# A Stateful Inspection of FireWall-1



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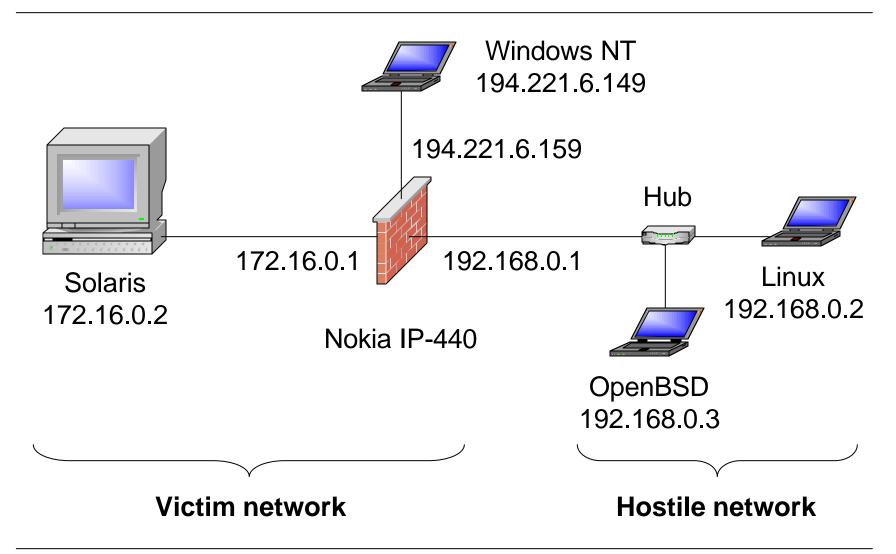


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#### Overview

- Architecture of FireWall-1
- Attacking the firewall's state I
- FWZ encapsulation
- Attacking the firewall's state II
- Attacking authentication between firewall modules
- Hardening FireWall-1
- The big picture

### Topology



### Problems in Inspection

- Unreliable / unauthenticated input
- Layering restrictions on inspection
- Layering violations in inspection
- Ambiguous end-to-end semantics

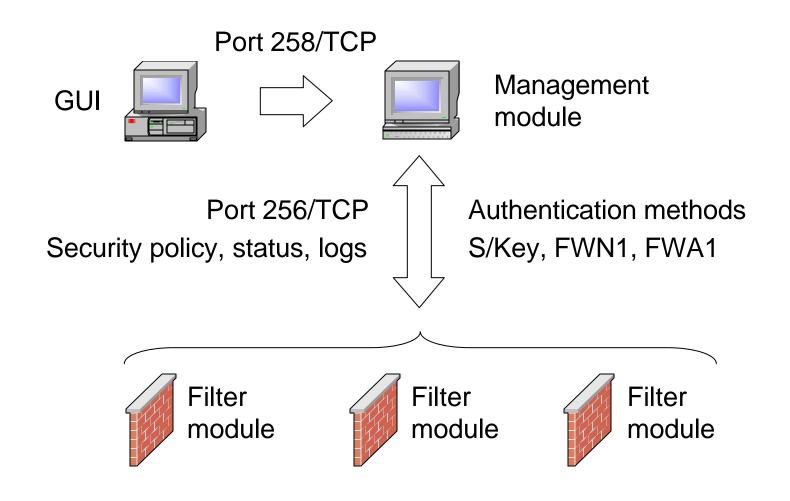
### **Example: Airport Security**

- Unreliable / unauthenticated input
   Examining baggage tags
- Layering restrictions on inspection
   Examining shape, size, weight
- Layering violations in inspection
   Parallelizing bag content inspection
- Ambiguous end-to-end semantics
   Checking for known contraband

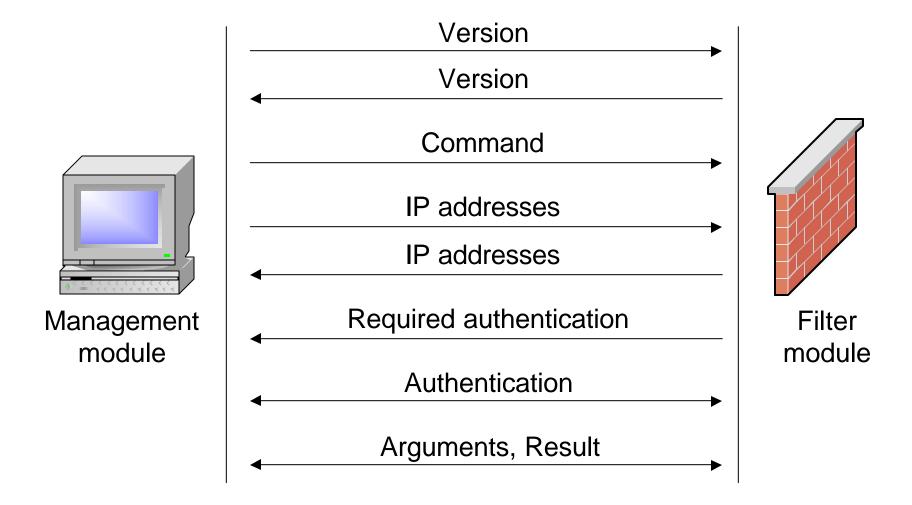
#### Classification of the Attacks

- Unreliable / unauthenticated input
  - TCP fastmode
- Layering restrictions on inspection
  - FWZ VPN encapsulation
- Layering violations in inspection
  - FTP data connection handling
  - unidirectional TCP data flow
  - RSH error connection handling
- Ambiguous end-to-end semantics
  - Parsing of FTP "PORT" commands

#### FireWall-1 Modules

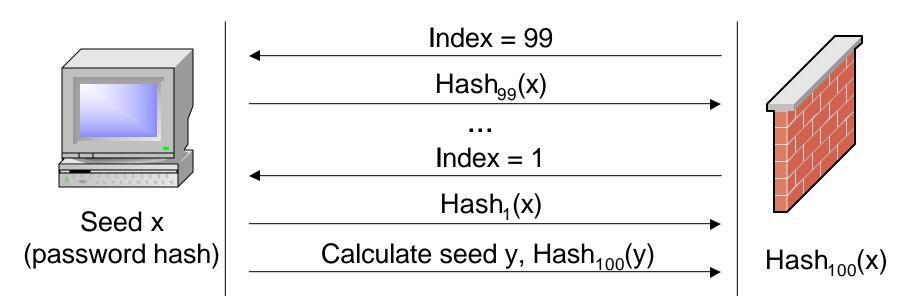


#### Inter-Module Protocol



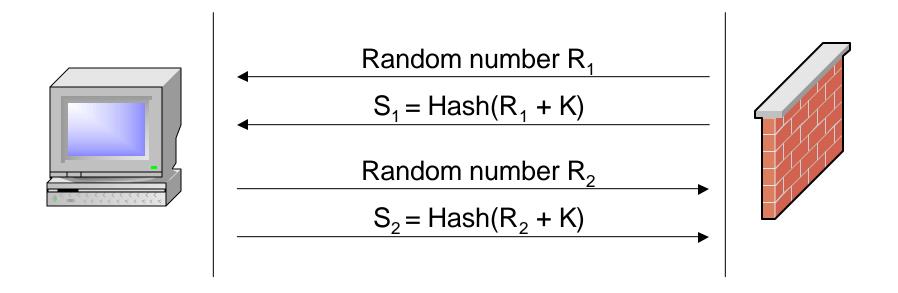
### S/Key Authentication

$$Hash_n(x) = Hash(Hash(... Hash(x))) = Hash(Hash_{n-1}(x))$$
  
n times



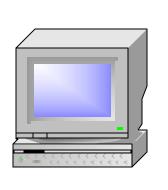
- "y = MakeSeed(time(NULL))"
- Attack: brute force

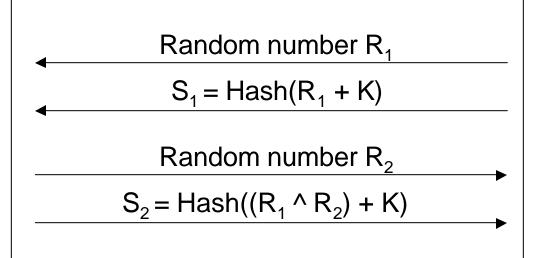
#### FWN1 Authentication

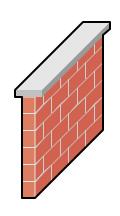


- Shared key K ("fw putkey")
- Attack: choose  $R_2 = R_1$ , so that  $S_2 = S_1$

#### FWA1 Authentication

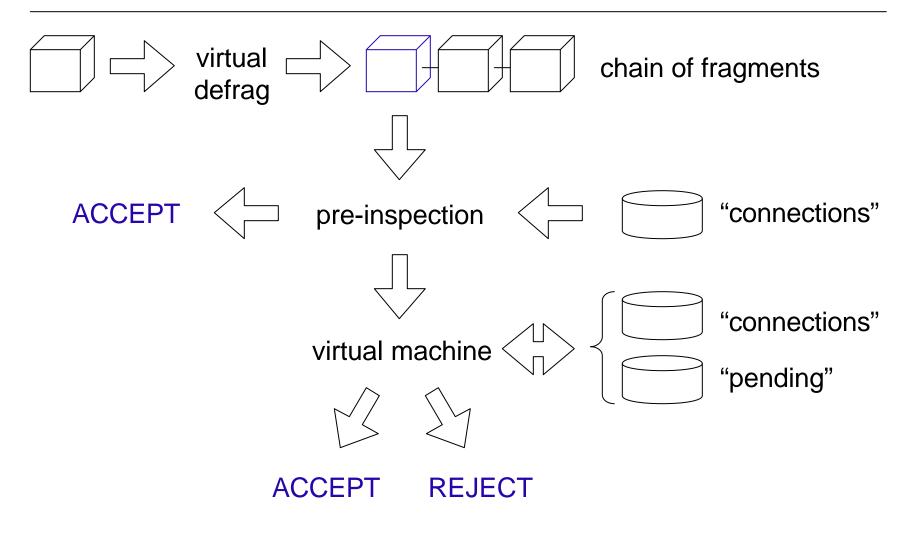




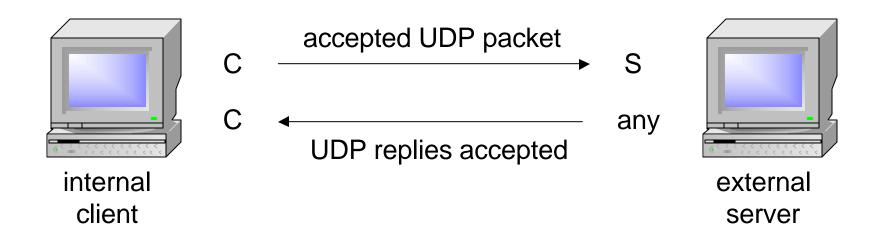


- Shared key K ("fw putkey")
- Attack: choose  $R_2 = 0$ , so that
  - $R_1 \land R_2 = R_1$  and
  - $S_2 = \text{Hash}((R_1 \land R_2) + K) = \text{Hash}(R_1 + K) = S_1$
- To be solved: encryption

### Stateful Inspection I

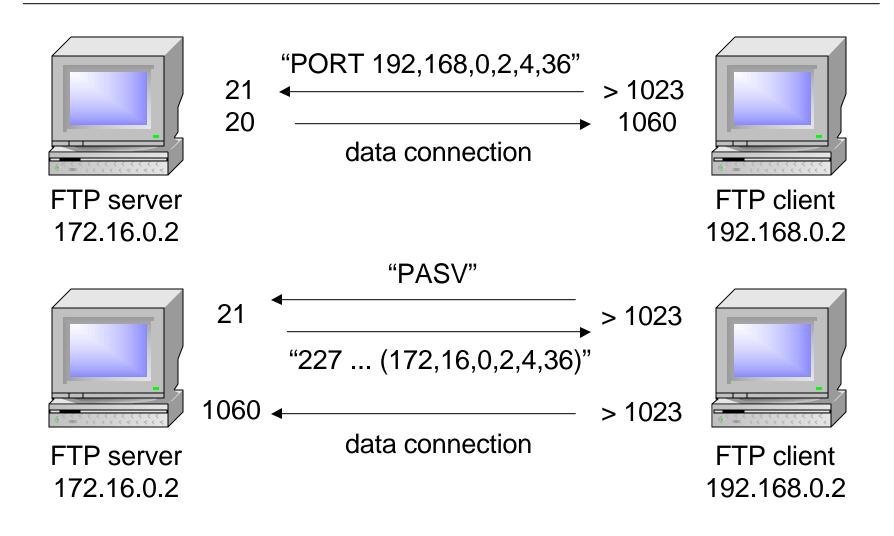


### Stateful Inspection II

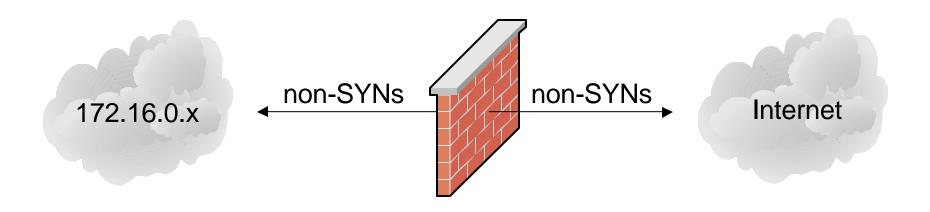


- UDP "connections"
  - from a client, port C
  - to a server, port S + wildcard port
- <s-address, s-port, d-address, d-port, protocol>

### Stateful Inspection III

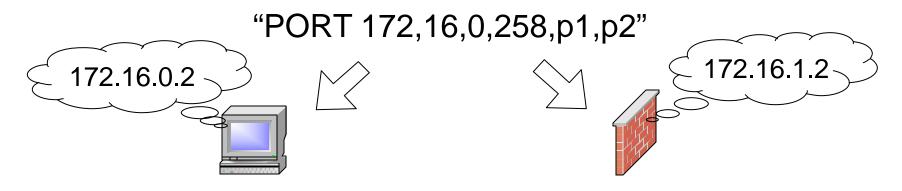


#### **Fastmode Services**

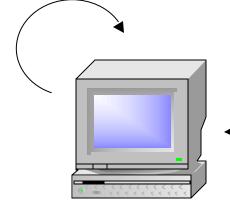


- non-SYN packets accepted
  - Source port = fastmode service
  - Destination port = fastmode service
- Stealth scanning (FINs, ...)

### FTP "PORT" Parsing



data connection



172.16.0.2

#### **Application:** bounce attack

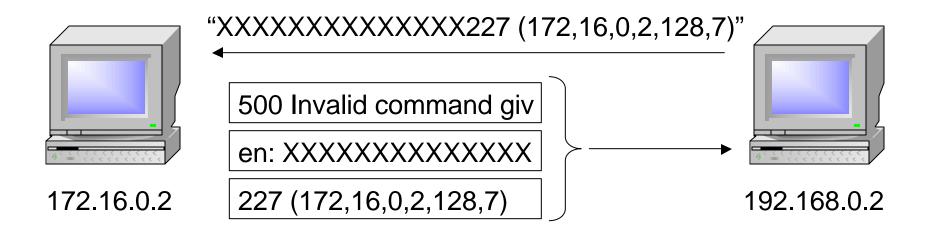
"PORT 172,16,1349632,2,p1,p2"

1349632 = 65536 \* (192 - 172) + 256 \* (168 - 16)



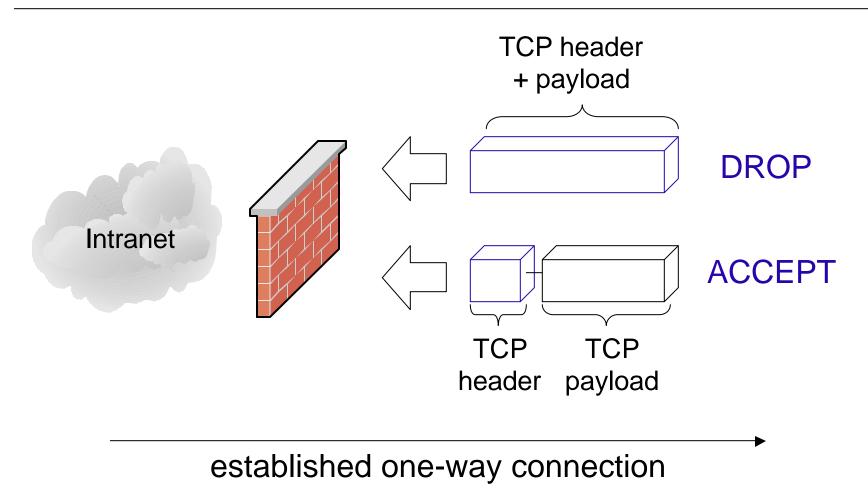
192.168.0.2

### FTP "PASV" Handling

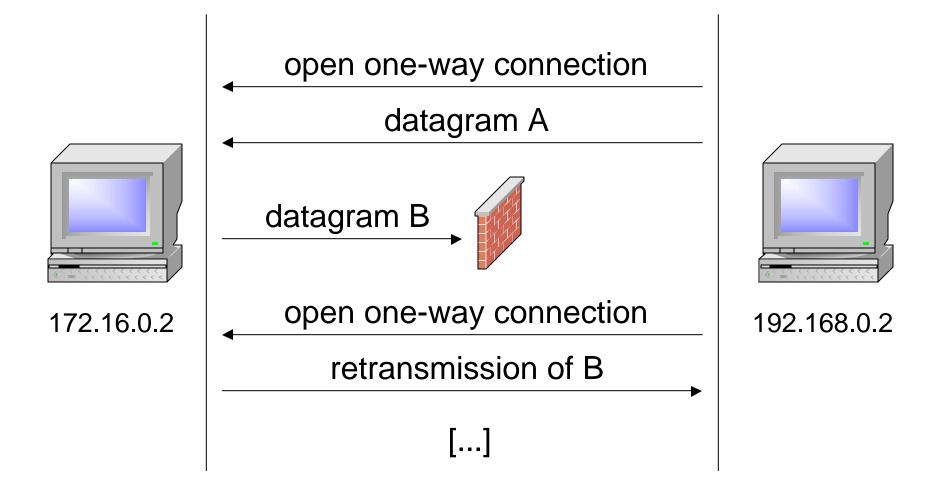


- Advertise small Maximal Segment Size
- Server replies split

### One-way Connections I

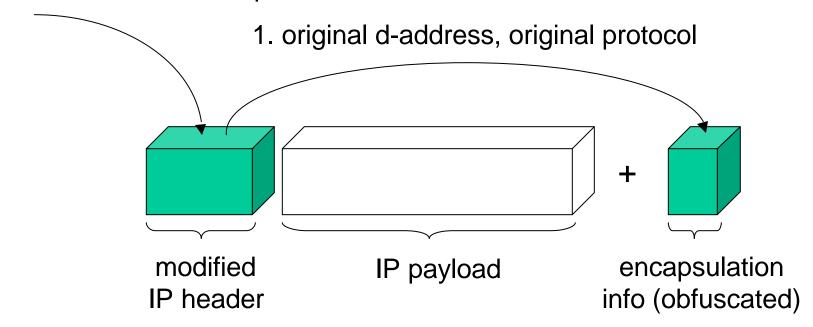


#### **One-way Connections II**



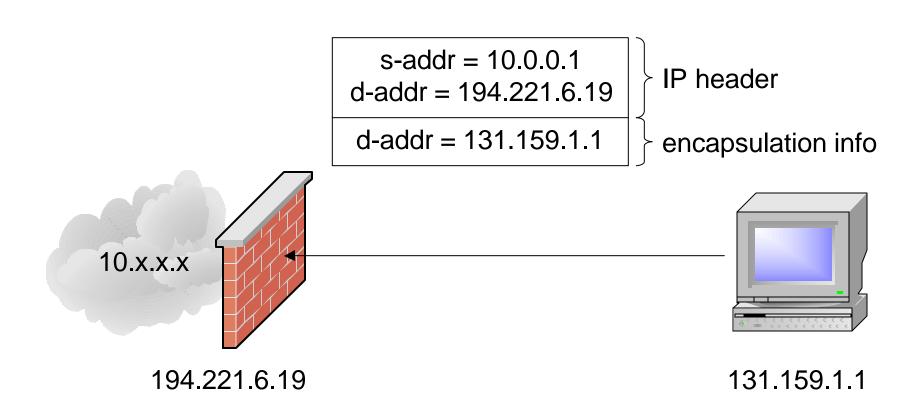
### FWZ Encapsulation I

2. d-address = firewall, protocol = 94



- VPN tunneling protocol
- Decapsulation without decryption or authentication
- Cannot be disabled

### FWZ Encapsulation II



#### **Key to spoofing attacks**

#### Fake "PORT" Commands

s-addr = 172.16.0.2 d-addr = 192.168.0.1

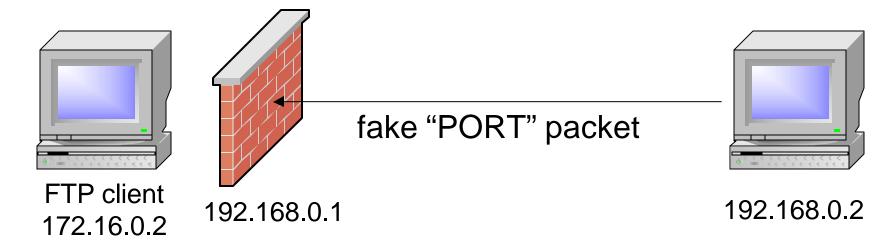
"PORT 172,16,0,2,128,7"

d-addr = 192.168.0.2

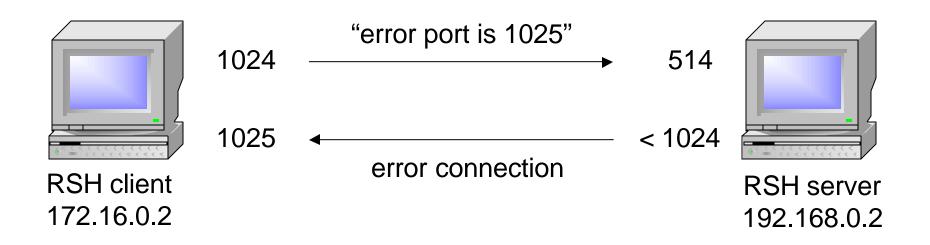
IP header

TCP header + payload

encapsulation info



#### RSH Error Connections I



- <172.16.0.2, 1024, 192.168.0.2, 514, 6> in "connections"
- <172.16.0.2, 1025, 192.168.0.2, magic, 6> in "pending"
- Reversed matching

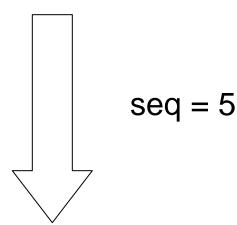
#### RSH Error Connections II

SYN

- s-addr:s-port
- d-addr:magic
- seq + 1
- 172.16.0.2:1024
- 192.168.0.2:magic
- 250001

- packet #2
  (port info)
- s-addr:error-port
- d-addr:magic
- protocol
- 172.16.0.2:1025
- 192.168.0.2:magic
- 6 (TCP)

- s-addr:s-port
- d-addr:magic
- seq + 1



- 172.16.0.2:32775
- 192.168.0.2:magic
- 6 = seq + 1 = TCP

### Fake UDP Requests

s-addr = 172.16.0.2 d-addr = 192.168.0.1

s-port = 161 d-port = 53

d-addr = 192.168.0.2

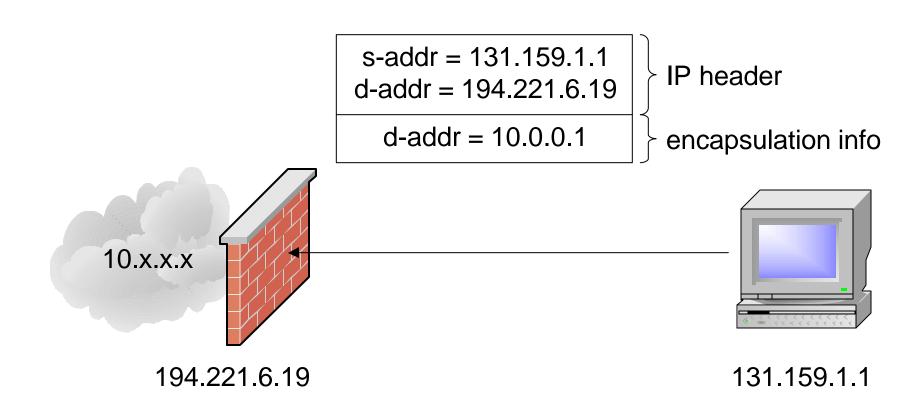
IP header

**UDP** header

encapsulation info



### FWZ Encapsulation III



#### Key to non-routable addresses

# Anti-Spoofing Protection I

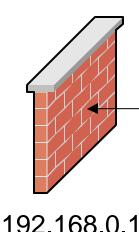
1. s-addr = 192.168.0.1 d-addr = 192.168.0.1

s-port = 161 d-port = 53

d-addr = 192.168.0.2

2. s-addr = 192.168.0.2 d-addr = 192.168.0.1

s-port = any d-port = 161



1. fake DNS request

2. tunnel to firewall



192.168.0.2

# Anti-Spoofing Protection II

1. s-addr = 224.0.0.1 d-addr = 192.168.0.1

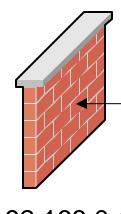
s-port = 161 d-port = 53

d-addr = 192.168.0.2

2. s-addr = 192.168.0.2 d-addr = 192.168.0.1

s-port = 53 d-port = 161

d-addr = 224.0.0.1



1. fake DNS request

2. tunnel to firewall



192.168.0.2

### Hardening I

- Disable implicit rules
  - DNS
  - control connections
  - ICMP
- Restrictive access rules
  - no "any" sources or destinations
  - deny broadcast / multicast addresses
  - "minimal privilege"
- Properly configure anti-spoofing mechanism
- Filter protocol 94 (e.g. IP Filter)

### Hardening II

- Different (virtual) IP addresses for public services
- Restrict control connections
  - FWA1 authentication
  - VPN technology
  - never use "127.0.0.1: \*/none"
- More than one line of defense!

### Fixes by Check Point

Solutions by Check Point available today at

http://www.checkpoint.com/techsupport

# Thanks.

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