

Backtracking

We write a function that determines whether the values of an `int` array can be placed into two bags whose sums are equal. For example, for `b` containing `{3, 1, 5, 4, 1, 2}` we can place them in two bags `{3, 5}` and `{1, 4, 1, 2}`, which both sum to 8. But putting the values `{4, 2, 1}` into two bags will always yield a bag with an even sum and a bag with an odd sum because there is only one odd value.

```
/** = "elements of b can be placed into
 * 2 bags whose sums are equal." */
static boolean split(int[] b)
```

Now, look at the specification of `split`. The spec doesn't require us to construct the bags in data structures! We need *only* calculate *the sum of the values in each bag*! This will simplify everything tremendously.

Function `split` has only array `b` as parameter, so we think of writing another function `split` that will have more parameters and be recursive. We envision a recursive function in which at each depth of recursion one more value is put into one of the bags. This leads us to the specification shown (below). Values in `b[0..k-1]` have been placed in the two bags and their sums are in parameters `s1` and `s2`. Return `true` if the rest of the array values, starting at `b[k]`, can be placed in the bags so that their sums are equal; otherwise, return `false`.

With 4-parameter method `split` specified, we can complete the original method `split` with the body that indicates that no values have been placed in the bags, so the values in the two bags sum to 0.

```
/** The values of b[0..k-1] have been placed into two
 * bags that sum to s1 and s2, respectively. Return true
 * iff the rest of the values b[k..] can be placed into the
 * two bags so that the two bags have the same sum. */
static boolean split(int[] b, int k, int s1, int s2)
```

```
/** = "elements of b can be placed into
 * 2 bags whose sums are equal." */
static boolean split(int[] b) {
    return split(b, 0, 0, 0);
}
```

At this point, stop and study the specification of the four-parameter `split` carefully. Make sure you understand it, as well as the call on it in the one-parameter function `split`.

We are now ready to write the body of 4-parameter method `split`.

The base case is easy. If `k` equals `b.length`, then all values have been placed in the bags, and `true` or `false` can be returned depending on whether the two bags have the same sum or not.

```
Base case:
if (k == b.length) return s1 == s2;
```

For the recursive case, we have two choices: place `b[k]` in the first bag and place `b[k]` in the second bag. (These are given by the two calls given to the right.) Which one do we do? It doesn't really matter.

```
Choice 1: split(b, k+1, s1 + b[k], s2)
Choice 2: split(b, k+1, s1, s2 + b[k])
```

Suppose we call the first choice: put `b[k]` in `s1`. If this choice returns `true`, `true` can be returned. But if this choice returns `false`, the other choice must be called and its value returned.

Thus, we write a single return statement, which returns the value of the first one OR the value of the second one. Note: If the first one returns `true`, short circuit evaluation means that its value will be returned, without calling the second one, but if the first one returns `false`, the value of the second one is returned.

Isn't that a neat little algorithm?

```
/** = The values of b[0..k-1] have been placed into two
 * bags that sum to s1 and s2, respectively. Return true
 * iff the rest of the values b[k..] can be placed into the
 * two bags so that the two bags have the same sum. */
static boolean split(int[] b, int k, int s1, int s2) {
    if (k == b.length) return s1 == s2;
    return split(b, k+1, s1 + b[k], s2)
        || split(b, k+1, s1, s2 + b[k]);
}
```