Data Visualization From a Category Theory Perspective

Davi Sales Barreira, Asla Medeiros e Sá

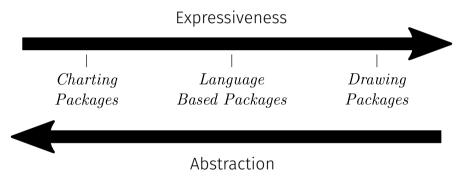
FGV - EMAp, IMPATech

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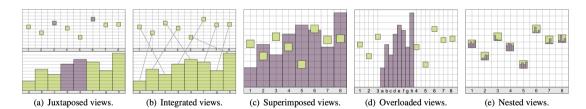
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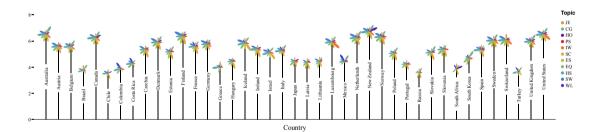
- 1. Category Theory behind Vizagrams
- 2. Vizagrams Hands-on

Balance expressiveness and abstraction in data visualization frameworks.

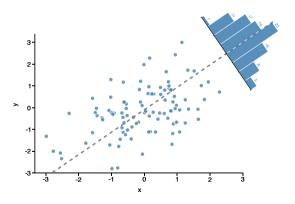


Visualization grammars often come short in terms of expressiveness, e.g., composite visualizations and visualizations with custom marks.





```
d = scatter_plot +
  T(x,y)R(θ)*histogram +
  S(:strokeDasharray=>2)*Line([p1,p2])
```



Introduction - Objective and Hypotheses

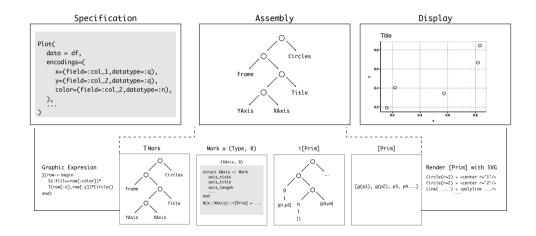
Main Objective: Develop a data visualization framework that extends the expressiveness of visualization grammars while preserving a high level of abstraction.

Hypothesis 1 (Cause): Limitations in expressiveness are due to a lack of integration between graphic specification and assembly.

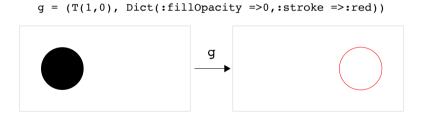
Hypothesis 2 (Solution): Integration of diagramming and graphic specification via a unified constructive framework.

Formalization Tool: Category Theory.

Introduction - Overview



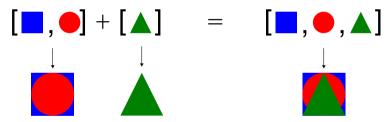
The basic building block for any diagram is the primitive. A primitive is any geometry that can be drawn and manipulated via geometric or stylistic transformations.



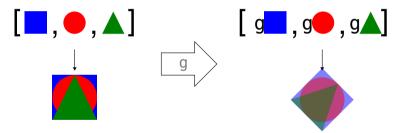
A diagram can be defined as an ordered list of primitives. ($[Prim], +, [\]$) can be interpreted as a monoid.

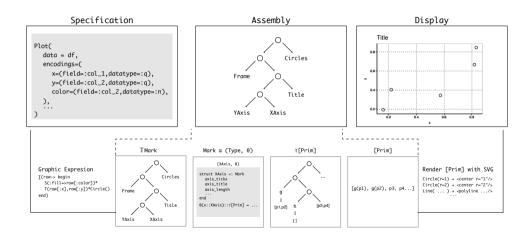
```
abstract type Prim end
struct Circle <: Prim
  radius::Real
  center::Tuple{Real, Real}
end
+(d1::Vector{Prim}, d2::Vector{Prim}) = vcat(d1,d2)</pre>
```

Diagrams are drawn by rendering the list of primitives on top of each other.



We can apply transformations to diagrams by applying the transformations to each primitive in the list.





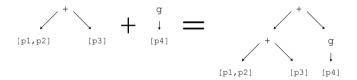
Data Visualization + Category Theory: Free Monads

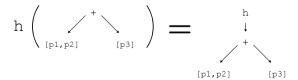
Let us improve our representation using expression trees. To do this, first we define a functor F that encodes diagram composition (i.e. +) and diagram transformations.

```
@data F{a} begin
Comp(::a, ::a)
Act(::H, ::a)
end
fmap(f::Function, x::Comp) = Comp(f(x._1),f(x._2))
fmap(f::Function, x::Act) = Act(x._1,f(x._2))
```

Data Visualization + Category Theory: Free Monads

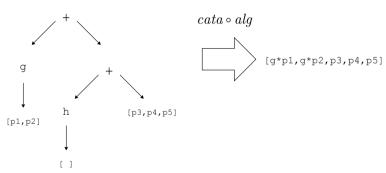
Using our functor F, we can define a free monad over such functor, i.e $\mathbb{T} := \mathsf{Free} F$. Thus, (\mathbb{T}, η, μ) is our monad, where \mathbb{T} is a parametric type representing trees with F as possible operations. To ease the process of expressing trees, we overload the + and the * operators to represent composition and transformation application respectively.

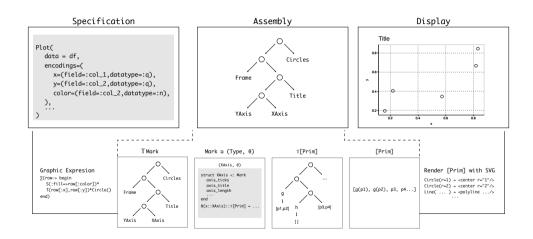




Data Visualization + Category Theory: Free Monads

A diagram tree is a value of type \mathbb{T} [Prim]. For this representation to be useful, we must be able to flatten it, i.e. we must define a function $f:\mathbb{T}[\mathsf{Prim}] \to [\mathsf{Prim}]$. We can do this using F-algebras and catamorphism.





The current definitions of marks equate them with primitives. We propose a new definition:

Def: A **Graphical Mark** is a tuple (A, θ_A) , where A is a type and θ_A is a function $\theta_A : A \to \mathbb{T}[Prim]$.

```
struct Face <: Mark
    size::Real
    smile::Real
end

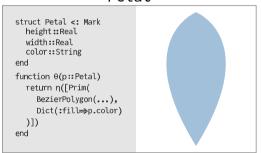
function θ(face::Face)
    eyes = Circle(...) + Circle(...)
    smile = Bezier(...)
    head = Circle(...)

    return head + eyes + smile
end</pre>
```

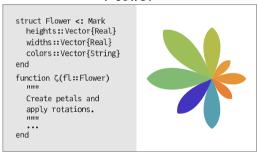


We want to be able to define marks using previously defined marks. In other words, instead of specifying $\theta_A:A\to \mathbb{T}[Prim]$, we want to define a function $\zeta_A:A\to \mathbb{T}Mark$. In other words, we want use ζ_A to infer θ_A .

Petal



Flower



Our ζ function can be formalized through the use of slice categories.

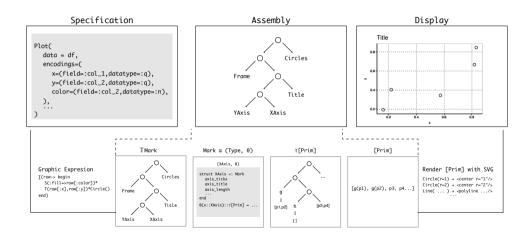
Def. The category of marks \mathcal{M} is a subcategory of $\mathbf{Set}_{\mathbb{T}}/[\mathsf{Prim}]$.

Assuming this category, a morphism $\zeta_A \in \operatorname{Hom}_{\mathcal{M}}(A,B)$ induces

$$\theta_{A} = \theta_{\mathbb{T}B} \circ \zeta_{A} = \mu \circ \mathbb{T}\theta_{B} \circ \zeta_{A}$$

This means that given an existing mark (B, θ_b) , we can define a new mark (A, θ_A) by defining ζ_A and inferring of θ_A .

Programmatically, we can define an abstract type 'Mark', such that every subtype of 'Mark' has a *theta* function.



In summary:

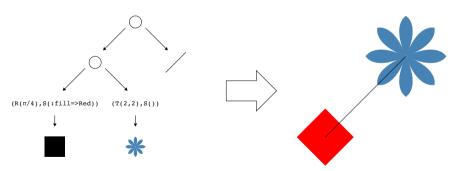
- A diagram is a value of [Prim];
- A diagram tree is a value of T[Prim];
- A mark is a value of type A <: Mark for which $\exists \theta: A \to \mathbb{T}[\mathsf{Prim}];$

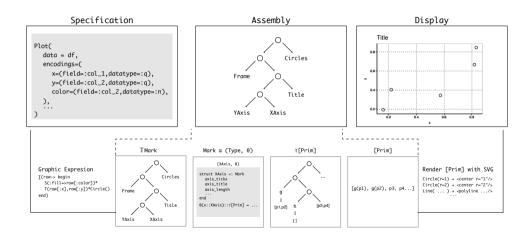
With all these concepts in mind, we can define **graphic** as a value of $\mathbb{T}Mark$, i.e. a tree where the leafs contain marks. Since each mark has a θ function, we can turn a $\mathbb{T}Mark$ into a $\mathbb{T}[Prim]$ using $\mu \circ \theta$.

Note that $\mathbb{T}\theta: \mathbb{T}\mathsf{Mark} \to \mathbb{T}\mathbb{T}\mathsf{Mark}$ and $\mu: \mathbb{T} \circ \mathbb{T} \to \mathbb{T}$.

Using + as diagram composition and * to apply transformations, we can construct values of \mathbb{T} Mark using a notation similar to mathematical expressions:

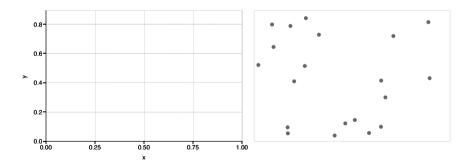
$$S(:fill=>:red)*R(\pi/4)*Square() + T(2,2)*Flower() + Line([0,0],[2,2])$$





Plot Specifications

 $\mathsf{Plot} = \mathsf{Guide} + \mathsf{Graphic} \,\, \mathsf{Expression} \circ \mathsf{Scale}(\mathsf{Data}, \, \mathsf{Encodings})$



References