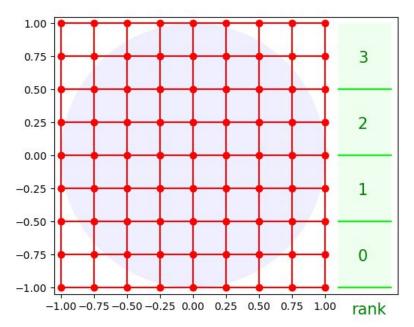
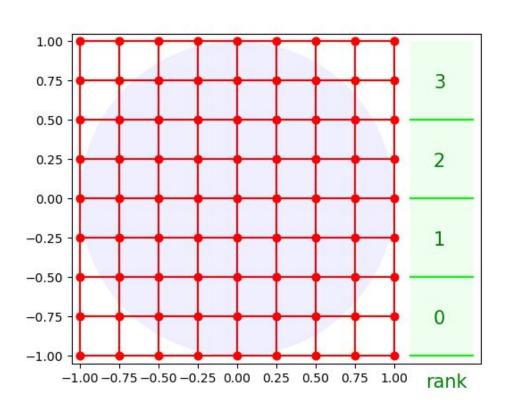
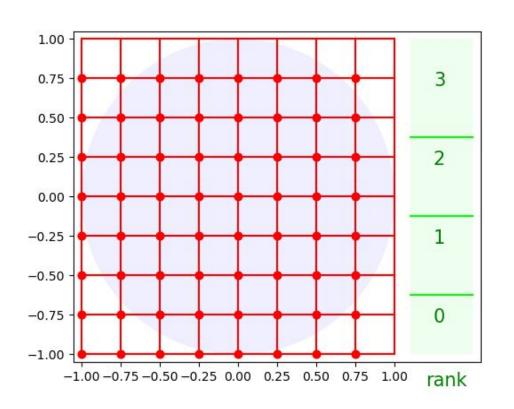
### Computação Paralela / Computação Avançada

Helmut Wolters cap4 - 2022-11-02

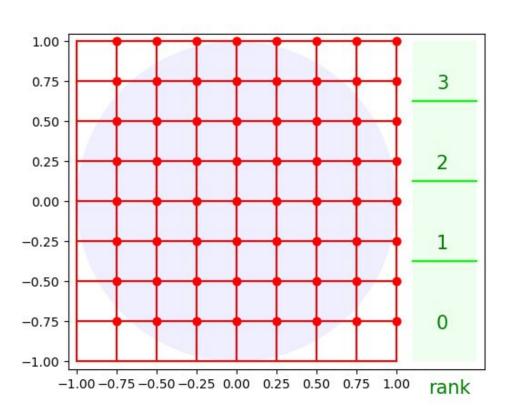




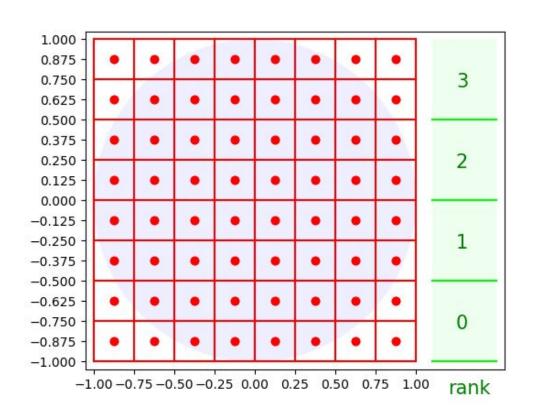
- We have an area of 1.0 x 1.0
- We want to divide it into a 8x8 grid
- This corresponds to a division into 64 0.25 x 0.25 squares
- Subdivide x in
   -1, -0.75, ..., 0.75, 1
- This gives 9 numbers, not 8!
- When we subdivide the y axis for distribution over 4 parallel processes, we are likely to repeat the y values on the intervals' borders
- This would be a **systematic error**



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- We want to divide it into a 8x8 grid
- This corresponds to a division into 64 0.25 x 0.25 squares
- Subdivide x in
  -1, -0.75, ..., 0.75, 1
  This gives 9 numbers, not 8!
- So we take away the highest values for x and y ...
- When dividing in subintervals, proceed the same way



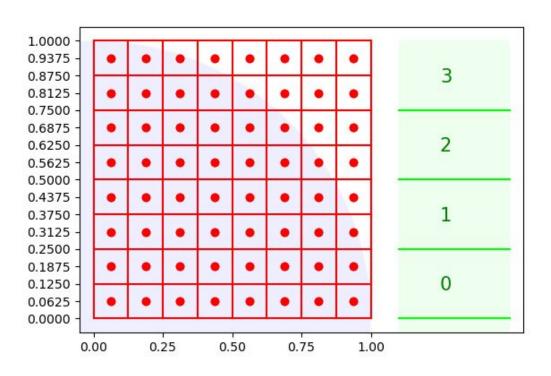
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- We want to divide it into a 8x8 grid
- This corresponds to a division into 64 0.25 x 0.25 squares
- Subdivide x in
  -1, -0.75, ..., 0.75, 1
  This gives 9 numbers, not 8!
- Or we take away the lowest values for x and y ...
- No repetition any more, but we still will make a systematic error shifting the values either up or down half of a square width/height



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- We want to divide it into a 8x8 grid
- This corresponds to a division into 64 0.25 x 0.25 squares
- Subdivide x in
  -1, -0.75, ..., 0.75, 1
  This gives 9 numbers, not 8!

#### • Correct solution:

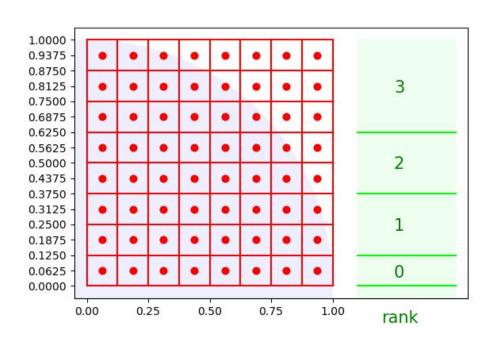
The discret x and y values are the intermediate points in the intervals



#### Correct solution:

The discret x and y values are the intermediate points in the intervals

- This also makes sure that we do not repeat values when we consider only the upper right quadrant of the square and the circle
- In the "optimized" version
   (sheet 3, problem 4)
   we make one run with small n and
   measure the time each process needs.
- We will observe that 0 is slowest,
   1 a bit faster, 2 still a bit faster,
   and 3 much faster.
- We redistribute the share of each process, 0 gets less until 3 that gets most

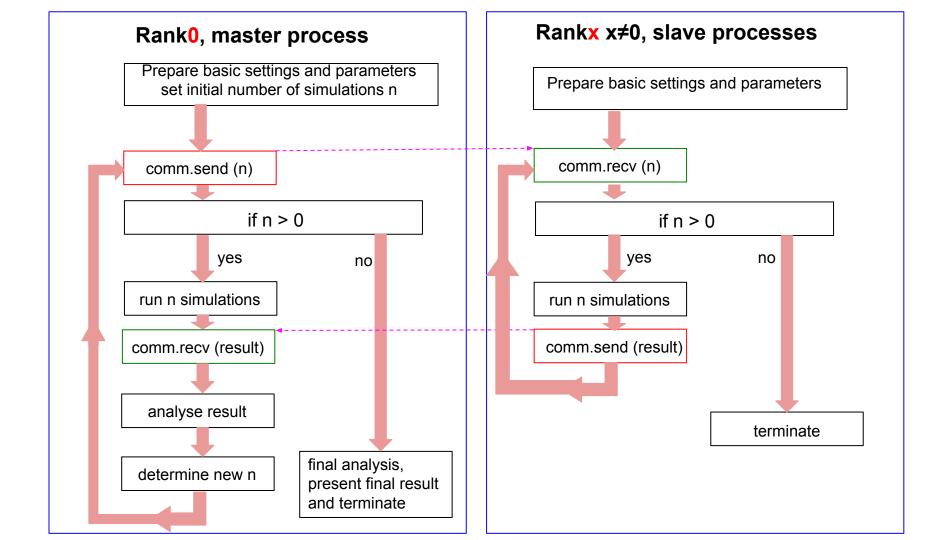


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   and 3 much faster.
- We redistribute the share of each process, 0 gets less until 3 that gets most
- Then we run with the final n.
- How to calculate the share?
   Be creative, just a simple formula:
   Δx<sub>p</sub> proportional to 1/Δt<sub>p</sub>

Where, for each processor p,  $\Delta t_{_{D}}$  is the time needed in the **first run**.

# Computação Paralela / Computação Avançada Basic flow of an iterative parallel process

- Idea: start with n simulations, distributed over m processes
- Verify quality of result
- While quality below expectations, substitute n ⇒ n' = 2\*n
- When result is satisfying, stop the whole process
- Whenever possible, accumulate the results we already obtained, i.e. do not throw away what we have already calculated
- The master controls the process
- The master sends the number of simulations to be done to all other processes
- All other processes wait for this message
- When the message arrives, each process starts the simulations, (as does the master)
- All other processes send their result to the master, and wait for a new message
- The master analyses the quality of result achieved
- If not satisfying, the master doubles n and sends the new number of simulations to be done to all other processes
- etc ...
- When the master sees that the analysis result is satisfying, it sends 0 to all other processes, and terminates
- When the message from the master is "0", the other processes terminate



# Computação Paralela / Computação Avançada Basic flow of an iterative parallel process

- You can use the "while" loop in python to repeat the sequences
- The best approach is a "permanent" while-loop, with a break-statement inside the loop at the point we want to finish. The break-statement continues with the next statement after the loop:

 Be careful that any message a process is waiting for in comm.recv, has a corresponding comm.send from the other process, otherwise you create a deadlock in the waiting process