

Due at the beginning of class on 30 January 2024

- Your answers should be neatly written and logically organized.
- You may collaborate on solving the problems, but the solutions you turn in should be your own.
- You may use any resource you find online (or elsewhere), but you must cite any resource you use.

Reading: Read §2.2 in [Rie14] and §1.5 in [Mal23].

- (1) Let \mathcal{C} be a homotopical category, and let \mathcal{J} be any category. The category $\text{Fun}(\mathcal{J}, \mathcal{C})$ of functors from \mathcal{J} to \mathcal{C} becomes a homotopical category with weak equivalences defined object-wise. By choosing a homotopical category \mathcal{C} and a category \mathcal{J} , show that the limit functor

$$\lim: \text{Fun}(\mathcal{J}, \mathcal{C}) \rightarrow \mathcal{C}, \quad F \mapsto \lim F$$

is *not* a homotopical functor.

SOLUTION: Plenty of examples. Homotopy pullbacks are not pullbacks. $S^1 \times_* S^1$ is a torus, but $S^1 \times_{D^2} S^1$ is homeomorphic to S^1 .

- (2) Let \mathcal{C} be a homotopical category. Prove that for any discrete (the only morphisms are identities) category \mathcal{J} with finitely many morphisms, $\text{ho}(\mathcal{C})^{\mathcal{J}}$ is equivalent to $\text{ho}(\mathcal{C}^{\mathcal{J}})$, where $\mathcal{C}^{\mathcal{J}}$ has weak equivalences defined pointwise. Prove that finite products in $\text{ho}(\mathcal{C})$ are homotopy products and finite coproducts in $\text{ho}(\mathcal{C})$ are homotopy coproducts.

SOLUTION:

You need to assume that \mathcal{J} is finite (or maybe a model category), otherwise you get bad behavior. For example, if \mathcal{C} has one object and two morphisms f and g and you invert g . Let \mathcal{J} be the integers. In $\text{ho}(\mathcal{C})^{\mathcal{J}}$ has morphisms of the form $(f, fg^{-1}f, fg^{-1}fg^{-1}f, \dots)$, but this is not a morphism in $\text{ho}(\mathcal{C}^{\mathcal{J}})$. The problem is that there are different lengths of zigzags at different objects of \mathcal{J} .

Let $L: \mathcal{C} \rightarrow \text{ho } \mathcal{C}$ be the localization functor. Then $L^{\mathcal{J}}: \mathcal{C}^{\mathcal{J}} \rightarrow (\text{ho } \mathcal{C})^{\mathcal{J}}$ is a functor which sends weak equivalences to isomorphisms. This yields a functor $\text{ho}(\mathcal{C}^{\mathcal{J}}) \rightarrow (\text{ho } \mathcal{C})^{\mathcal{J}}$. By construction, this functor is bijective on objects. Further, it is faithful, for it sends a zig-zag of tuples of morphisms to the corresponding tuple of zig-zags of morphisms.

We need additional assumptions to guarantee that this functor is full. If \mathcal{J} is finite, then the supremum of lengths of zig-zags accuring in an \mathcal{J} -indexed tuple is finite, and thus arises from a zig-zag of tuples. Otherwise, if \mathcal{C} is a model category (or more generally has an n -step calculus of fractions), the supremum of lengths accuring in a zig-zag of \mathcal{J} -indexed tuples is again finite, and thus arises from a zig-zag of tuples.

Assume this functor is full. Then, $\text{ho}(\mathcal{C}^{\mathcal{J}}) \rightarrow (\text{ho } \mathcal{C})^{\mathcal{J}}$ is an equivalence (even an isomorphism, although it is forbidden to acknowledge that some functors are isomorphisms).

Assume that \mathcal{C} has homotopy products and $\text{ho } \mathcal{C}$ has products. Then, the homotopy product functor $\text{ho}(\mathcal{C}^{\mathcal{J}}) \rightarrow \text{ho } \mathcal{C}$ is the right derived functor of the product $\mathcal{C}^{\mathcal{J}} \rightarrow \mathcal{C}$. By the previous problem set, this is the adjoint of the left derived functor of the diagonal, $\mathcal{C} \rightarrow \mathcal{C}^{\mathcal{J}}$. Since the diagonal is already homotopical, its left derived functor is again the diagonal $\text{ho}(\mathcal{C})^{\mathcal{J}} \rightarrow \text{ho } \mathcal{C}$ after identifying $\text{ho}(\mathcal{C}^{\mathcal{J}})$ and $(\text{ho } \mathcal{C})^{\mathcal{J}}$. But the adjoint to this diagonal is exactly the homotopy product in $\text{ho } \mathcal{C}$. See also [Rie14, Remark 6.3.1 and footnote 3 therein].

(3) A *coequalizer* is the colimit of a diagram of shape $\bullet \rightrightarrows \bullet$ in a category.

- (a) Prove that the data of the coequalizer of two parallel morphisms $A \xrightarrow[f]{g} B$ is equivalent to the data of the pushout of the diagram

$$A \xleftarrow{\nabla} A \amalg A \xrightarrow{(f,g)} B,$$

where $\nabla: A \amalg A \rightarrow A$ is the fold map.

SOLUTION: We will show that each object has the universal property of the other. First, assume that P is a pushout in the diagram below

$$\begin{array}{ccc} A \amalg A & \xrightarrow{(f,g)} & B \\ \downarrow \nabla & & \downarrow p \\ A & \xrightarrow{q} & P \end{array}$$

First, claim that p equalizes f and g . To see this, consider the composite $p(f, g)i_1$, where $i_1: A \rightarrow A \amalg A$ is the left injection into the coproduct. We have $(f, g)i_1 = f$, and $\nabla i_1 = \text{id}_A$. Then:

$$pf = p(f, g)i_1 = q\nabla i_1 = q$$

Similarly, $(f, g)i_2 = g$ and $\nabla i_2 = \text{id}_A$, so:

$$pg = p(f, g)i_2 = q\nabla i_2 = q$$

So $pf = pg$. So p equalizes f and g . It remains to be seen that it is the universal morphism that does so.

Let $x: B \rightarrow X$ be any morphism such that $xf = xg$. Then $x(f, g)i_1 = xf = xg = x(f, g)i_2$, which is the same as $x(f, g) = x\nabla = xq$. Then we can draw a commuting diagram

$$\begin{array}{ccc} A \amalg A & \xrightarrow{(f,g)} & B \\ \downarrow \nabla & & \downarrow p \\ A & \xrightarrow{q} & P \end{array} \quad \begin{array}{c} \searrow x \\ \downarrow \\ X \end{array}$$

$xf = xg$

and fill it in with the dashed arrow, which shows that there is a unique morphism $P \rightarrow X$ exhibiting P as the coequalizer of f and g .

The converse is similar.

- (b) Use part (a) to describe the homotopy coequalizer of two maps in the category **Top** of (unpointed) topological spaces¹.

SOLUTION: The homotopy coequalizer of two parallel maps $f, g: X \rightarrow Y$ is the mapping torus

$$X \times [0, 1] \sqcup Y / \sim$$

where \sim identifies

$$(x, 0) \sim f(x) \text{ and } (x, 1) \sim g(x).$$

¹To be precise, we assume all spaces are compactly generated and weakly Hausdorff.

(4) A pointed space X is *well-based* if the inclusion of the basepoint is a cofibration. Let $f: X \rightarrow Y$ be a pointed map of well-based spaces.

- (a) Let $\text{cof}(f)$ be the homotopy cofiber of f . Prove that the homotopy cofiber of $Y \rightarrow \text{cof}(f)$ is homotopy equivalent to ΣX .
- (b) Prove the dual statement: if $\text{fib}(f)$ is the homotopy fiber of f , then the homotopy fiber of $\text{fib}(f) \rightarrow X$ is homotopy equivalent to ΩY .

SOLUTION:

- (a) Consider the diagram

$$\begin{array}{ccccc} X & \xrightarrow{f} & Y & \longrightarrow & * \\ \downarrow & & \downarrow & & \downarrow \\ * & \longrightarrow & \text{cof}(f) & \longrightarrow & \text{cof}(Y \rightarrow \text{cof}(f)) \end{array}$$

Both squares are individually homotopy pushouts. Thus, by a version of the pasting lemma, the rectangle is a homotopy pushout.

(To see this explicitly, recall that homotopy pushouts are computed by cofibrant replacement and literal pushout. Cofibrantly replace the top left cospan to compute $\text{cof}(f)$ by literal pushout. Then, cofibrantly replace the entire square, along with the map to the point. The square is still a pushout after cofibrant replacement and now $\text{cof}(Y \rightarrow \text{cof}(f))$ can be computed by literal pushout. Now the literal pasting lemma yields that the entire diagram is a literal pushout, which proves the homotopy pasting lemma.)

Since the pushout of X mapping to points is the suspension of X , this demonstrates that $\text{cof}(Y \rightarrow \text{cof}(f))$ is homotopy equivalent to ΣX . The well-pointedness assumption guarantees that this homotopy pushout is the same as the smash product $S^1 \wedge X$, but if you take the pushout as the definition of suspension, the well-pointedness assumption is not needed.

- (b) Prove the dual statement: if $\text{fib}(f)$ is the homotopy fiber of f , then the homotopy fiber of $\text{fib}(f) \rightarrow X$ is homotopy equivalent to ΩY .

Dually, consider the diagram

$$\begin{array}{ccccc} \text{fib}(\text{fib}(f) \rightarrow X) & \longrightarrow & \text{fib}(f) & \longrightarrow & * \\ \downarrow & & \downarrow & & \downarrow \\ * & \longrightarrow & X & \xrightarrow{f} & Y \end{array}$$

Both squares are homotopy pullbacks, so the rectangle is a homotopy pullback by the same reasoning to the above. But since the pushout of points along Y is the loop-space of Y , this demonstrates that $\text{fib}(\text{fib}(f) \rightarrow X)$ is homotopy equivalent to ΩY .

This problem is [Mal23, Exercise 16 in Chapter 1]. You can probably also find it (without any category theory) in Hatcher.

REFERENCES

- [Mal23] Cary Malkiewich. Spectra and stable homotopy theory. http://people.math.binghamton.edu/malkiewich/spectra_book_draft.pdf, October 2023.
- [Rie14] Emily Riehl. *Categorical homotopy theory*, volume 24 of *New Mathematical Monographs*. Cambridge University Press, Cambridge, 2014.