Formal Semantics

The structure of a language defines its syntax, but what defines <u>semantics</u> or meaning?

⇒ Behavior!

The most straight forward way to define semantics is to provide a <u>simple interpreter</u> for the programming language that highlights the behavior of the language,

⇒ Operational Semantics

Operational Semantics

Let's develop an operational semantics for a simple programming language called *ONE*;

```
ONE: \langle exp \rangle^* ::= \langle exp \rangle + \langle mulexp \rangle \mid \langle mulexp \rangle

\langle mulexp \rangle ::= \langle mulexp \rangle * \langle rootexp \rangle \mid \langle rootexp \rangle

\langle rootexp \rangle ::= (\langle exp \rangle) \mid \langle constant \rangle

\langle constant \rangle ::= \text{all valid integer constants}
```

Note: The grammar is unambiguous, both precedence and associativity rules of "standard" arithmetic are observed.

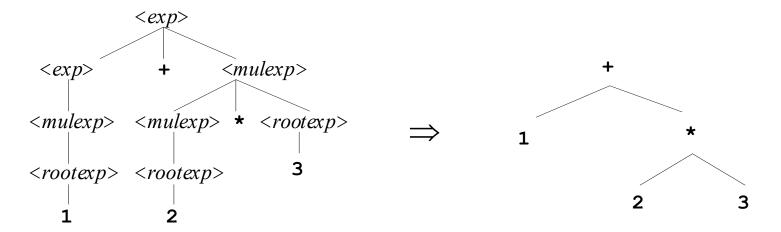
Do the following sentences belong to L(ONE)? Why? Why not?

$$s = 1 + 2 * 3$$

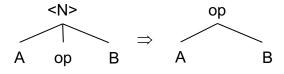
 $s = (1 + 2) * 3$
 $s = a + 3$

Abstract Syntax Trees

We want to define an operational semantics, i.e., an abstract interpreter for the language, but parse trees are not very convenient, too many non-terminal symbols ⇒ <u>Abstract Syntax Tree</u> (AST)



Transformation Rules:



Note: This rule also applies to unary operators and operators with arity > 2.

Observations

<u>Definition</u>: An abstract syntax tree is a finite, labeled, directed tree, where the internal nodes are labeled by operators, and the leaf nodes represent the operands of the node operators. -Wikipedia, 2006

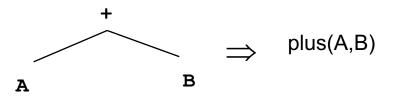
Observation: The abstract syntax tree is a simplified form of the parse tree: same order as the parse tree, but no non-terminals.

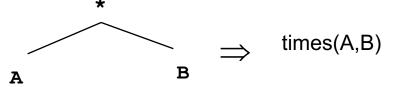
ASTs & Parentheses

- What happens to parentheses in the AST representation of a program?
- They are <u>not needed!</u>
- ASTs naturally represent associativity and precedence relations.
- Consider: (1 + 2) * 3
- Parentheses do not contribute to computations, therefore the following tree transformations can be applied:

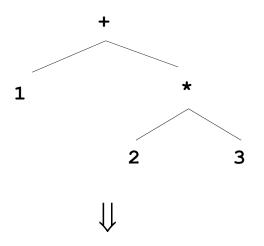
Prolog ASTs

We can represent ASTs in Prolog:





c (constant) \Longrightarrow const(c)



plus(const(1),times(const(2),const(3)))

ONE: Prolog Interpreter

A simple interpreter that computes a semantic value for syntactic constructs, the computation of this semantic value can be interpreted as the behavior: val1 / 2, AST input and semantic value as output.

```
?- val1(const(1),X).

X = 1

Yes
?- val1(plus(const(1),const(2)),X).

X = 3

Yes
?- val1(plus(const(1),times(const(2),const(3))),X).

X = 7

Yes
```

Exercises

- Extend the grammar for language ONE with the subtraction operator
- Extend the operational semantics appropriately, e.g.,
 - 6 3 should give the value 3

Assume that the abstract syntax of this operator is sub(x,y).

- Compute the semantic value for the following expressions:
 - sub(3,1)
 - sub(4,2)

Exercises

- Extend the grammar for language ONE with the '!' factorial operator
- Extend the operational semantics appropriately, e.g.,
 - 3! should give the value 6

Assume that the abstract syntax of this operator is fact(x).

- Compute the semantic value for the following expressions:
 - fact(3)
 - fact(4)