

**EFFICIENT DATA STRUCTURE
FOR
REPRESENTING & SIMPLIFYING
SIMPLICIAL COMPLEXES
IN
HIGH DIMENSIONS**

D. ATTALI

CNRS, GIPSA-LAB
GRENOBLE

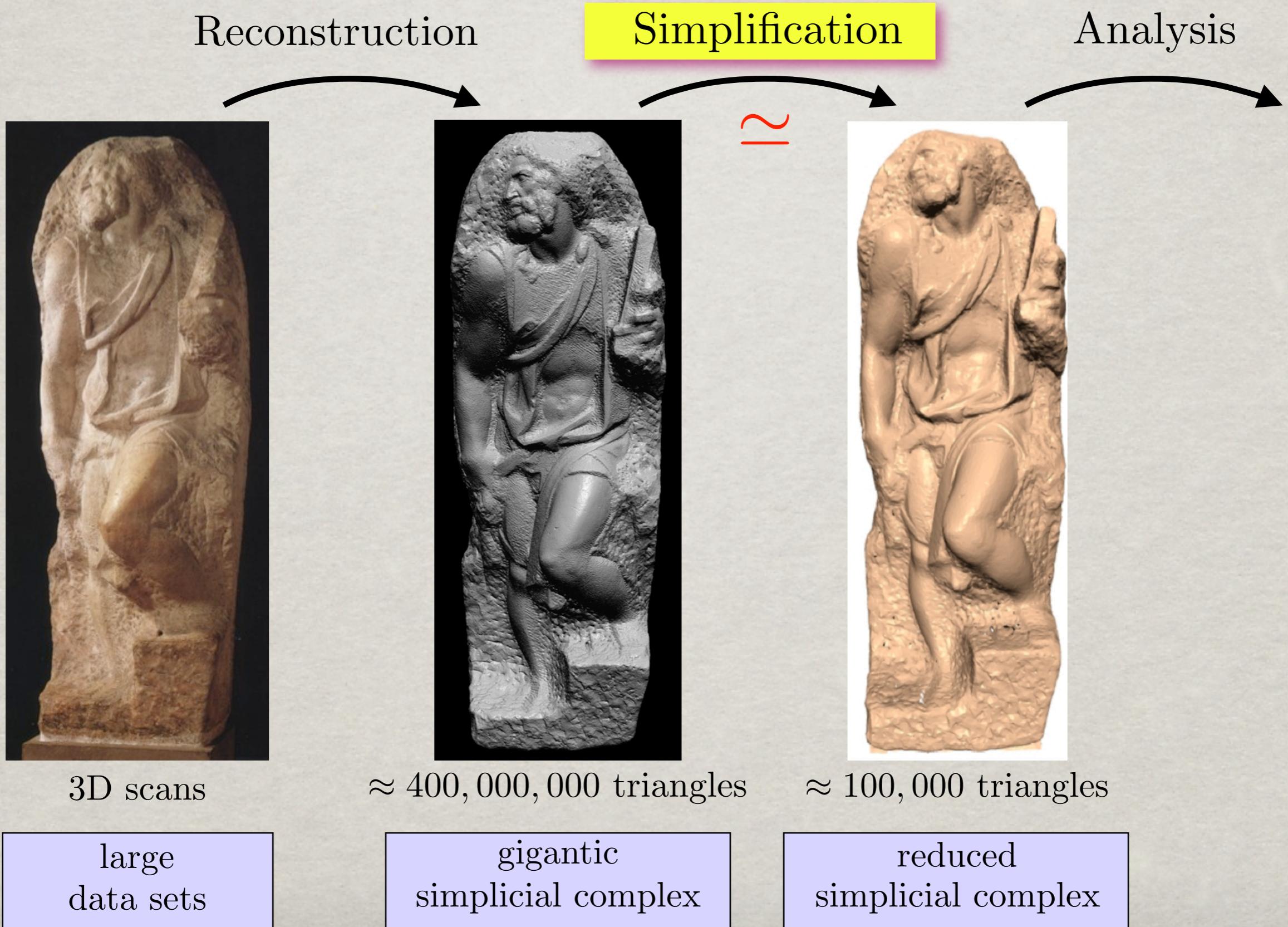
A. LIEUTIER

DASSAULT SYSTÈME
AIX-EN-PROVENCE

D. SALINAS

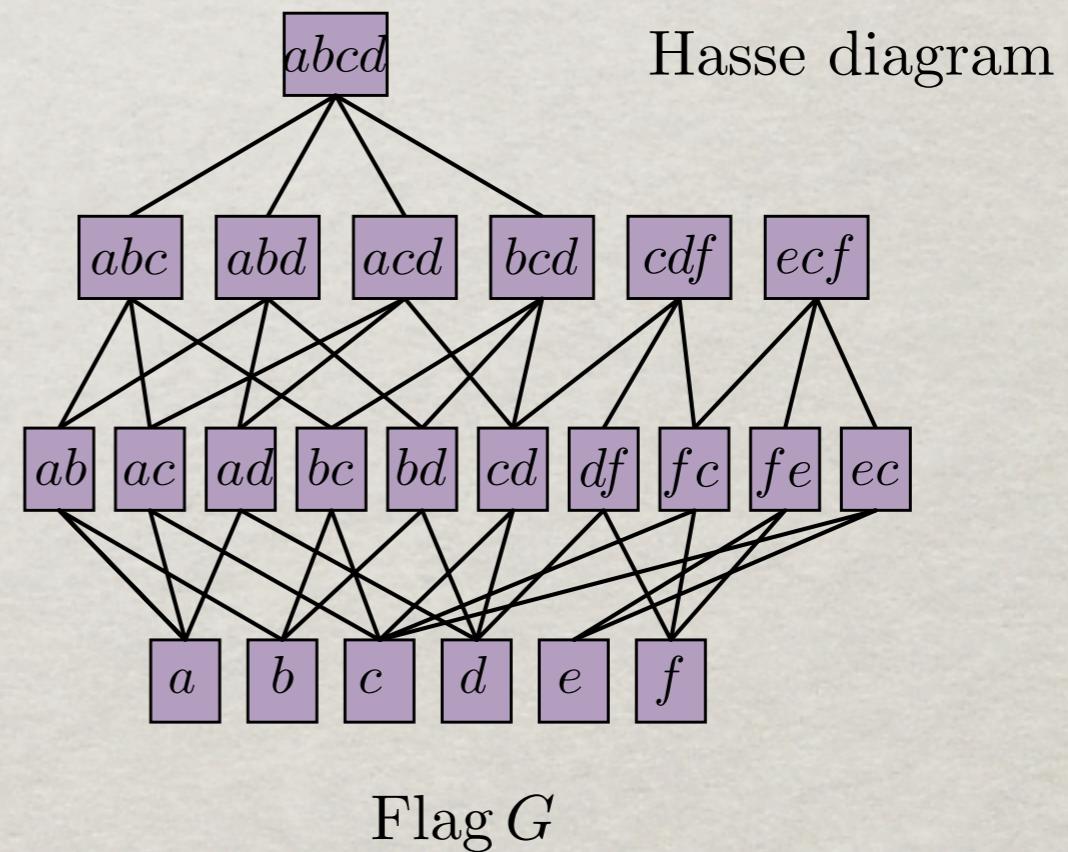
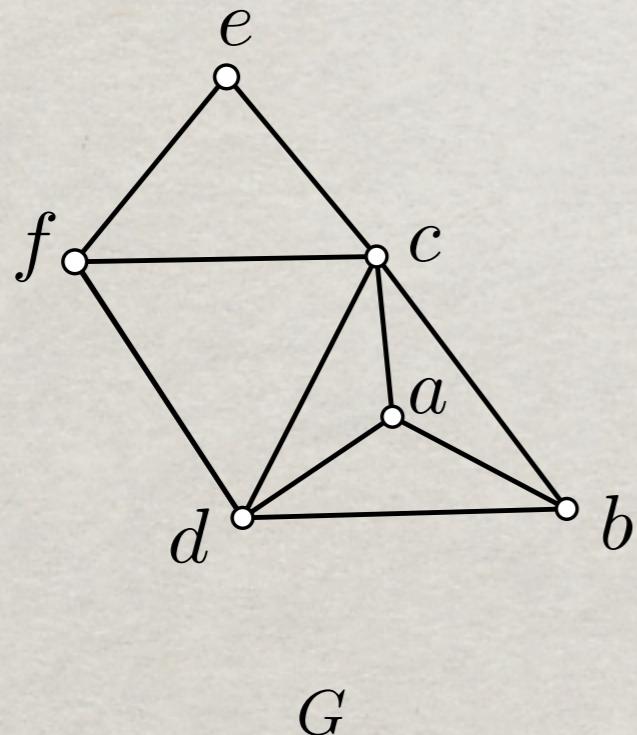
GIPSA-LAB,
GRENOBLE

MOTIVATION



FLAG COMPLEXES

- * Flag G = largest simplicial complex whose 1-skeleton is G .
- * $\{v_0, \dots, v_k\} \in \text{Flag } G \iff v_i v_j \in G \ \forall i, j$

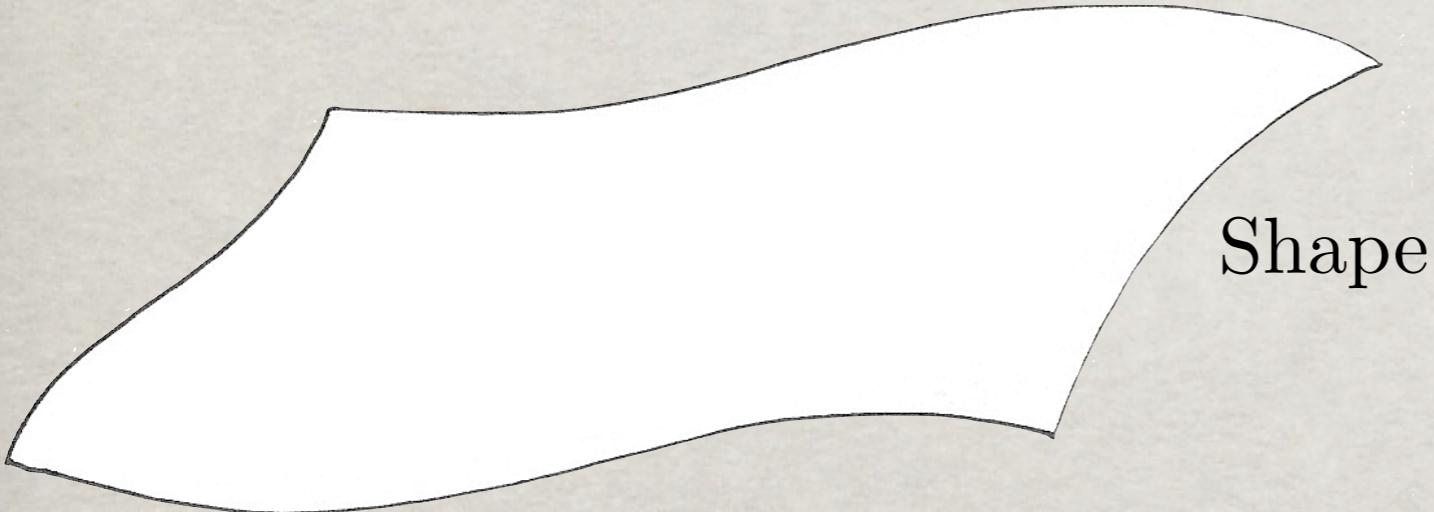


Flag complexes have a compressed form of storage

SHAPE RECONSTRUCTION

INPUT

Point cloud P



OUTPUT

Flag $G_\alpha(P) = \text{Rips complex}$

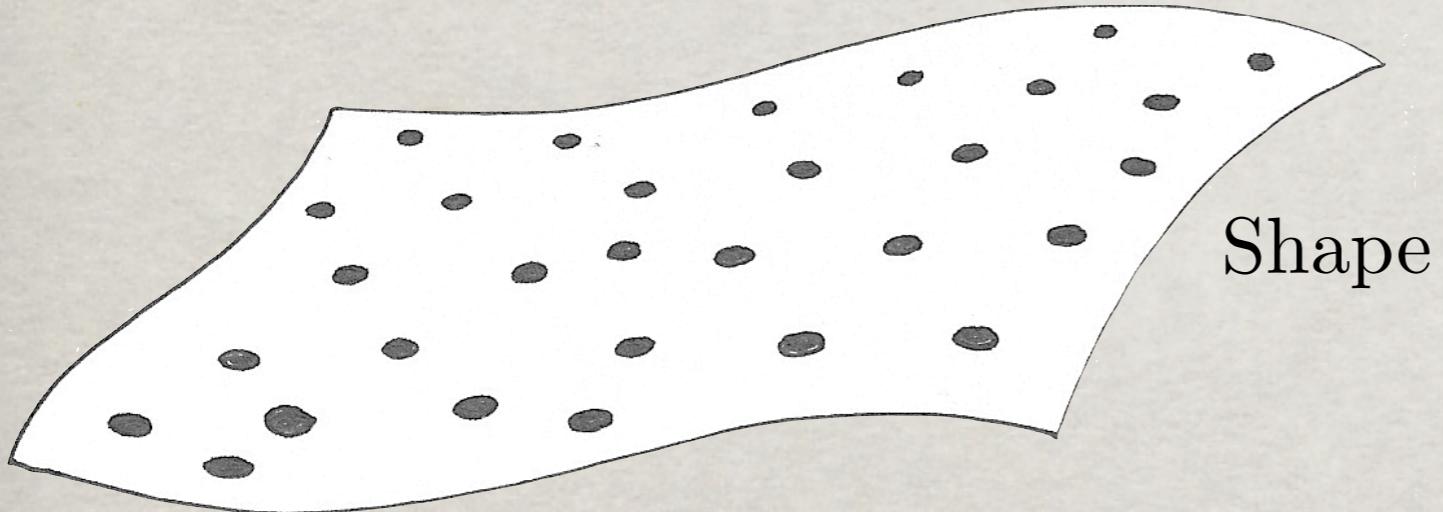
$G_\alpha(P) = \text{proximity graph}$

$pq \in G_\alpha(P) \iff d(p, q) \leq 2\alpha$

SHAPE RECONSTRUCTION

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Point cloud P



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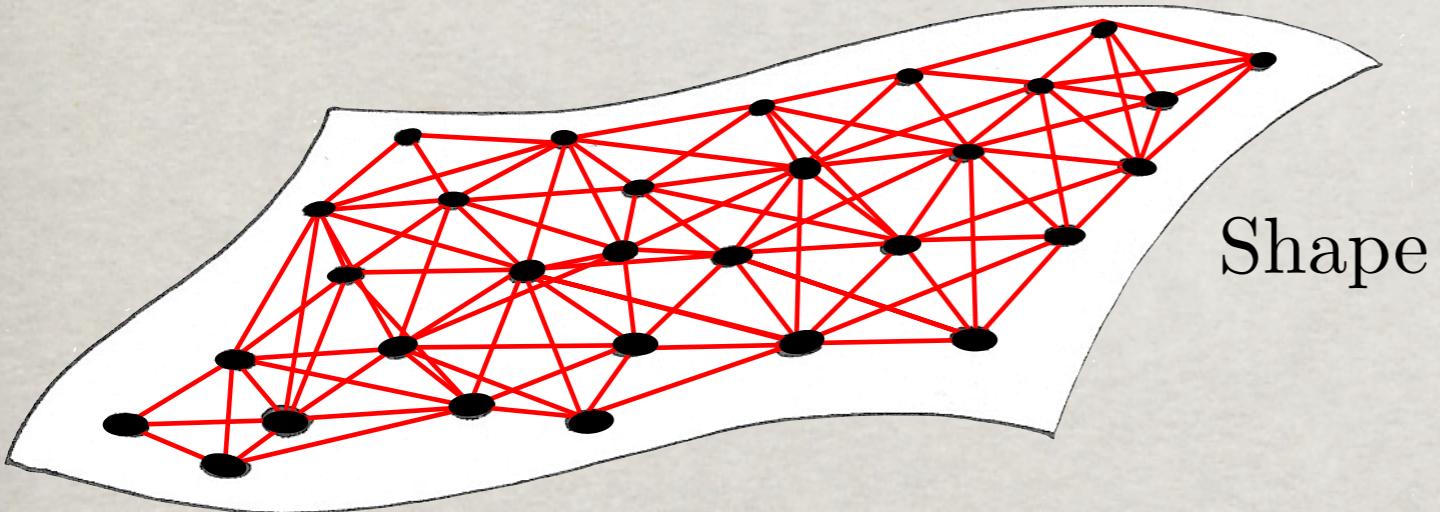
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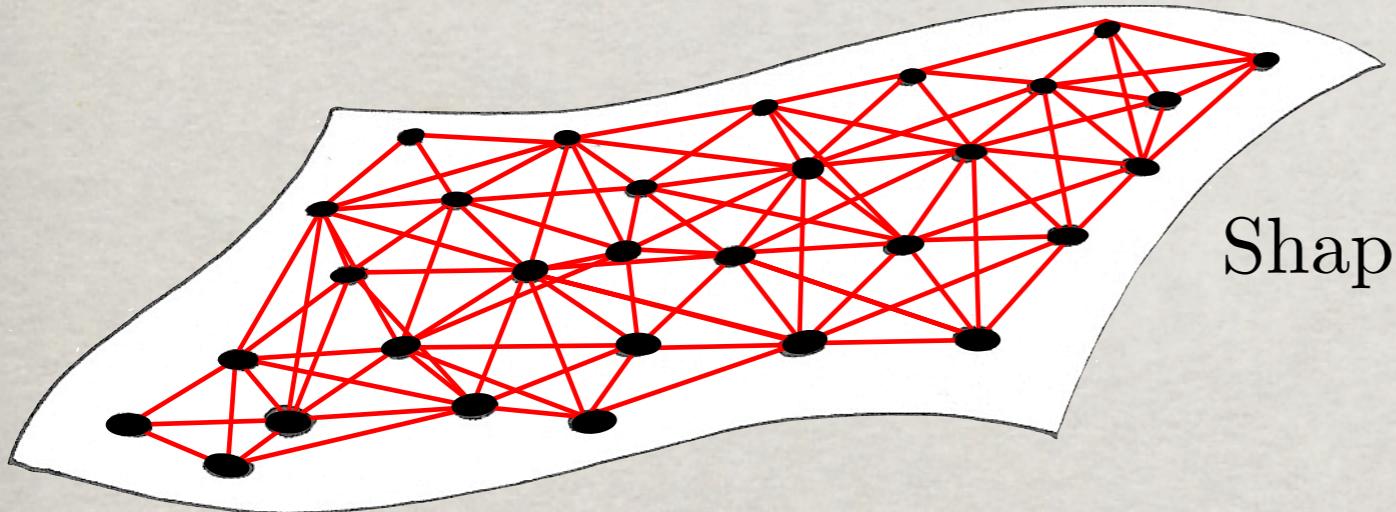
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SHAPE RECONSTRUCTION

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for sampling conditions stated in

[AL 2010] when $d = d_\infty$

[ALS 2011] when $d = d_2$

$G_\alpha(P) = \text{proximity graph}$

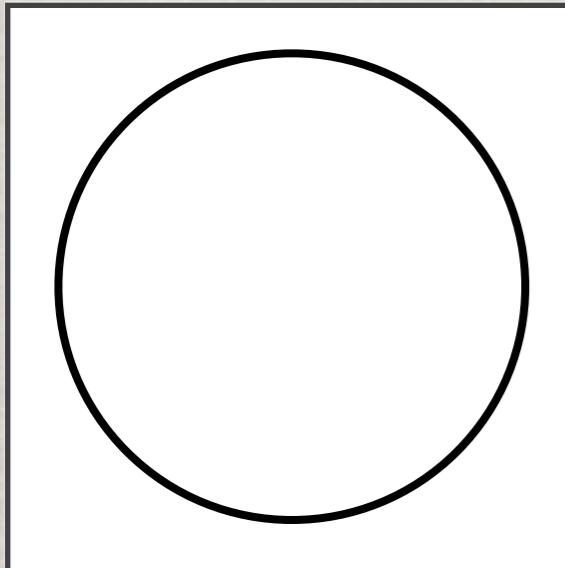
$$pq \in G_\alpha(P) \iff d(p, q) \leq 2\alpha$$

OVERVIEW

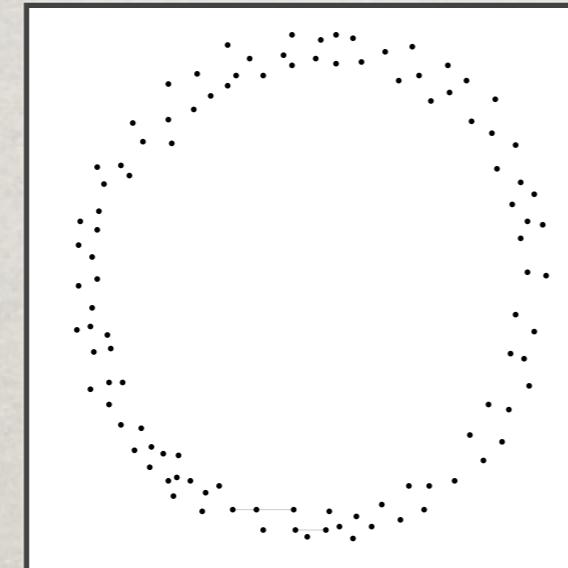
Part I:
Guarantees

Part II:
Simplification

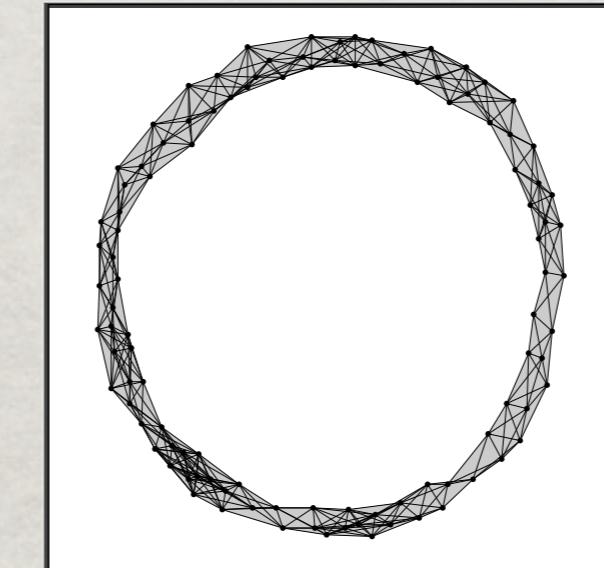
under some “good” sampling conditions



Shape



Point cloud



Flag complex

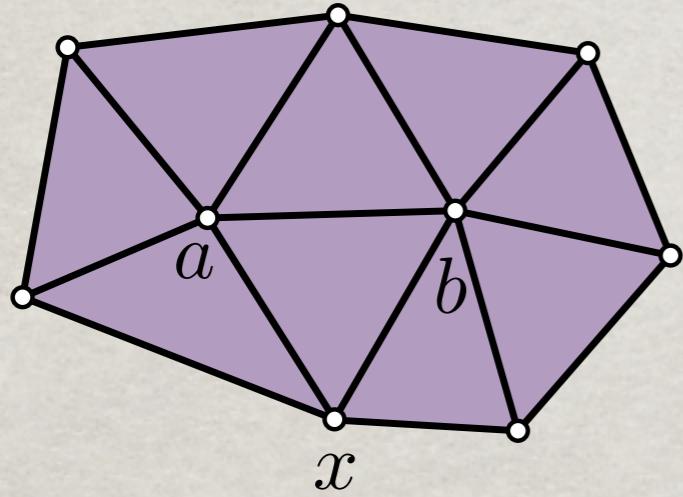
can be high dimensional !



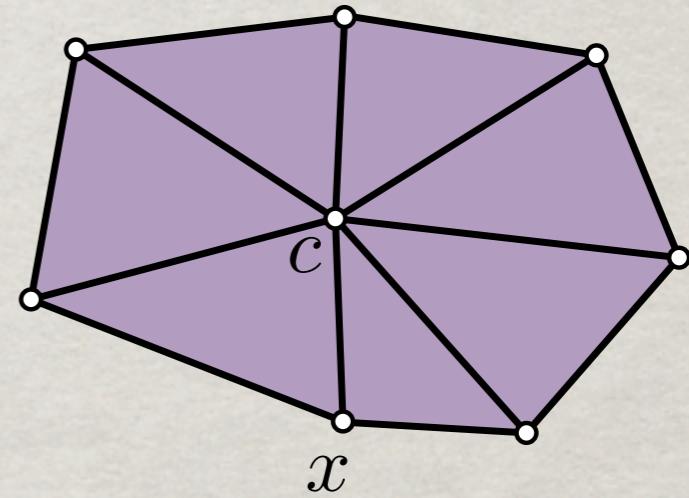
EDGE CONTRACTION

operation that identifies vertices a and b to vertex c

$$K \xrightarrow{ab \mapsto c} K'$$



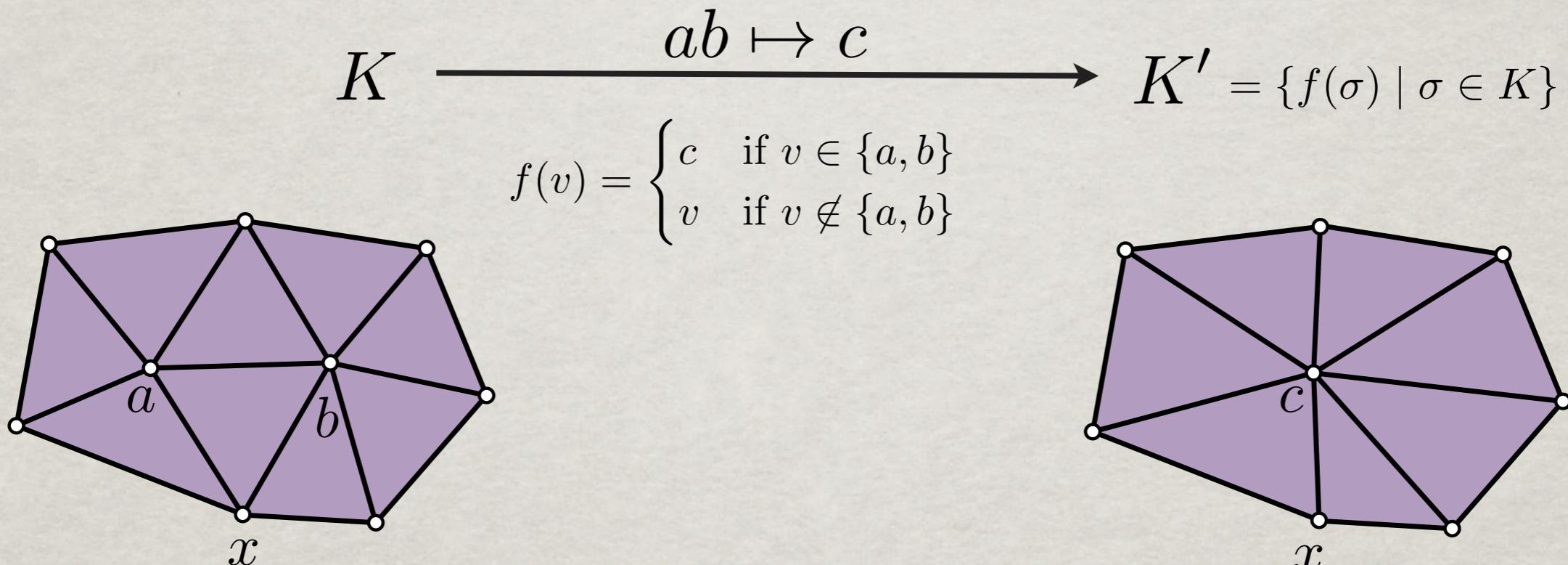
$$\{\cdots, a, b, x, ab, ax, bx, abx, \cdots\}$$



$$\{\cdots, c, x, cx, \cdots\}$$

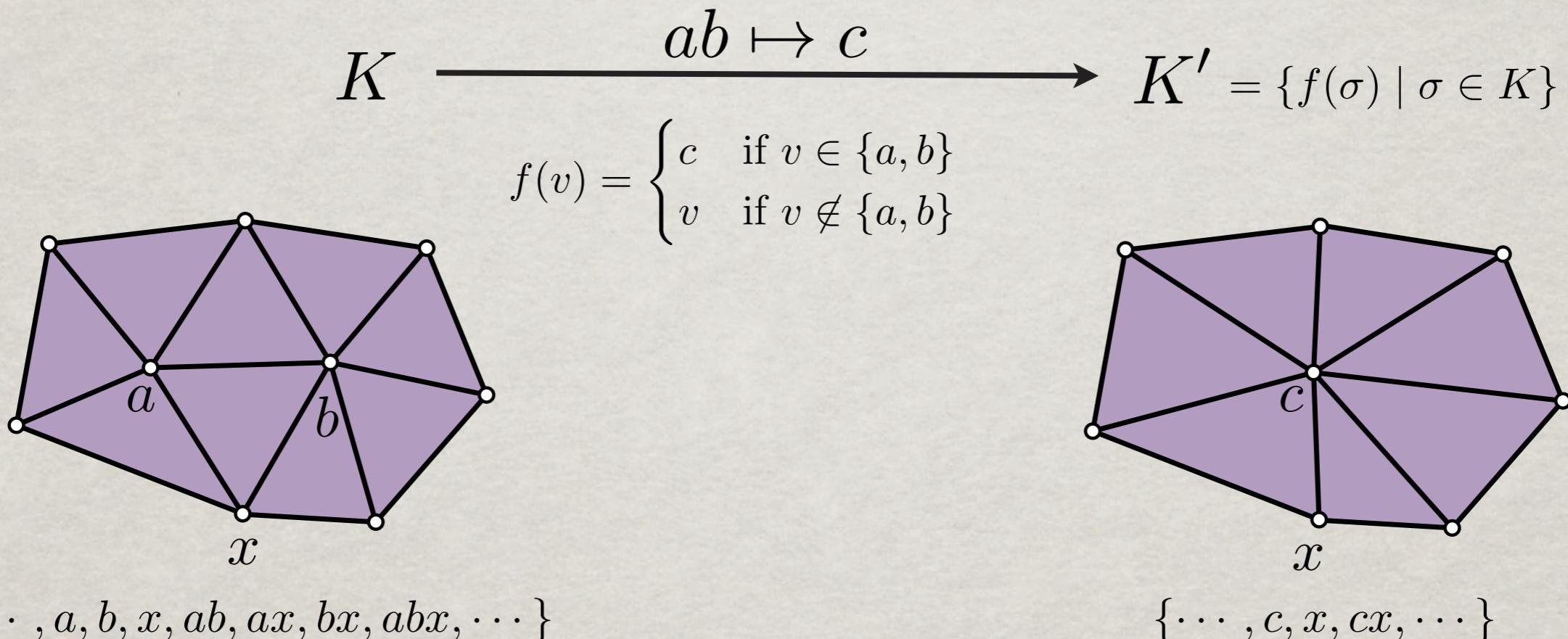
EDGE CONTRACTION

operation that identifies vertices a and b to vertex c



EDGE CONTRACTION

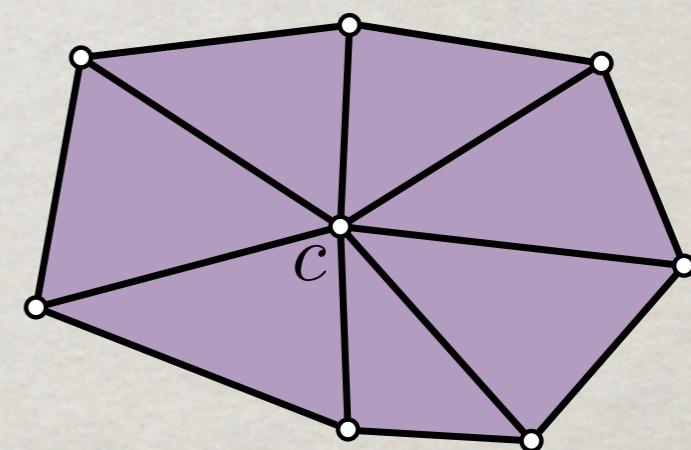
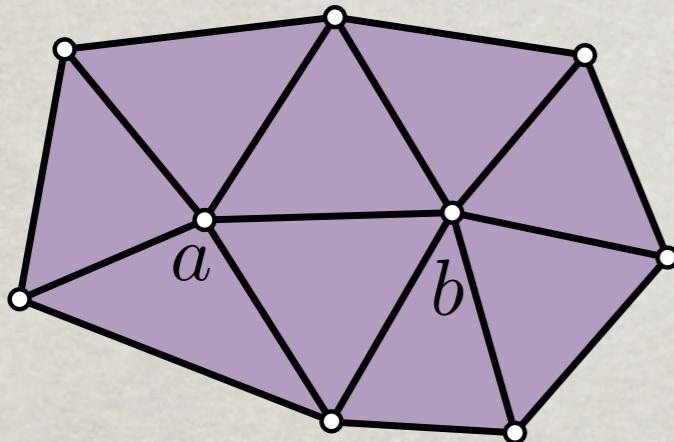
operation that identifies vertices a and b to vertex c



- * What if the result is not a flag complex?
- * How to preserve homotopy type?

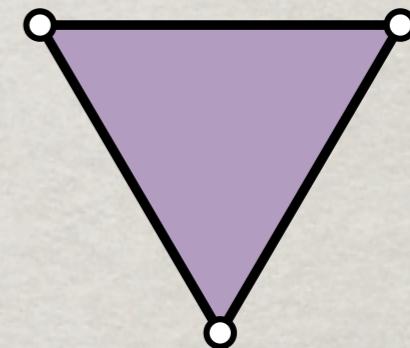
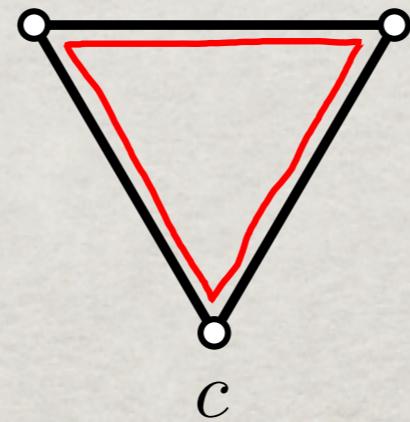
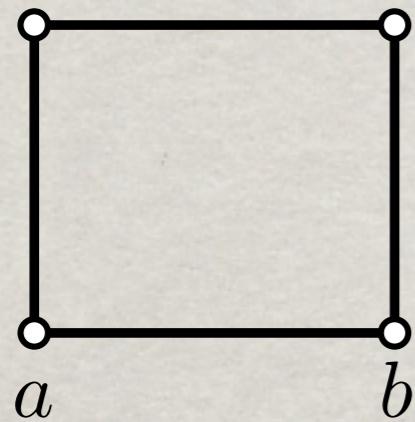
EDGE CONTRACTION

$$K = \text{Flag } K^{(1)} \xrightarrow{ab \mapsto c} K' = \text{Flag } K'^{(1)}$$



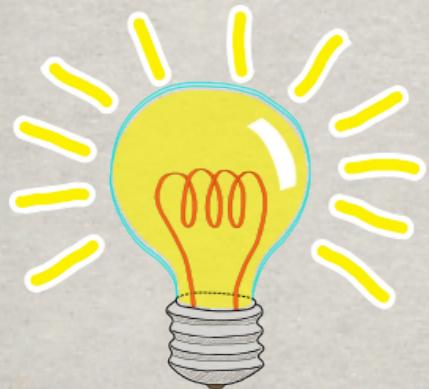
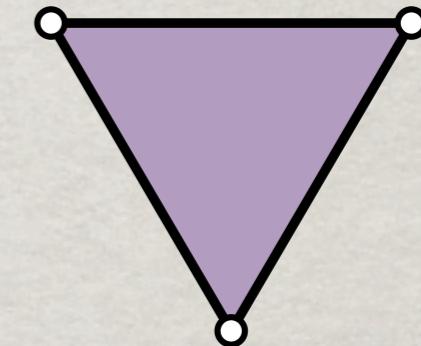
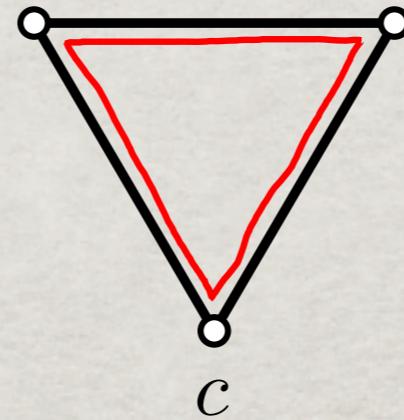
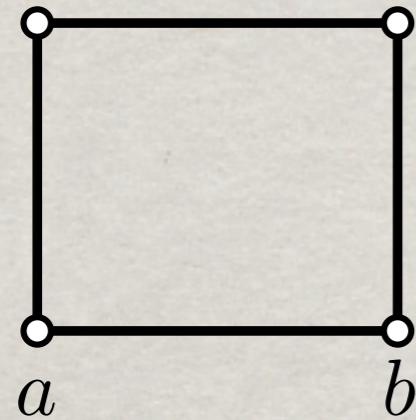
EDGE CONTRACTION

$$K = \text{Flag } K^{(1)} \xrightarrow{ab \mapsto c} K' \neq \text{Flag } K'^{(1)}$$



EDGE CONTRACTION

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Encode a simplicial complex K by storing the pair:

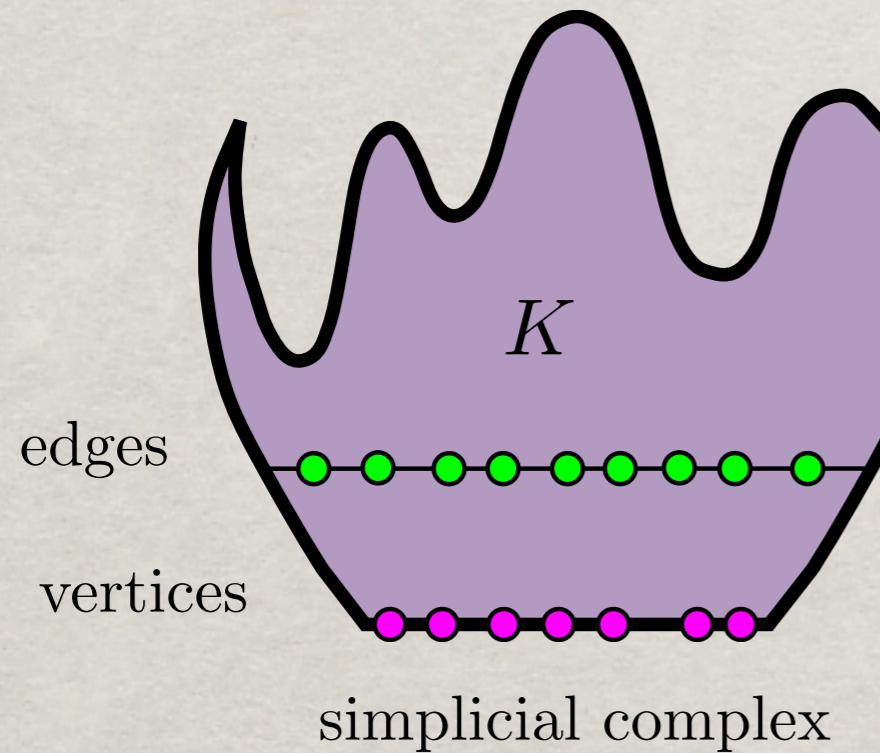
($K^{(1)}$, Blockers(K))

vertices and edges

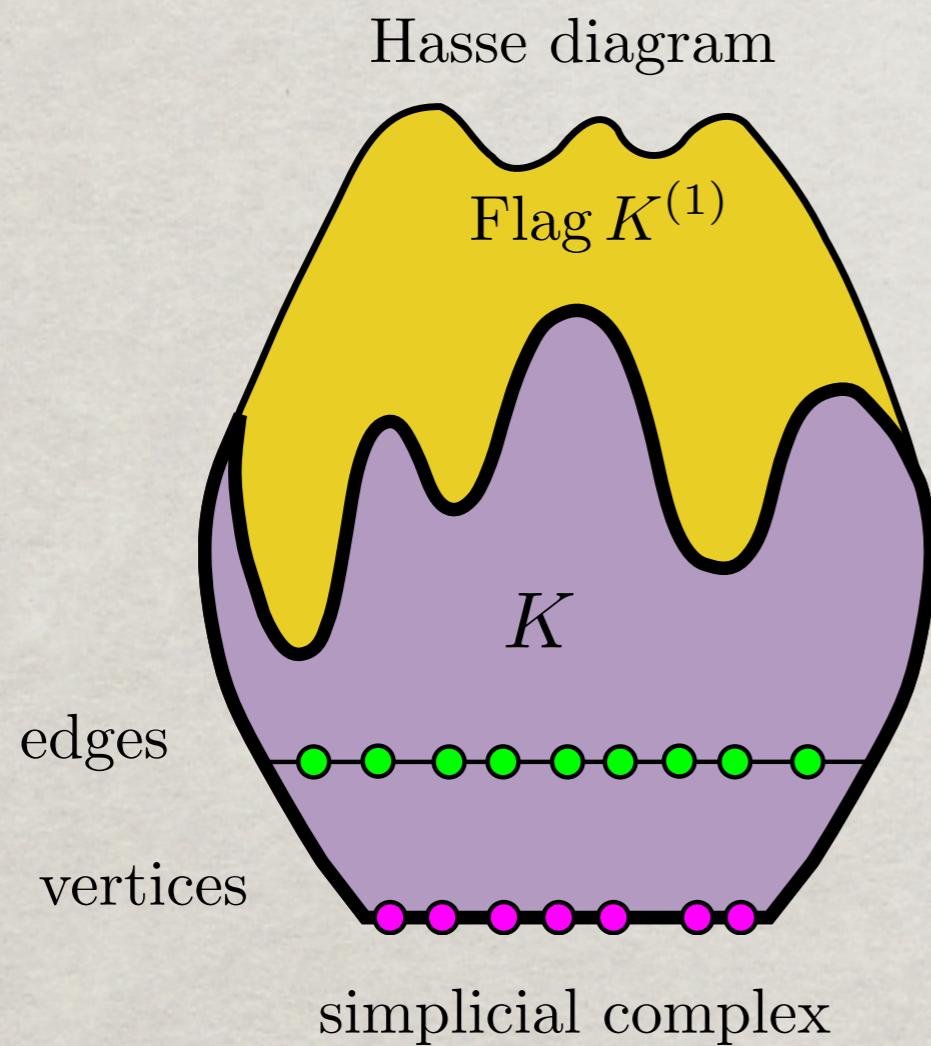
indicates how much
 K differs from $\text{Flag } K^{(1)}$

DATA STRUCTURE FOR SIMPLICIAL COMPLEXES

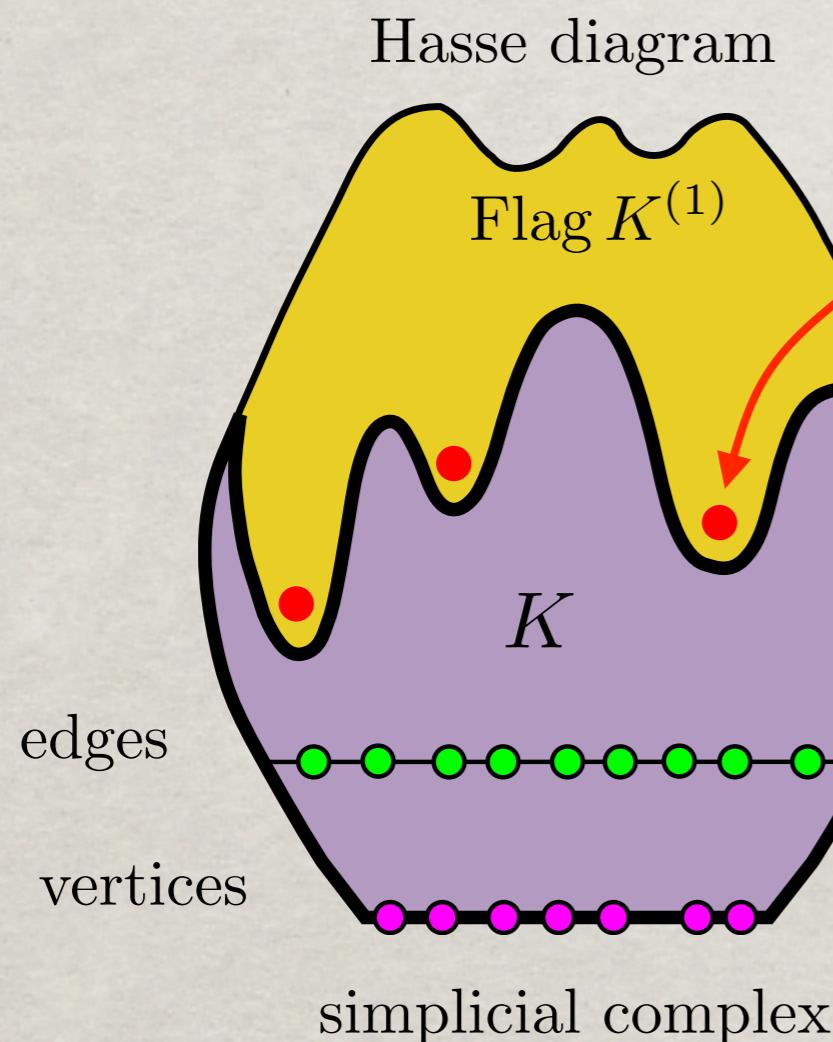
Hasse diagram



DATA STRUCTURE FOR SIMPLICIAL COMPLEXES

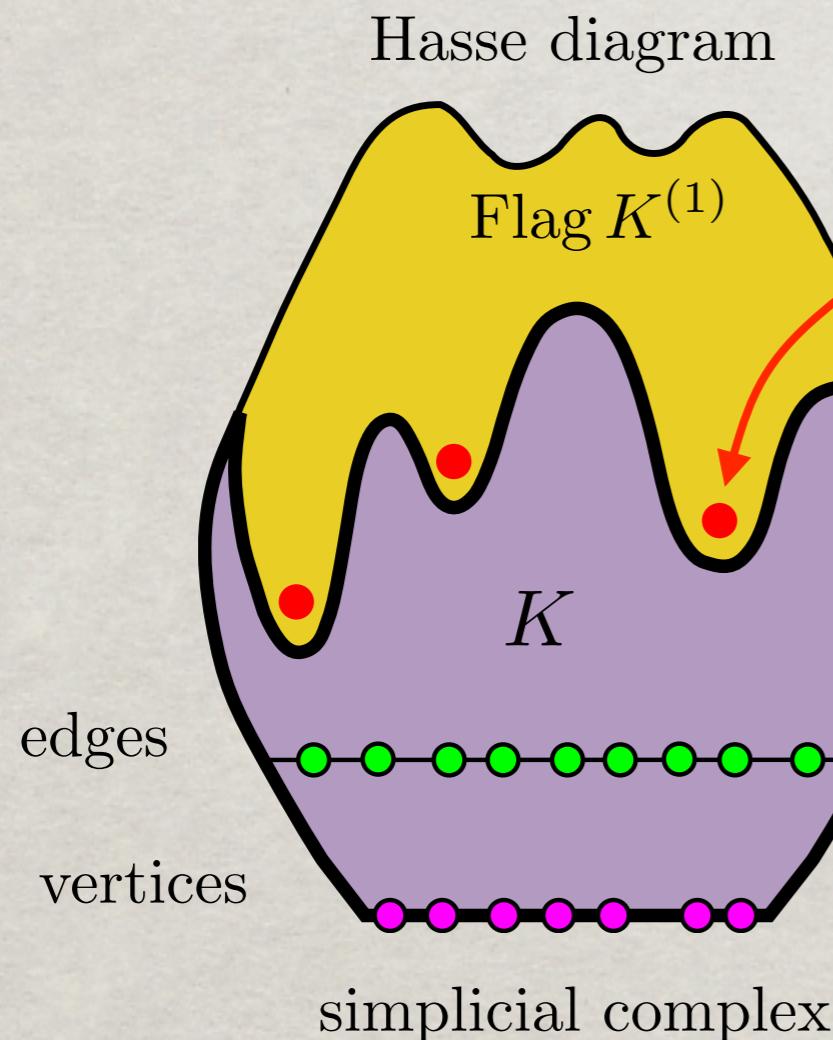


DATA STRUCTURE FOR SIMPLICIAL COMPLEXES



Blockers of K are inclusion-minimal simplices of $\text{Flag } K^{(1)} \setminus K$

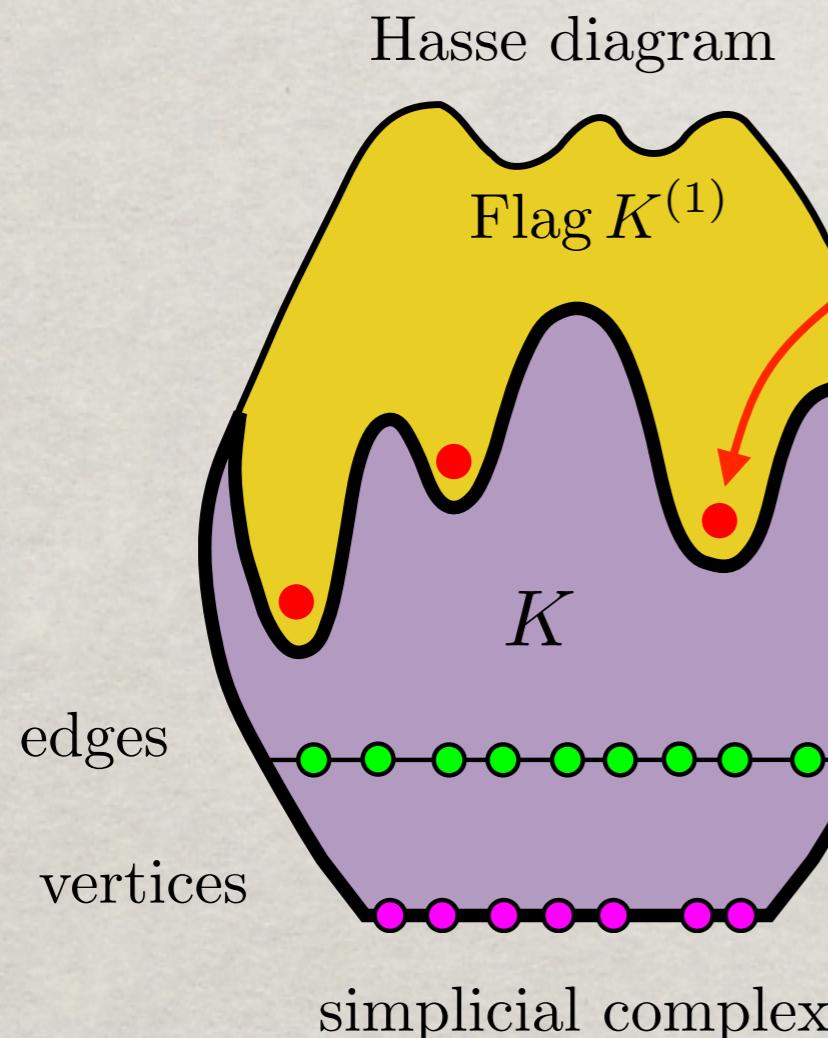
DATA STRUCTURE FOR SIMPLICIAL COMPLEXES



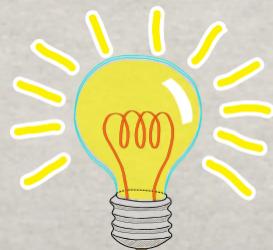
Blockers of K are inclusion-minimal simplices of $\text{Flag } K^{(1)} \setminus K$

σ blocker of K
 \iff
 $\dim \sigma \geq 2$
 $\sigma \notin K$
 $\forall \tau \subsetneq \sigma, \tau \in K$

DATA STRUCTURE FOR SIMPLICIAL COMPLEXES

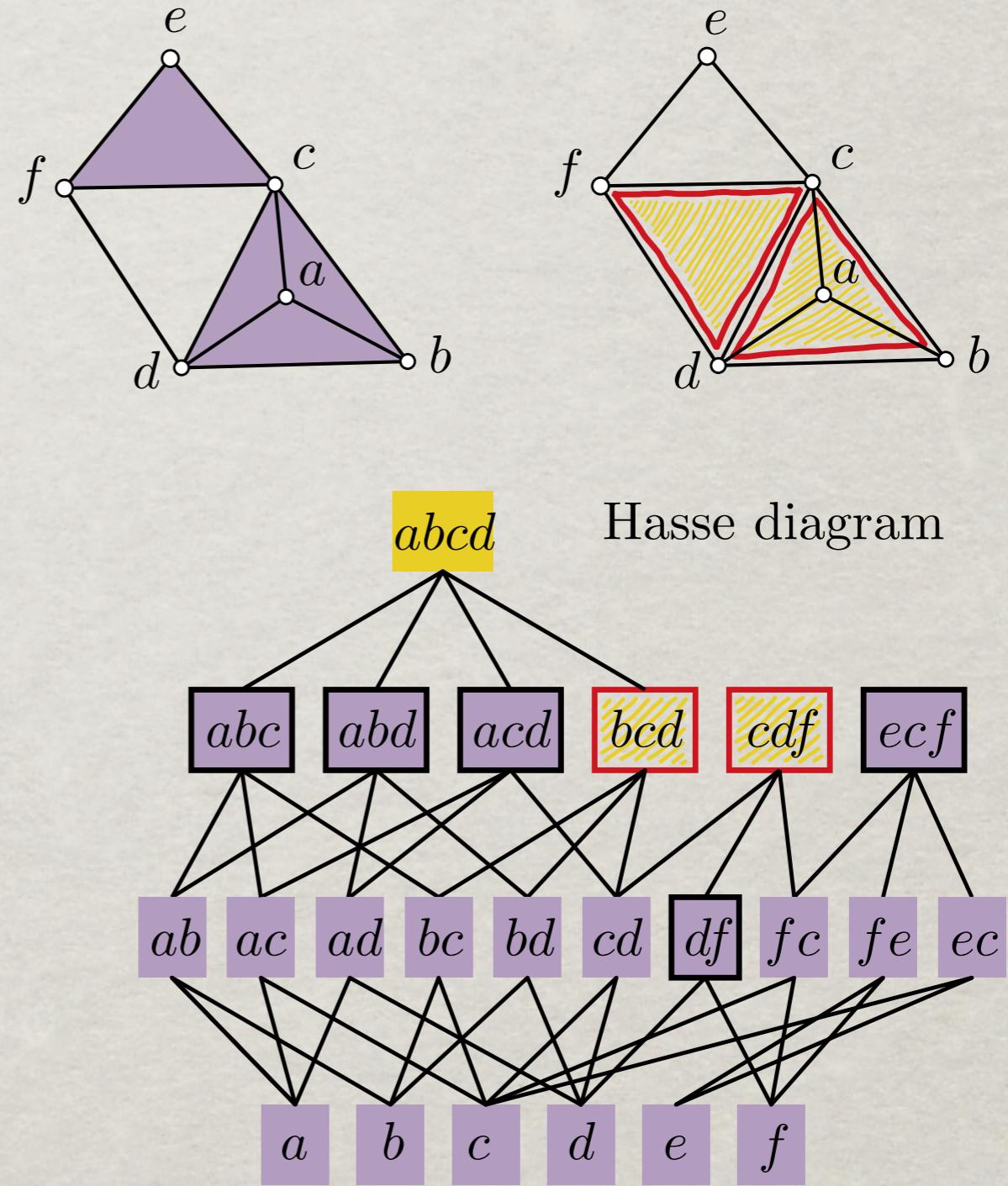
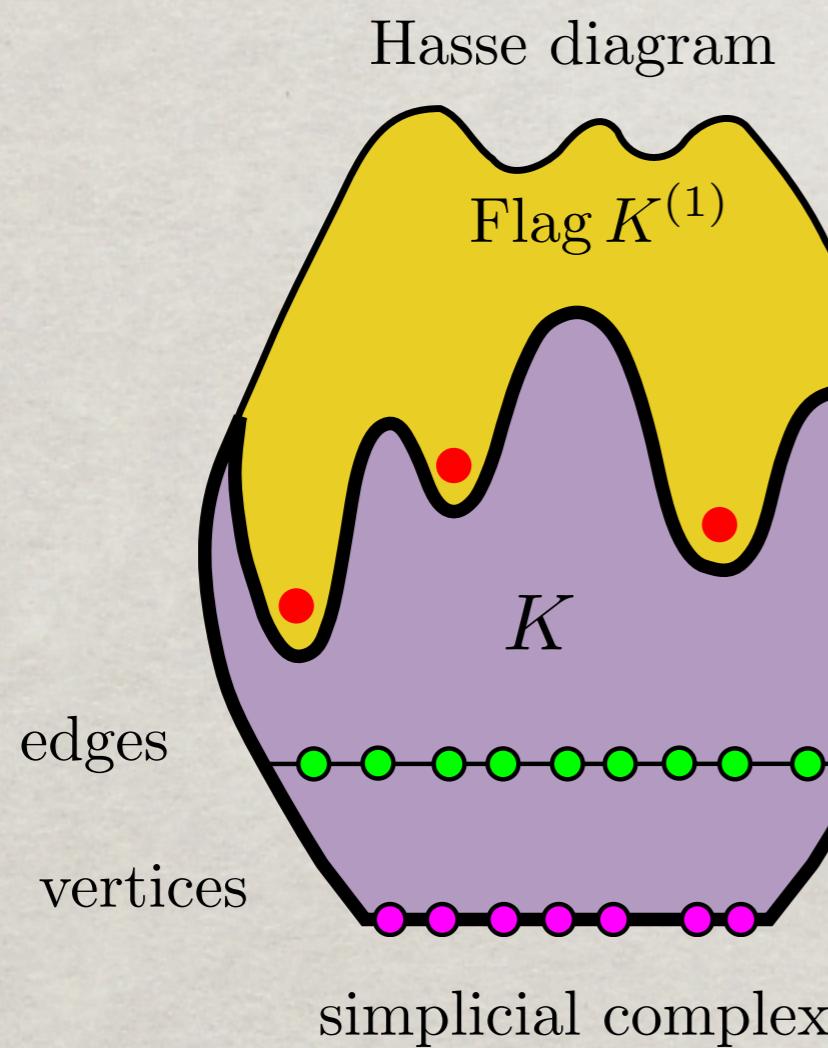


Blockers of K are inclusion-minimal simplices of $\text{Flag } K^{(1)} \setminus K$

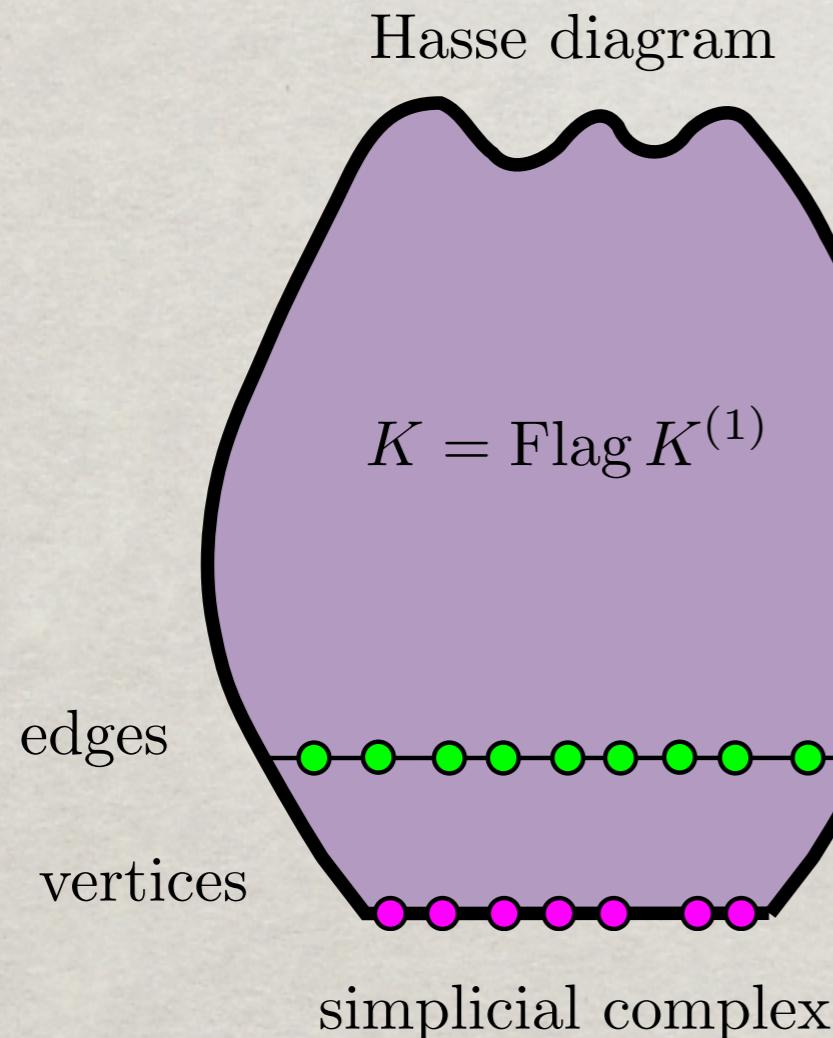
$$\begin{aligned} & \sigma \text{ blocker of } K \\ \iff & \\ & \dim \sigma \geq 2 \\ & \sigma \notin K \\ & \forall \tau \subsetneq \sigma, \tau \in K \end{aligned}$$


Encode a simplicial complex K by storing the pair:
 $(K^{(1)}, \text{Blockers}(K))$

A SMALL EXAMPLE

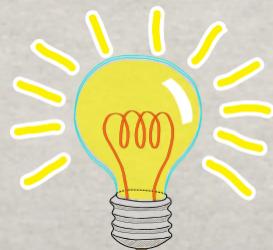


DATA STRUCTURE FOR SIMPLICIAL COMPLEXES



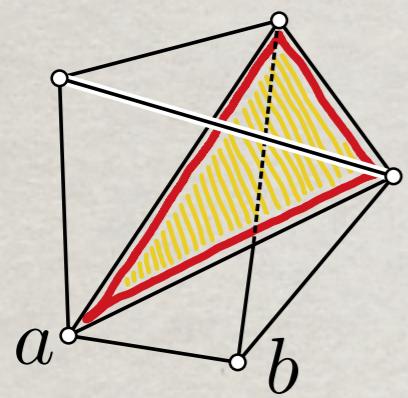
Blockers of K are inclusion-minimal simplices of $\text{Flag } K^{(1)} \setminus K$

If K is a flag complex
 $\text{Blockers}(K) = \emptyset$

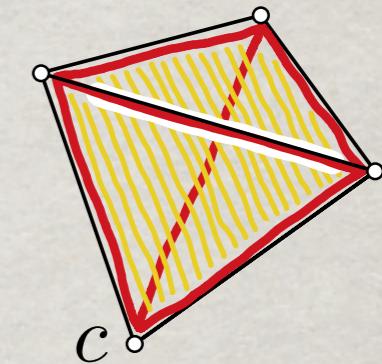


For a flag complex K , the pair reduces to:
 $(K^{(1)}, \emptyset)$

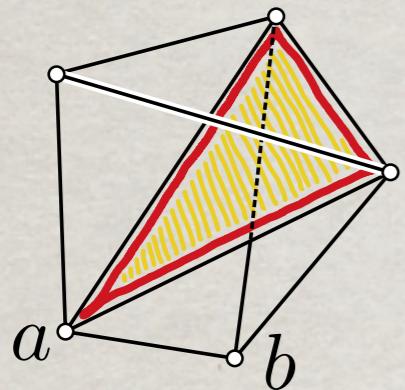
UPDATING DATA STRUCTURE



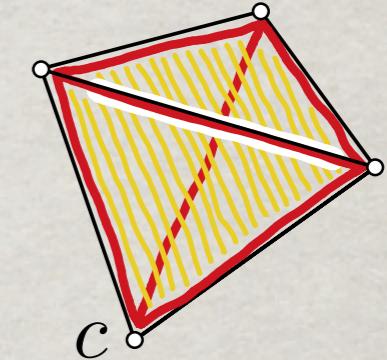
$$K = (K^{(1)}, \text{Blockers}(K)) \xrightarrow{ab \mapsto c} K' = (K'^{(1)}, \text{Blockers}(K'))$$



UPDATING DATA STRUCTURE



$$K = (K^{(1)}, \text{Blockers}(K)) \xrightarrow{ab \mapsto c} K' = (K'^{(1)}, \text{Blockers}(K'))$$

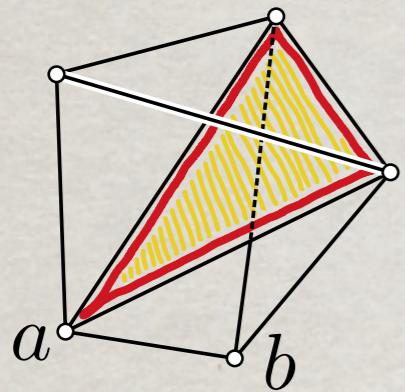


Lemma 1. $c\sigma \in \text{Blockers}(K')$ with $\sigma \subset \text{Vert}(K) \setminus \{a, b\}$ and $\dim \sigma \geq 1$ iff:

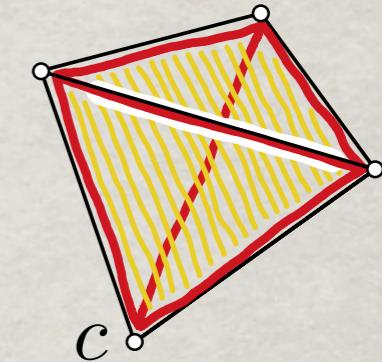
- (i) $\sigma \in K$; for all $\tau \subsetneq \sigma$, $\tau \in \text{Lk}(a) \cup \text{Lk}(b)$;
- (ii) $\sigma = \alpha\beta$ with $a\beta \in \text{Blockers}_0(K)$ and $b\alpha \in \text{Blockers}_0(K)$,

where $\text{Blockers}_0(K) = \text{Blockers}(K) \cup$ complement of $K^{(1)}$

UPDATING DATA STRUCTURE



$$K = (K^{(1)}, \text{Blockers}(K)) \xrightarrow{ab \mapsto c} K' = (K'^{(1)}, \text{Blockers}(K'))$$



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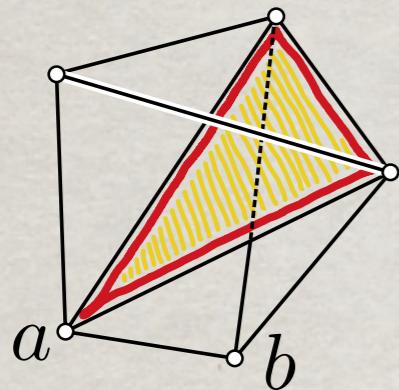
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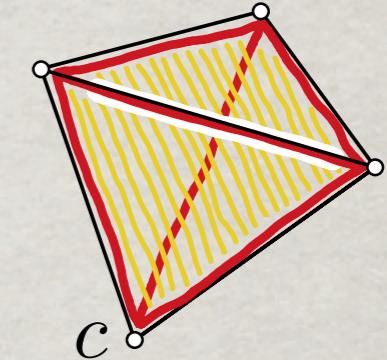


$\text{Lk } v = ((\text{Lk } v)^{(1)}, \text{Blockers}(\text{Lk } v))$

UPDATING DATA STRUCTURE



$$K = (K^{(1)}, \text{Blockers}(K)) \xrightarrow{ab \mapsto c} K' = (K'^{(1)}, \text{Blockers}(K'))$$



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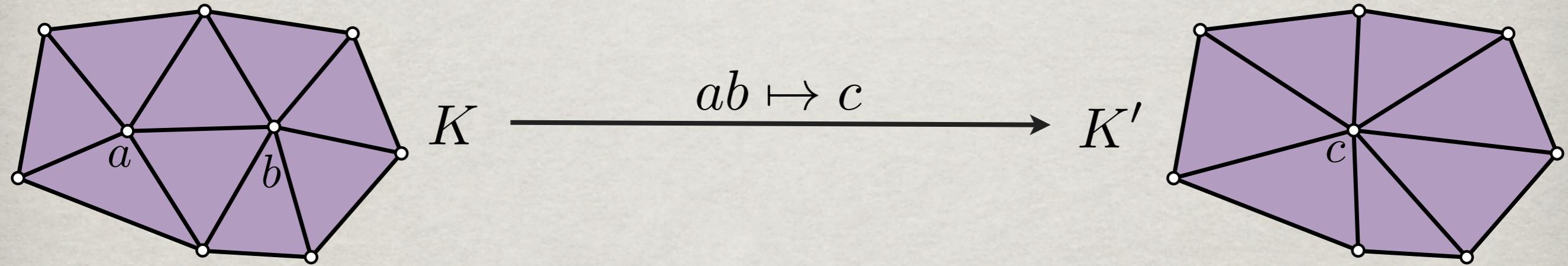


$\text{Lk } v = ((\text{Lk } v)^{(1)}, \text{Blockers}(\text{Lk } v))$

If no blockers “around” a and b , costs in $\tilde{O}(\#\text{neighbors}(a) \times \#\text{neighbors}(b))$

EDGE CONTRACTION

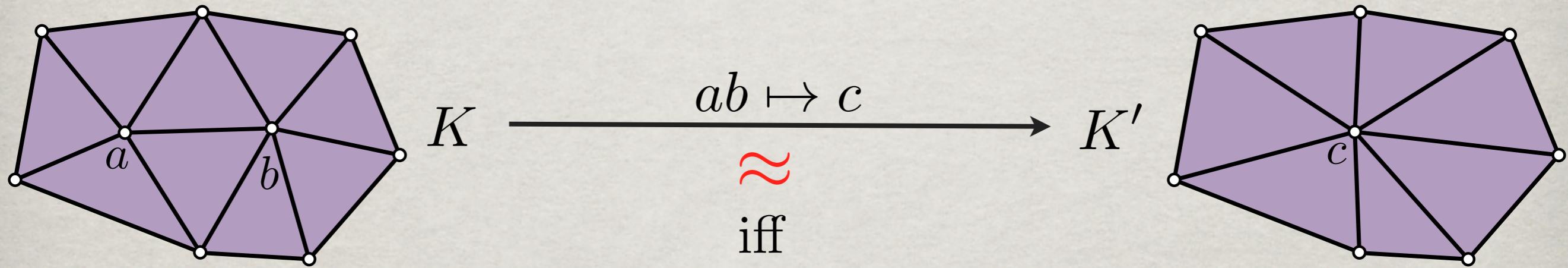
How to preserve homotopy type?



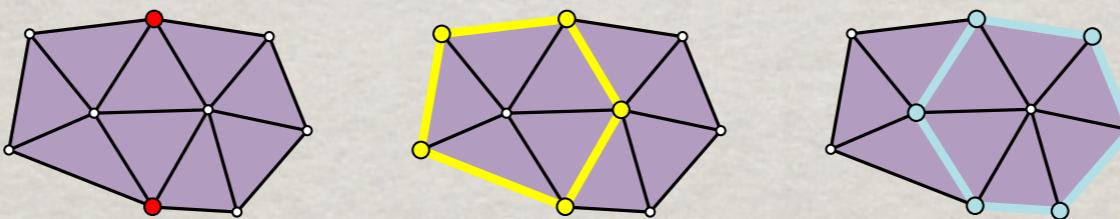
EDGE CONTRACTION

[Dey, Edelsbrunner, Guha & Nekhayev 1999]

If K triangulates a 2- or 3-manifold

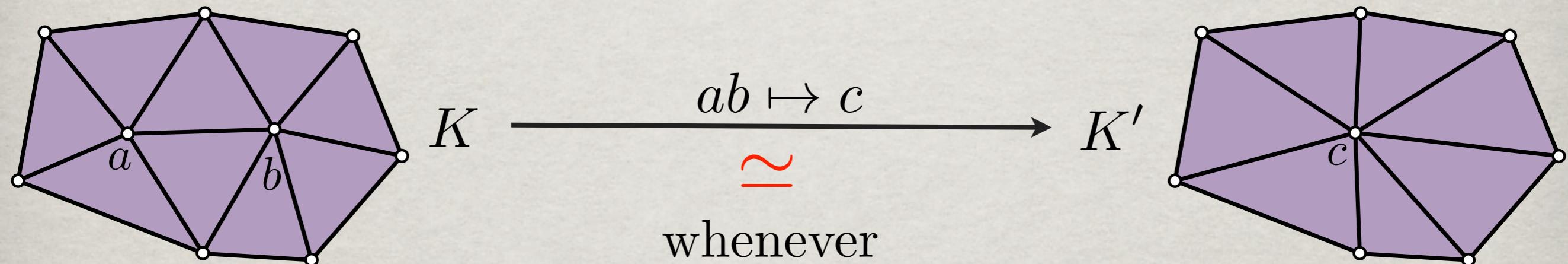


$$\text{Lk } ab = \text{Lk } a \cap \text{Lk } b$$

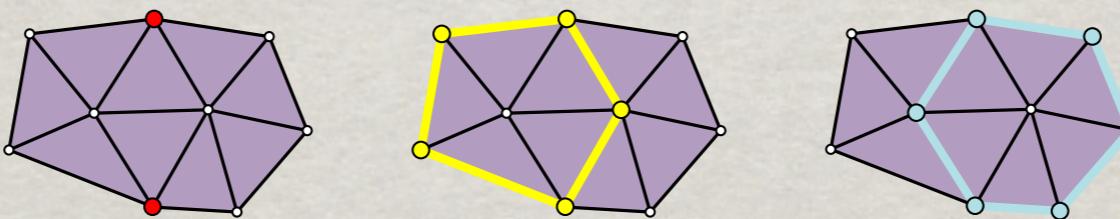


EDGE CONTRACTION

For arbitrary simplicial complexes, we established that:

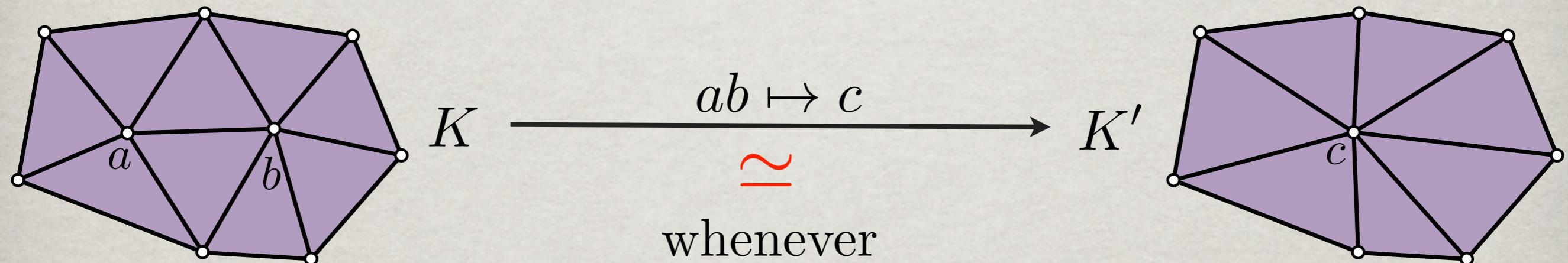


$$\text{Lk } ab = \text{Lk } a \cap \text{Lk } b$$

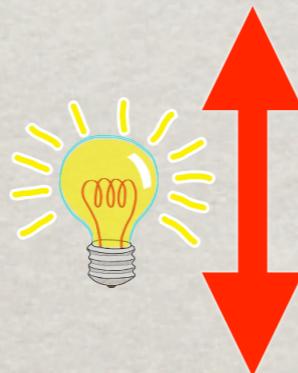


EDGE CONTRACTION

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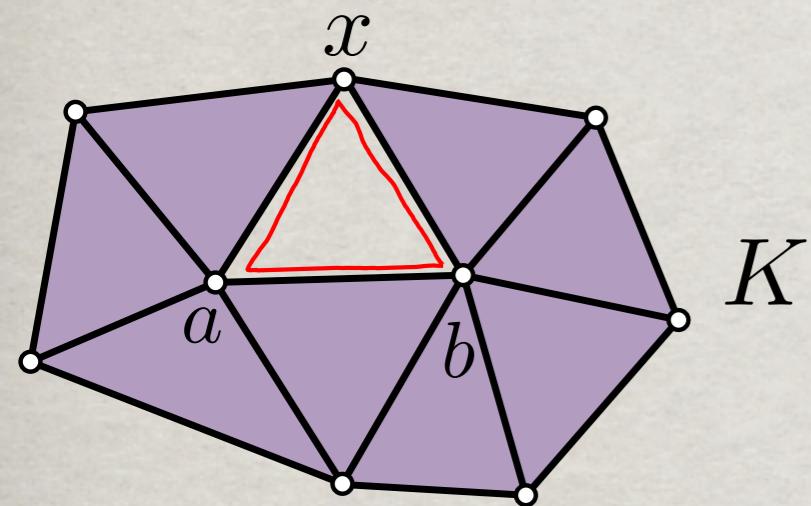
$$\text{Lk } ab = \text{Lk } a \cap \text{Lk } b$$



No blocker of K contains ab

EDGE CONTRACTION

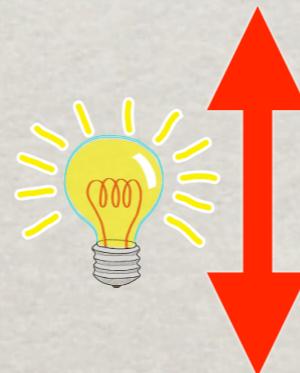
For arbitrary simplicial complexes, we established that:



$$K \xrightarrow{ab \mapsto c} K'$$

implies that

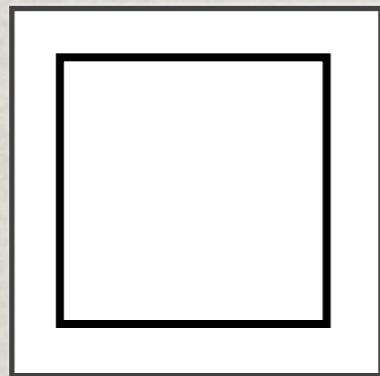
$$\text{Lk } ab \neq \text{Lk } a \cap \text{Lk } b$$



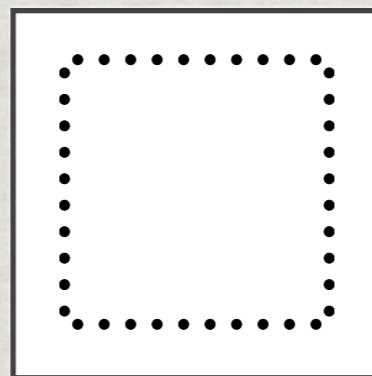
\exists a blocker of K containing ab

EXPERIMENTS

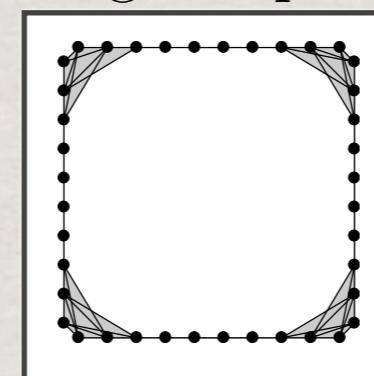
$$C_d = \partial[-1, 1]^{d+1}$$



Point cloud

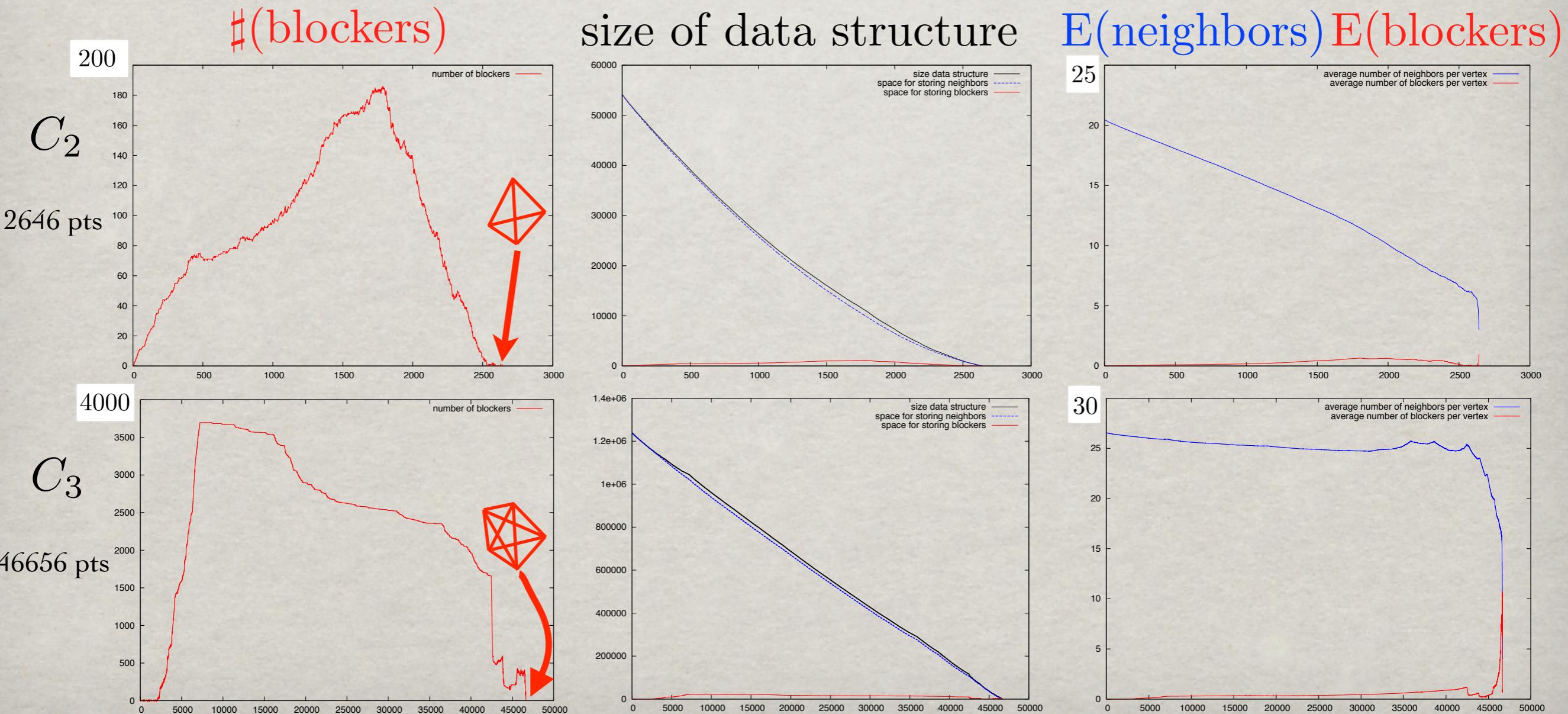


Flag complex



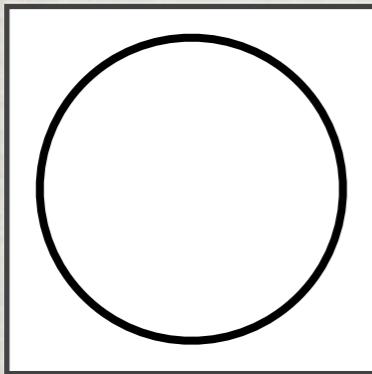
\approx

We keep contracting
shortest edge with
no blocker through it

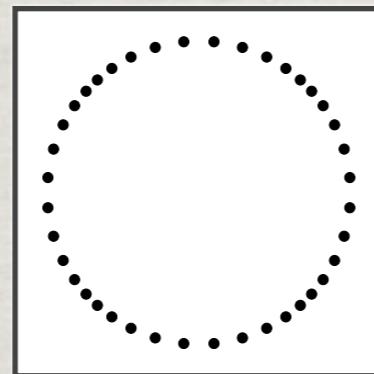


EXPERIMENTS

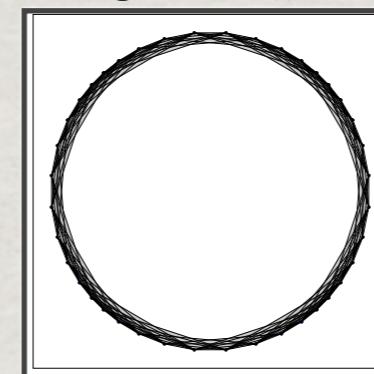
$S_d = d\text{-sphere}$



Point cloud

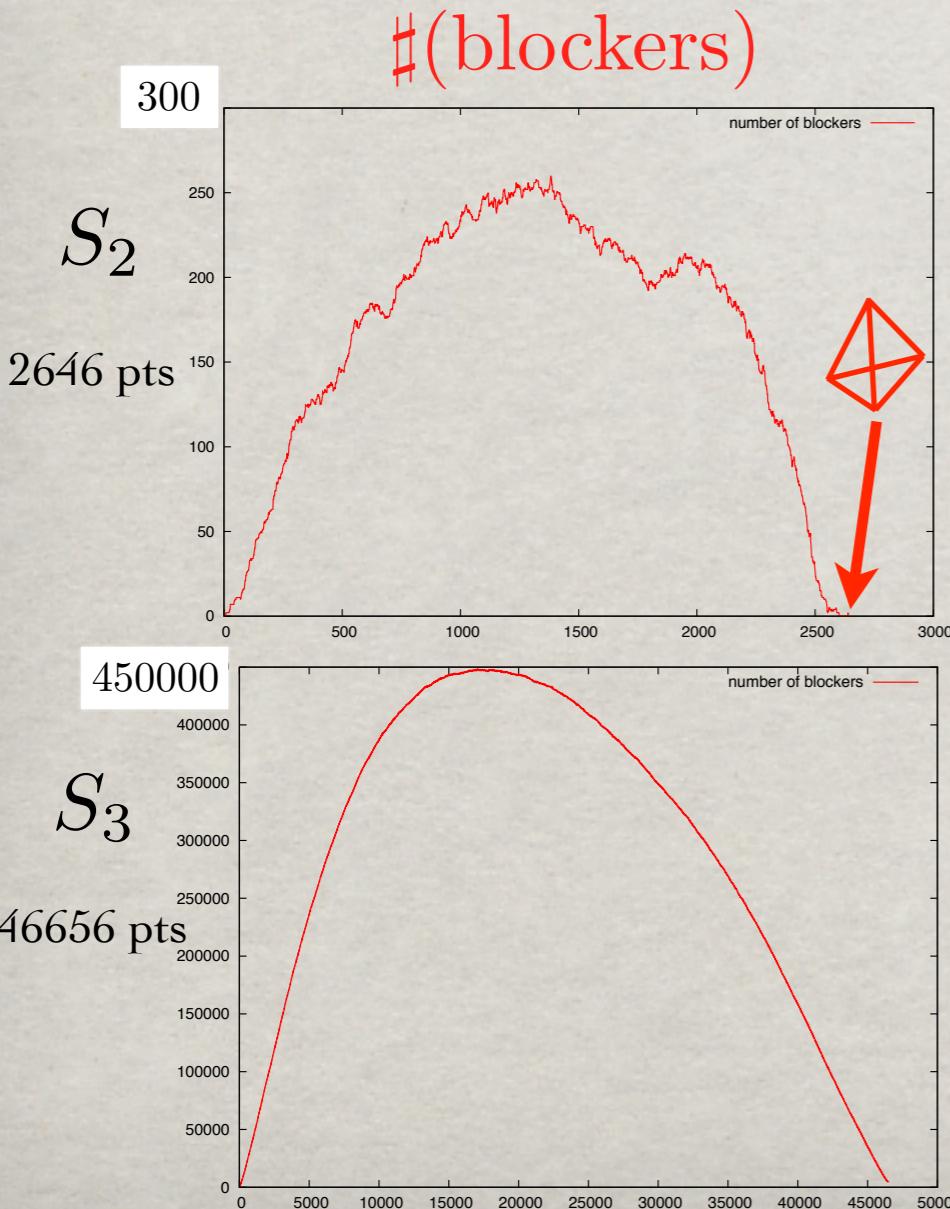


Flag complex

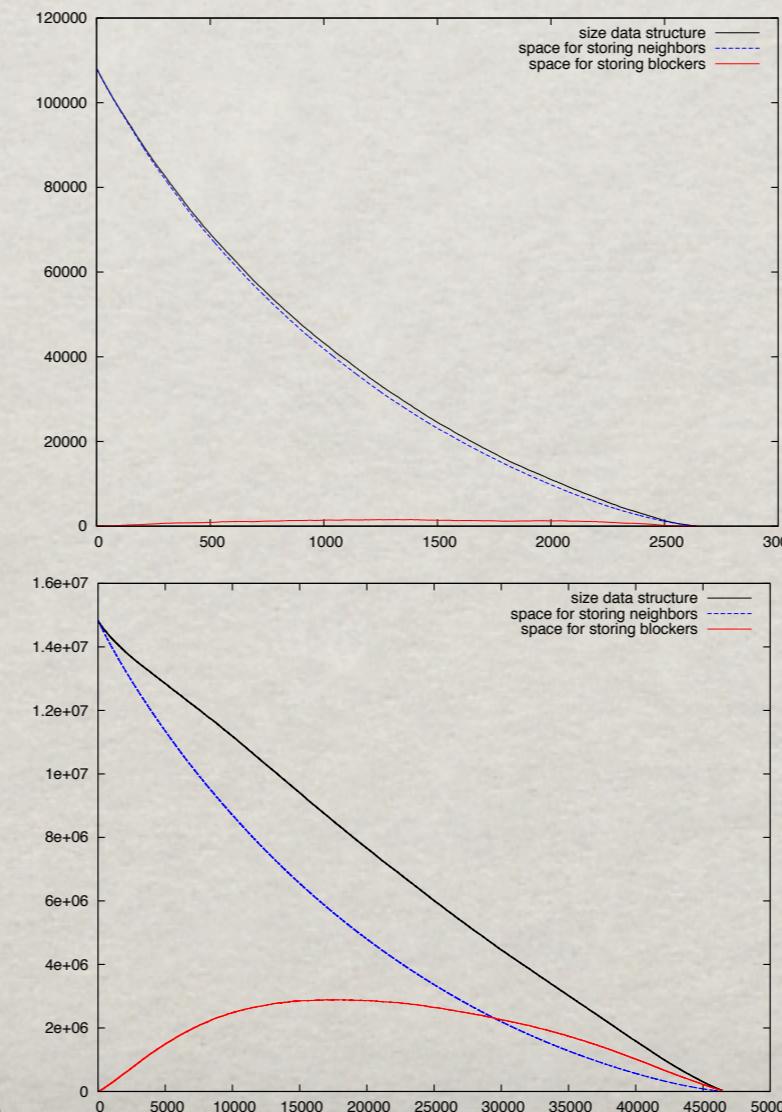


\approx

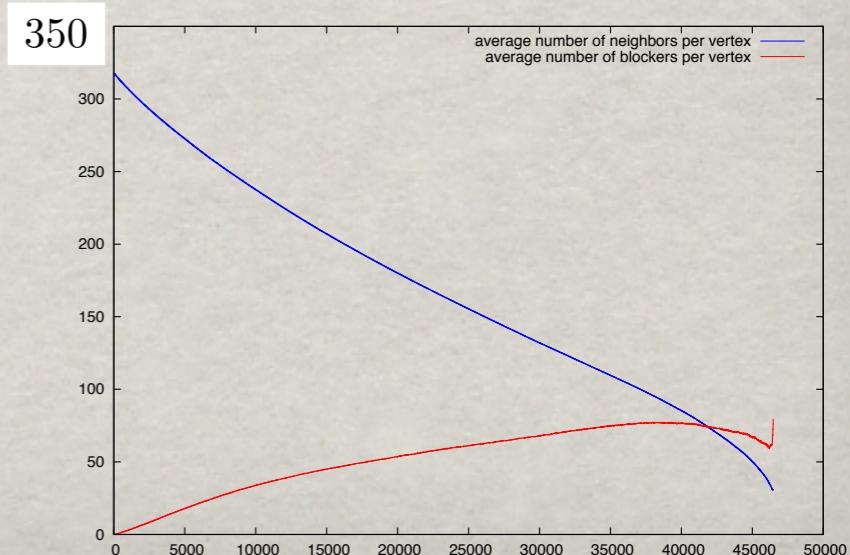
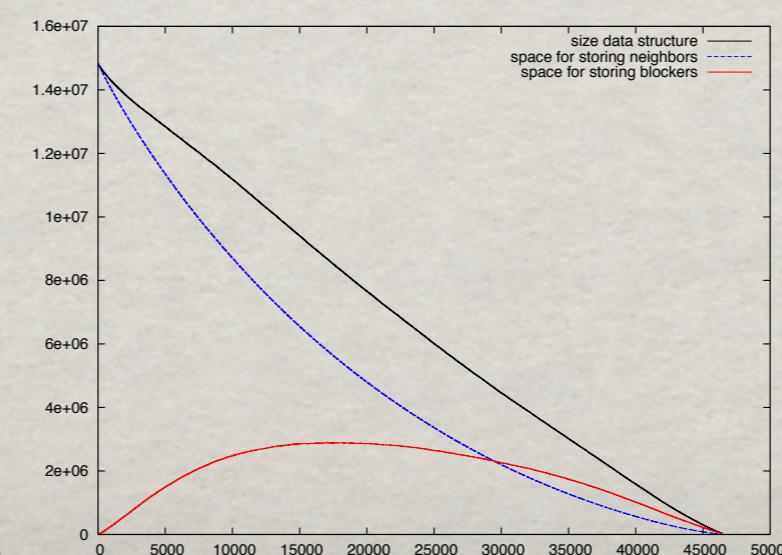
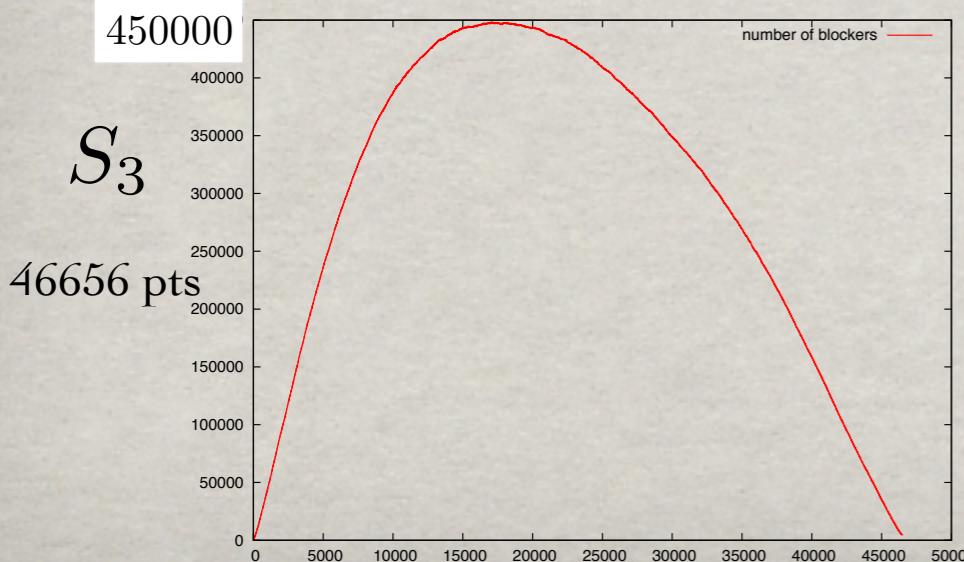
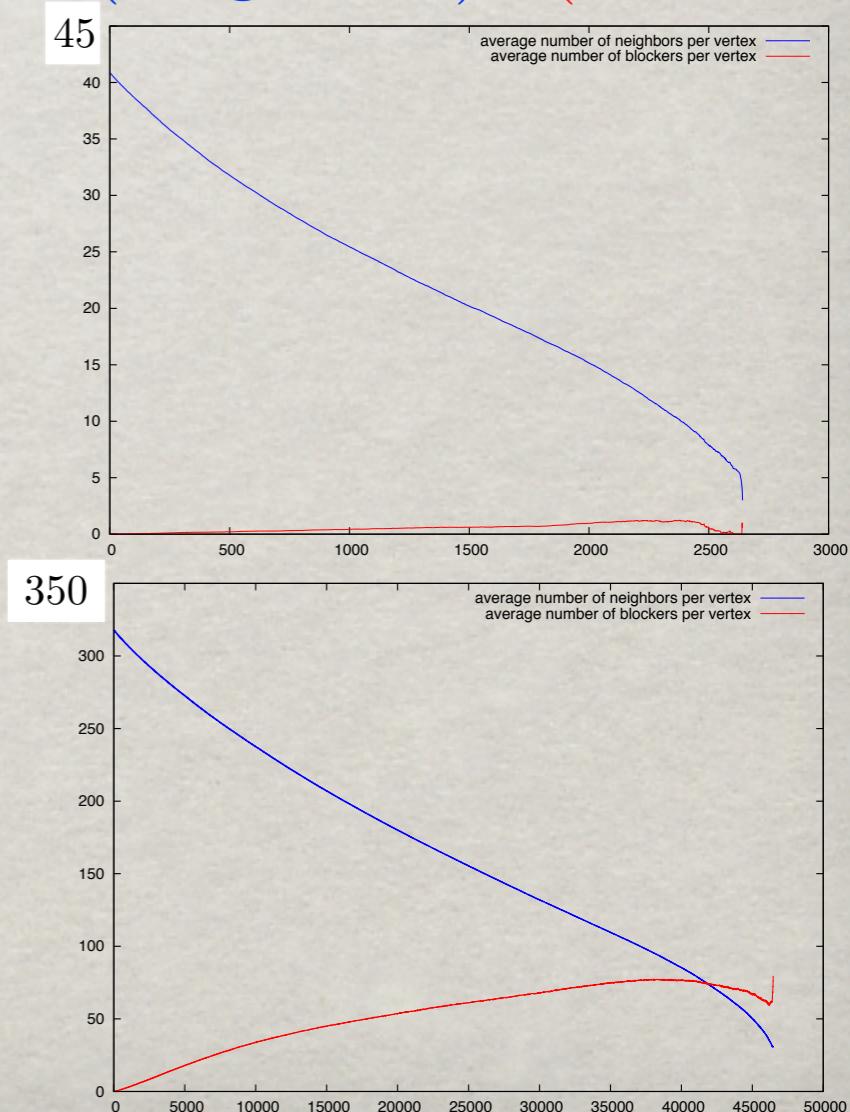
We keep contracting
shortest edge with
no blocker through it



size of data structure



E(neighbors) E(blockers)



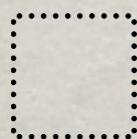
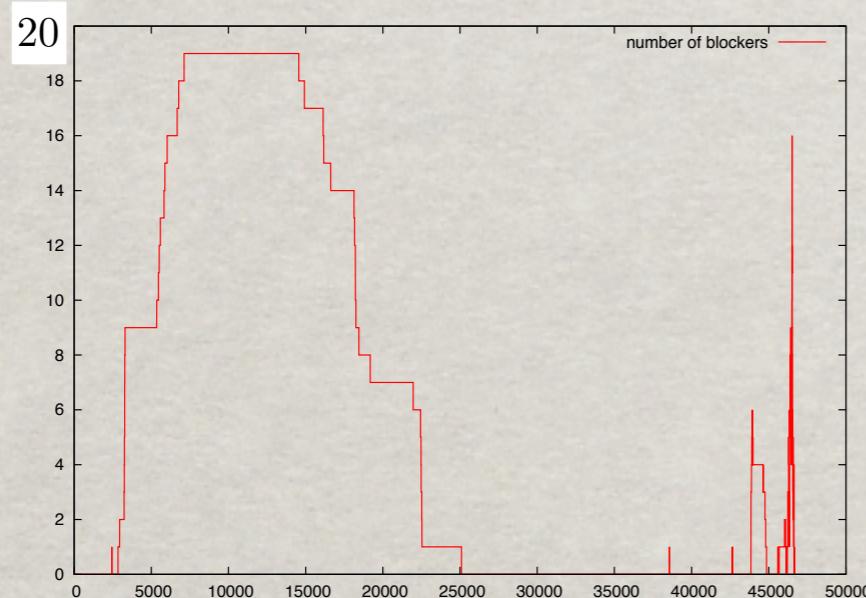
EXPERIMENTS

$K = (G, B) \xrightarrow{\text{extended anticollapse}} K' = (G, B \setminus \{b\})$

\approx

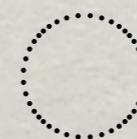
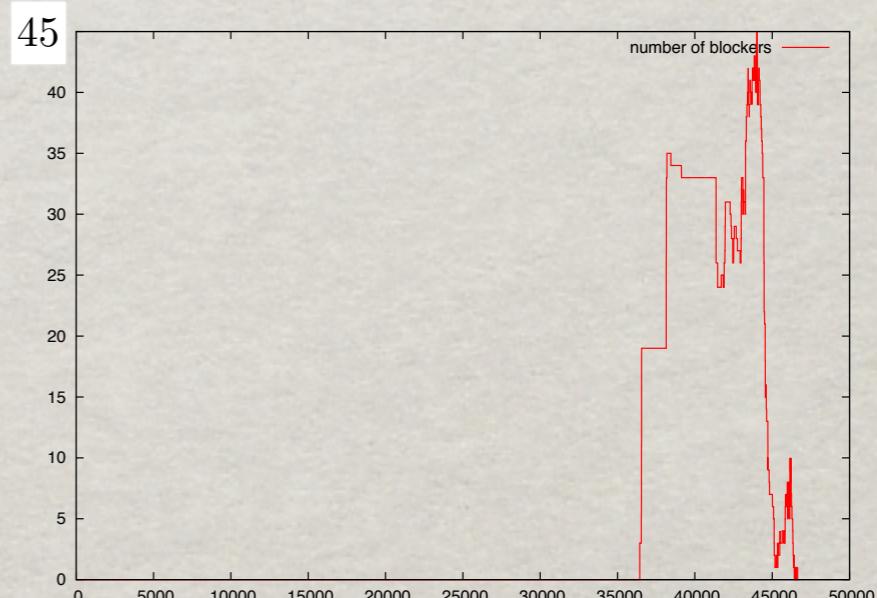
whenever

Lk _{K'} b is a cone



C3

46656 pts



S3

46656 pts

CONCLUSION

- ✿ Promising data structure for encoding high dimensional simplicial complexes.
- ✿ Need further investigations:
 - ✿ study other strategies to prioritize contractions besides edge length
 - ✿ study other simplification operations besides edge contractions (collapses, ...)

