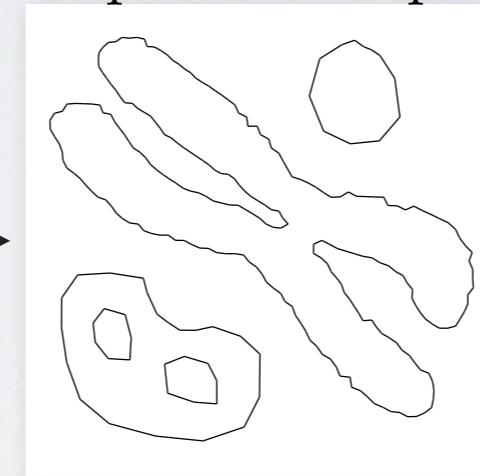


n points in \mathbb{R}^2



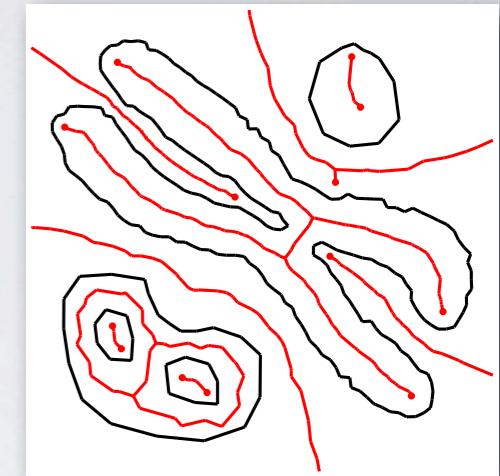
RECONSTRUCTION

Simplicial complex



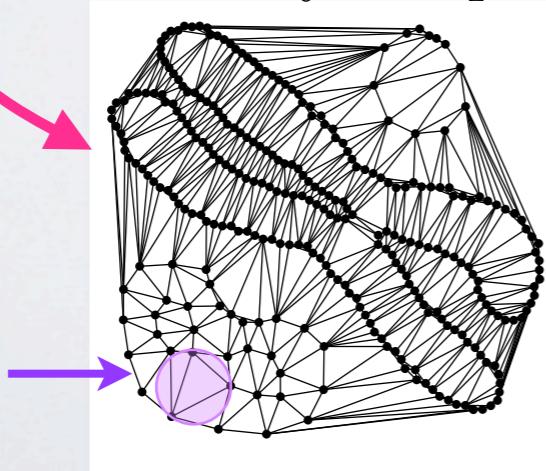
PROCESSING

Medial axis



BUILDING

Delaunay complex



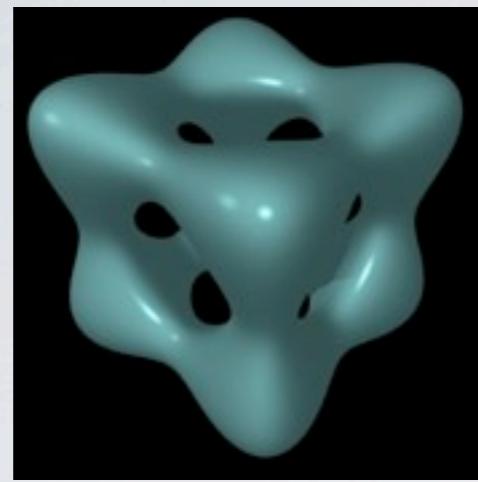
Empty circle property

* In \mathbb{R}^2 , has size $O(n)$

(1995 – 2005) HEURISTICS
(Crust, Power crust, Co-cone, Wrap, ...)

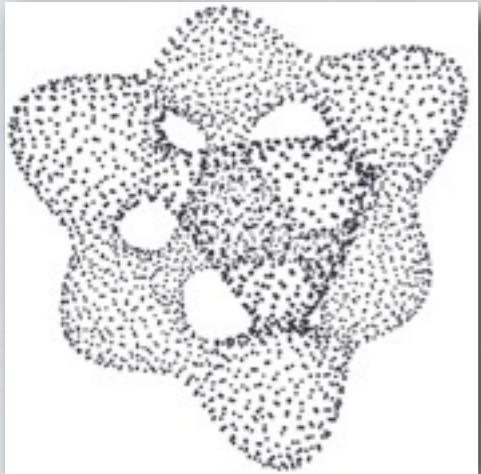


Delaunay of 10M points in 2D ≈ 10 s



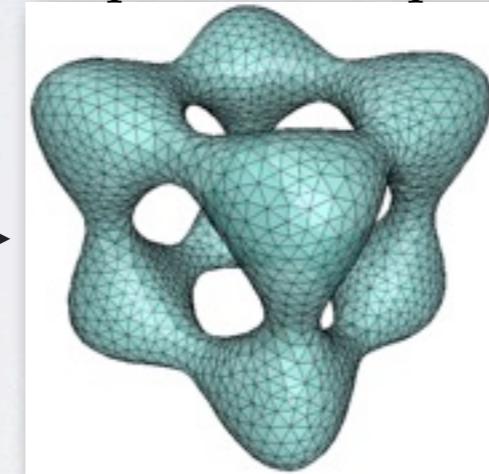
in 3D

n points in \mathbb{R}^3



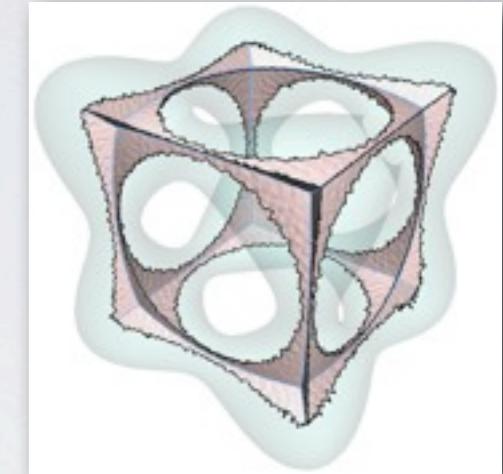
RECONSTRUCTION

Simplicial complex



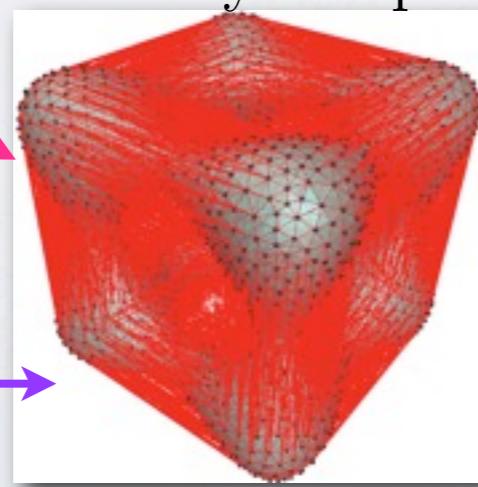
PROCESSING

Medial axis



BUILDING

Delaunay complex



(1995 – 2005)

(Crust, Power crust, Co-cone, Wrap, ...)

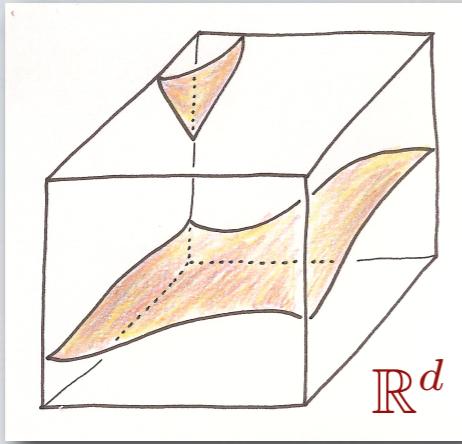
- * In \mathbb{R}^3 , has size $O(n^2)$
- * In practice, has size $O(n)$

Empty sphere property



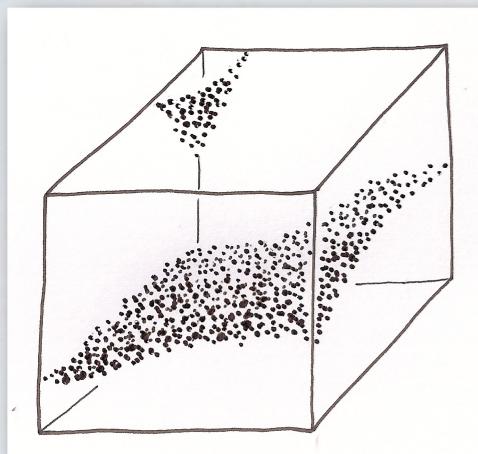
Delaunay of 10M points in 3D ≈ 80 s

Shape

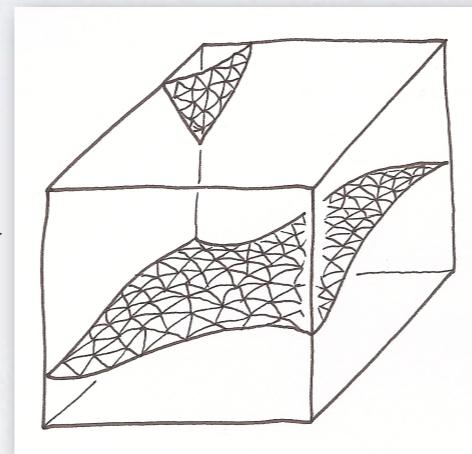


in dD

n points in \mathbb{R}^d



Simplicial complex



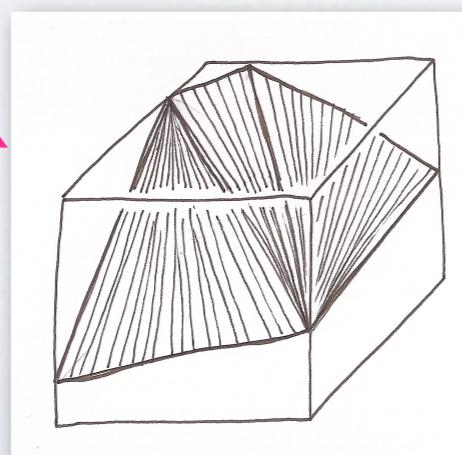
RECONSTRUCTION

PROCESSING

Betti numbers
Volume
Medial axis
Signatures
...

BUILDING

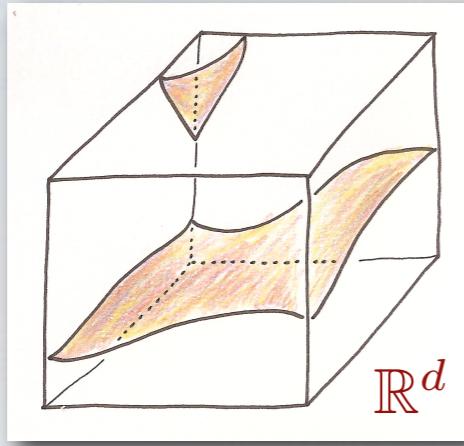
Delaunay complex



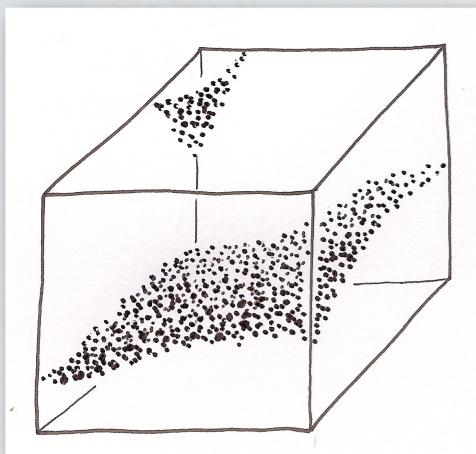
curse of dimensionality

- * In \mathbb{R}^d , has size $O(n^{\lceil d/2 \rceil})$
- * The bound is tight (and achieved for points that sample curves).

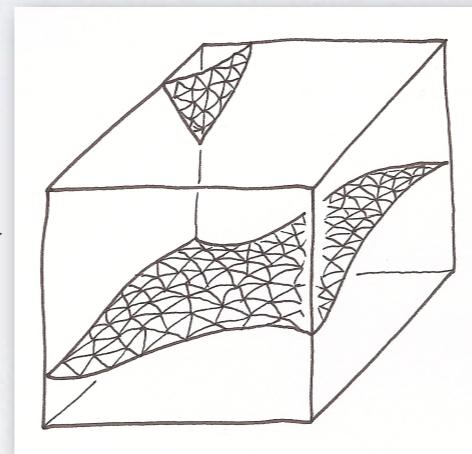
Shape



n points in \mathbb{R}^d



Simplicial complex



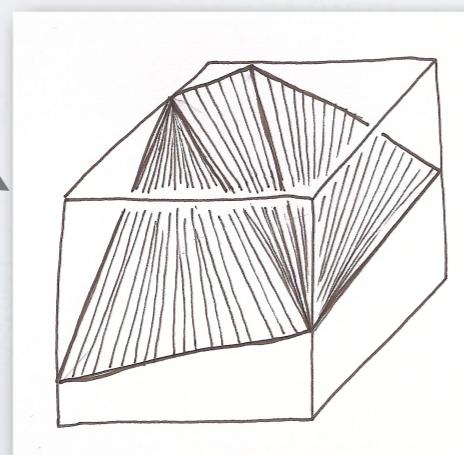
RECONSTRUCTION

PROCESSING

Betti numbers
Volume
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...

~~Building~~

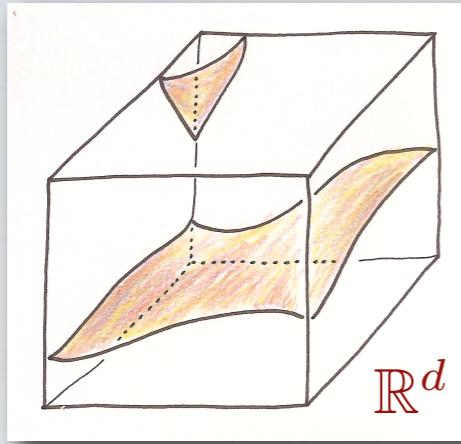
Delaunay complex



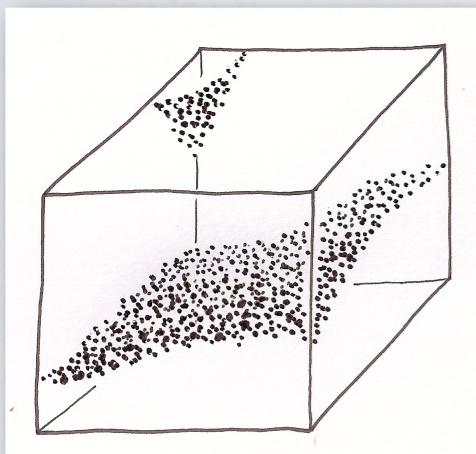
How to reconstruct without Delaunay?

Shape

in dD

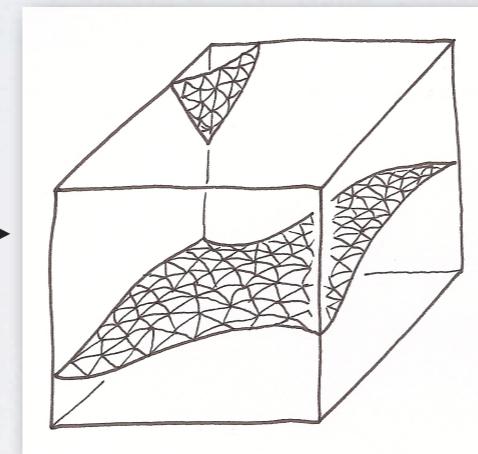


n points in \mathbb{R}^d



RECONSTRUCTION

Simplicial complex

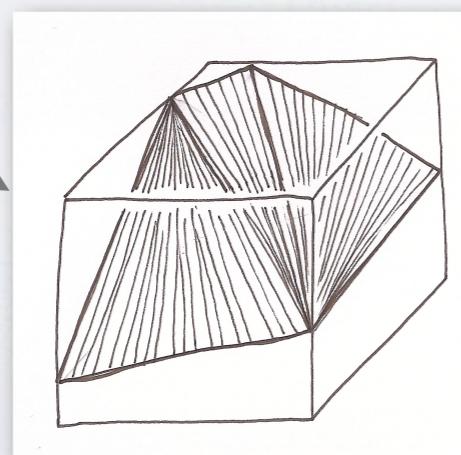


PROCESSING

Betti numbers
Volume
Medial axis
Signatures
...

~~BUILDING~~

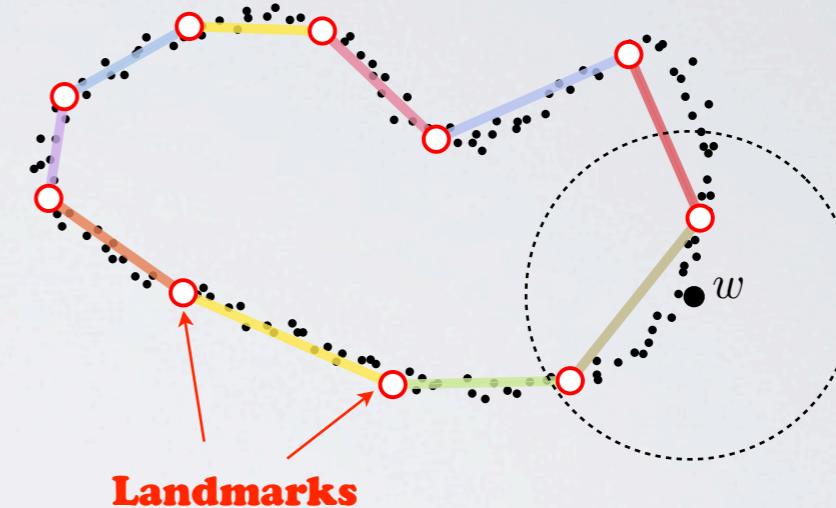
Delaunay complex



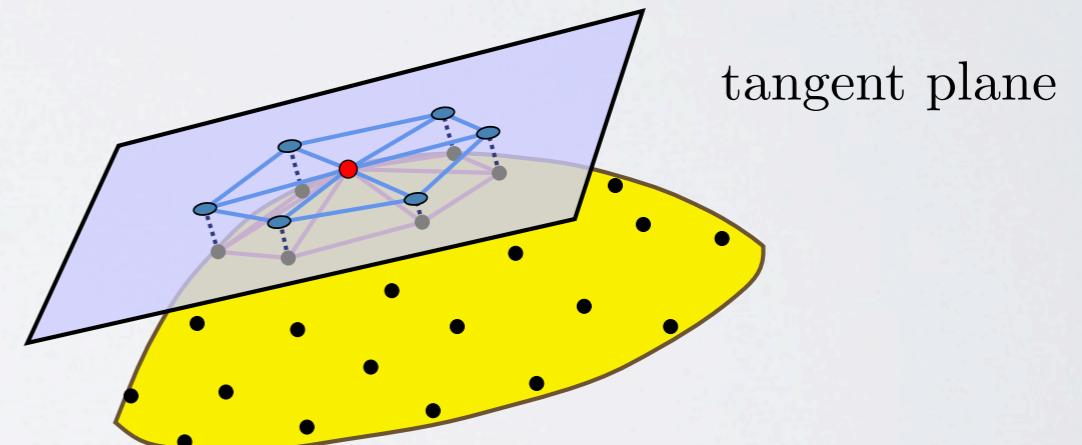
How to reconstruct without Delaunay?

How to reconstruct without building the whole Delaunay complex?

- * weak Delaunay triangulation
[V. de Silva 2008]

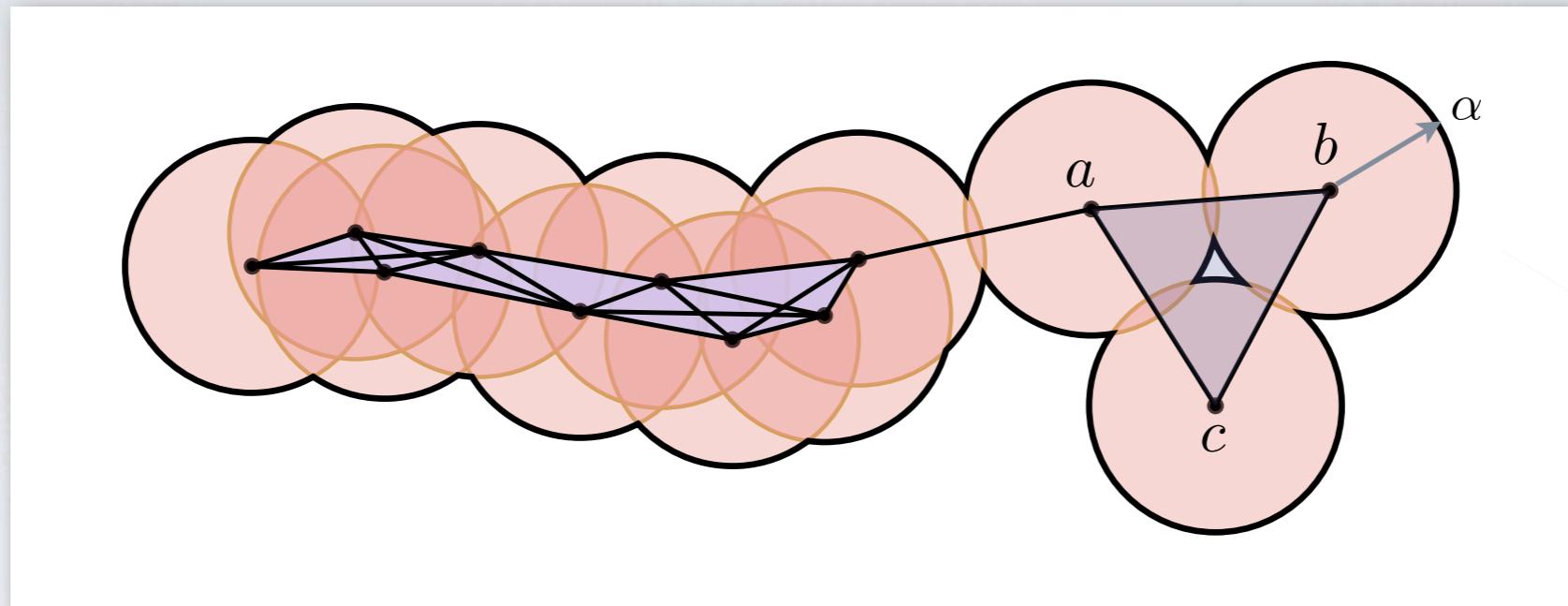


- * tangential Delaunay complexes
[J. D. Boissonnat & A. Ghosh 2010]



- * Rips complexes
our approach with André Lieutier and David Salinas

Rips complexes

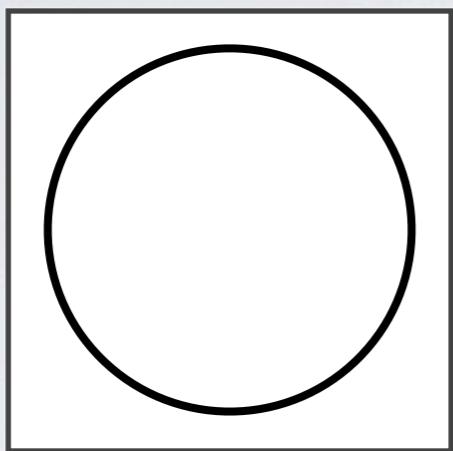


$$\text{Rips}(P, \alpha) = \{\sigma \subset P \mid \text{Diameter}(\sigma) \leq 2\alpha\}$$

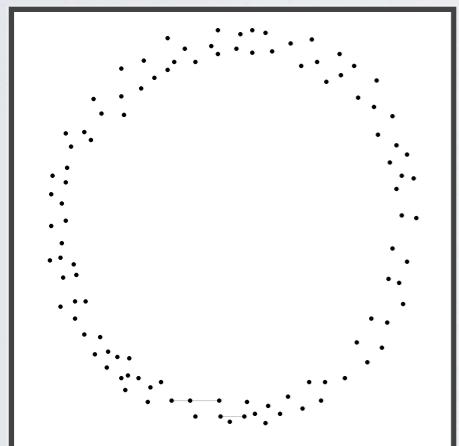
- ✳ proximity graph G_α connects every pair of points within 2α
- ✳ $\text{Rips}(P, \alpha) = \text{Flag } G_\alpha$ [Flag G = largest complex whose 1-skeleton is G]
- ✳ compressed form of storage through the 1-skeleton
- ✳ easy to compute

SHAPE RECONSTRUCTION

Shape *A*

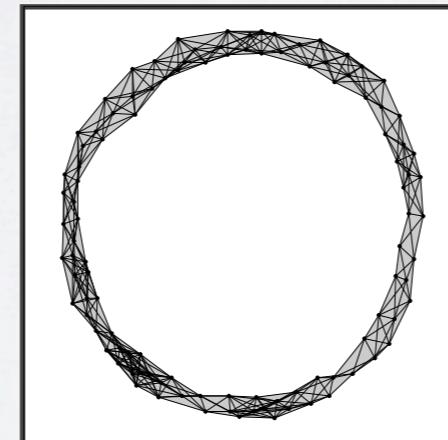


$P \subset \mathbb{R}^d$



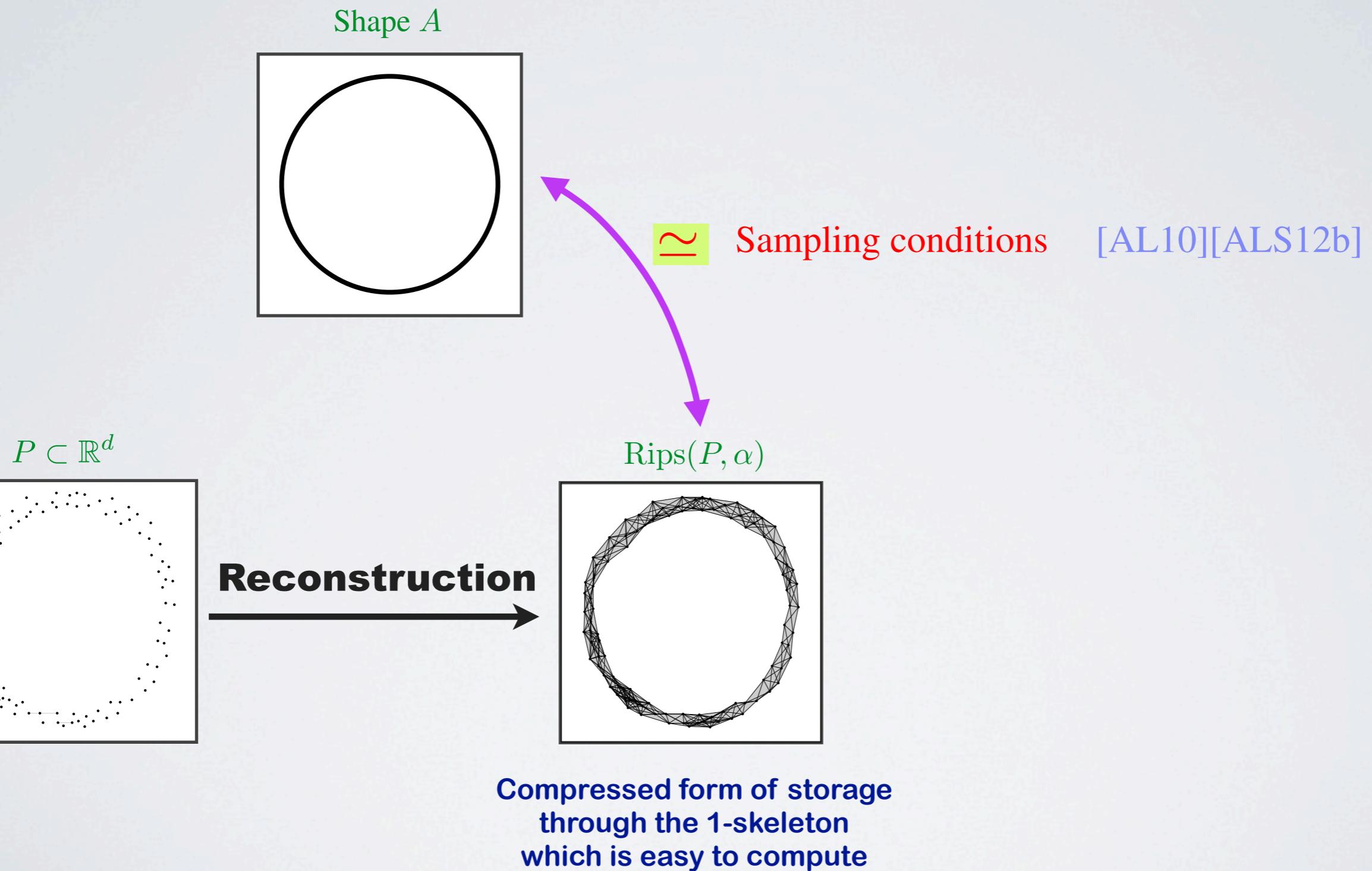
Reconstruction

Rips(P, α)

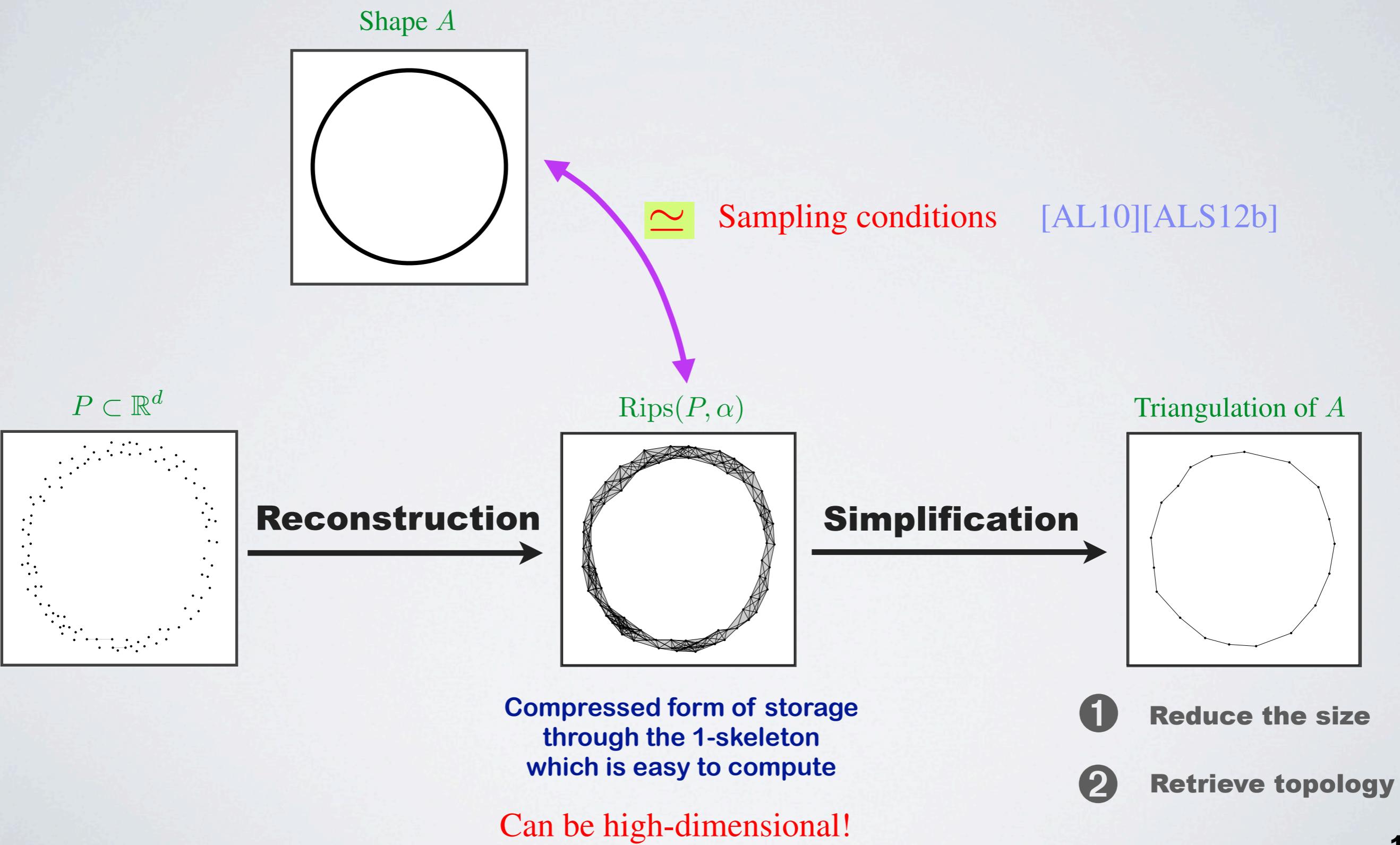


Compressed form of storage
through the 1-skeleton
which is easy to compute

SHAPE RECONSTRUCTION

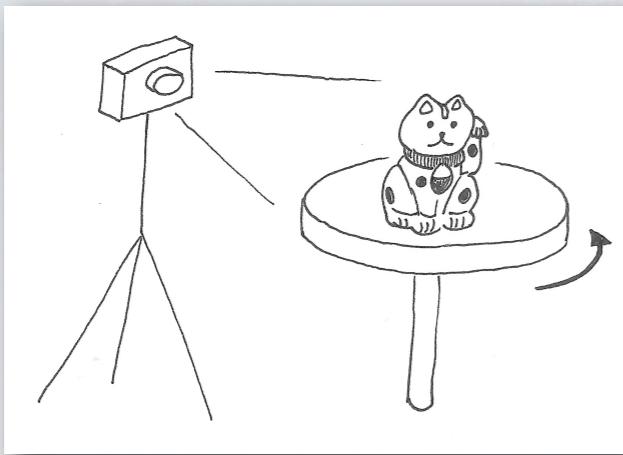


SHAPE RECONSTRUCTION



Example

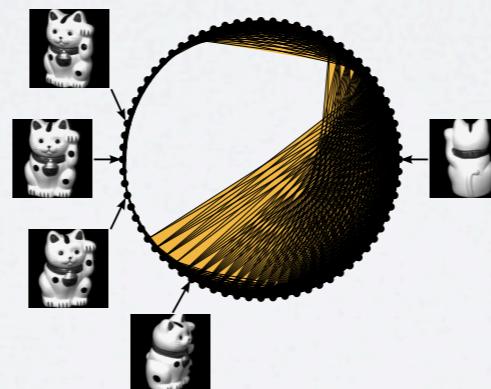
Physical system



Point cloud in \mathbb{R}^{128^2}



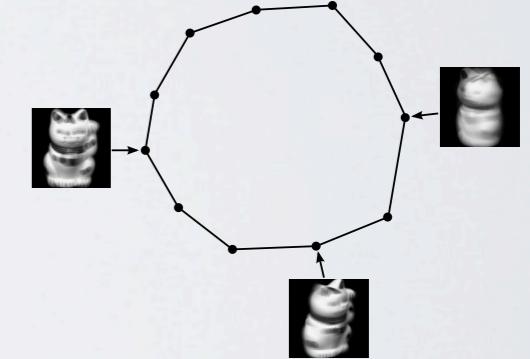
Rips complex



Correct homotopy type

Is high-dimensional!

Polygonal curve



Correct intrinsic dimension

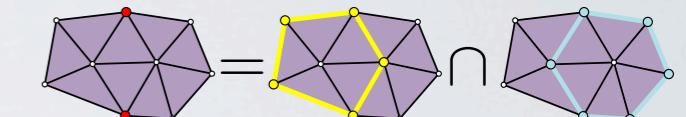
Simplification by iteratively applying elementary operations

- * Edge contraction $ab \mapsto c$



- * Identifies vertices a and b to vertex c

- * Preserves homotopy type if $\text{Lk}_K(ab) = \text{Lk}_K(a) \cap \text{Lk}_K(b)$

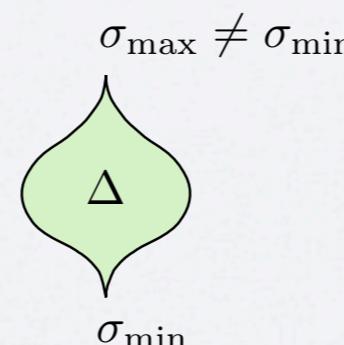
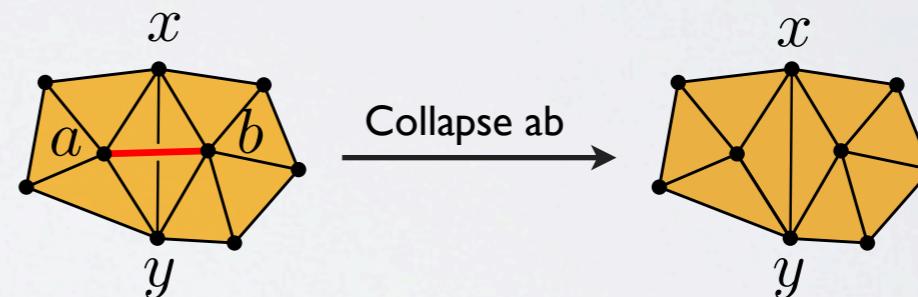


$$\text{Lk}_K(\sigma) = \{\tau \mid \tau \cap \sigma = \emptyset, \tau \cup \sigma \in K\}$$

- * Collapse of a simplex σ_{\min}

- * Removes σ_{\min} and its cofaces Δ

- * Preserves homotopy type if Δ has a unique maximal element $\sigma_{\max} \neq \sigma_{\min}$



Does a simplification exist?

- ✿ Different strategies:
 - ✿ Edge contractions;
 - ✿ Vertex and edge collapses;
 - ✿ Seems to work well in practice ...

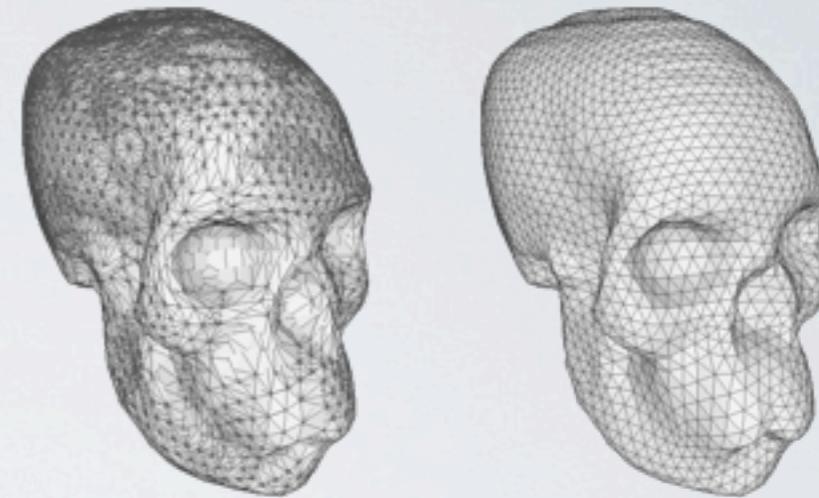
- ✿ And yet, not at all obvious that the Rips complex whose vertices sample a shape contains a subcomplex homeomorphic to that shape.

- ✿ A triangulated Bing's house is contractible but not collapsible



The diagram shows a sequence of geometric shapes illustrating the construction of a triangulated Bing's house. It starts with a cube, followed by a sequence of operations: a subtraction operation (indicated by a minus sign) followed by two union operations (indicated by a plus sign with a 'u' superscript), and finally another union operation. The resulting shape is a complex, symmetric structure with multiple interconnected cubes and rectangular prisms, representing the triangulated version of the Bing's house.

- ✿ Geometry has to play a key role.

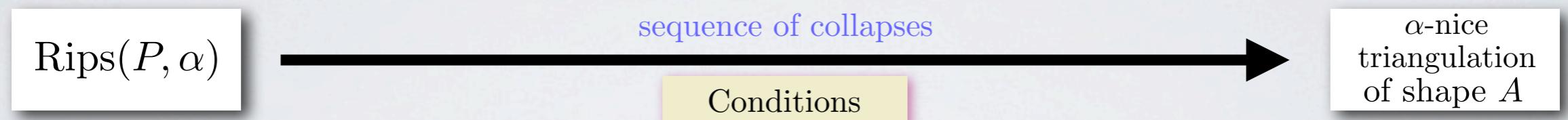


Simplifying Rips complexes

$A \subset \mathbb{R}^d$ is a compact set

$P \subset \mathbb{R}^d$ is a finite point set

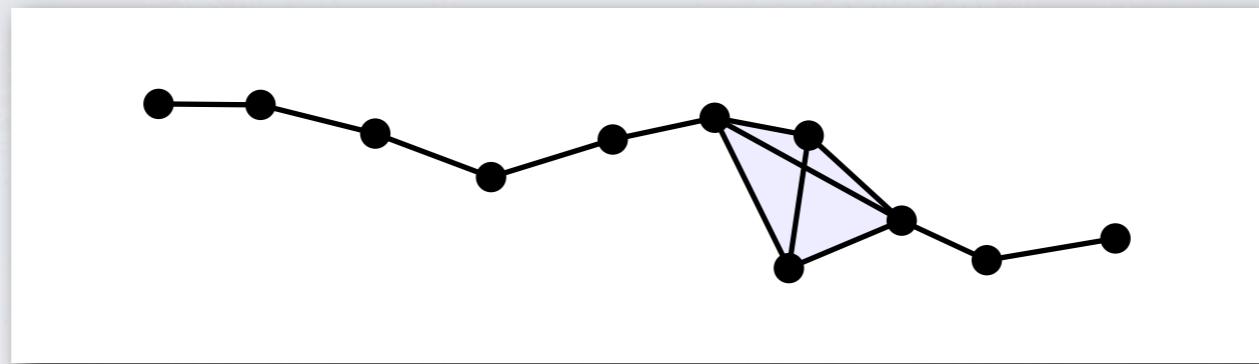
$\alpha > 0$



- ✳️ Unfortunately:
 - ✳️ the proof is not constructive (no algorithm!);
 - ✳️ it only works for shapes that have an α -nice triangulation!

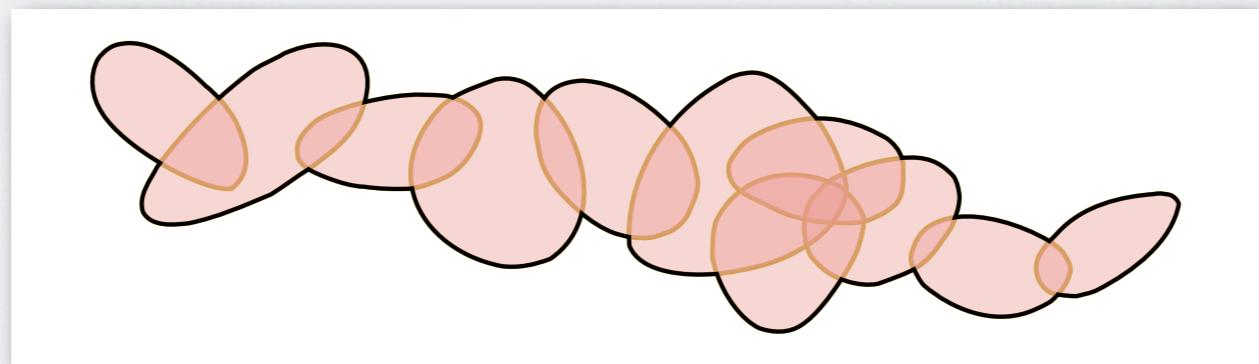
A key tool

$$\text{Nerve } \mathcal{C} = \{\sigma \subset P \mid \sigma \neq \emptyset \text{ and } \bigcap_{p \in \sigma} C_p \neq \emptyset\}$$



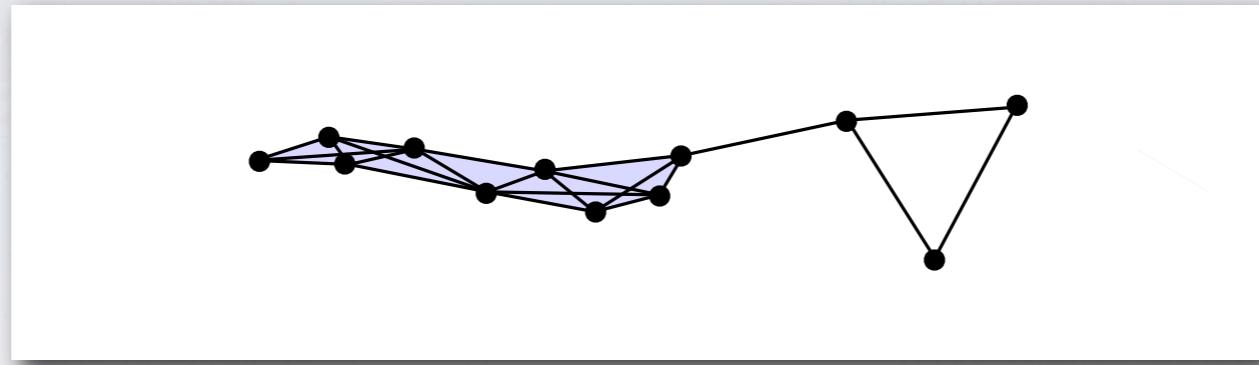
If $\bigcap_{z \in \sigma} C_z$ is either empty or contractible
Nerve Lemma.

$\bigcup \mathcal{C}$, where $\mathcal{C} = \{C_p \mid p \in P\}$ finite collection of closed sets



Čech complexes

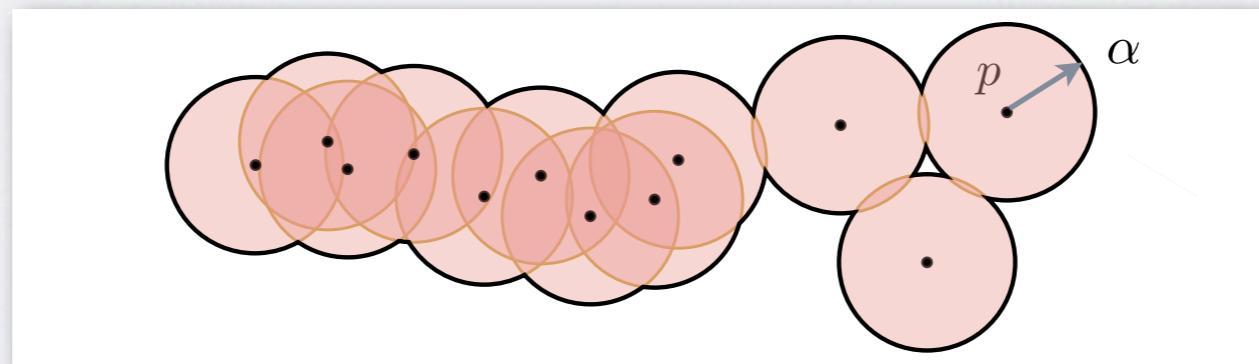
$$\text{Cech}(P, \alpha) = \text{Nerve}\{B(p, \alpha) \mid p \in P\}$$



↑
 \simeq

Nerve Lemma.

$$P^{\oplus \alpha} = \bigcup_{p \in P} B(p, \alpha) \quad \text{α-offset of } P$$

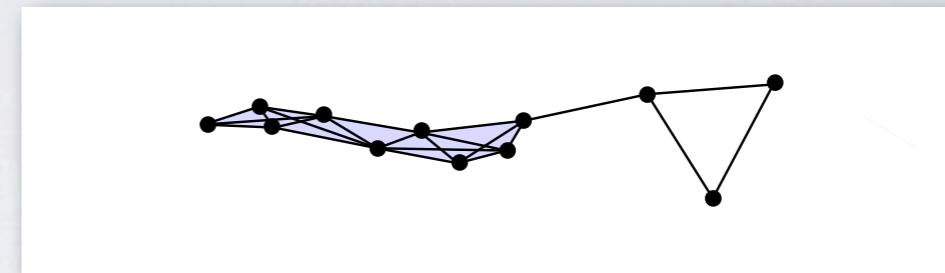
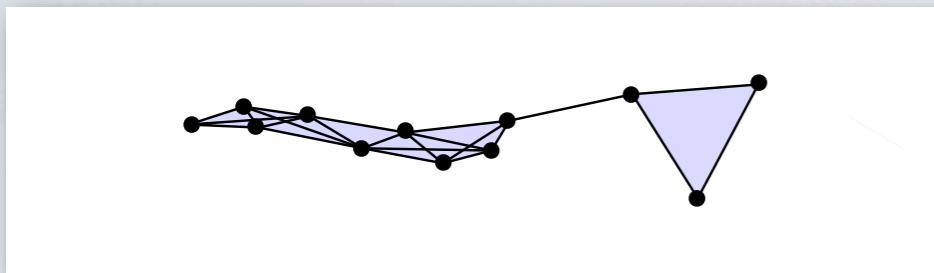


Čech complexes

$\text{Rips}(P, \alpha)$

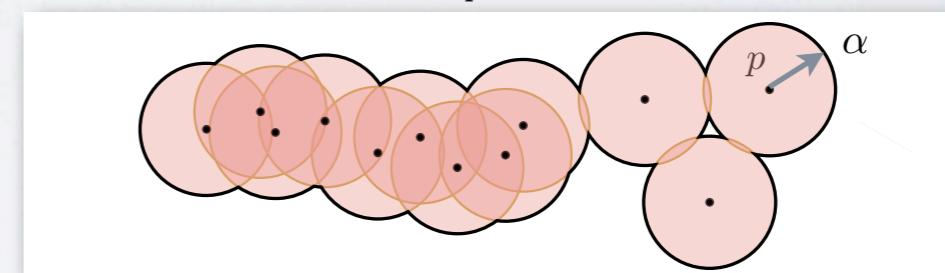
\supset

$\text{Cech}(P, \alpha) = \text{Nerve}\{B(p, \alpha) \mid p \in P\}$

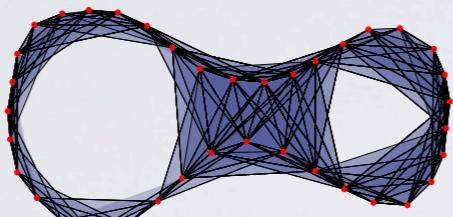
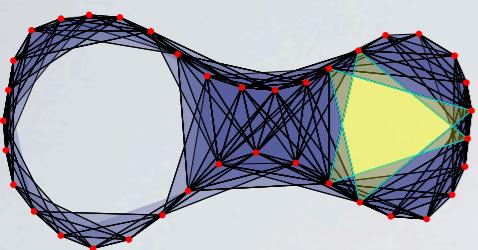


$\uparrow \downarrow \simeq$ **Nerve Lemma.**

$$P^{\oplus \alpha} = \bigcup_{p \in P} B(p, \alpha) \quad \text{α-offset of P}$$



Overview of what we knew!



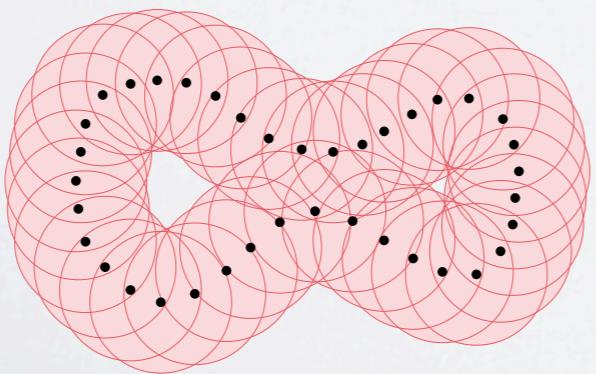
$\text{Rips}(P, \alpha)$

\supset

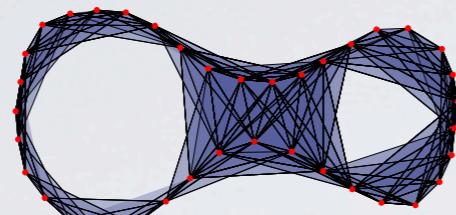
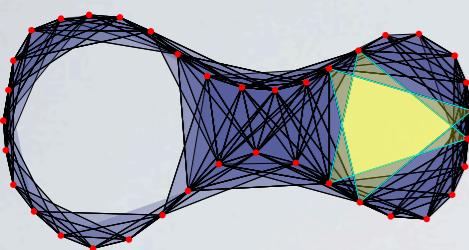
$\text{Cech}(P, \alpha)$

$$\begin{array}{c} \uparrow \\ \simeq \text{ Nerve Lemma} \\ \downarrow \end{array}$$

$P^{\oplus \alpha}$



Overview of what we knew!



$\text{Rips}(P, \alpha)$

\supset

$\text{Cech}(P, \alpha)$

$$\begin{array}{c} \uparrow \\ \simeq \text{ Nerve Lemma} \\ \downarrow \end{array}$$

$P^{\oplus \alpha}$

deform. retracts

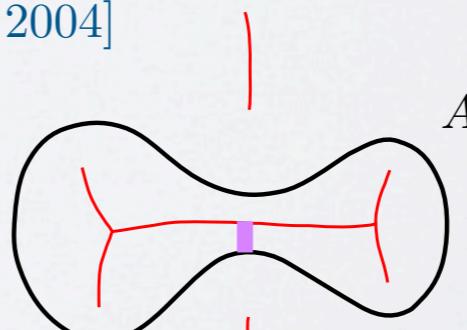
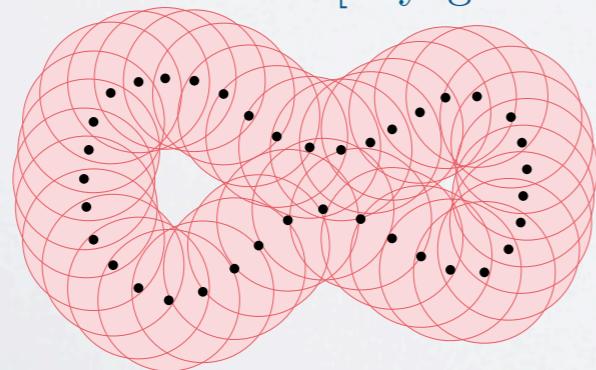
Shape A

(SC1)

$$d_H(A, P) \leq \varepsilon < (3 - \sqrt{8}) \text{Reach } A$$

$$\alpha = (2 + \sqrt{2})\varepsilon$$

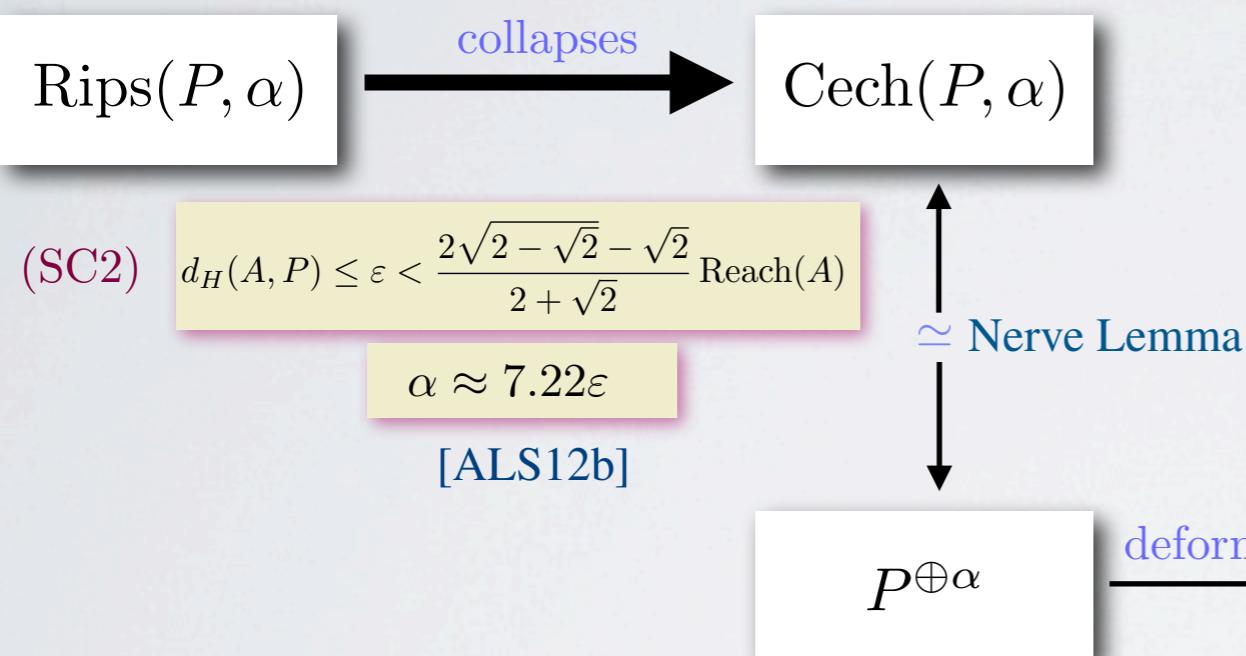
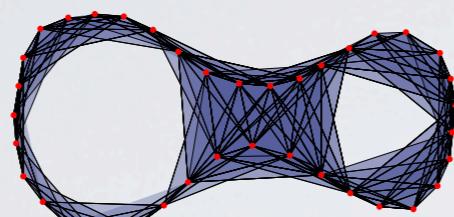
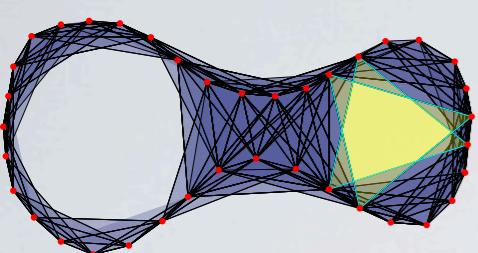
[Niyogi Smale Weinberger 2004]



$$\text{Reach } A = d(A, \text{MedialAxis}(A))$$

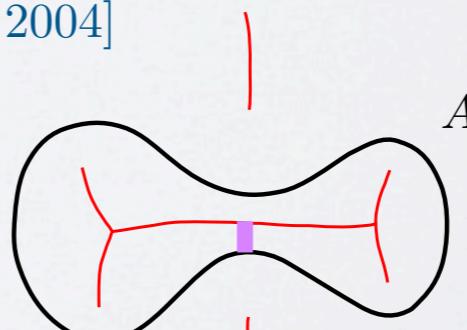
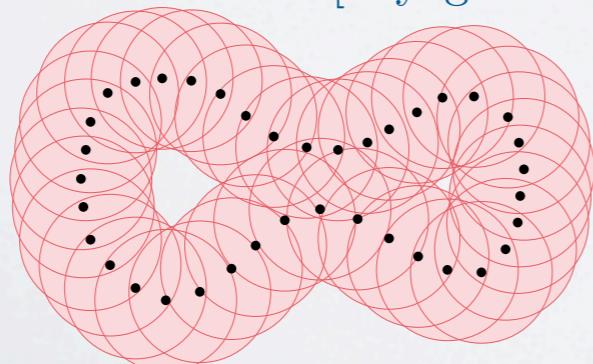
$$\text{MedialAxis}(A) = \{m \in \mathbb{R}^d \mid m \text{ has at least two closest points in } A\}$$

Overview of what we knew!



$(\text{SC1}) \quad d_H(A, P) \leq \varepsilon < (3 - \sqrt{8}) \text{Reach } A$
 $\alpha = (2 + \sqrt{2})\varepsilon$

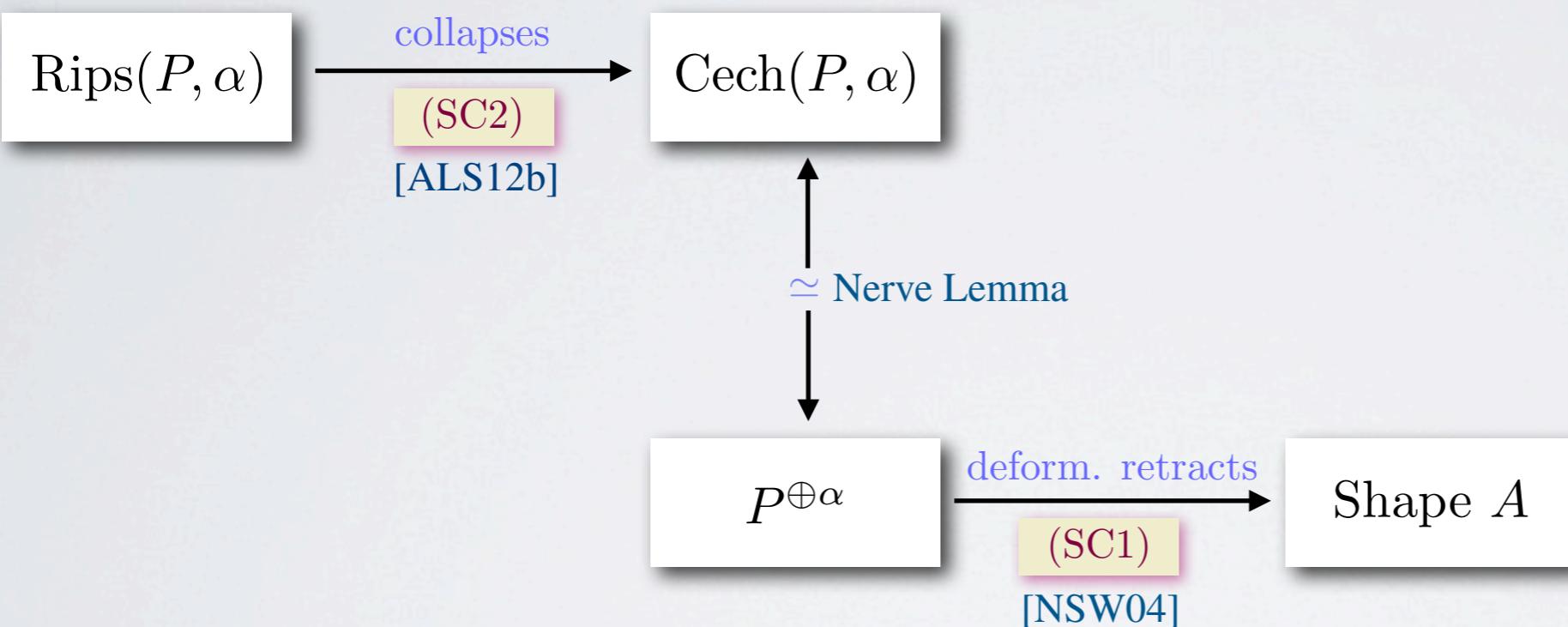
[Niyogi Smale Weinberger 2004]



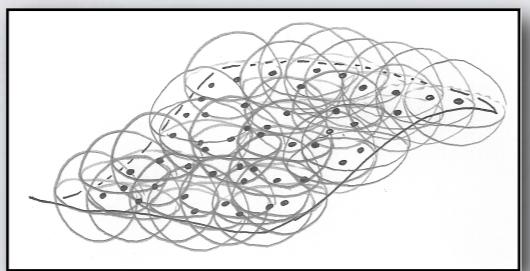
$\text{MedialAxis}(A) = \{m \in \mathbb{R}^d \mid m \text{ has at least two closest points in } A\}$ 21

$\text{Reach } A = d(A, \text{MedialAxis}(A))$

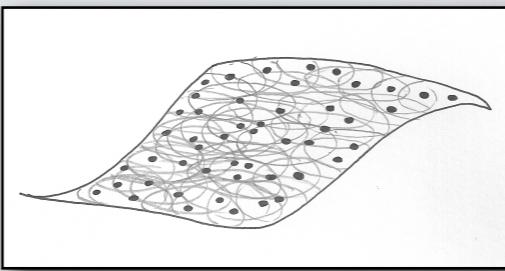
Overview of what we knew!



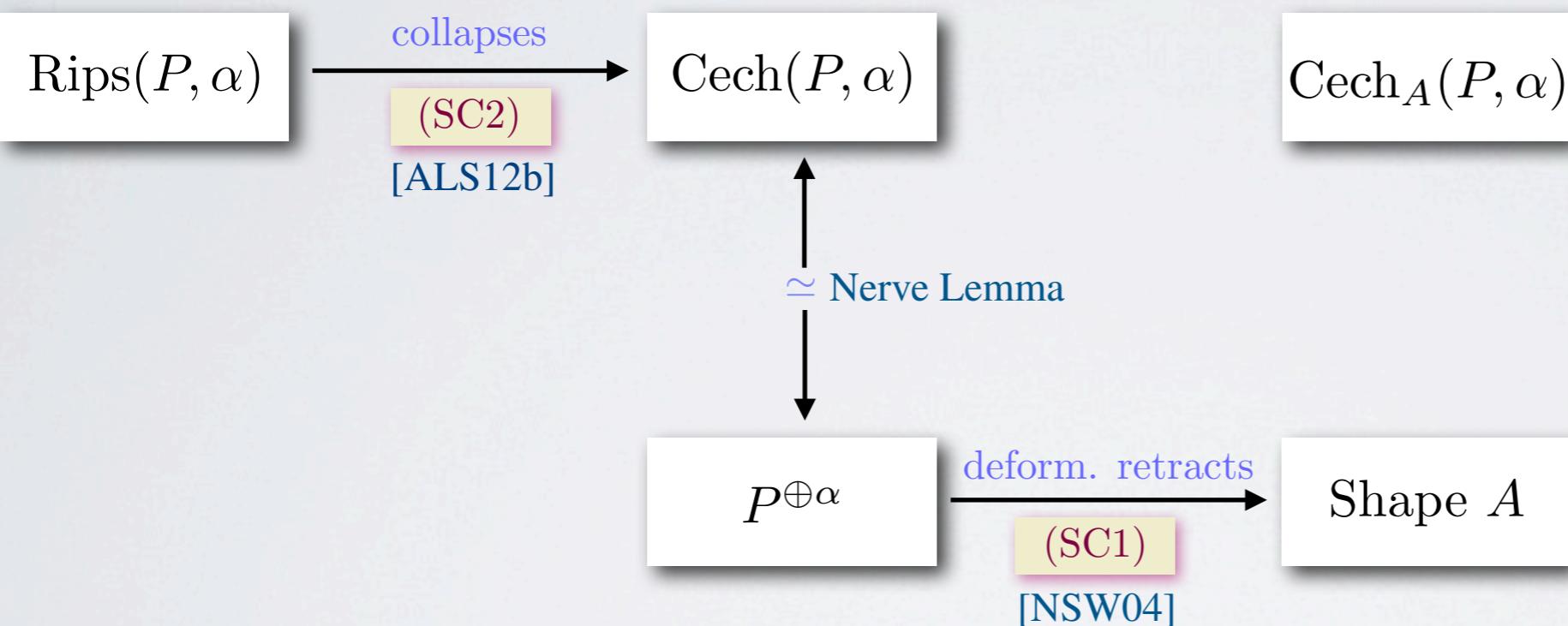
Overview of what is new



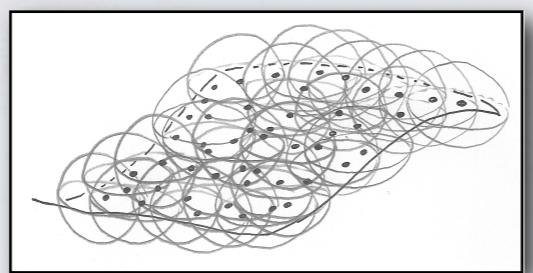
$\text{Nerve}\{B(p, \alpha) \mid p \in P\}$



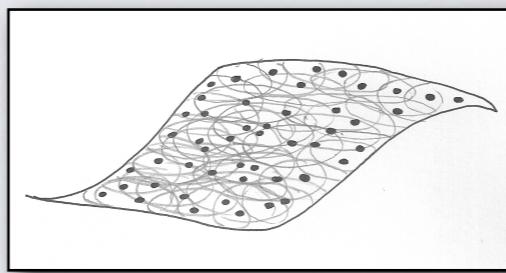
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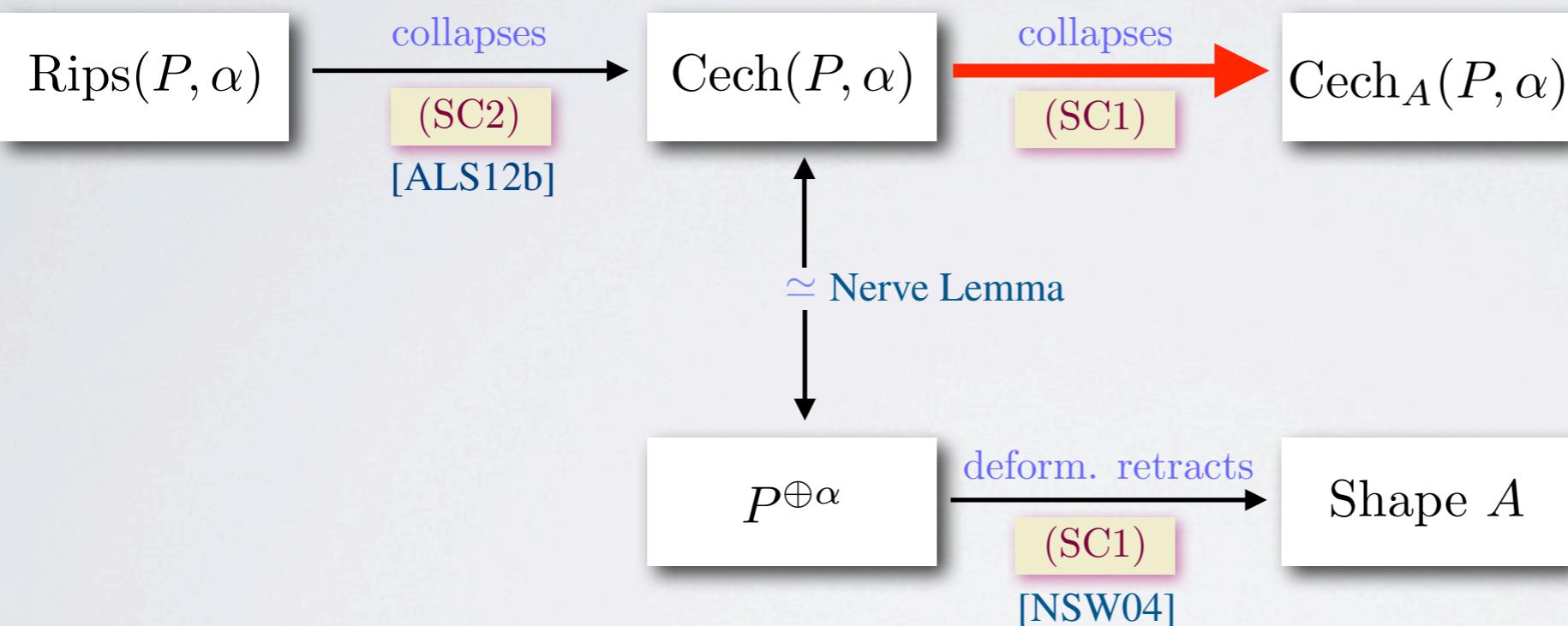
Overview of what is new



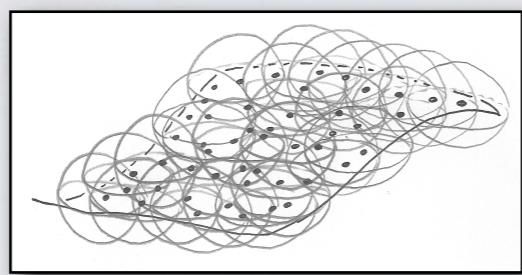
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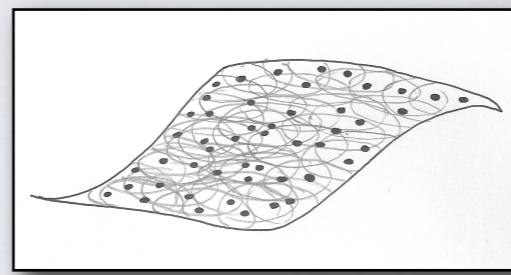
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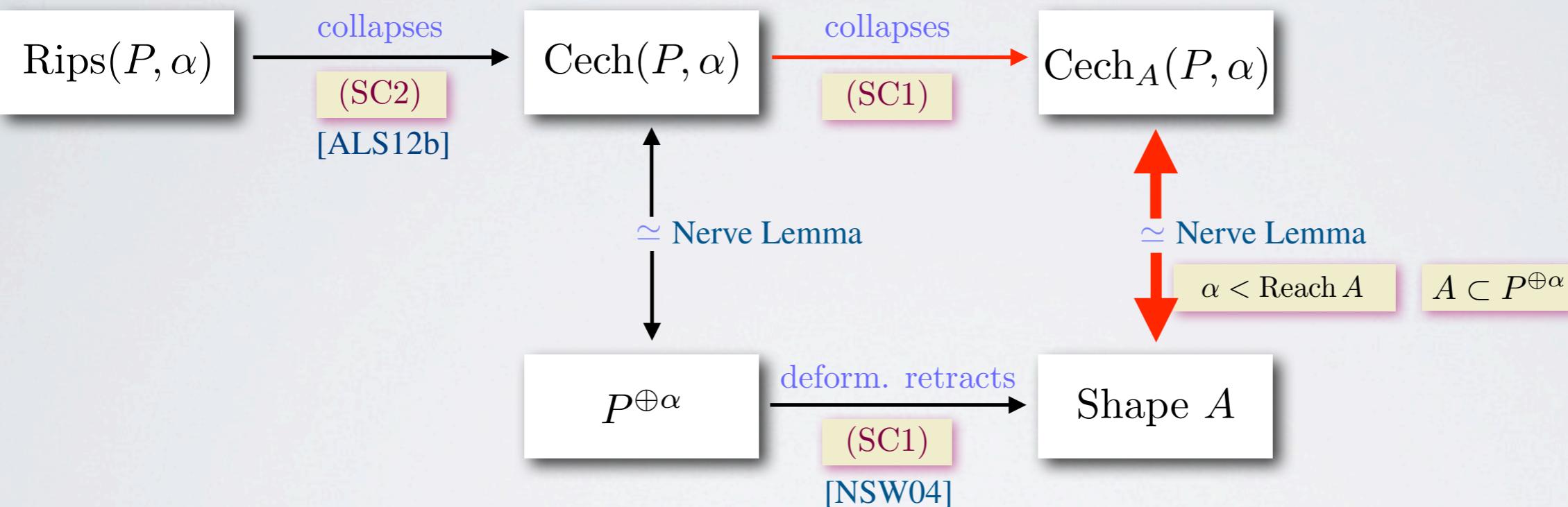
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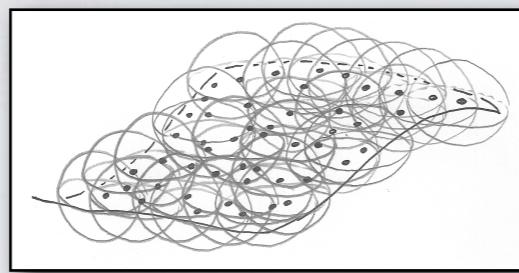
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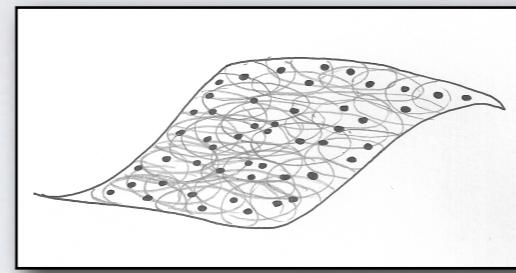
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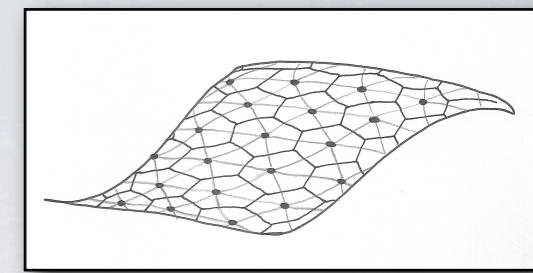
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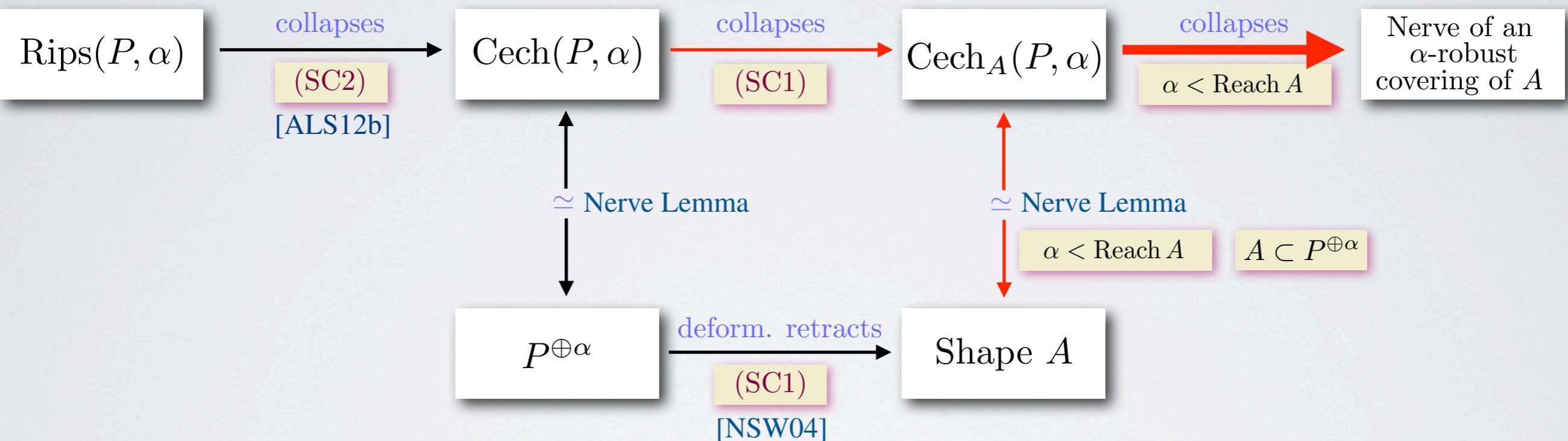
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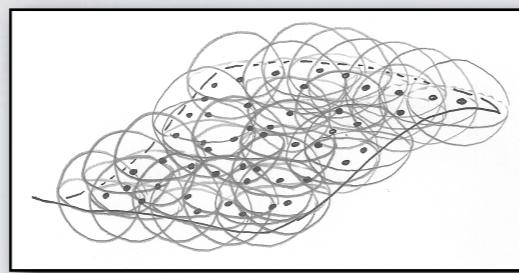
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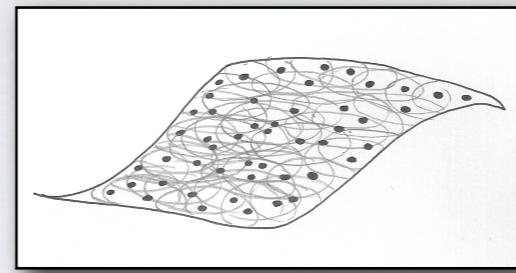
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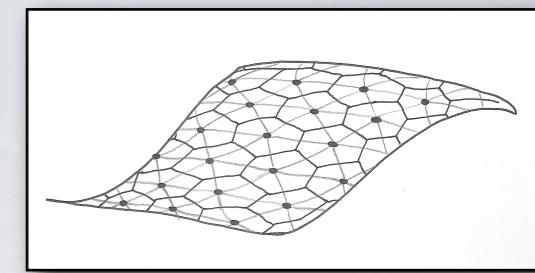
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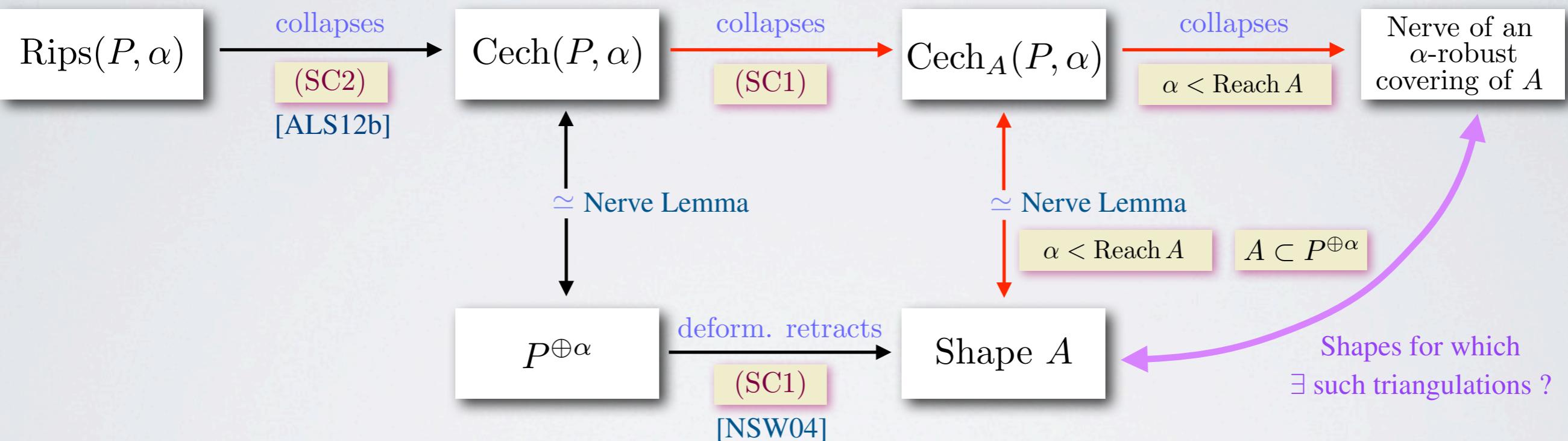
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$\text{Nerve}\{A \cap B(p, \alpha) \mid p \in P\}$

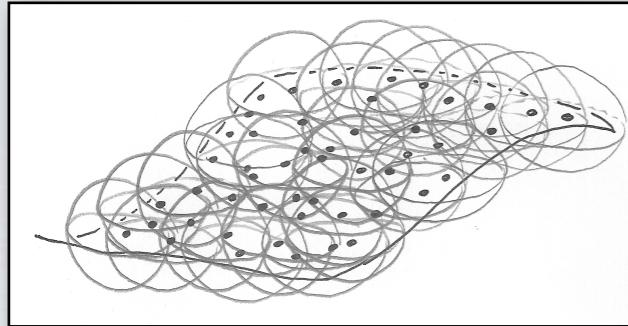


$\text{Nerve}\{A \cap \text{Hull}_\alpha(C_v) \mid v \in V\}$



Restricting the Čech complex

Theorem 2 If $d_H(A, P) \leq \varepsilon < (3 - \sqrt{8}) \text{Reach}(A)$ and $\alpha = (2 + \sqrt{2})\varepsilon$, then there exists a sequence of collapses from $\text{Cech}(P, \alpha)$ to $\text{Cech}_A(P, \alpha)$.



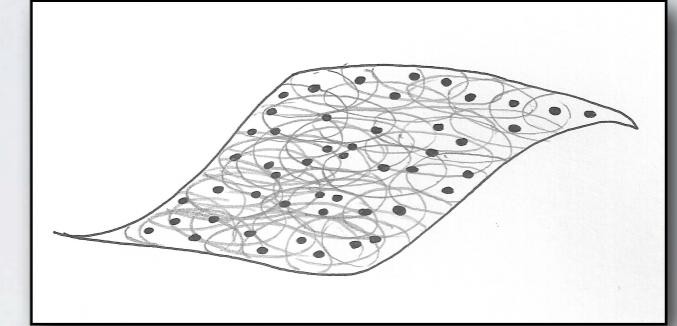
$\text{Nerve}\{B(p, \alpha) \mid p \in P\}$

\parallel
 $K(+\infty)$

$\text{Cech}(P, \alpha)$

Define collapses?

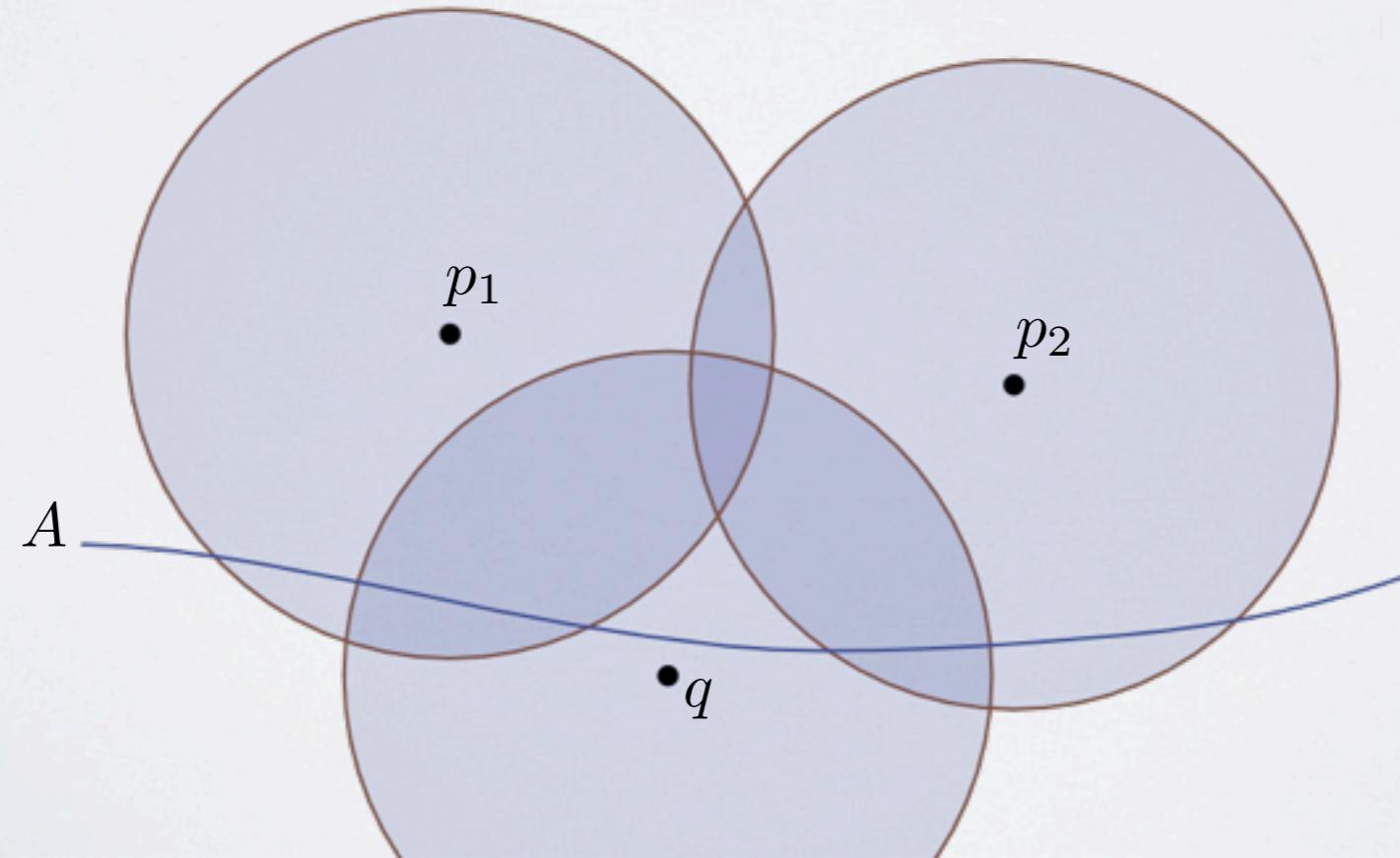
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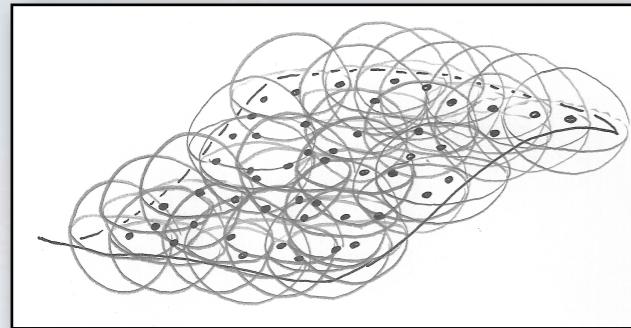
\parallel
 $K(0)$

$K(t) = \text{Nerve}\{A^t \cap B(p, \alpha) \mid p \in P\}$



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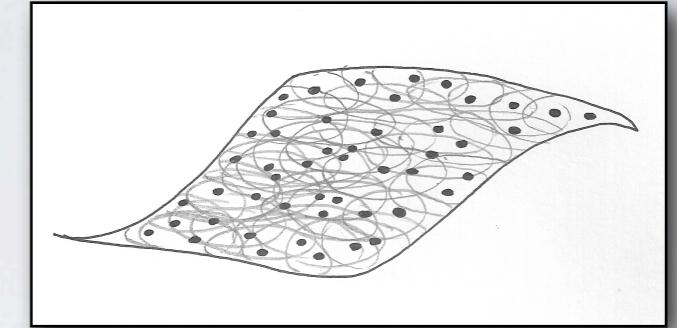
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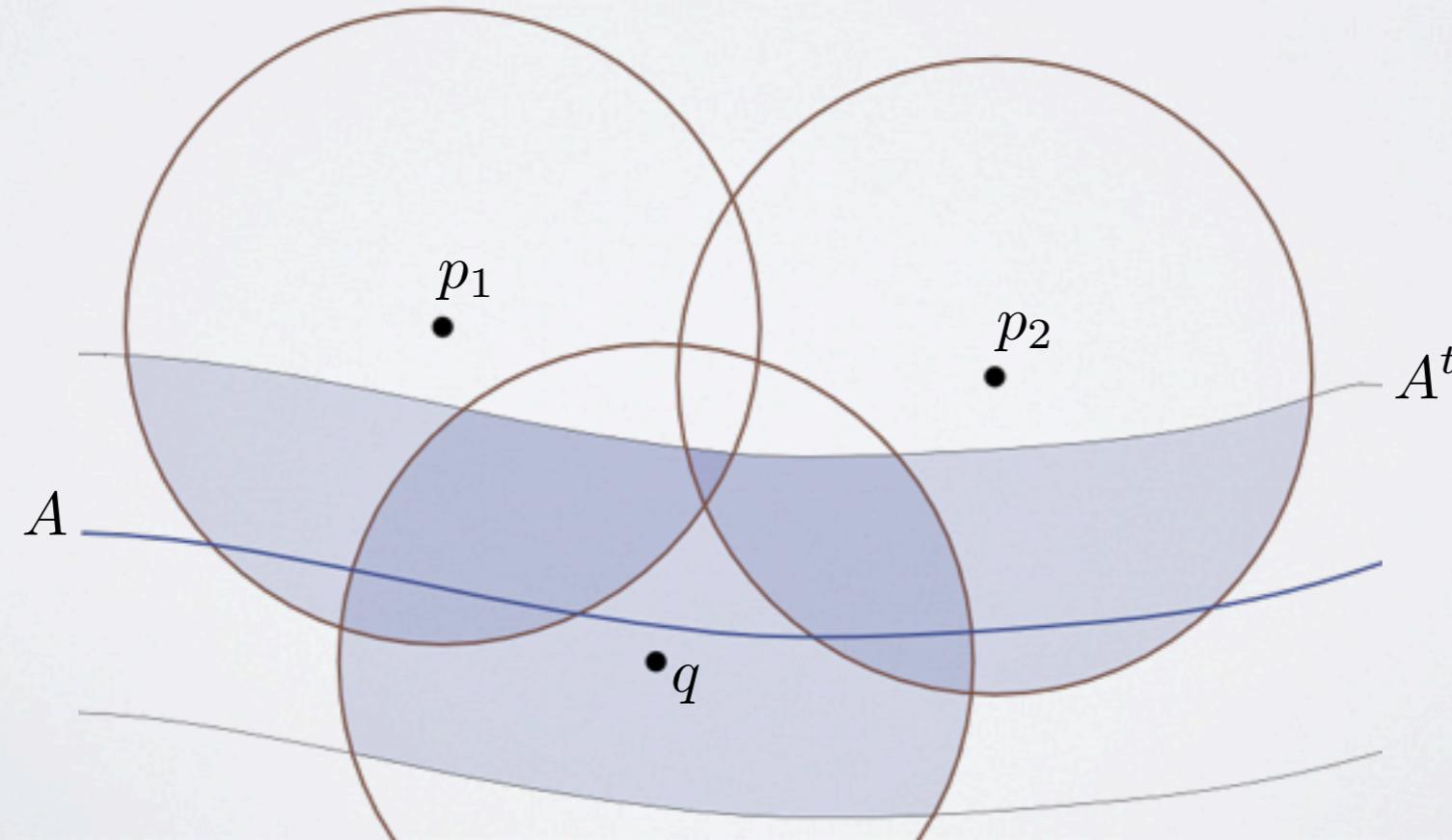
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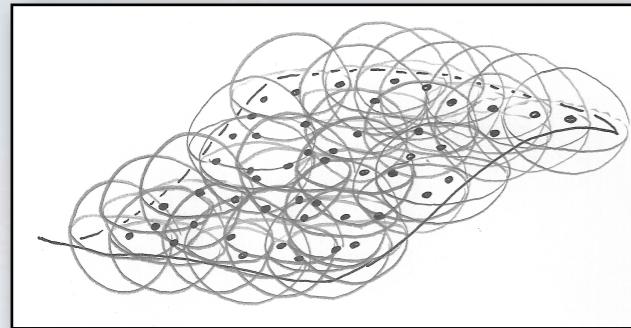
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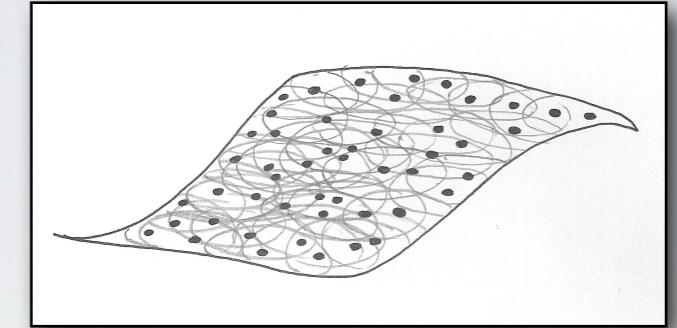
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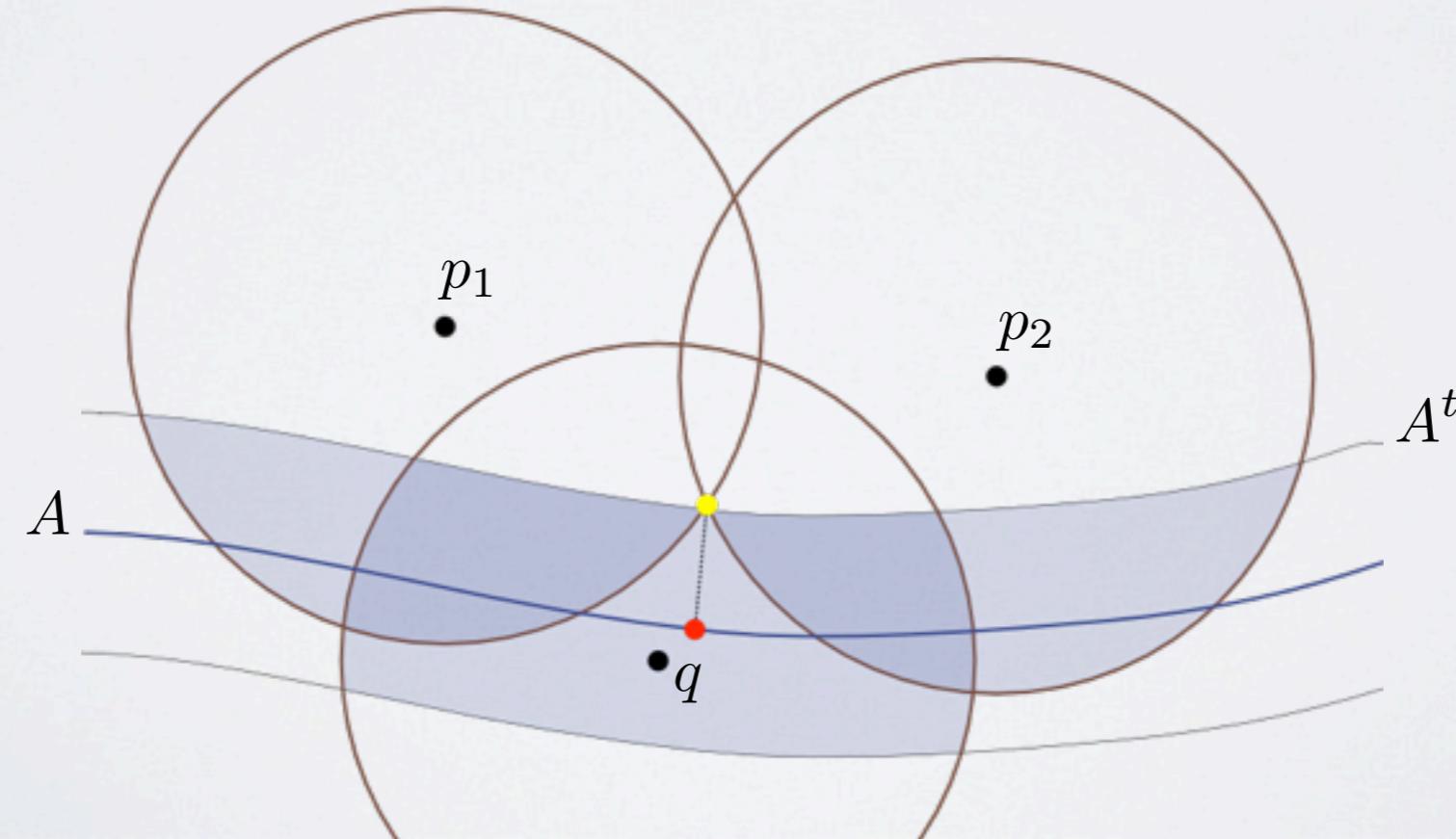
$\text{Cech}_A(P, \alpha)$



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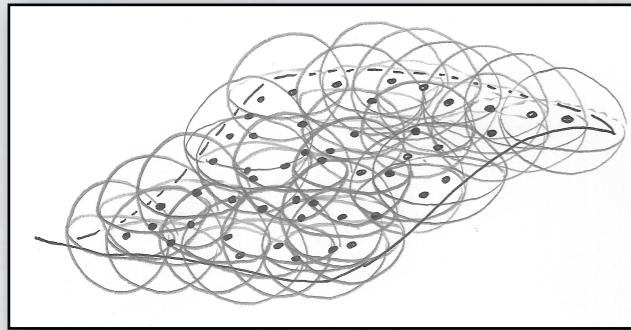
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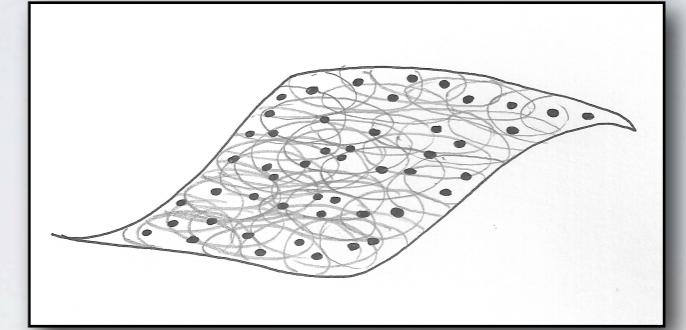
Restricting the Čech complex

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$\text{Cech}(P, \alpha)$

Define collapses?



$\text{Cech}_A(P, \alpha)$

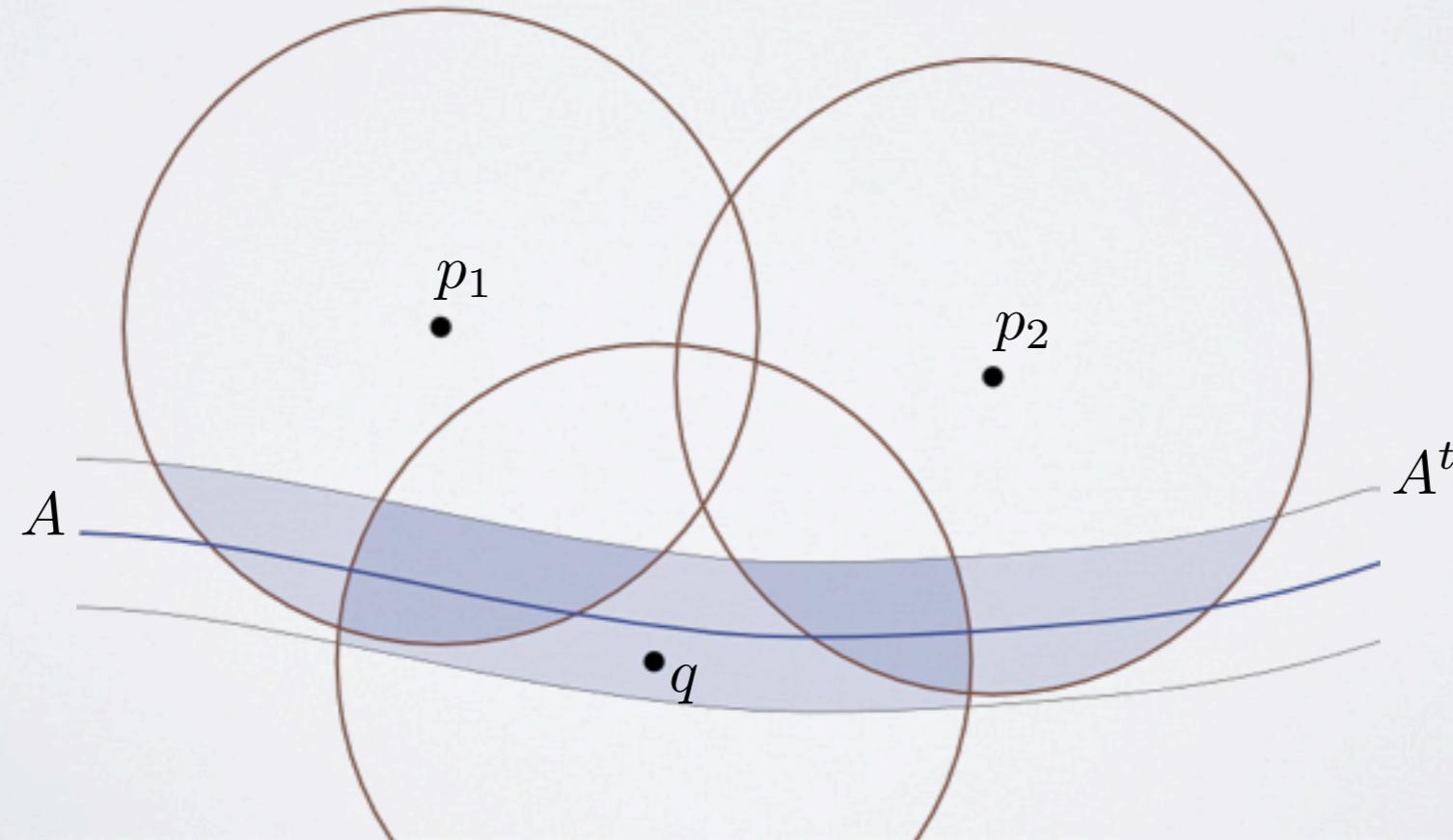
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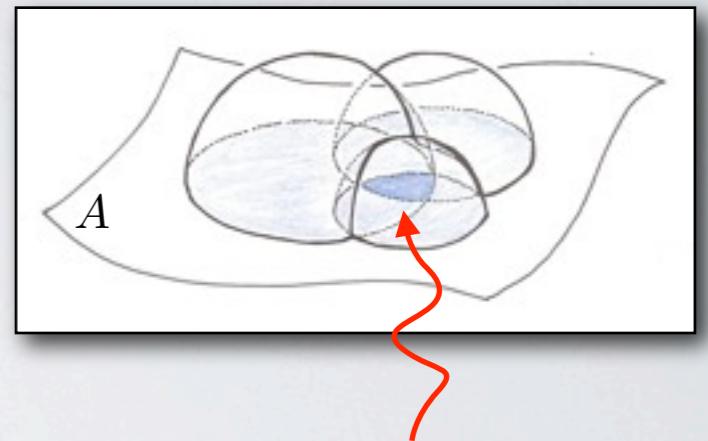
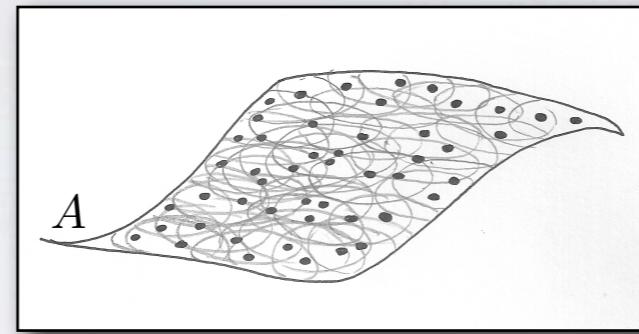
The restricted Čech complex

Theorem 1 If $\alpha < \text{Reach}(A)$ and $A \subset P^{\oplus\alpha}$, then $\text{Čech}_A(P, \alpha) \simeq A$.

$$\text{Čech}_A(P, \alpha) = \text{Nrv}\{A \cap B(p, \alpha) \mid p \in P\}$$

\uparrow \downarrow
 $\simeq?$ Apply the Nerve Lemma

$$A = \bigcup_{p \in P} [A \cap B(p, \alpha)]$$



contractible?



Technical Lemma. $A \cap \bigcap_{z \in \text{compact subset } \sigma} B(z, \alpha)$ is either empty or contractible if $\alpha < \text{Reach}(A)$

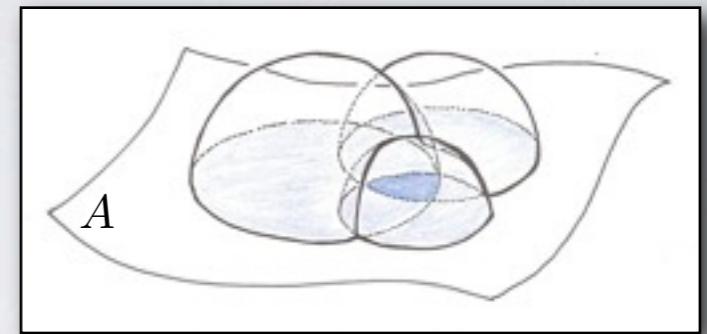
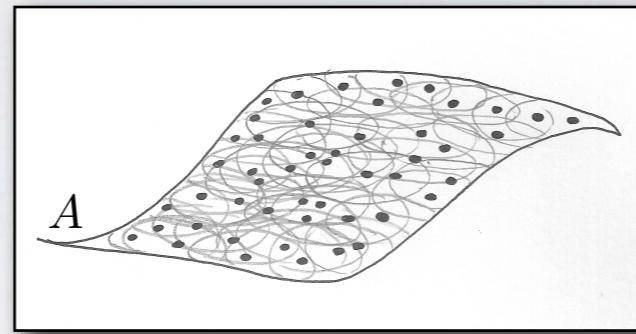
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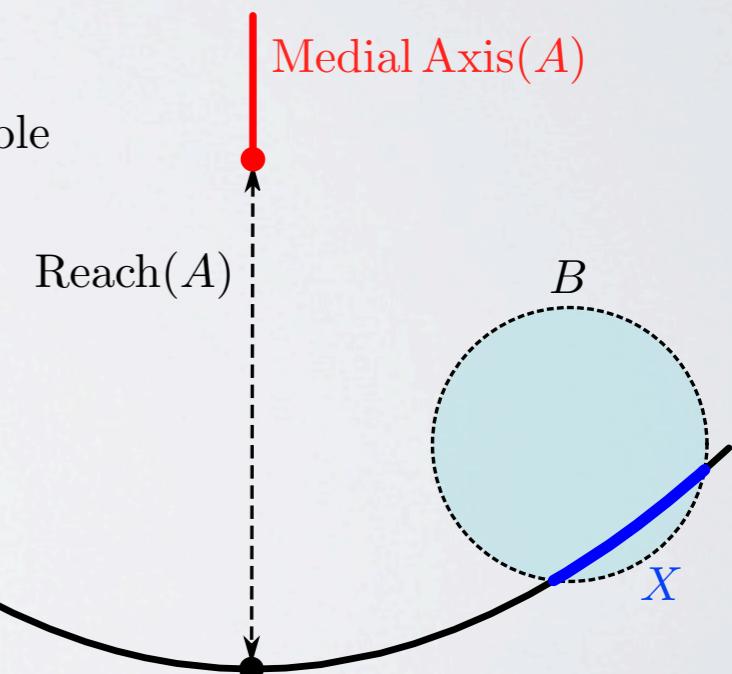
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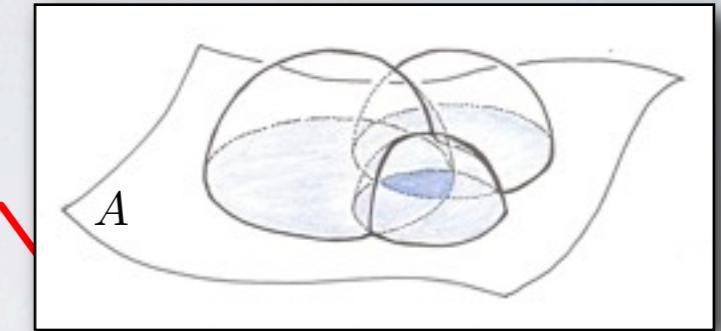
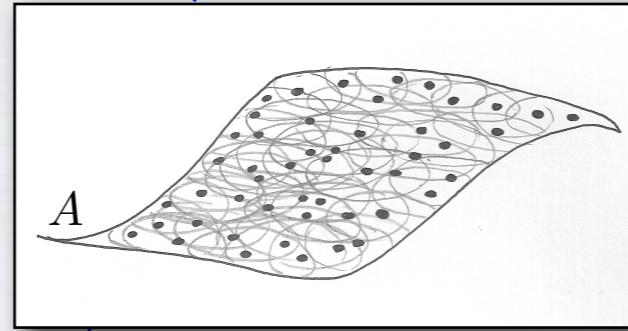
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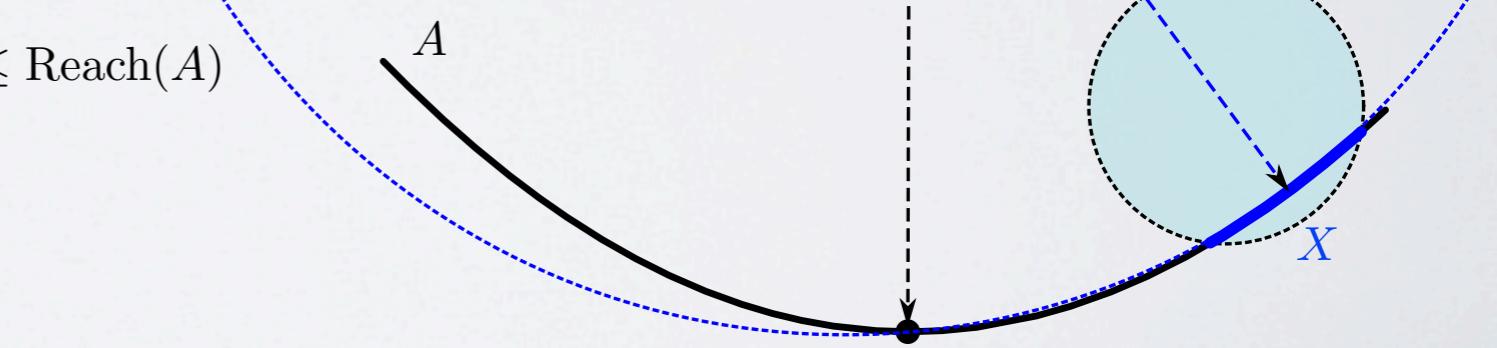
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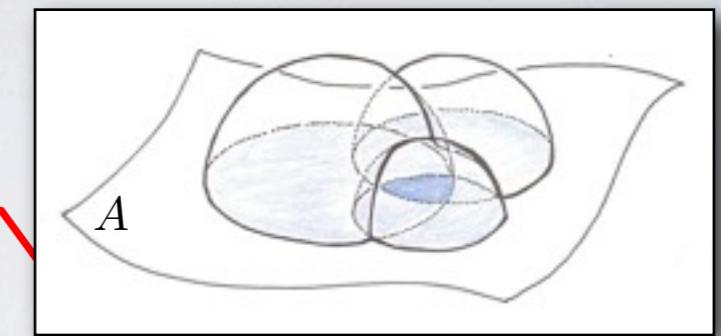
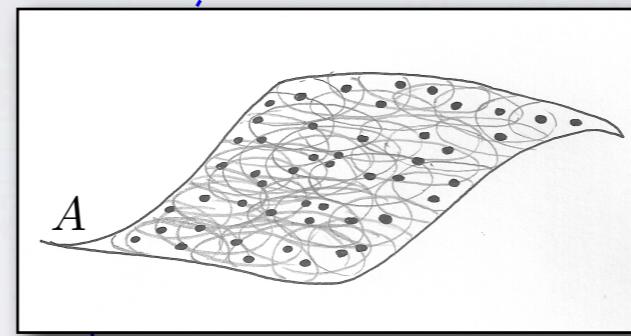
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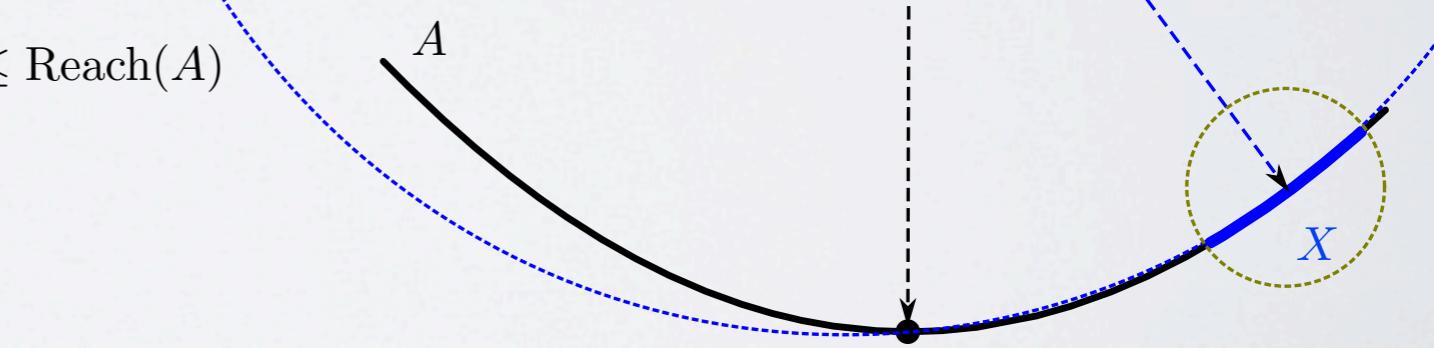
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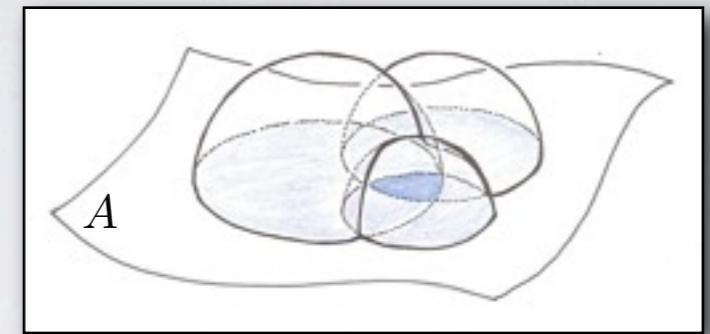
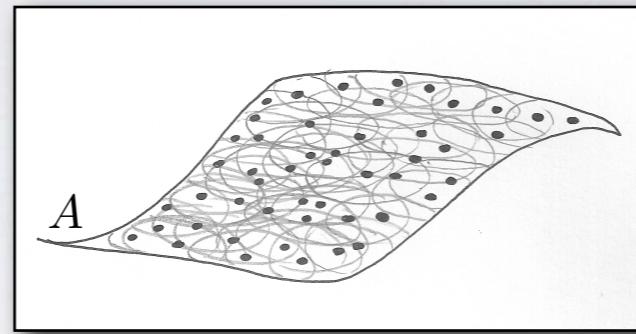
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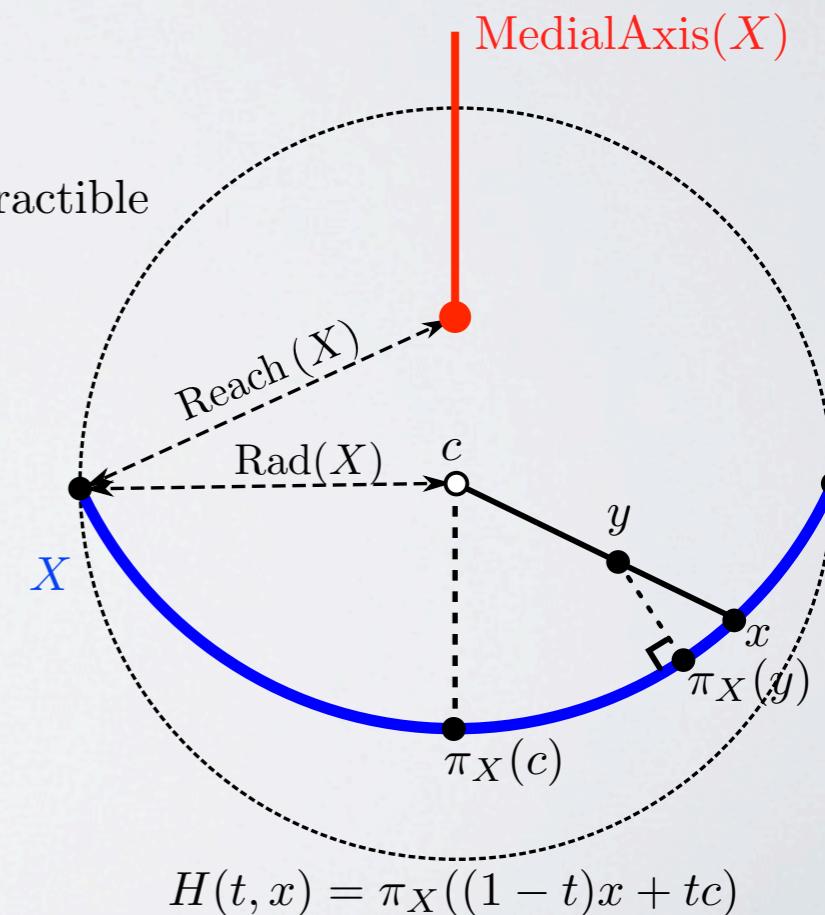
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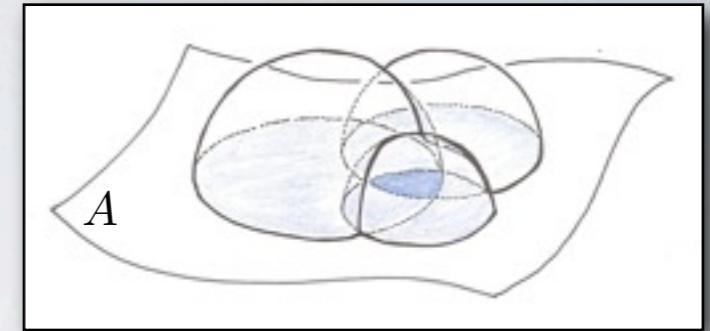
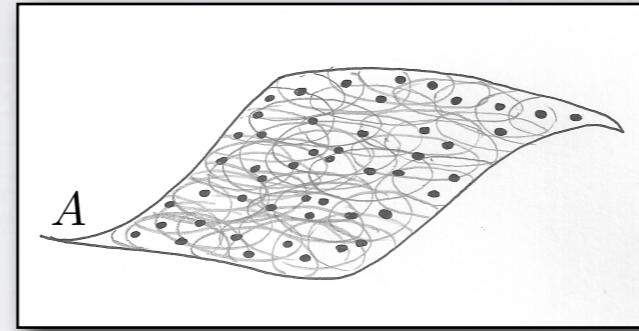
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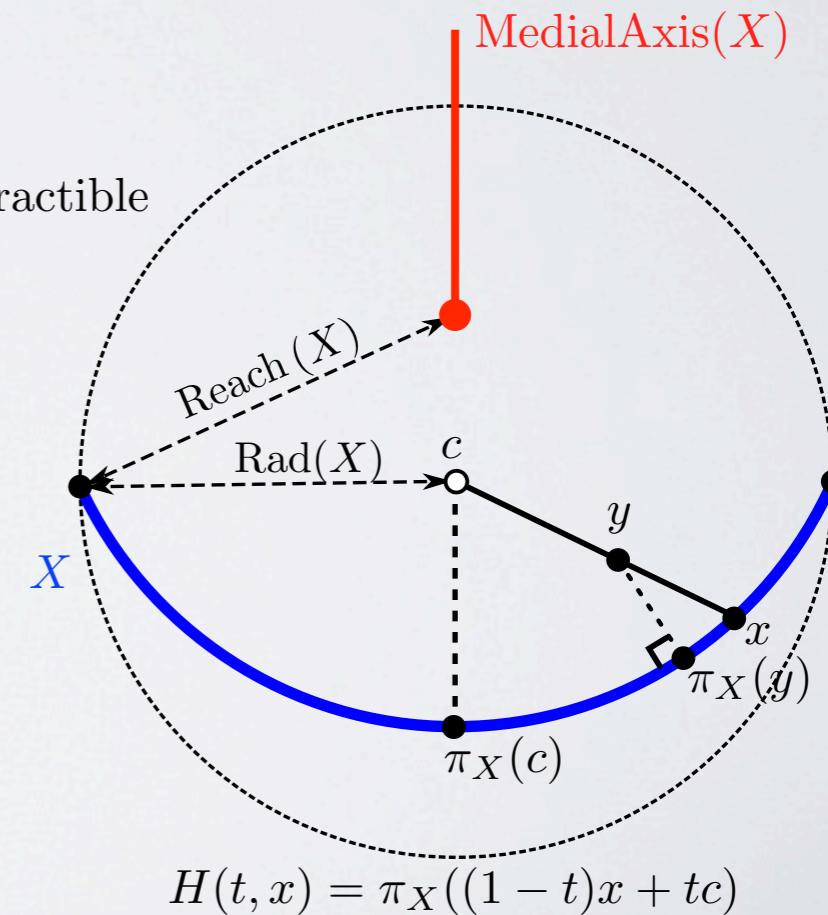
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$$\text{Rad}(A \cap \bigcap_{z \in \sigma} B(z, \alpha)) \leq \alpha < \text{Reach}(A) \leq \text{Reach}(A \cap \bigcap_{z \in \sigma} B(z, \alpha))$$

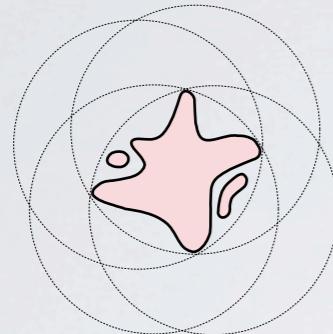


α -robust coverings



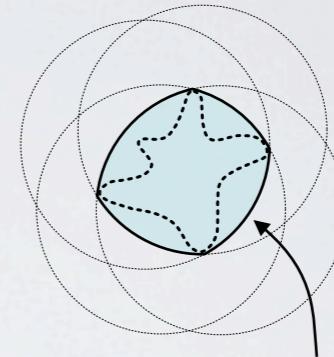
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$X \subset A$ with $\text{Rad}(X) < \alpha$



not necessarily contractible!

$A \cap \text{Hull}_\alpha(X)$



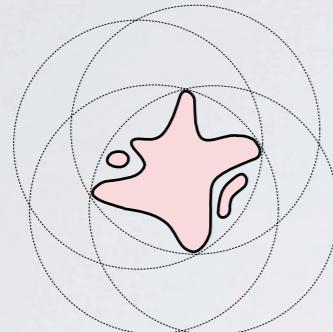
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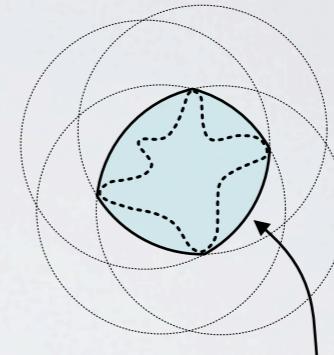
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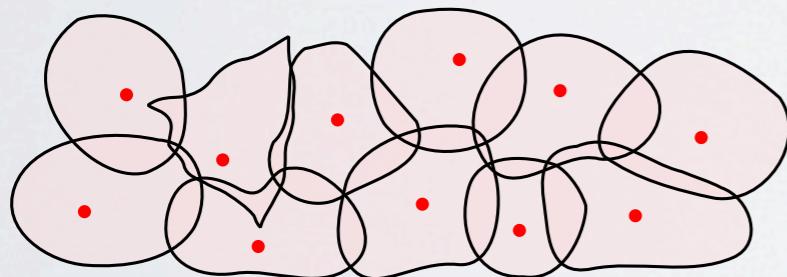
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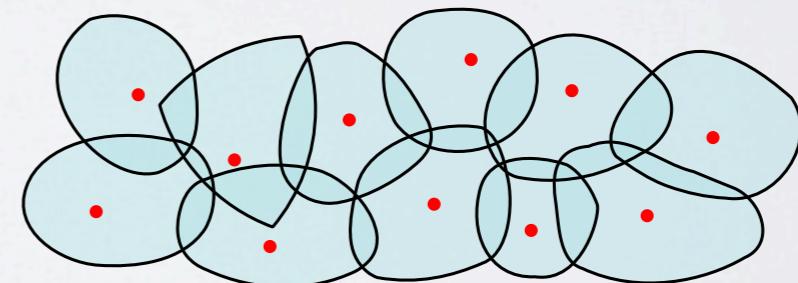
contractible
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$$\mathcal{C} = \{C_v \mid v \in V\}$$



\mathcal{C} α -robust covering of A
if $\text{Nrv } \mathcal{C} = \text{Nrv } \mathcal{C}'$

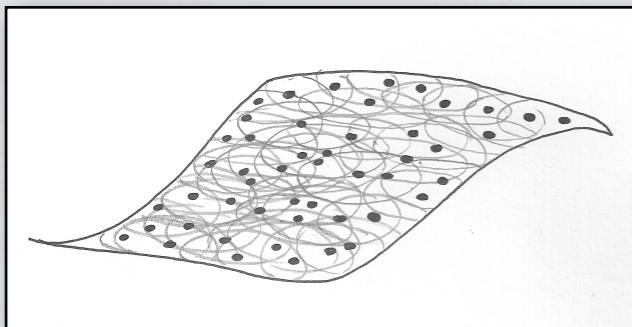
$$\mathcal{C}' = \{A \cap \text{Hull}_\alpha(C_v) \mid v \in V\}$$



$\text{Nrv } \mathcal{C}' \simeq A$
if $\alpha < \text{Reach}(A)$

Collapsing restricted Čech complex

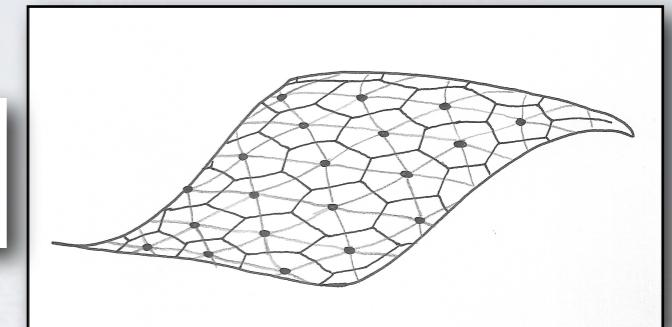
Theorem 3 Let $\mathcal{C} = \{C_v \mid v \in V\}$ an α -robust covering of A with $V \subset P$. Suppose there exists $f : V \rightarrow P$ injective such that $C_v \subset B^\circ(f(v), \alpha)$. If $\alpha < \text{Reach}(A)$, then there is a sequence of collapses from $\text{Cech}_A(P, \alpha)$ to $\text{Nrv } \mathcal{C}$.



$\text{Cech}_A(P, \alpha)$

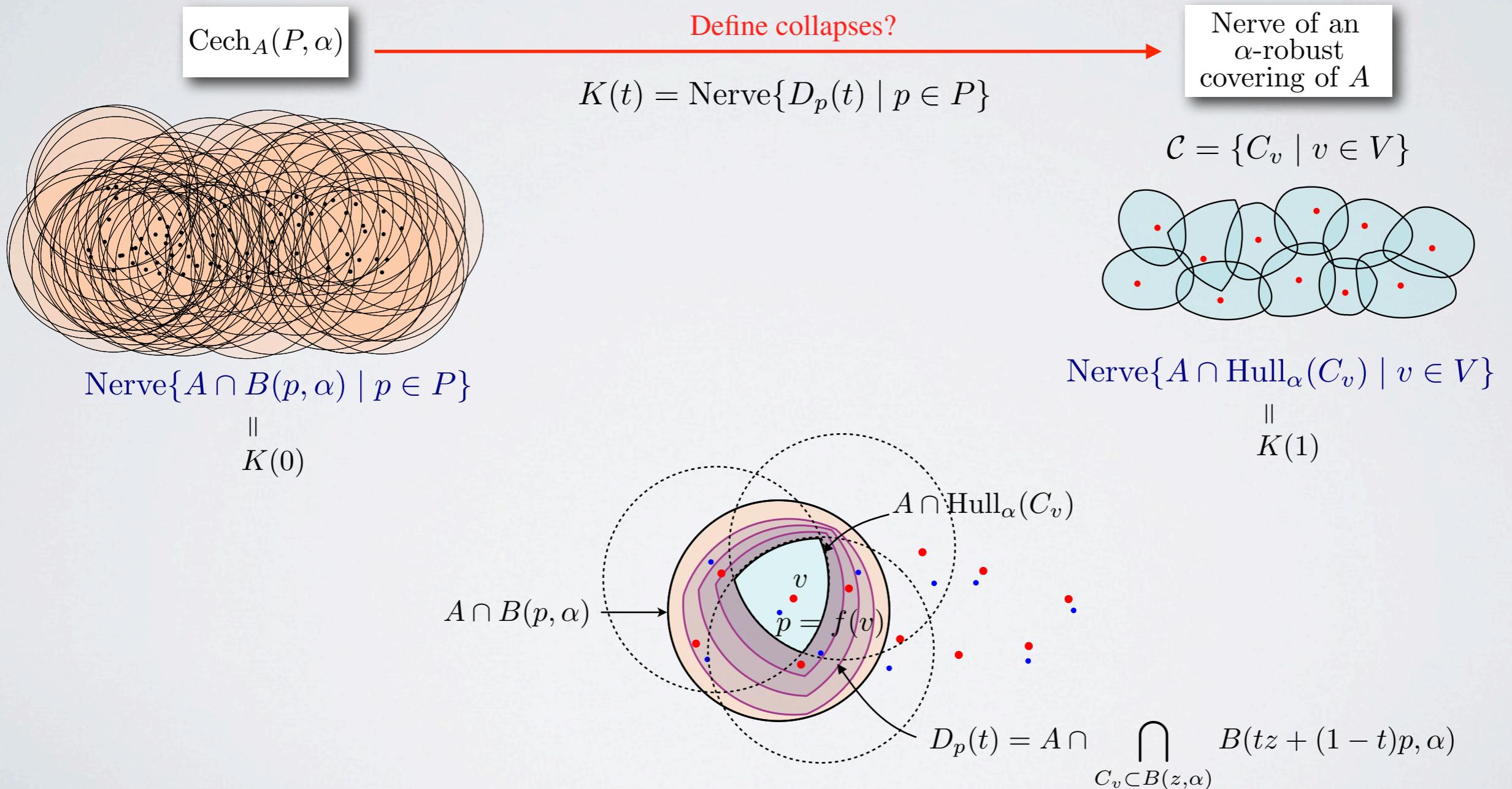
Define collapses?

Nerve of an
 α -robust
covering of A



Collapsing restricted Čech complex

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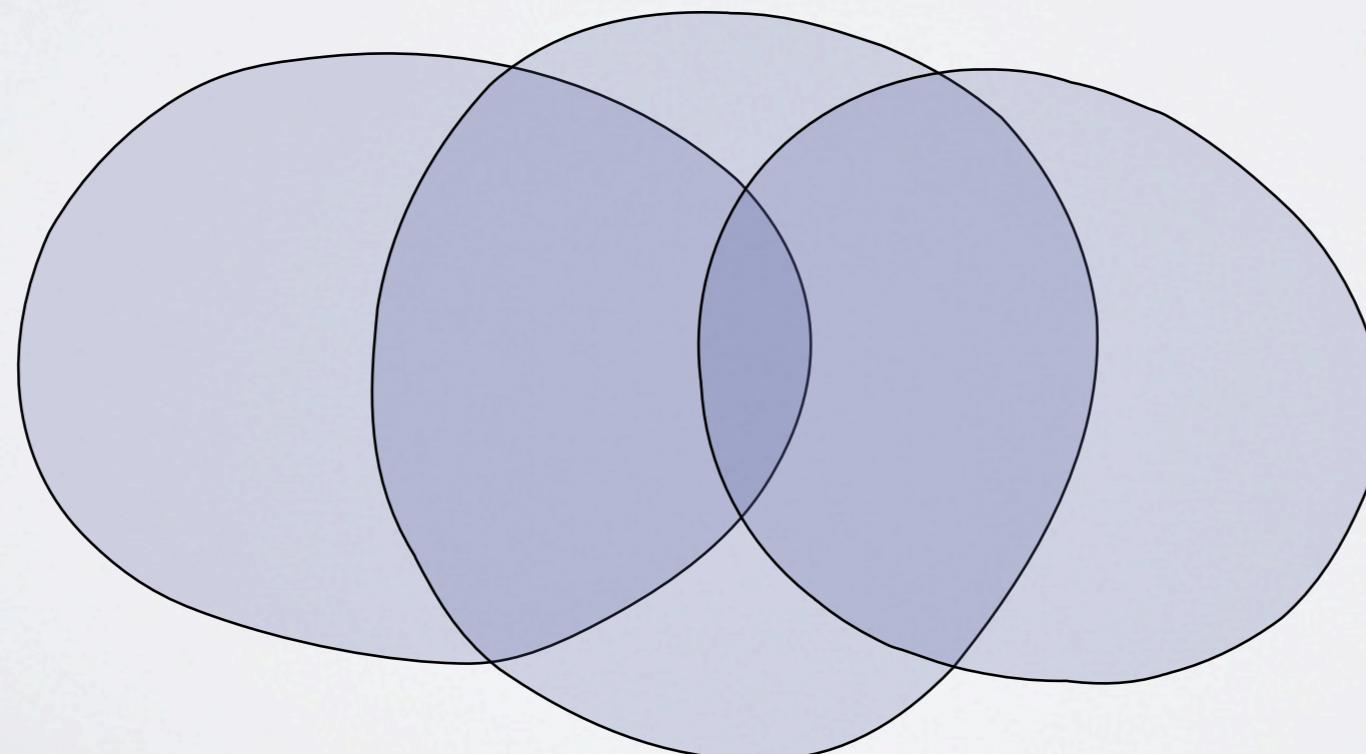


Evolving family of compact sets

Suppose $K(t) = \text{Nerve}\{D_p(t) \mid p \in P\}$ such that $\forall t_1 < t_2, \forall \sigma \subset P, \forall t$

- (a) $\bigcap_{p \in \sigma} D_p(t)$ empty or connected ;
- (b) before disappearing $\bigcap_{p \in \sigma} D_p(t)$ is reduced to a single point.
- (c) $D_p(t_2) \subset D_p(t_1)^\circ$;

Then, generically $K(t)$ undergoes collapses as t increases.

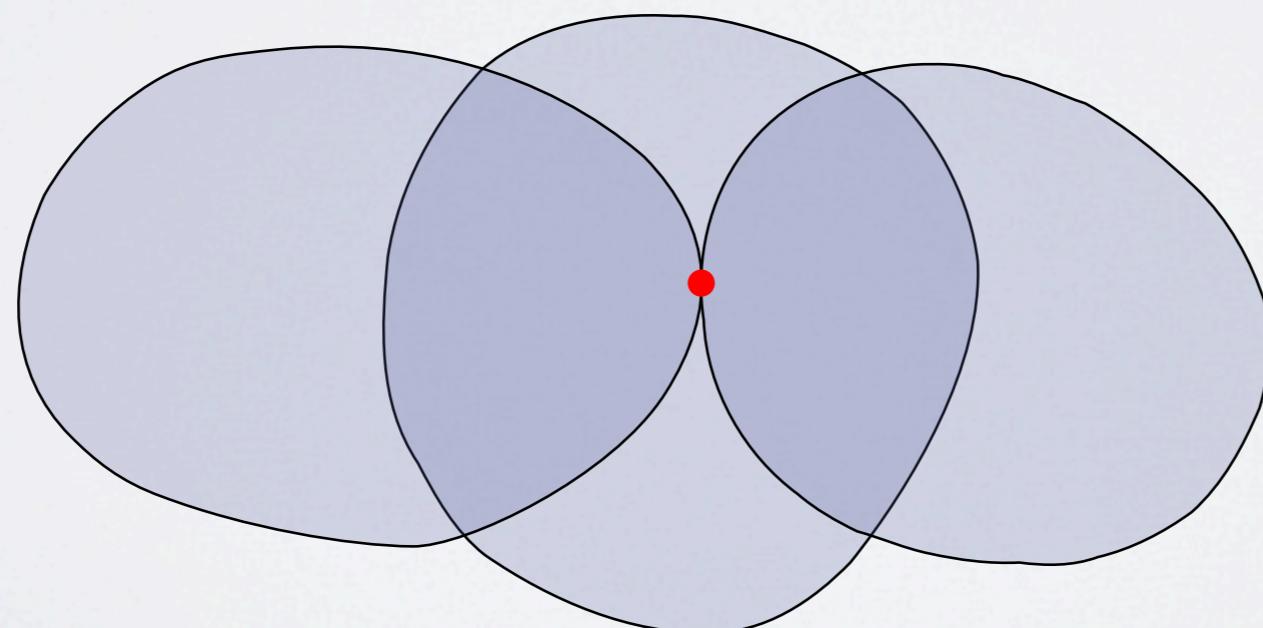


Evolving family of compact sets

Suppose $K(t) = \text{Nerve}\{D_p(t) \mid p \in P\}$ such that $\forall t_1 < t_2, \forall \sigma \subset P, \forall t$

- (a) $\bigcap_{p \in \sigma} D_p(t)$ empty or connected ;
- (b) before disappearing $\bigcap_{p \in \sigma} D_p(t)$ is reduced to a single point.
- (c) $D_p(t_2) \subset D_p(t_1)^\circ$;

Then, generically $K(t)$ undergoes collapses as t increases.

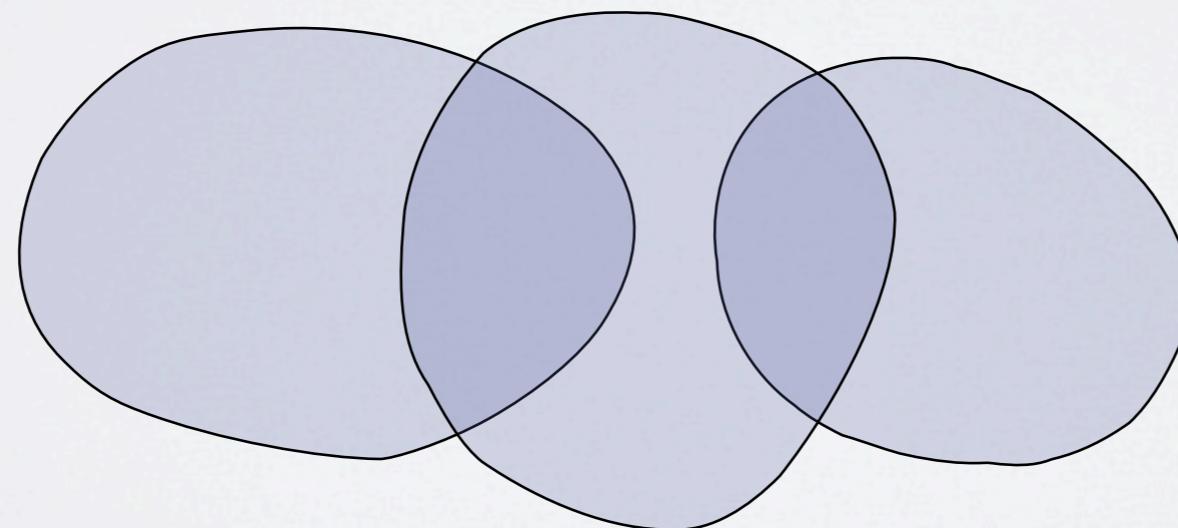


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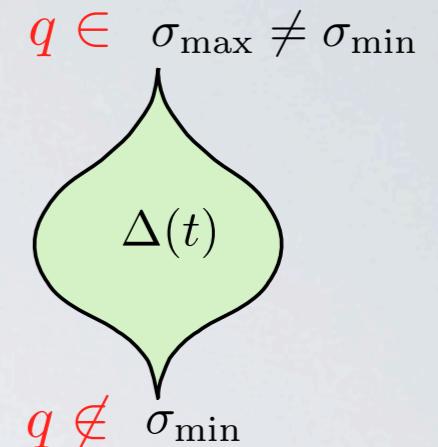
Then, generically $K(t)$ undergoes collapses as t increases.



Steps for proving that collapses

$$K(t) = \text{Nerve}\{D_p(t) \mid p \in P\}$$

$\Delta(t)$ = set of simplices that disappear at time t



Does the operation that removes $\Delta(t)$ from $K(t)$ a collapse?

(1) Generically, $\Delta(t)$ has a unique minimal element σ_{\min}

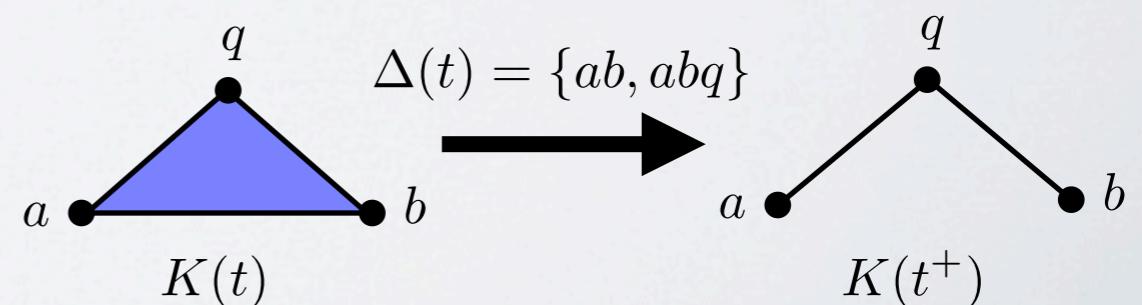
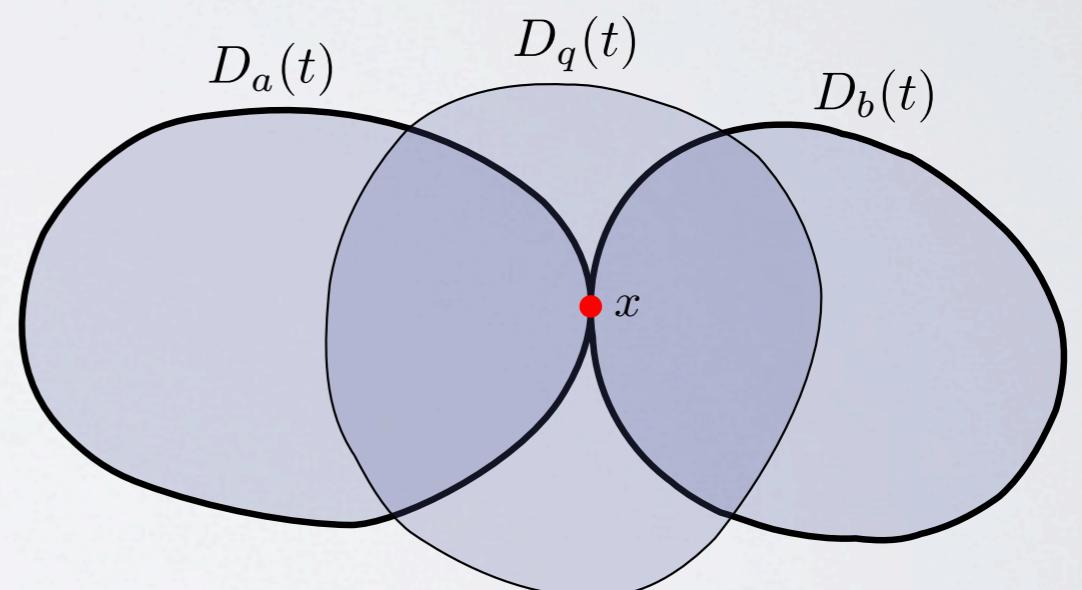
(2) $\bigcap_{p \in \sigma_{\min}} D_p(t) = \{x\}$

(3) $\sigma_{\max} = \{p \in P \mid x \in D_p(t)\}$

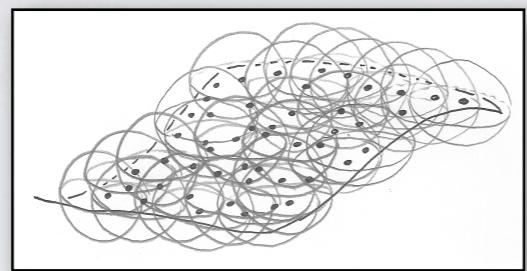
(4) $x \in \partial D_p(t), \quad \forall p \in \sigma_{\min}$

(5) $\exists q \in P$ such that $x \in D_q(t)^\circ$

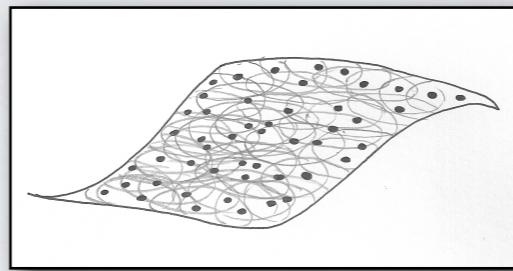
(6) $\sigma_{\min} \neq \sigma_{\max} \implies$ removing $\Delta(t)$ is a collapse



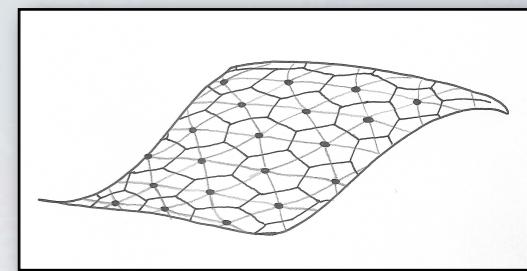
Summary



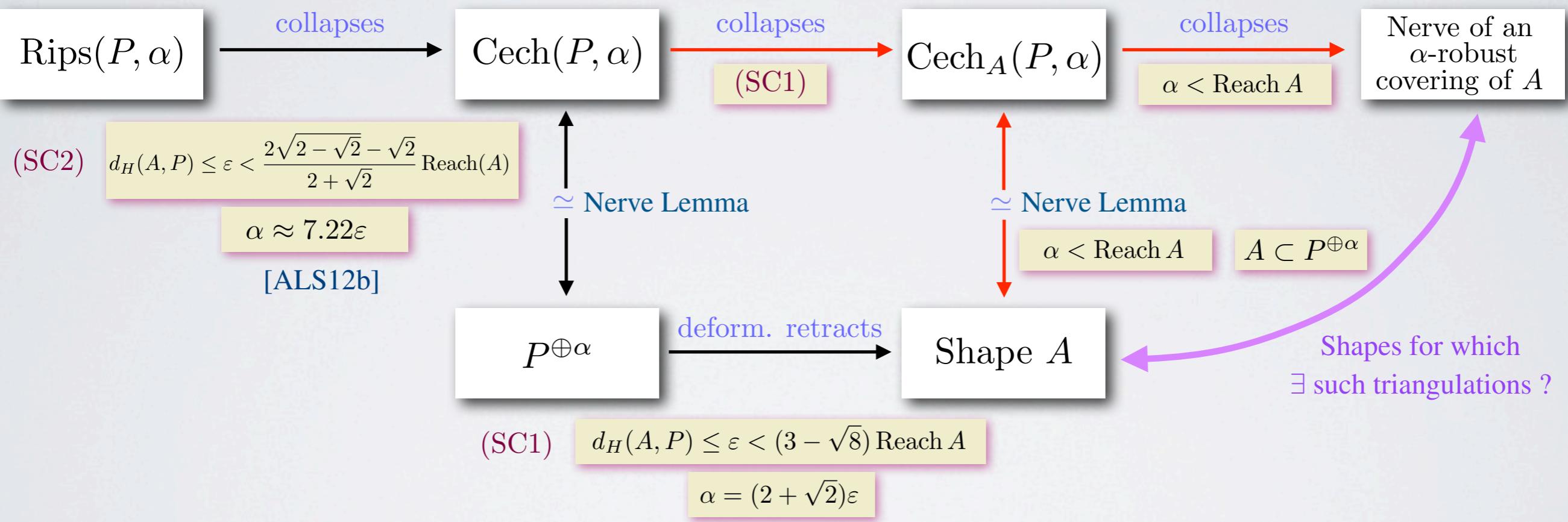
$\text{Nerve}\{B(p, \alpha) \mid p \in P\}$



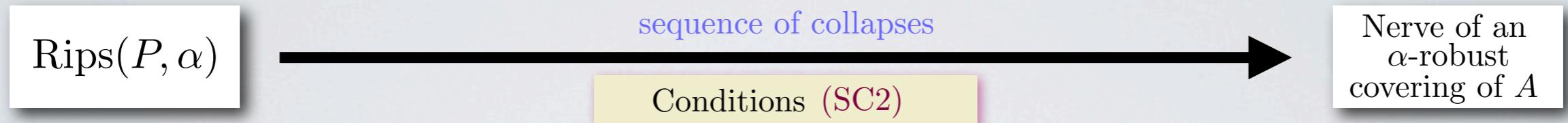
$\text{Nerve}\{A \cap B(p, \alpha) \mid p \in P\}$



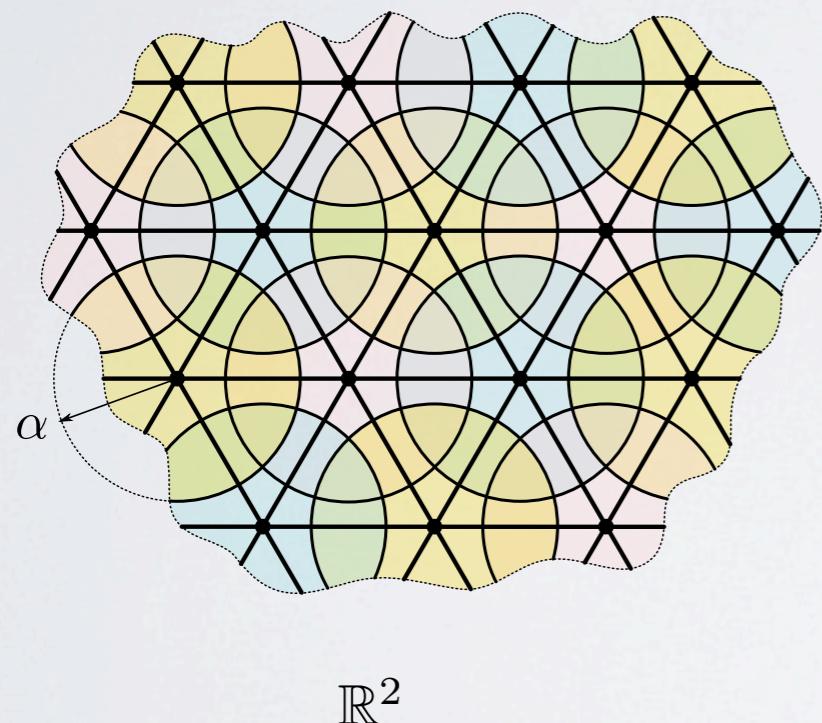
$\text{Nerve}\{A \cap \text{Hull}_\alpha(C_v) \mid v \in V\}$



α -Nice triangulations



A triangulation of A is α -nice if nerve of an α -robust covering of A



$T =$ triangulation of \mathbb{R}^2 with equilateral triangles

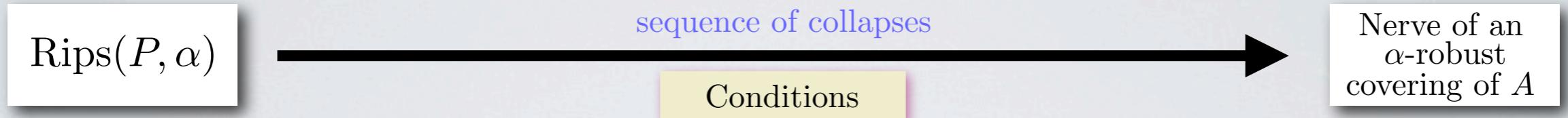
$\mathcal{C} = \{B(v, \alpha) \mid v \in \text{Vertices}(T)\} \setminus B(v, \alpha) \subset \text{St}_T(v)$

Then, $T = \text{Nerve}(\mathcal{C})$

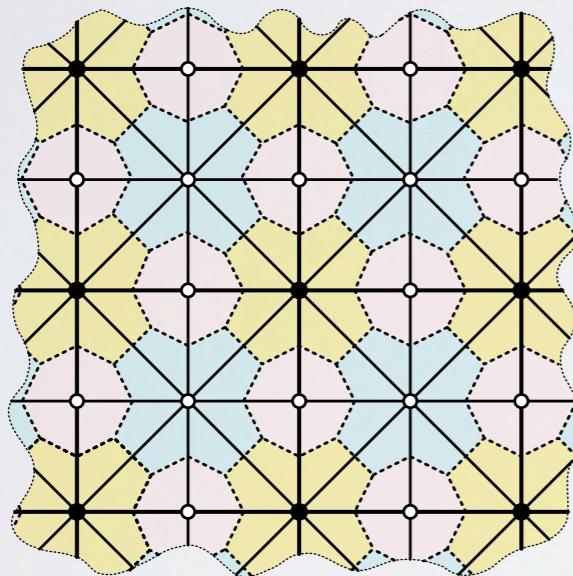
\mathcal{C} : α -robust

T : α -nice

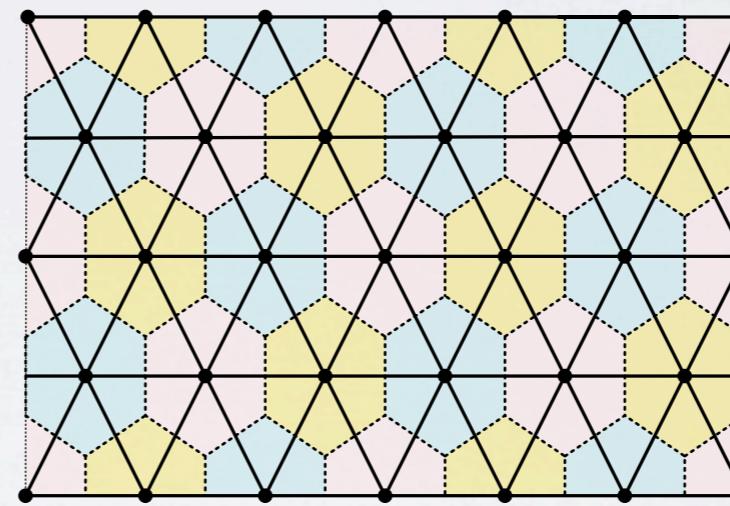
Nicely triangulable spaces



A space is “nicely triangulable” if it has an α -nice triangulation for all α



\mathbb{R}^m



The flat torus $\mathbb{T}^2 \subset \mathbb{R}^4$

Can we find other spaces that are “nicely triangulable”?

Can we turn all this into a practical algorithm?

