Runtime Verification for Trustworthy Secure Shell Deployment

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Cryptographic protocol concerns

- Standard practice to establish provable security guarantees of a protocol (e.g., Dolev-Yao model)
- However, the execution of the protocol might still be insecure
 - Protocol implementation might not adhere to the protocol specifications
 - Incorrect implementation of protocol steps
 - Missing checks
 - Malware might interfere with the execution of security-critical code

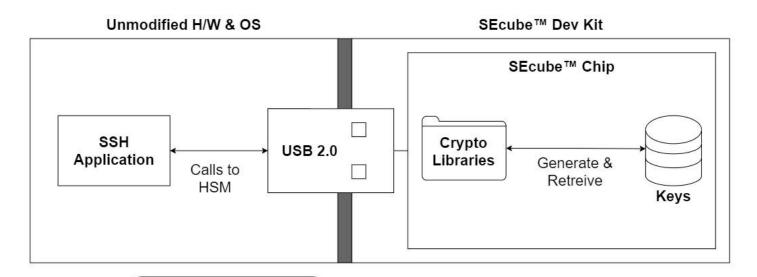
Secure Shell (SSH)

- Adapted as an internet standard in 2006
- Complements existing work on the TLS 1.3 protocol
- Widely popular in the industry
- Some features:
 - Remote login and command execution
 - Virtual private networks (VPNs)
 - Port forwarding
 - Secure file transfer

Trusted Execution Environments

- Isolates security-critical code from untrusted code (potentially compromised by malware)
- In practice, TEEs are implemented as **CPU modes** offering non-addressable, reserved, and encrypted memory pages
 - Intel SGX
 - AMD SEV/SME
 - ARM TrustZone

The proposed solution: RV-TEE





RV Monitor

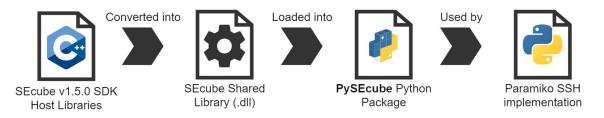
Paramiko SSHv2 protocol code function call tracing and memory leak protection

Protocol implementation

- The chosen SSH implementation is Paramiko, having as of April 2021:
 - 11.4K dependant repositories
 - 898 dependant packages (Docker Python SDK, Ansible and Apache Airflow)
- Uses cryptography Python library for cryptographic operations
 - Based on the OpenSSL implementation

Utilisation of HSM in protocol implementation

- The protocol implementation source code is modified such that:
 - The HSM is utilised for cryptographic operations, rather than *cryptography*
 - Ephemeral cryptographic keys are stored on the HSM, rather than kept in memory
- Changes involving cryptographic operations:
 - Hashing during key exchange and key derivation (SHA256)
 - Encryption and decryption of protocol messages (AES256)
 - Message authentication code generation (HMACSHA256)



Property derivation and specification

- 17 properties were systematically derived from RFCs
 - Focused on RFC standard keywords, e.g. "MUST" and "SHOULD"
- Focused on the client side of the protocol
 - Weaker security
 - SEcube device used is aimed towards end-users.

Property Category	Number of Properties
Temporal	11
Point assertion	5
Real-time	1

Example temporal property

"When a **KEXDH_REPLY** message is received from the server, the client **must verify** the **public host key** with the **signature of the hash obtained**"

Violation of the property leads to:

- Vulnerability to active man-in-the-middle attacks
- Client-to-attacker and attacker-to-server connections are established

High-level RV deployment

- Assert the executed protocol steps conform with the specification
- Instrumentation
 - Monkey-patching used for instrumenting Python code
 - In-line hooking used for compiled code
- Properties are manually modelled into RV monitors using LARVA
 - Automata-based approach
- Initial offline RV configuration
 - Instrumentation is limited to function call tracing
 - Events are replayed and monitored after the execution of the protocol

Low-level RV deployment

- Monitors data flows across the trust boundary
- Detect malicious interactions with the HSM
- Assert sensitive data is not leaked from the HSM

Future work

- Analyse the impact of the introduced overheads, in the context of SSH
- Deploy online RV consisting of
 - Synchronous monitors for basic protocol sequence properties
 - Asynchronous monitors for memory-based and processor-intensive properties
- Experiment with other software and hardware configurations