3.1 (16) Algorithm to find the smallest integer in a finite sequence of numbers.

Procedure smallest (a, ..., an integers)

min := a,

for i = 2 to n

if a < min then min = a;

return min

3.1 (42) S'election-Sort in pseudocode

procedure selection (a, a, a, a, a)

for i = 1 to n-1

min-index := i // place where remaining minimum element is found is found is found if a camin-index then min-index:= j

exchange amin\_index and a:

of the list is now sorted }

Growth Functions

3.2 (10) Show that  $x^3$  is  $O(x^4)$  but that  $x^4$  is NOT  $O(x^3)$ .

· oc³ is O(oc4) because:

oc3 < oc4 whenever oc> 1 take as witnesses C = 1, K = 1

recall definition  $|f(\infty)| \leq C g(\infty)$   $|f(\infty)| \leq C g(\infty)$   $|f(\infty)| \leq C g(\infty)$   $|f(\infty)| \leq C g(\infty)$ 

· oc4 is not O(xe3) because:

Suppose it is true that

oc4 & C oc3 for some witnesses CK

But then it follows that (since oc is assumed positive)

oc & C // divide Both sides By x3

But this can not Rold since C is a constant and oc grows without Bound ( : x > k).

This contradiction completes proof.

## 3.2 (12) Show that exclose or is $O(x^2)$ but that $x^2$ is not $O(x \log x)$

## • $\frac{\text{oclog}}{\text{oc}}$ is $O(\alpha^2)$ because:

 $\propto \log c \leq 0c \cdot c$  for all positive oc take C=1, K=0 as witnesses

Aside: The fact that log oc < oc can be seen by looking at the graphs of the two functions, for example. Or it can be proved formally, by induction: It follows from the fact that oc < 2° and that log oc is a monotonically increasing function in oc.

## • $x^2$ is not $O(x \log x)$ because:

oc & C

But this can not hold because  $\frac{x}{\log x}$  is unbounded:

This can be seen by noting, for example, that log oc < 150 for oc>16

11 book out
11 graphs or
11 by ealculus

And therefore

 $\frac{\partial C}{\log \partial c} > \frac{\partial C}{\sqrt{\partial c}} = \sqrt{\partial c}$ , and  $\sqrt{\partial c}$  is clearly unbounded: