

## Practice with the second loop question

### Introduction

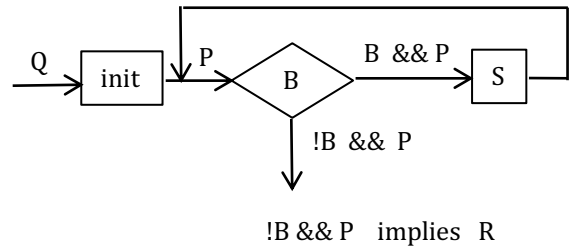
We practice finding a loop condition B by using the second loop question: Is  $!B \ \&\& \ P \Rightarrow R$  true? Thus, we look for  $!B$  that makes  $!B \ \&\& \ P \Rightarrow R$  true and complement  $!B$  to get B.

Here's the invariant P and postcondition for our first example.

P: s is the sum of  $m..k-1$  and  $m \leq k \leq n$   
R: s is the sum of  $m..n-1$

Knowing that P is true, and doing some pattern matching with P and R, we see that R will be true if  $k = n$ . Therefore,  $!B$  is  $k = n$ , so the loop condition B is  $k \neq n$ . Looking at the restriction on k in invariant P, we can write the loop condition at  $k < n$  if we want. Thus, we use either

**while** ( $k \neq n$ ) { ... }      or      **while** ( $k < n$ ) { ... }



### A second example

Here are the invariant and postcondition for a loop to calculate the minimum value in array segment  $b[0..n-1]$ :

P:  $v = \text{minimum of } b[0..k-1]$  and  $0 \leq k \leq n$   
R:  $v = \text{minimum of } b[0..n-1]$

Using reasoning like we did the first example, you can see that we get the same answer for B as in the previous example.

**while** ( $k \neq n$ ) { ... }      or      **while** ( $k < n$ ) { ... }

### Computing $z = b^c$

Here are the invariant and postcondition for a loop to store  $b^c$  in z, given  $c \geq 0$ :

P:  $b^c = z * x^y$  and  $y \geq 0$   
R:  $z = b^c$

Again doing pattern matching, we see that R will be true when P is true and  $x^y = 1$ . That last formula,  $x^y = 1$ , is true, when  $y = 0$ . So our loop condition is  $y \neq 0$ :

**while** ( $y \neq 0$ ) { ... }

### Exercises

In the two examples below, find the loop condition. Answers are at the end of the pdf script for this video.

1. P: s is the sum of  $k..n-1$  and  $m \leq k \leq n$   
R: s is the sum of  $m..n-1$
2. P:  $v = \text{minimum of } b[k..n]$  and  $0 \leq k \leq n$   
R:  $v = \text{minimum of } b[0..n]$  and  $0 \leq k \leq n$

### Answers

In the first exercise, doing pattern matching on P and R, we see that  $k = m$  is needed. Therefore the loop condition is  $k \neq m$ . This can be written as  $m < k$  if you want, since  $m \leq k \leq n$ :

**while** ( $k \neq m$ ) { ... }      or      **while** ( $m < k$ ) { ... }

In the second exercise, pattern matching on P and R, we see that  $k = 0$  is needed. Therefore, the loop condition is  $k \neq 0$ . This can be written as  $0 < k$  if you want, since  $0 \leq k \leq n$ :

**while** ( $k \neq 0$ ) { ... }      or      **while** ( $0 < k$ ) { ... }