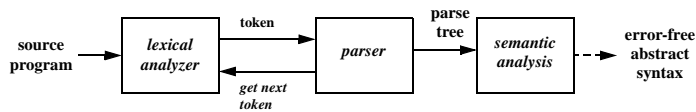


Compiler Front-End

- Almost all compilers and interpreters contain the same **front-end** --- it consists of three components:

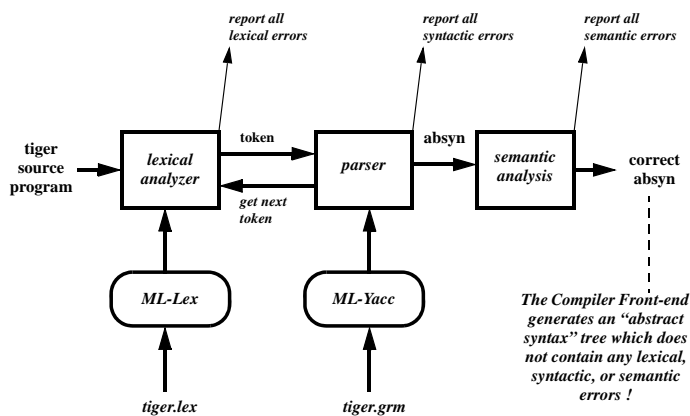
- Lexical Analysis** --- report lexical errors, output a list of tokens
- Syntax Analysis** --- report syntactic errors, output a parse tree
- Semantic Analysis** --- report semantic errors (e.g., type-errors, undefined identifiers, ...) --- generate a clean and error-free “abstract syntax tree”



“Concrete” vs. “Abstract” Syntax

- The grammar specified in “tiger.grm” (for Yacc) is mainly used for parsing only ----- the key is to resolve all ambiguities. This grammar is called **Concrete Syntax**.
- Abstract Syntax (Absyn)** is used to characterize the essential structure of the program ----- the key is to be as simple as possible; **Absyn** may contain ambiguities.
- The grammar for **Abstract Syntax** is defined using **ML datatypes**.
- Traditional Compilers:** do semantic analysis on Concrete Syntax --- implemented as “actions” in Section 3 of “tiger.grm” file (for Yacc)
- Modern Compilers:** “tiger.grm” constructs the Abstract Syntax tree; the semantic analysis is performed on the Absyn later after parsing!

Tiger Compiler Front End



Tiger Program and Expression

- A Tiger program `prog` is just an expression `exp`
- An expression can be any of the following:

l-value	<code>foo, foo.bar, foo[1]</code>
Nil	<code>nil</code>
Integer literal	<code>34</code>
String literal	<code>"Hello, World\n"</code>
Sequencing	<code>(exp; exp; ...; exp)</code>
Function call	<code>id(), id(exp{,exp})</code>
Arithmetic expression	<code>exp arith-op exp</code>
Comparison expression	<code>exp comp-op exp</code>
Boolean operators	<code>exp & exp, exp exp</code>
Record creation	<code>ty-id {id = exp, ...}, {}</code>
Array creation	<code>ty-id [exp₁] of exp₂</code>
Assignment	<code>lvalue := exp</code>

Tiger Expression and Declaration

- *More Tiger expressions:*

If-then-else	if exp_1 then exp_2 else exp_3
If-then	if exp_1 then exp_2
While-expression	while exp_1 do exp_2
For-expression	for $id := exp_1$ to exp_2 do exp_3
Break-expression	break
Let-expression	let $decsq$ in $\{exp\}$ end

- *A Tiger declaration sequence is a sequence of type, variable, and function declarations:*

```
dec -> tydec | vardec | fundec
decsq -> decsq dec | ε
```

Tiger Type Declaration

- *Tiger Type declarations:*

```
tydec -> type id = ty
ty -> id | { tyfields } | array of id
tyfields -> ε | id : type-id { ,id: type-id }
```

- *You can define mutually-recursive types using a consecutive sequence of type declarations*

```
type tree = {key : int, children : treelist}
type treelist = {hd : tree, tl : treelist}
```

recursion cycle must pass through a record or array type !

Variable and Function Declaration

- *Tiger Variable declarations:*

```
short-form:   vardec -> var id := exp
long-form:   vardec -> var id : type-id := exp
```

"var x := 3" in Tiger is equivalent to **"val x = ref 3"** in ML

- *Tiger Function declarations:*

```
procedure:   fundec -> function id (tyfields) := exp
function:    fundec -> function id (tyfields):type-id := exp
```

- *Function declarations may be mutually recursive --- must be declared in a sequence of consecutive function declarations! Variable declarations cannot be mutually recursive !*

Tiger Absyn "Hack"

- *When translating from Concrete Syntax to Abstract Syntax, we can do certain syntactic transformations*

```
MINUS exp    ==>    0 MINUS exp
exp1 & exp2 ==>    if exp1 then exp2 else 0
exp1 | exp2 ==>    if exp1 then 1 else exp2
```

This can make **Abstract Syntax** even simpler.

Toy does not support Macros. If the source language supports macros, they can be processed here.

Tiger Semantics

- **nil** --- a value belong to every record type.
- Scope rule --- similar to PASCAL, Algol ---- support nested scope for types, variables, and functions; redeclaration will hide the same name.

```
function f(v : int) =
  let var v := 6
  in print(v);
  let var v := 7 in print(v) end;
  print(v);
  let var v := 8 in print(v) end;
  print(v)
end
```

- Support two different **name space**: one for types, and one for functions and variables. You can have a type called `foo` and a variable `foo` in scope at same time.

An Example

```
(* A program to solve the 8-queens problem, see Appel's book *)

let
  var N := 8

  type intArray = array of int
  var row := intArray [ N ] of 0
  var col := intArray [ N ] of 0
  var diag1 := intArray [N+N-1] of 0
  var diag2 := intArray [N+N-1] of 0

  function printboard() =
    (for i := 0 to N-1
     do (for j := 0 to N-1
        do print(if col[i]=j then " O" else " .");
        print("\n"));
    print("\n"))

  function try(c:int) =
    (* for i:= 0 to c do print("."); print("\n"); flush(); *)
    if c=N then printboard()
    else for r := 0 to N-1
        do if row[r]=0 & diag1[r+c]=0 & diag2[r+7-c]=0
            then (row[r]:=1; diag1[r+c]:=1; diag2[r+7-c]:=1;
                col[c]:=r; try(c+1);
                row[r]:=0; diag1[r+c]:=0; diag2[r+7-c]:=0)

in try(0)
end
```