Université d'Ottawa Faculté de génie

University of Ottawa Faculty of Engineering



École de science informatique et de génie électrique (SIGE)

School of Electrical Engineering and Computer Science (EECS)

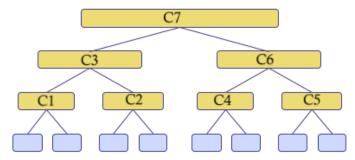
CSI2110 (Fall 2022) **Assignment 11 (2.5%) – 10 points**

Due: Wednesday Dec 7, 11:59PM

Late assignment policy: for A11 there is a grace period until Sunday Dec 10, 11:59 with no penalty.

Question 1 (Mergesort) (3.5 points)

Consider the code for mergesort given in the appendix. Below is its recursion tree for array of length 8:



Consider the mergesort algorithm applied to the following array of length 8:

12	2	9	4	10	6	15	8	

The array of 8 numbers is constinuously updated with the recursive calls. Show what will be the array at the **END** of each recursive call numbered C1 to C7 in the recursion tree above.

C1:	2	12	9	4	10	6	15	8
C2:	2	12	4	9	10	6	15	8
C3:	2	4	9	12	10	6	15	8
C4:	2	4	9	12	6	10	15	8
C5:	2	4	9	12	6	10	8	15
C6:	2	4	9	12	6	8	10	15
C7:	2	4	6	8	9	10	12	15

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Question 2 (Quicksort) (3.5 points=2.5+1)

Consider the code for quicksort, taken from the textbook by Goodrich, Tamassia and Goldwasser 6th ed.

```
/** Sort the subarray S[a..b] inclusive. */
 2
      private static <K > void quickSortInPlace(K[] S, Comparator<K > comp,
                                                                                               JOHANA
 3
                                                                          int a, int b) {
 4
        if (a >= b) return;
                                  // subarray is trivially sorted
 5
        int left = a;
 6
        int right = b-1;
        K pivot = S[b];
 8
                                  // temp object used for swapping
        K temp;
        while (left <= right) {
10
          // scan until reaching value equal or larger than pivot (or right marker)
11
          while (left \leq right && comp.compare(S[left], pivot) \leq 0) left++;
12
          // scan until reaching value equal or smaller than pivot (or left marker)
13
          while (left \leq right && comp.compare(S[right], pivot) > 0) right—;
          if (left <= right) { // indices did not strictly cross
14
15
            // so swap values and shrink range
16
            temp = S[left]; S[left] = S[right]; S[right] = temp;
17
            left++; right--;
18
19
        // put pivot into its final place (currently marked by left index)
20
21
        temp = S[left]; S[left] = S[b]; S[b] = temp;
22
        // make recursive calls
23
        quickSortInPlace(S, comp, a, left -1);
24
        quickSortInPlace(S, comp, left + 1, b);
25
    Code Fragment 12.6: In-place quick-sort for an array S. The entire array can be
```

Consider the following array S in the call quickSortInPlace (S, comp, 0, 7)

```
S: 12 2 9 4 10 6 15 8
```

sorted as quickSortInPlace(S, comp, 0, S.length-1).

Show what will be the state of the array at the end of the first partition (at line 22). Before showing this final array, show intermediate arrays after each swap of S[right] and S[left] in the partition (line 16).

a) arrays after each swap:

```
6,2,9,4,10,12,15,8
6,2,4,9,10,12,15,8
6,2,4,8,10,12,15,9
```

b) array after partition:

6	2	4	8	10	12	15	9

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Question 3 Bucketsort and Radixsort (3 points)

Consider the following variation of radix-sort

```
Algorithm radicchioSort(S)
        Input sequence S of b-bit integers
        Output sequence S sorted
       for i \leftarrow 0 to b - 1
               // use as the key k of each item x of S the bit x_i of x=(x_{b-1}...x_i x_0)
                bucketSort(S, 2,i); /***/
```

Conisder b=3 and the following array: [7, 4, 1, 5, 3, 2]

Show the array after each call of bucketSort.

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Appendix: Mergesort algorithm

```
public class MergeSort{
                                                          Ressorts Old Ottawa
 private static void merge(Comparable[] a, Comparable[] aux, int lo, int mid, int hi){
    for (int k = lo; k \le hi; k++)
                                    aux[k] = a[k];
    int i = lo, j = mid+1;
    for (int k = lo; k \le hi; k++)
          (i > mid)
     if
                            a[k] = aux[j++];
     else if (j > hi)
                            a[k] = aux[i++];
     else if (less(aux[i], aux[i])) a[k] = aux[i++];
     else
                          a[k] = aux[i++];
 private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi)
   if (hi <= lo) return;
   int mid = lo + (hi - lo) / 2;
   sort(a, aux, lo, mid);
   sort(a, aux, mid+1, hi);
   merge(a, aux, lo, mid, hi);
 public static void sort(Comparable[] a)
   Comparable[] aux = new Comparable[a.length];
   sort(a, aux, 0, a.length - 1);
```

Code from: "Algorithms", Robert Sedgewick and Kevin Wayne https://www.cs.princeton.edu/courses/archive/spr14/cos226/lectures/22Mergesort.pdf