# ECE 440 - Introduction to Computer Networks

Spring 2020 Lecture 3

### Packet switching versus circuit switching

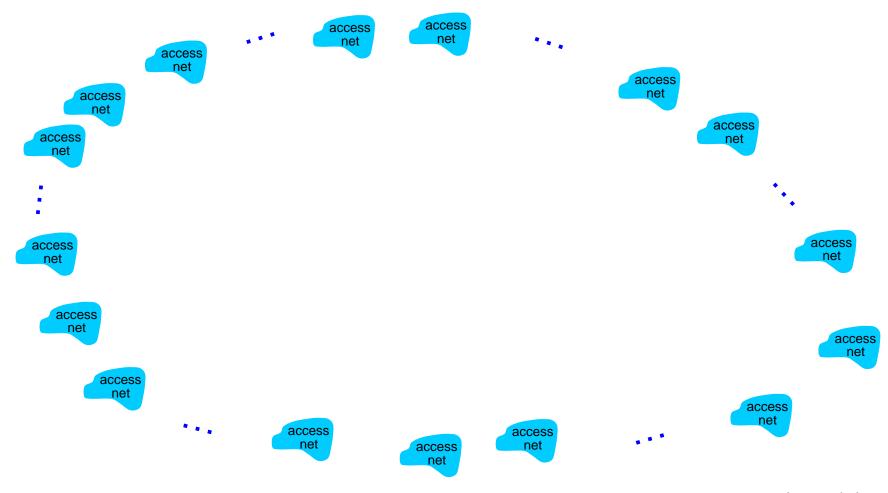
### is packet switching a "slam dunk winner?"

- great for bursty data
  - resource sharing
  - simpler, no call setup
- excessive congestion possible: packet delay and loss
  - protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior?
  - bandwidth guarantees needed for audio/video apps
  - still an unsolved problem (chapter 7)

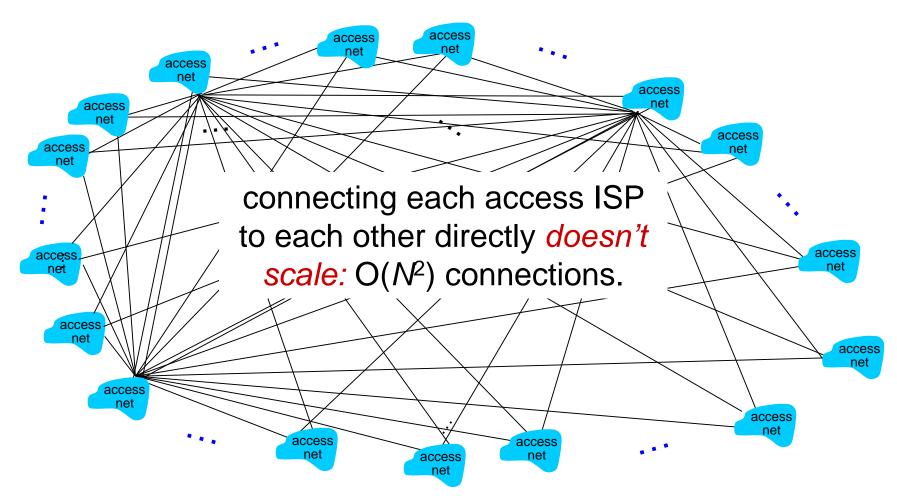
Q: human analogies of reserved resources (circuit switching) versus on-demand allocation (packet-switching)?

- End systems connect to Internet via access ISPs (Internet Service Providers)
  - residential, company and university ISPs
- Access ISPs in turn must be interconnected.
  - so that any two hosts can send packets to each other
- Resulting network of networks is very complex
  - evolution was driven by economics and national policies
- Let's take a stepwise approach to describe current Internet structure

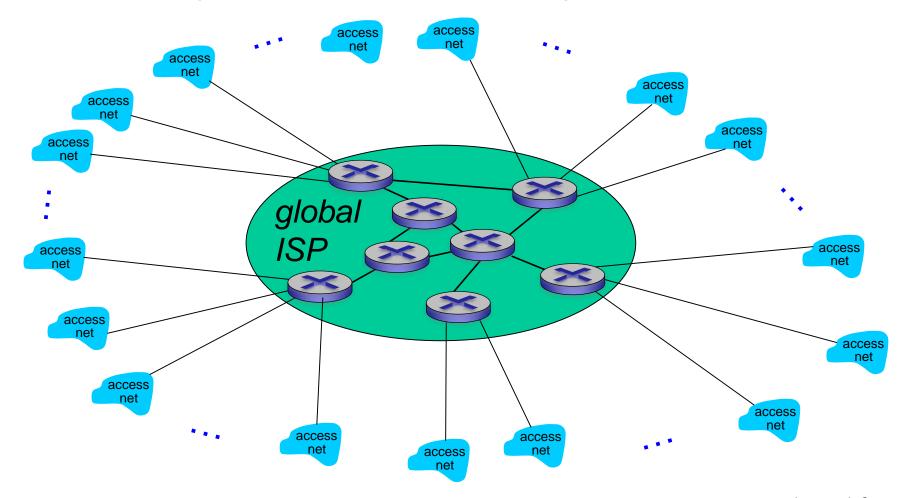
Question: given millions of access ISPs, how to connect them together?



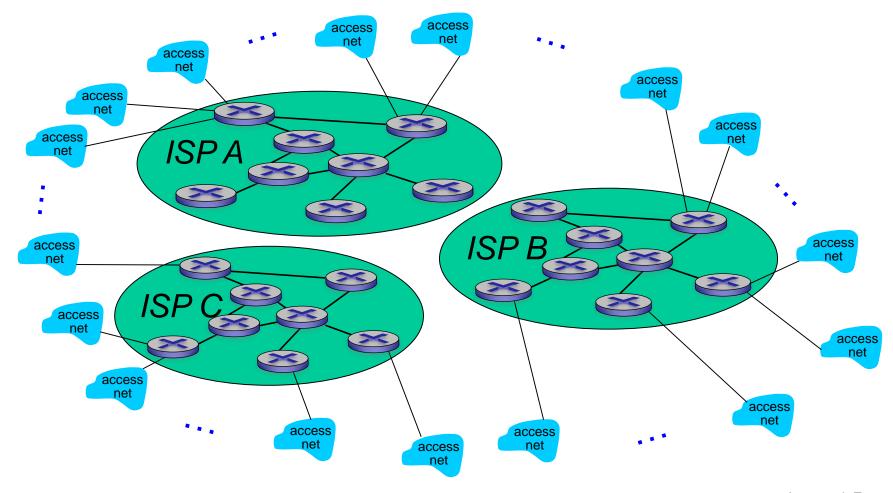
Option: connect each access ISP to every other access ISP?



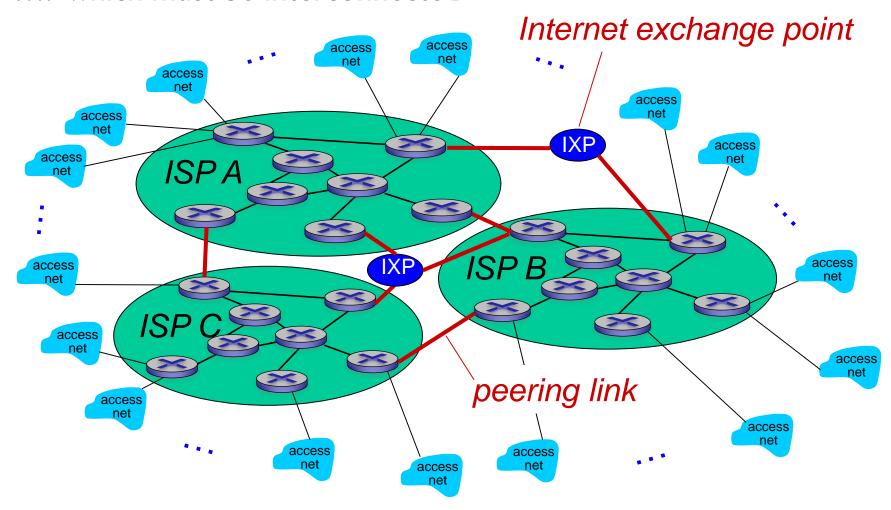
Option: connect each access ISP to one global transit ISP? Customer and provider ISPs have economic agreement.



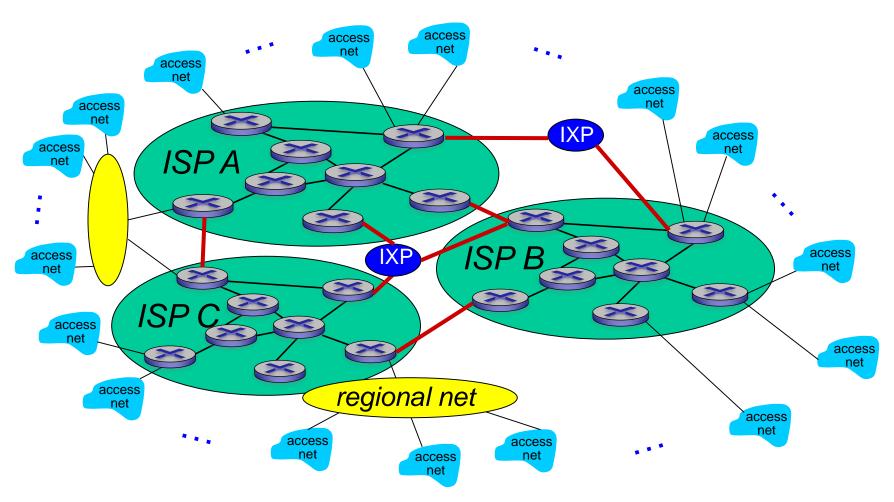
But if one global ISP is viable business, there will be competitors ....



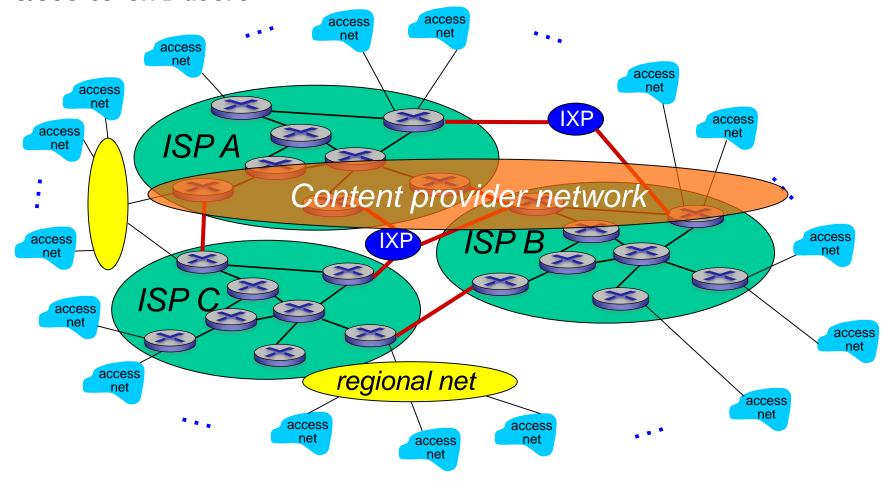
But if one global ISP is viable business, there will be competitors .... which must be interconnected

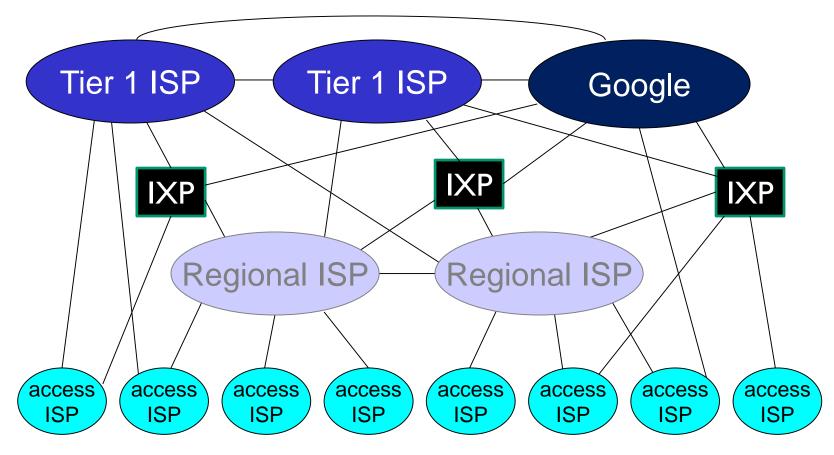


... and regional networks may arise to connect access nets to ISPs



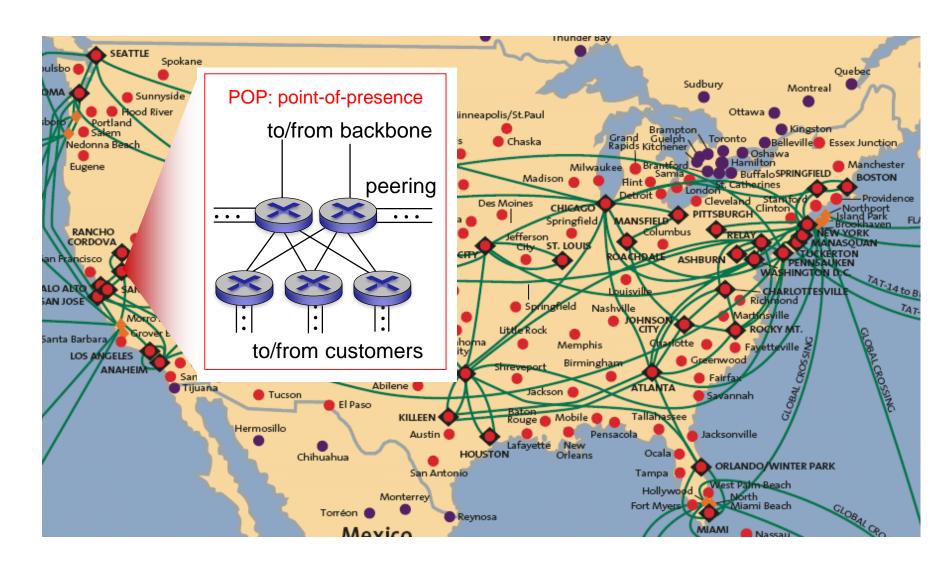
... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users





- at center: small # of well-connected large networks
  - "tier-I" commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
  - content provider network (e.g., Google): private network that connects it data centers to Internet, often bypassing tier-I, regional ISPs  $_{\rm Introduction\ 1-11}$

## Tier-I ISP: e.g., Sprint



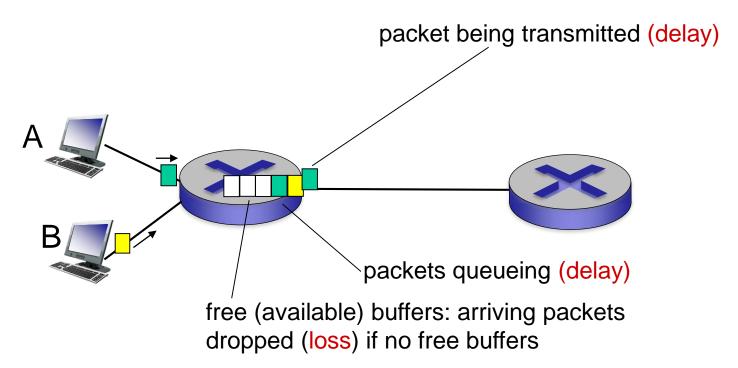
## Chapter 1: roadmap

- I.I what is the Internet?
- 1.2 network edge
  - end systems, access networks, links
- 1.3 network core
  - packet switching, circuit switching, network structure
- 1.4 delay, loss, throughput in networks
- 1.5 protocol layers, service models
- 1.6 networks under attack: security
- 1.7 history

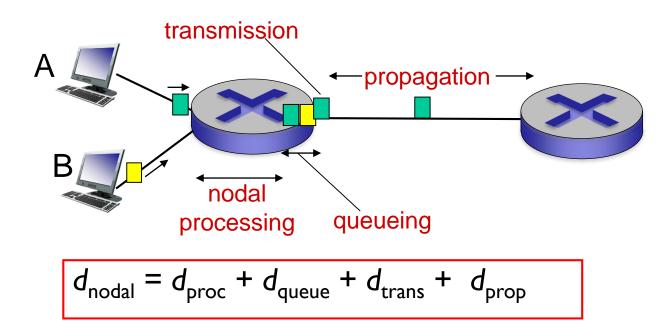
### How do loss and delay occur?

#### packets queue in router buffers

- packet arrival rate to link (temporarily) exceeds output link capacity
- packets queue, wait for turn



### Four sources of packet delay



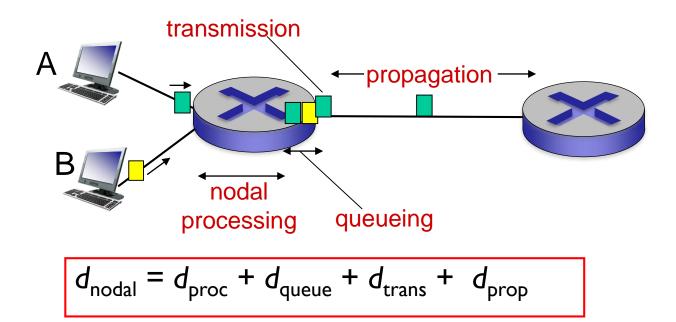
### $d_{proc}$ : nodal processing

- check bit errors
- determine output link
- typically < msec</p>

### d<sub>queue</sub>: queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

### Four sources of packet delay



#### $d_{\text{trans}}$ : transmission delay:

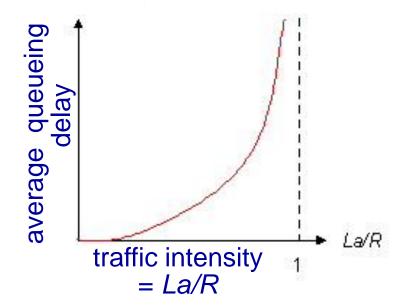
- L: packet length (bits)
- R: link bandwidth (bps)
- $d_{trans} = L/R \leftarrow d_{trans}$  and  $d_{prop} \rightarrow d_{prop} = d/s$ very different

#### $d_{prop}$ : propagation delay:

- d: length of physical link
- s: propagation speed (~2×10<sup>8</sup> m/sec)

### Queueing delay (revisited)

- R: link bandwidth (bps)
- L: packet length (bits)
- a: average packet arrival rate



- La/R ~ 0: avg. queueing delay small
- La/R -> I: avg. queueing delay large
- La/R > I: more "work" arriving than can be serviced, average delay infinite!

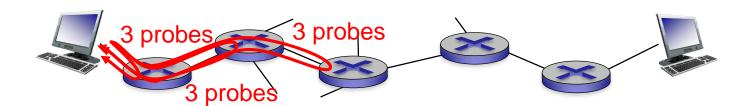


 $La/R \sim 0$ 

La/R ->

### "Real" Internet delays and routes

- what do "real" Internet delay & loss look like?
- traceroute program: provides delay measurement from source to router along endend Internet path towards destination. For all i:
  - sends three packets that will reach router *i* on path towards destination
  - router i will return packets to sender
  - sender times interval between transmission and reply.



### "Real" Internet delays, routes

traceroute: gaia.cs.umass.edu to www.eurecom.fr

```
3 delay measurements from
                                             gaia.cs.umass.edu to cs-gw.cs.umass.edu
1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms
2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
5 in1-so7-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms
                                                                          trans-oceanic
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms 4 9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms 10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
                                                                            link
11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms
12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms
14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms 16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
                       means no response (probe lost, router not replying)
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms
```

<sup>\*</sup> Do some traceroutes from exotic countries at www.traceroute.org