

STL Notes

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July 2, 2019

1 STL for Hybrid Systems

1.1 Hybrid Systems Preliminaries

Definition 1.1 (Hybrid System). A hybrid system $\mathcal{H} = (C, F, D, G)$ is defined by a flow set C , flow map F , jump set D , and jump map G , following [3]. $x \in \mathcal{X}$ is the system's state where $\mathcal{X} \subset \mathbb{R}^n$ is the state space. Continuous dynamics are defined on flow set C and are captured by the set valued map $F : \mathcal{X} \rightrightarrows \mathcal{X}$. Discrete dynamics are defined on jump set D and are captured by the set-valued map $G : \mathcal{X} \rightrightarrows \mathcal{X}$. Taken together, \mathcal{H} can be written as:

$$\begin{aligned} \dot{x} &\in F(x) & x &\in C \\ x^+ &\in F(x) & x &\in D \end{aligned} \tag{1}$$

Solutions to a hybrid system \mathcal{H} are defined on ordinary and discrete time:

1. Ordinary Time: $t \in \mathbb{R}_{\geq 0} := [0, \infty)$
2. Discrete Time: $j \in \mathbb{N} := \{0, 1, 2, \dots\}$

Definition 1.2 (Hybrid Time Instance). A *hybrid time instance* is given by

$$(t, j) \in \mathbb{R}_{\geq 0} \times \mathbb{N}$$

Definition 1.3 (Compact Hybrid Domain). A set $E \subset \mathbb{R}_{\geq 0} \times \mathbb{N}$ is a *compact hybrid domain* if it can be written as

$$E := \bigcup_{j=0}^J ([t_j, t_{j+1}] \times \{j\}) \tag{2}$$

for a finite sequence of times $0 = t_0 \leq t_1 \leq t_2 \leq \dots \leq t_{J+1}$.

Definition 1.4 (Hybrid Domain). A set $E \subset \mathbb{R}_{\geq 0} \times \mathbb{N}$ is a *hybrid domain* if for each $(T, J) \in E$

$$E \cap ([0, T] \times \{0, 1, \dots, J\})$$

is a compact hybrid time domain.

Definition 1.5 (Compact Hybrid Shift). Given a compact hybrid domain E and (t^*, j^*) , the forward compact hybrid shift of E by (t^*, j^*) is

$$(t^*, j^*) + E = \bigcup_{j=0}^J ([t_j + t^*, t_{j+1} + t^*] \times \{j + j^*\}) \tag{3}$$

and the backward hybrid shift of E by (t^*, j^*) is

$$E - (t^*, j^*) = \bigcup_{j=0}^J ([t_j - t^*, t_{j+1} - t^*] \times \{j - j^*\}) \tag{4}$$

for a finite sequence of times $0 = t_0 \leq t_1 \leq t_2 \leq \dots \leq t_{J+1}$ and E satisfying (2).

Definition 1.6 (Hybrid Interval). A *hybrid interval* is defined by two hybrid time instances (t_A^*, j_A^*) and (t_B^*, j_B^*) , where $t_A^* + j_A^* \leq t_B^* + j_B^*$, and an hybrid domain E , over the range $[t_A^*, t_B^*] \times \{j_A^*, \dots, j_B^*\}$.

Definition 1.7 (Compact Hybrid Interval). A hybrid interval is a *compact hybrid interval* if it has the form

$$\mathcal{I} := (t_A^*, j_A^*) + E = \bigcup_{j=0}^J ([t_j + t_A^*, t_{j+1} + t_A^*] \times \{j + j_A^*\}) \quad (5)$$

- where $J = j_B^* - j_A^*$ and $t_0 = 0 \leq t_1 \leq t_2 \leq \dots \leq t_{J+1} = t_B^* - t_A^*$.

Definition 1.8 (Open Hybrid Interval). A hybrid time interval is an *open hybrid interval* if it has the form of the following cases:

1. Left open interval

$$\begin{aligned} \mathcal{I} &:= \left((t_A^*, t_1 + t_A^*] \times \{j_A^*\} \right) \cup \left\{ (t_A^* + t_1, j_A^* + 1) + E \right\} \\ &= \left((t_A^*, t_1 + t_A^*] \times \{j_A^*\} \right) \cup \left\{ \bigcup_{j=0}^J ([t_j + t_A^* + t_1, t_{j+1} + t_A^* + t_1] \times \{j + j_A^* + 1\}) \right\} \end{aligned} \quad (6)$$

- where $J = j_B^* - j_A^* - 1$ and $t_0 = 0 \leq t_1 \leq t_2 \leq \dots \leq t_{J+1} = t_B^* - t_A^* - t_1$;

2. Right open interval

$$\begin{aligned} \mathcal{I} &:= \left\{ (t_A^*, j_A^*) + E \right\} \cup \left([t_{j_B^*}^*, t_B^*) \times \{j_B^*\} \right) \\ &= \left\{ \bigcup_{j=0}^J ([t_j + t_A^*, t_{j+1} + t_A^*] \times \{j + j_A^*\}) \right\} \cup \left([t_{j_B^*}^*, t_B^*) \times \{j_B^*\} \right) \end{aligned} \quad (7)$$

- where $J = j_B^* - j_A^* - 1$ and $t_0 = 0 \leq t_1 \leq t_2 \leq \dots \leq t_{J+1} = t_{j_B^*}^* - t_A^*$;

3. Open interval

$$\begin{aligned} \mathcal{I} &:= \left((t_A^*, t_1 + t_A^*] \times \{j_A^*\} \right) \cup \left\{ (t_A^* + t_1, j_A^* + 1) + E \right\} \cup \left([t_{j_B^*}^*, t_B^*) \times \{j_B^*\} \right) \\ &= \left((t_A^*, t_1 + t_A^*] \times \{j_A^*\} \right) \cup \left\{ \bigcup_{j=0}^J ([t_j + t_A^* + t_1, t_{j+1} + t_A^*] \times \{j + j_A^* + 1\}) \right\} \cup \left([t_{j_B^*}^*, t_B^*) \times \{j_B^*\} \right) \end{aligned} \quad (8)$$

- where $J = j_B^* - j_A^* - 2$ and $t_0 = 0 \leq t_1 \leq t_2 \leq \dots \leq t_{J+1} = t_{j_B^*}^* - t_A^*$.

Definition 1.9 (Hybrid Arc). A function $\phi : E \mapsto \mathcal{X}$ is a *hybrid arc* if E is a hybrid time domain and for each $j \in \mathbb{N}$, the function $t \mapsto \phi(t, j)$ is absolutely continuous on the interval $I_j := \{t : (t, j) \in E\}$. The interior of I^j is denoted $\text{int } I^j$.

Definition 1.10 (Solution of a Hybrid System). A hybrid arc ϕ is a solution to hybrid system $\mathcal{H} = (C, F, D, G)$ if

1. $\phi(0, 0) \in \overline{C} \cup D$
2. for all $j \in \mathbb{N}$ such that I^j has a nonempty interior
 - $\phi(t, j) \in C$ for all $t \in \text{int } I^j$
 - $\dot{\phi}(t, j) \in F(\phi(t, j))$ for almost all $t \in I^j$
3. for all $(t, j) \in \text{dom } \phi$ such that $(t, j + 1) \in \text{dom } \phi$
 - $\phi(t, j) \in D$
 - $\phi(t, j + 1) \in G(\phi(t, j))$

1.2 STL Definitions

Definition 1.11. (Atomic Proposition)

An atomic proposition $p^\mu : \mathcal{X} \mapsto \mathbb{B}$ is a function that maps the state space of the system to a boolean value. The set of all atomic propositions is denoted by \mathcal{P} .

The function $\mu : \mathcal{X} \mapsto \mathbb{R}$ represents a robustness measure of the proposition and has following relation to an atomic proposition.

$$\begin{aligned}\mu > 0 &\Leftrightarrow p^\mu = 1 \\ \mu \leq 0 &\Leftrightarrow p^\mu = 0\end{aligned}$$

Definition 1.12 (STL Grammar). STL formulas are defined recursively by the following grammar:

$$\varphi ::= p^\mu \mid \neg\varphi \mid \varphi_1 \wedge \varphi_2 \mid \Diamond_{\mathcal{I}} \varphi \mid \varphi_1 \mathcal{U}_{\mathcal{I}} \varphi_2$$

Where \mathcal{I} is a open or closed hybrid interval and $\varphi, \varphi_1, \varphi_2$ are STL formulae.

The fact that a hybrid arc satisfies STL formula φ at hybrid time instance (t, j) is given by $\varphi(\phi(t, j)) = 1$. When a hybrid arc does not satisfy a proposition then $\varphi(\phi(t, j)) = 0$.

The validity of a formula φ with respect to hybrid arc ϕ at at time (t, j) is defined inductively as

$$\begin{aligned}(\phi, (t, j)) \models p^\mu &\iff \mu(\phi(t, j)) > 0 \\ (\phi, (t, j)) \models \varphi &\iff \varphi(\phi(t, j)) = 1 \\ (\phi, (t, j)) \models \neg\varphi &\iff \neg((\phi, (t, j)) \models \varphi) \\ (\phi, (t, j)) \models \varphi \wedge \psi &\iff (\phi, (t, j)) \models \varphi \wedge (\phi, (t, j)) \models \psi \\ (\phi, (t, j)) \models \Diamond_{\mathcal{I}} \varphi &\iff \exists(t', j') \in (t, j) + \mathcal{I} \text{ s.t. } (\phi, (t', j')) \models \varphi \\ (\phi, (t, j)) \models \psi \mathcal{U}_{\mathcal{I}} \varphi &\iff \exists(t', j') \in (t, j) + \mathcal{I} \text{ s.t. } (\phi, (t', j')) \models \varphi \wedge \forall(t'', j'') \in ((t, j) + \mathcal{I}) \cap ([t, t'] \times [j, \dots, j']), (\phi, (t'', j'')) \models \psi\end{aligned}$$

References

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