

STL Notes

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1 Autonomous Vehicles

1.1 Intersection Management

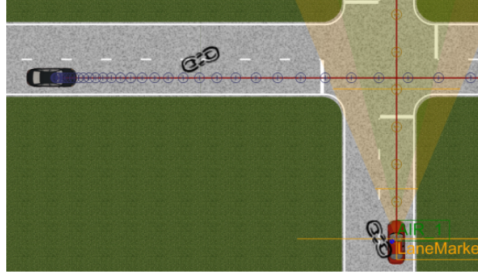


Figure 1: Intersection Scenario

In [?] the ego-vehicle's behavior at an intersection is defined using STL. The system has a state:

$$\mathbf{x}_t = [x_t^{ego} \ y_t^{ego} \ x_t^{adv} \ y_t^{adv} \ v_t^{ego} \ v_t^{adv}]^T \quad (1)$$

Figure 1 shows the ego-vehicle in blue and the adversary-vehicle in red. The authors of [?] use the following specification to describe the behavior: ϕ_s : "When the ego-vehicle is within 2 meters of the adversary-vehicle the ego-vehicle will stop for two seconds."

$$\phi_s = \Box(|x_t^{ego} - x_t^{adv}| < 2) \implies \Box_{[0,2]}(|v_t^{ego}| < 0.1) \quad (2)$$

1.1.1 Alternate Formulation

The above STL specification relies on the adversary-vehicle to stop at the intersection, but a car should always stop at an intersection. Inspired from [?] the following specification describes: "When the ego-vehicle reaches an intersection it will stop for one second and only cross when safe (ϕ_{cross}) otherwise it will stay stopped. ($\phi_{stopped} \mathcal{U}_{[1,3]} \phi_{stopped}$)."

$$\phi_{stopped} = (|v_t^{ego}| == 0) \quad (3)$$

$$\phi_{cross} = \Box_{[1,3]}(|y_t^{adv} - \mathbf{Y}^{\text{INT}}| > 0.1) \wedge (x_t^{ego} > 5) \quad (4)$$

$$\phi_{int} = \Box(|x_t^{ego} - \mathbf{X}^{\text{INT}}| < 0.1) \implies \Box[0.5, 1](\phi_{stopped}) \wedge (\phi_{stopped} \mathcal{U}_{[1,3]} \phi_{cross} \vee \phi_{stopped} \mathcal{U}_{[1,3]} \phi_{stopped}) \quad (5)$$

Where $\mathbf{X}^{\text{INT}} = -5$ and $\mathbf{Y}^{\text{INT}} = -5$ are the positional coordinates of the intersection.

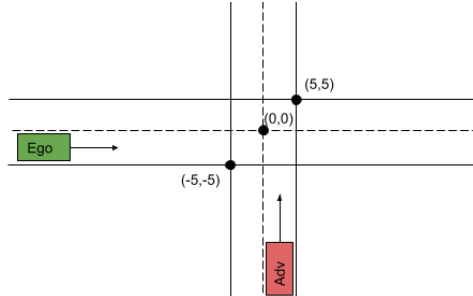


Figure 2: Alternate Scenario

2 STL for Hybrid Systems

2.1 Hybrid Time STL

2.1.1 Hybrid System Solutions

Solutions to hybrid systems are defined on ordinary and discrete time:

1. Ordinary Time: $t \in \mathbb{R}_{\geq 0} := [0, \infty)$
2. Discrete Time: $j \in \mathbb{N} := \{0, 1, 2, \dots\}$

Definition 2.1 (Hybrid Time Instance). A *hybrid time instance* is given by

$$(t, j) \in \mathbb{R}_{\geq 0} \times \mathbb{N}$$

Definition 2.2 (Compact Hybrid Domain). A set $E \subset \mathbb{R}_{\geq 0} \times \mathbb{N}$ is a *compact hybrid domain* if it can be written as

$$E = \bigcup_{j=0}^{J-1} ([t_j, t_{j+1}] \times \{j\}) \quad (6)$$

for a finite sequence of times $0 = t_0 \leq t_1 \leq t_2 \leq \dots \leq t_J$.

Definition 2.3 (Hybrid Domain). A set $E \subset \mathbb{R}_{\geq 0} \times \mathbb{N}$ is a *hybrid domain* if for each $(T, J) \in E$

$$E \cap ([0, T] \times \{0, 1, \dots, J\})$$

is a compact hybrid time domain.

Definition 2.4 (Compact Hybrid Shift). Given a compact hybrid domain E and (t^*, j^*) , the forward compact hybrid shift of E by (t^*, j^*) is

$$(t^*, j^*) + E = \bigcup_{j=0}^{J-1} ([t_j + t^*, t_{j+1} + t^*] \times \{j + j^*\}) \quad (7)$$

and the backward hybrid shift of E by (t^*, j^*) is

$$E - (t^*, j^*) = \bigcup_{j=0}^{J-1} ([t_j - t^*, t_{j+1} - t^*] \times \{j - j^*\}) \quad (8)$$

for a finite sequence of times $0 = t_0 \leq t_1 \leq t_2 \leq \dots \leq t_J$ and E satisfying (??).

Definition 2.5 (Compact hybrid Interval). A *compact hybrid interval* is defined by two hybrid time instances (t_A^*, j_A^*) and (t_B^*, j_B^*) , where $t_A^* + j_A^* \leq t_B^* + j_B^*$, and an hybrid domain E , over the range $[t_A^*, t_B^*] \times \{j_A^*, \dots, j_B^*\}$, by

$$\mathcal{I} := (t_A^*, j_A^*) + E = \bigcup_{j=0}^{J-1} ([t_j + t_A^*, t_{j+1} + t_A^*] \times \{j + j_A^*\}) \quad (9)$$

Where $J = j_B^* - j_A^*$ and $t_0 = 0 \leq t_1 \leq t_2 \leq \dots \leq t_J = t_B^* - t_A^*$

Definition 2.6 (Open Hybrid Interval). An open hybrid interval is defined by two hybrid time instances, (t_A^*, j_A^*) and (t_B^*, j_B^*) , on hybrid domain E , over the range $[t_A^*, t_B^*] \times \{j_A^*, \dots, j_B^*\}$, by the following cases:

1. Left open interval

$$\mathcal{I} := (t_A^*, t_1] \times \{j_A^*\} \cup \bigcup_{j=1}^{J-1} ([t_j + t_A^*, t_{j+1} + t_A^*] \times \{j + j_A^*\}) \quad (10)$$

Where $J = j_B^* - j_A^*$ and $t_0 = 0 \leq t_1 \leq t_2 \leq \dots \leq t_J = t_B^* - t_A^*$

2. Right open interval

$$\mathcal{I} := \bigcup_{j=0}^{J-2} ([t_j + t_A^*, t_{j+1} + t_A^*] \times \{j + j_A^*\}) \cup [t_{J-1}, t_B^*) \times \{j_B^*\} \quad (11)$$

Where $J = j_B^* - j_A^*$ and $t_0 = 0 \leq t_1 \leq t_2 \leq \dots \leq t_J = t_B^* - t_A^*$

Remark: Authors of [?] denote a ordinary time interval \mathcal{I} , relative to ordinary time t as:

$$t + \mathcal{I} := \{t + t' : t' \in \mathcal{I}\}$$

The same notation can be used denoting a hybrid temporal interval, \mathcal{I} relative to hybrid time instance (t, j) :

$$(t, j) + I := \{(t + t', j + j') : (t', j') \in I\}$$

2.2 Continuous Time STL

A STL formula ϕ is defined recursively as:

$$\psi ::= p^\mu \mid \neg\psi \mid \psi_1 \wedge \psi_2 \mid \Diamond_{\mathcal{I}}\psi \mid \psi_1 U_{\mathcal{I}}\psi_2 \quad (12)$$

Where ψ_1, ψ_2 are STL formula and \mathcal{I} is an ordinary time interval of form $\mathcal{I} = (t_a, t_b), (t_a, t_b], [t_a, t_b)$, or $[t_a, t_b]$ with $t_a < t_b$ and $\mathcal{I} \subset \mathbb{R}_{\geq 0}$.

Inductive STL Definition

$$\begin{aligned} \phi &\models \psi && \Leftrightarrow (\phi, t) \models p \\ (\phi, t) &\models p^\mu && \Leftrightarrow \mu(\phi(t)) > 0 \\ (\phi, t) &\models \Box_{[a,b]}\psi && \Leftrightarrow \forall t' \in [t+a, t+b] \text{ s.t. } (\phi, t') \models \psi \end{aligned}$$

2.2.1 Possible Notion of Hybrid STL

Recursive STL Definition

$$\psi ::= p^\mu \mid \neg\psi \mid \psi_1 \wedge \psi_2 \mid \Diamond_I \psi \mid \psi_1 U_I \psi_2 \quad (13)$$

Where ψ_1, ψ_2 are STL formula and I is a hybrid time interval.

Inductive STL Definition

$$\begin{aligned} \phi &\models \psi && \Leftrightarrow (\phi, (t, j)) \models p \\ (\phi, (t, j)) &\models p^\mu && \Leftrightarrow \mu(\phi(t)) > 0 \\ ((\phi, (t, j)) &\models \Box_I \psi && \Leftrightarrow \forall (t', j') \in (t, j) + I \text{ s.t. } (\phi, (t', j')) \models \psi \end{aligned}$$

2.2.2 Previous Notions of Hybrid STL

From Atreya - Modeling Hybrid Systems using Signal Temporal Logic

Inductive STL Definition

$$((\phi, (t, j)) \models \Box_{(t, j)} \psi \Leftrightarrow \forall (t', j') \in \text{dom } \phi, t' + j' > t + j \text{ s.t. } (\phi, (t', j')) \models \psi \quad (14)$$

.1 Lane Changing

The authors of [?] use STL and Hamilton Jacobi reachability to combine temporal specification with quantification of a reachable set: STL temporal specifications are related to reachability operators.

.1.1 Example

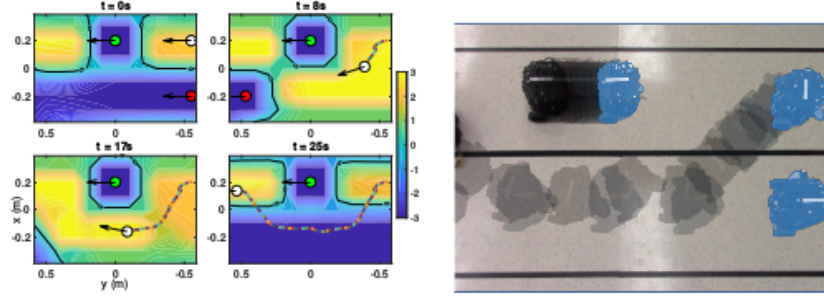


Figure 3: Lane Change Scenario

In Figure 3 white represents the autonomous car, green a slow car, and red an adjacent car. The white car's lane change specification is given as: "Within 25 seconds, satisfy general safety ϕ until you can pass. Otherwise stay within the lane. Also ensure you are always within 5 seconds from re-entering the lane. (ψ^{lane})."

$$\phi = \phi^{off-road} \wedge \phi^{on-road} \wedge \phi^{avoid} \quad (15)$$

$$\Box_{[0,25]} \Diamond_{[0,5]} \psi^{lane} \wedge (\phi \mathcal{U}_{[0,25]} \psi^{pass} \vee \phi \mathcal{U}_{[0,25]} \psi^{stay}) \quad (16)$$

.2 Solutions for Discrete Systems

Consider a discrete system of the form:

$$x^+ \in G(x, \gamma) \quad (17)$$

Let $x \in \mathbb{R}^n$ be states of the discrete system.

Let $\gamma \in \mathbb{R}^m$ be input to the discrete system.

Discrete time is denoted: $k \in \mathbb{N}$

Then, given an initial state, x_0 and an input $k \mapsto \gamma(k)$, a state trajectory(solution) can be constructed. The domain of γ is given: $dom \gamma = \{0, 1, 2, \dots, K\} \cap \mathbb{N}$ where $K \in \mathbb{N} \cup \{\infty\}$. Solutions are defined as: $k \mapsto \phi(k)$ such that the following are satisfied:

1. $\phi(0) = x_0$
2. $dom \phi = dom \gamma$ if $K = \infty$
3. $dom \phi \setminus \{max dom \phi\} = dom \gamma$ if K is finite
4. $\forall k \in dom \gamma: \phi(k+1) = G(\phi(k), \gamma(k))$

.2.1 Discrete Time STL

$$\psi ::= p^\mu \mid \neg \psi \mid \psi_1 \wedge \psi_2 \mid \Diamond_I \psi \mid \psi_1 U_I \psi_2 \quad (18)$$

Where ϕ_1, ϕ_2 are STL formula and I is an interval on $dom \phi$.

Inductive STL Definition

$$\begin{aligned} \sigma &\models \psi && \Leftrightarrow (\sigma, k) \models p \\ (\sigma, k) &\models p^\mu && \Leftrightarrow \mu(x_k, \gamma_k) > 0 \\ (\sigma, k) &\models \Box_{[n,m]} \psi && \Leftrightarrow \forall k' \in [k+n, k+m] \text{ s.t } (\sigma, k') \models \psi \end{aligned}$$