# August 22: Exploring causal assumptions with string diagrams

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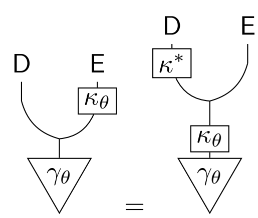
## 1 Recoverability

A natural assumption suggested by the notion of a CSDP is that of *recoverability* - that a causal theory $\mathcal{T} : E \times D \rightarrowtriangle E$ permits some decision function that reproduces the distribution of the observed data. That is, we assume that for every $(\kappa_\theta, \mu_\theta) := \theta \in \mathcal{T}$ there exists $\gamma_\theta \in \Delta(\mathcal{D})$ such that

$$\gamma_\theta \kappa_\theta = \mu_\theta \tag{1}$$

Suppose also that we have some $\kappa^*$ that, for all $\theta \in \mathcal{T}$, is a Bayesian inversion of $\gamma_\theta$ and $\kappa_\theta$; that is:




(2)

A sufficient condition for the existence of such a $\kappa^*$ is the assumption that decisions correspond to *variable setting* - that is, there is some variable $\mathsf{X} : E \to X$ such that for all $a \in D, \theta \in \mathcal{T}$ we have $\delta_a \kappa_\theta F_{\mathsf{X}} = \delta_a$ (such an assumption arises in graphical models as hard interventions, and in potential outcomes as "potential-outcome identifiers"). Indeed $F_{\mathsf{X}}$ is in this case a candidate for $\kappa^*$. It is not necessary that $\kappa^*$ be deterministic, however - suppose every $\kappa$ ignores $D$. Then choose $\gamma_\theta = \gamma$ for arbitrary $\gamma \in \Delta(\mathcal{D})$ and it can be verified that $\kappa^* : b \mapsto \gamma$ satisfies 2.

I believe a weaker sufficient condition for the existence of a universal $\kappa^*$ is that every $\kappa_\theta$ factorises as $\kappa_\theta = hj_\theta$ for some fixed $h$, but I have not yet shown this.

We will proceed somewhat rashly: suppose that by defining $\gamma : \mathcal{T} \to \Delta(\mathcal{D})$, $\mu : \mathcal{T} \to \Delta(\mathcal{E})$ and $\kappa : \mathcal{T} \times D \to \Delta(\mathcal{E}$ by $\gamma : \theta \to \gamma_\theta$, $\mu : \theta \to \mu_\theta$ and $\kappa : (\theta, d) \to \kappa_\theta(d; \cdot)$ that all resulting objects are Markov kernels, and that $\mathcal{T}$ is a standard measurable space.

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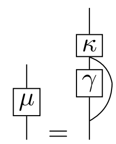
By previous assumptions, we have the following properties:

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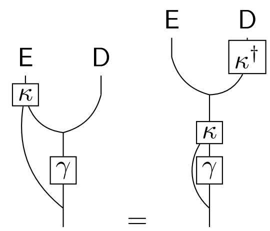


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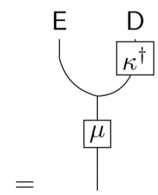
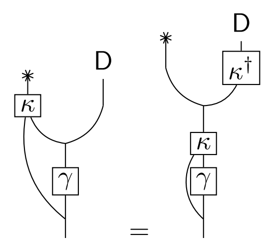



(5)

From 4 we also have

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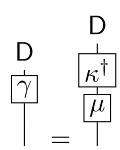
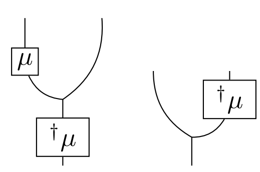



(7)

Where 7 follows from 3.

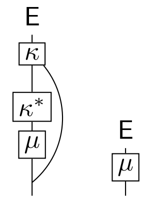
Now, suppose we have a *left-inverse* of $\mu$, denoted $^\dagger\mu$. This means that from observations E we can distinguish any two states $\theta$ and $\theta'$ unless $\mu_\theta = \mu_{\theta'}$; this could be true, for example, if we have an infinite sequence of samples.

A corollary of Lemma 2.5 is that left inverses have the following property:

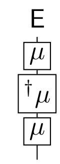



(8)

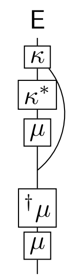
We then have



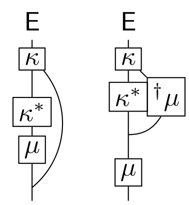

$$= \tag{9}$$




$$= \tag{10}$$

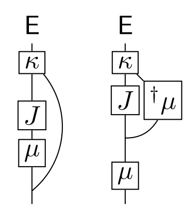



$$= \tag{11}$$




$$= \tag{12}$$

A key question is does 12 imply anything non-trivial regarding the following "identifiability" condition for arbitrary $J : E \to \Delta(\mathcal{D})$:




$$= \tag{13}$$

I call this identifiability because the left hand side is the kernel $\mathcal{T} \to \Delta(\mathcal{E})$ that computes the "result" of a given state and decision function, while the right hand side implies it is possible to find a $J$ that minimises the expected utility $\kappa u$ independent of the causal state (this is because we see one of the input wires and control the other).

We may be able to get some insight into this by asking, given matrices $A, B, C, D$ of appropriate shapes, if $BA = CA$ when does $BDA = CDA$?

## 2 Notes on category theoretic probability and string diagrams

Category theoretic treatments of probability theory often start with *probability monads* (for a good overview, see [Jacobs, 2018]). A monad on some category $C$ is a functor $T : C \to C$ along with natural transformations called the unit $\eta : 1_C \to T$ and multiplication $\mu : T^2 \to T$. Roughly, functors are maps between categories that preserve identity and composition structure and natural transformations are "maps" between functors that also preserve composition structure. The monad unit is similar to the identity element of a monoid in that application of the identity followed by multiplication yields the identity transformation. The multiplication transformation is also (roughly speaking) associative.

An example of a probability monad is the discrete probability monad given by the functor $\mathcal{D}$ : **Set** $\to$ **Set** which maps a countable set $X$ to the set of functions from $X \to [0,1]$ that are probability measures on $X$, denoted $\mathcal{D}(X)$. $\mathcal{D}$ maps a measurable function $f$ to $\mathcal{D}f : X \to \mathcal{D}(X)$ given by $\mathcal{D}f : x \mapsto \delta_{f(x)}$. The unit of this monad is the map $\eta_X : X \to \mathcal{D}(X)$ given by $\eta_X : x \mapsto \delta_x$ (which is equivalent to $\mathcal{D}1_X$) and multiplication is $\mu_X : \mathcal{D}^2(X) \to \mathcal{D}(X)$ where $\mu_X : \Omega \mapsto \sum_\phi \Omega(\phi)\phi$.

For continuous distributions we have the Giry monad on the category **Meas** of mesurable spaces given by the functor $\mathcal{G}$ which maps a measurable space $X$ to the set of probability measures on $X$, denoted $\mathcal{G}(X)$. Other elements of the monad (unit, multiplication and map between morphisms) are the “continuous” version of the above.

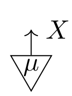
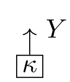
Of particular interest is the Kleisli category of the monads above. The Kleisli $C_T$ category of a monad $T$ on category $C$ is the category with the same objects and the morphisms $X \to Y$ in $C_T$ is the set of morphisms $X \to TY$ in $C$. Thus the morphisms $X \to Y$ in the Kleisli category **Set**$_\mathcal{D}$ are morphisms $X \to \mathcal{D}(Y)$ in **Set**, i.e. stochastic matrices, and in the Kleisli category **Meas**$_\mathcal{G}$ we have Markov kernels. Composition of arrows in the Kleisli categories correspond to Matrix products and “kernel products” respectively.

Both $\mathcal{D}$ and $\mathcal{G}$ are known to be *commutative* monads, and the Kleisli category of a commutative monad is a symmetric monoidal category.

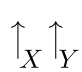
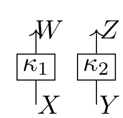
Diagrams for symmetric monoidal categories consist of wires with arrows, boxes and a couple of special symbols. The identity object (which we identify with the set $\{*\}$) is drawn as nothing at all $\{*\} :=$ and identity maps are drawn as bare wires:

$\mathrm{Id}_X :=$  (14)

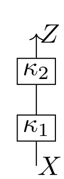
We draw Kleisli arrows from the unit (i.e. probability distributions) $\mu : \{*\} \to X$ as triangles and Kleisli arrows $\kappa : X \to Y$ (i.e. Markov kernels $X \to \Delta(\mathcal{Y})$) as boxes. We draw the Kleisli arrow $\mathbb{1}_X : X \to \{*\}$ (which is unique for each $X$) as below

$\mu :=$  $\kappa :=$  (15)

The product of objects in **Meas** is given by $(X, \mathcal{X}) \cdot (Y, \mathcal{Y}) = (X \times Y, \mathcal{X} \otimes \mathcal{Y})$, which we will often write as just $X \times Y$. Horizontal juxtaposition of wires indicates this product, and horizontal juxtaposition also indicates the tensor product of Kleisli arrows. Let $\kappa_1 : X \to W$ and $\kappa_2 : Y \to Z$:

$(X \times Y, \mathcal{X} \otimes \mathcal{Y}) :=$  $\kappa_1 \otimes \kappa_2 :=$  (16)

Composition of arrows is achieved by “wiring” boxes together. For $\kappa_1 : X \to Y$ and $\kappa_2 : Y \to Z$ we have

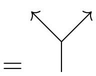
$\kappa_1\kappa_2(x; A) = \int_Y \kappa_2(y; A)\kappa_1(x; dy) :=$  (17)

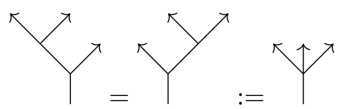
Symmetric monoidal categoris have the following coherence theorem[Selinger, 2010]:

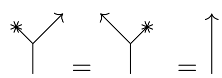
**Theorem 2.1** (Coherence (symmetric monoidal))**.** *A well-formed equation between morphisms in the language of symmetric monoidal categories follows from the axioms of symmetric monoidal categories ifand only if it holds, up to isomorphism of diagrams, in the graphical language.*

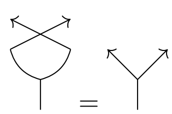
Isomorphism of diagrams for symmetric monoidal categories (somewhat informally) is any planar deformation of a diagram including deformations that cause wires to cross. We consider a diagram for a symmetric monoidal category to be well formed only if all wires point upwards.

In fact the Kleisli categories of the probability monads above have (for each object) unique *copy*: $X \to X \times X$ and *erase*: $X \to \{*\}$ maps that satisfy the *commutative comonoid axioms* that (thanks to the coherence theorem above) can be stated graphically. These differ from the copy and erase maps of *finite product* or *cartesian* categories in that they do not necessarily respect composition of morphisms.

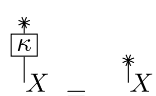
Erase $= \mathbb{1}_X :=$  Copy $= x \mapsto \delta_{x,x} :=$  (18)

 (19)

 (20)

 (21)

Finally, $\{*\}$ is a terminal object in the Kleisli categories of either probability monad. This means that the map $X \to \{*\}$ is unique for all objects $X$, and as a consequence for all objects $X, Y$ and all $\kappa : X \to Y$ we have

 (22)

This is equivalent to requiring for all $x \in X$ $\int_Y \kappa(x; dy) = 1$. In the case of $\mathbf{Set}_{\mathcal{D}}$, this condition is what differentiates a stochastic matrix from a general positive matrix (which live in a larger category than $\mathbf{Set}_{\mathcal{D}}$).

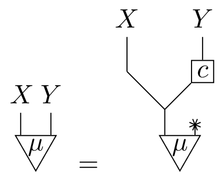
Thus when manipulating diagrams representing Markov kernels in particular (and, importantly, not more general symmetric monoidal categories) diagram isomorphism also includes applications of 19, 20, 21 and 22.

A particular property of the copy map in $\mathbf{Meas}_{G}$ (and probably $\mathbf{Set}_{\mathcal{D}}$ as well) is that it commutes with Markov kernels iff the markov kernels are deterministic [Fong, 2013].

## 2.1 Disintegration and Bayesian inversion

*Disintegration* is a key operation on probability distributions (equivalently arrows $\{*\} \to X$) in the categories under discussion. It corresponds to "finding the conditional probability" (though conditional probability is usually formalised in a slightly different way).

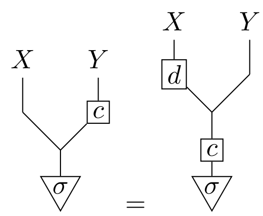
Given a distribution $\mu : \{*\} \to X \otimes Y$, a disintegration $c : X \to Y$ is a Markov kernel that satisfies




(23)

Disintegrations always exist in $\mathbf{Set}_{\mathcal{D}}$ but not in $\mathbf{Meas}_{\mathcal{G}}$. The do exist in the latter if we restrict ourselves to standard measurable spaces. If $c_1$ and $c_2$ are disintegrations $X \to Y$ of $\mu$, they are equal $\mu$-A.S. In fact, this equality can be strengthened somewhat - they are equal almost surely with respect to any distribution that shares the "$X$-marginal" of $\mu$.

Given $\sigma : \{*\} \to X$ and a channel $c : X \to Y$, a Bayesian inversion of $(\sigma, c)$ is a channel $d : Y \to X$ such that



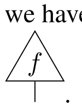

(24)

We can obtain disintegrations from Bayesian inversions and vise-versa.

Clerc et al. [2017] offer an alternative view of Bayesian inversion which they claim doesn't depend on standard measurability conditions, but there is a step in their proof I didn't follow.

## 2.2 Generalisations

Cho and Jacobs [2019] make use of a larger "CD" category by dropping 22. I'm not completely clear whether you end up with arrows being "Markov kernels for general measures" or something else (can we have negative arrows?). This allows for the introduction of "observables" or "effects" of the form



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Jacobs et al. [2019] make use of an embedding of $\mathbf{Set}_{\mathcal{D}}$ in $\mathbf{Mat}(\mathbb{R}^+)$ with morphisms all positive matrices (I'm not totally clear on the objects, or how they are self-dual - this doesn't seem to be exactly the same as the category of finite dimensional vector spaces). This latter category is compact closed, which - informally speaking - supports the same diagrams as symmetric monoidal categories with the addition of "upside down" wires.

## 2.3 Key questions for Causal Theories

We will first define *labeled diagrams*. Rather than labelling the wires of our diagrams with *spaces* (as is typical [Selinger, 2010]), we assign a unique label to each "wire segment" (with some qualifications). That is, we assign a unique label to each bare wire in the diagram with the following additonal qualifications:

I'm sure one of the papers I read mentioned labeled diagrams, I just couldn't find it when I looked for it

- If we have a box in the diagram representing the identity map, the incoming and outgoing wires are given the same label
- If we have a wire crossing in the diagram, the diagonally opposite wires are given the same label
- The input wire and the *two* output wires of the copy map are given the same label

Given two diagrams $G_1$ and $G_2$ that are isomorphic under transformations licenced by the axioms of symmetric monoidal categories and commutative comonoid axioms, suppose we have a labelling of $G_1$. We can label $G_2$ using the following translation rule:

Since writing this, I found Kissinger [2014] as an example of a diagrammatic system with labeled wires, I will follow it up

- For each box in $G_2$, we can identify a corresponding box in $G_1$ via labels on each box. For each such pair of boxes, we label the incoming wires of the $G_2$ box with the labels of the $G_1$ box preserving the left-right order. We do likewise for outgoing wires.

These rules will lead to a unique labelling of $G_2$ with all wire segments are labelled. We would like for these rules to yield the following:

- For any sequence of diagram isomorphisms beginning with $G_1$ and ending with $G_2$, we end up with the same set of labels
- If we label $G_2$ according to the rules above then relabel $G_1$ from $G_2$ according to the same rules we retrieve the original labels of $G_1$

We do not prove these properties here, but motivate them via the following considerations:

- These properties obviously hold for the wire segments into and out of boxes
- The only features a diagram may have apart from boxes and wires are wire crossings, copy maps and erase maps
- The labeling rule for wire crossings respects the symmetry of the swap map
- The labeling rule for copy maps respects the symmetry of the copy map and the property described in Equation 21

We will follow the convention whereby "internal" wire labels are omitted from diagrams.

Note also that each wire that terminates in a free end can be associated with a random variable. Suppose for $N \in \mathbb{N}$ we have a kernel $\kappa : A \to \Delta(\times_{i\in N} X_i)$. Define by $p_j$ ($j \in [N]$) the projection map $p_j : \times_{i\in N} X_i \to X_j$ defined by $p_j : (x_0, ..., x_N) \mapsto x_j$. $p_j$ is a measurable function, hence a random variable. Define by $\pi_j$ the projection kernel $\mathcal{G}(\pi_j)$ (that is, $\pi_j : \mathbf{x} \mapsto \delta_{p_j(\mathbf{x})}$). Note that $\kappa(y; p_j^{-1}(A)) = \int_{X_j} \delta_{p_j(\mathbf{x})}(A)\kappa(y; d\mathbf{x}) = \kappa\pi_j$. Diagrammatically, $\pi_j$ is the identity map on the $j$-th wire tensored with the erase map on every other wire. Thus the $j$-th wire carries the distribution associated with the random variable $p_j$. We will therefore consider the labels of the "outgoing" wires of a diagram to denote random varaibles (though there are obviously many random variables not represented by such wires). We will additionally distinguish wire labels from spaces by font - wire labels are sans serif $\mathsf{A}, \mathsf{B}, \mathsf{C}, \mathsf{X}, \mathsf{Y}, \mathsf{Z}$ while spaces are serif $A, B, C, X, Y, Z$.

Wire labels appear to have a key advantage over random variables: they allow us to "forget" the sample space as the correct typing is handled automatically by composition and erasure of wires

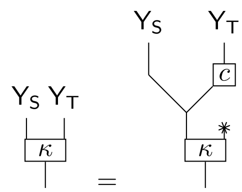
**generalised disintegrations** : Of key importance to our work is generalising the notion of disintegration (and possibly Bayesian inversion) to general kernels $X \to Y$ rather than restricting ourselves to probability distributions $\{*\} \to Y$. We will define generalised disintegrations as a straightforward analogy regular disintegrations, but the conditions under which such disintegrations exist are more restrictive than for regular disintegraions.

**Definition 2.2** (Label signatures)**.** If a kernel $\kappa : X \to \Delta(Y)$ can be represented by a diagram $G$ with incoming wires $\mathsf{X}_1, ...\mathsf{X}_\mathsf{n}$ and outgoing wires $\mathsf{Y}_1, ..., \mathsf{Y}_\mathsf{m}$, we can assign the kernel a "label signature" $\kappa : \mathsf{X}_1 \otimes ... \otimes \mathsf{X}_\mathsf{n} \dashrightarrow \mathsf{Y}_1 \otimes ... \otimes \mathsf{Y}_m$ or, for short, $\kappa : \mathsf{X}_{[n]} \dashrightarrow \mathsf{Y}_{[m]}$. Note that this signature associates each label with a unique space - the space of $\mathsf{X}_1$ is the space associated with the left-most wire of $G$ and so forth. We will implicitly leverage this correspondence and write with $X_1$ the space associated with $\mathsf{X}_1$ and so forth. Note that while $\mathsf{X}_1$ is by construction always different from $\mathsf{X}_2$ (or any other label), the space $X_1$ may coincide with $X_2$ - the fact that labels always maintain distinctions between wires is the fundamental reason for introducing them in the first place.

There might actually be some sensible way to consider $\kappa$ to be transforming the measurable functions of a type similar to $\otimes_{i\in[n]}\mathsf{X}_i$ to functions of a type simlar to $\otimes_{i\in[m]}\mathsf{Y}_i$ (or vise versa - perhaps related to Clerc et al. [2017]), but wire labels are all we need at this point

**Definition 2.3** (Generalised disintegration)**.** Given a kernel $\kappa : X \to \Delta(Y)$ with label signature $\kappa : \mathsf{X}_{[\mathsf{n}]} \dashrightarrow \mathsf{Y}_{[\mathsf{m}]}$ and disjoint subsets $S, T \subset [m]$ such that $S \cup T = [m]$, a kernel $c$ is a $g$-

*disintigration from $S$ to $T$* if it's type is compatible with the label signature $c : \mathsf{Y}_{\mathsf{S}} \dashrightarrow \mathsf{Y}_{\mathsf{T}}$ and we have the identity (omitting incoming wire labels):




(25)

I have introduced without definition additional labeling operations here: first, each label has a particular space associated with it (in order to license the notion of "type compatible with label signature"), and we have supposed labels can be "bundled".

In contrast to regular disintegrations, generalised disintegrations "usually" do not exist. Consider $X = \{0,1\}$, $Y = \{0,1\}^2$ and $\kappa$ has label signature $\mathsf{X}_1 \dashrightarrow \mathsf{Y}_{\{1,2\}}$ with

$$\kappa : \begin{cases} 1 \mapsto \delta_1 \otimes \delta_1 \\ 0 \mapsto \delta_1 \otimes \delta_0 \end{cases} \tag{26}$$

$\kappa$ imposes contradictory requirements for any disintegration $c : \{0,1\} \to \{0,1\}$ from $\{1\}$ to $\{2\}$: equality for $\mathsf{X}_1 = 1$ requires $c(1;\cdot) = \delta_1$ while equality for $\mathsf{X}_1 = 0$ requires $c(1;\cdot) = \delta_0$. Subject to some regularity conditions (similar to standard Borel conditions for regular disintegrations), we can define g-disintegrations of a canonically related kernel that do generally exist; intuitively, g-disintegrations exist if they take the "input wires" of $\kappa$ as input wires themselves.

**Lemma 2.4.** *Given $\kappa : X \to \Delta(Y)$, a kernel $\kappa^\dagger$ is a right inverse iff we have for all $x \in X$ $\kappa^\dagger(y;A) = \delta_x(A)$, $\kappa(x;\cdot)$-almost surely.*

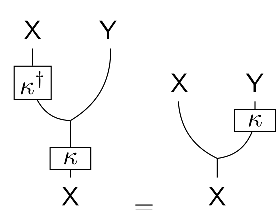
*Proof.* Suppose $\kappa^\dagger$ satisfies the almost sure equality for all $x \in X$. Then for all $x \in X$, $A \in \mathcal{Y}$ we have $\kappa\kappa^\dagger(x;A) = \int_Y \kappa^\dagger(y;A)\kappa(x;dy) = \int_Y \delta_x(A)\kappa(x;dy) = \delta_x(A)$; that is, $\kappa\kappa^\dagger = \mathrm{Id}_X$, so $\kappa^\dagger$ is a right inverse of $\kappa$.

Suppose we have a right inverse $\kappa^\dagger$. By definition, for all $x \in X$ and $A \in \mathcal{Y}$ we have $\int_Y \kappa^\dagger(y;A)\kappa(x;dy) = \delta_x(A)$. Suppose $x \notin A$ and let $B_\epsilon = \kappa_A^{\dagger-1}((\epsilon,1])$ for some $\epsilon > 0$. We have $\int_Y \kappa^\dagger(y;A)\kappa(x;dy) \geq \epsilon\kappa(x;B)$. For any $\epsilon > 0$ we have $\kappa(x;B_\epsilon) = 0$. Consider the set $B_0 = \kappa_A^{\dagger-1}((0,1])$. For some sequence $\{\epsilon_i\}_{i\in\mathbb{N}}$ such that $\lim_{i\to\infty}\epsilon_i = 0$ we have $B_0 = \cup_{i\in\mathbb{N}} B_{\epsilon_i}$. By countable additivity, $kappa(x;B_0) = 0$. Suppose $x \in A$ and let $B^{1-\epsilon} = \kappa_A^{\dagger-1}([0,1-\epsilon))$. By an argument analogous to the above, we have $\kappa(x;B^1) = 0$. Thus the $\kappa(x;\cdot)$ measure of the set on which $\kappa^\dagger(y;A)$ disagrees with $\delta_x(A)$ is $\kappa(x;B_0) + \kappa(x;B^1) = 0$ and hence $\kappa^\dagger(y;A) = \delta_x(A)$ $\kappa(x;\cdot)$-almost surely. □

I haven't shown that any map inverting $\kappa$ implies the existence of a Markov kernel that does so

I am using countable sets below to get my general argument in order without getting too hung up on measurability; I will try to lift it to standard measurable once it's all there

**Lemma 2.5.** *Given $\kappa : X \to \Delta(Y)$ and a right inverse $\kappa^\dagger$, we have*




(27)

*Proof.* Let the diagram on the left hand side be $L$ and the diagram on the right hand side be $R$.

$$L(x; A \times B) = \int_Y \int_{Y \times Y} \mathrm{Id}_Y \otimes \kappa_S^\dagger(y, y'; A \times B)\delta_{(z,z)}(dy \times dy')\kappa\pi_S(x; dz) \tag{28}$$
$$= \int \mathrm{Id}_Y \otimes \kappa^\dagger(z, z; A \times B)\kappa\pi_S(x; dz) \tag{29}$$
$$= \int \delta_z(A)\kappa_S^\dagger(z; B)\kappa\pi_S(x; dz) \tag{30}$$
$$= \int_A \kappa_S^\dagger(z; B)\kappa\pi_S(x; dz) \tag{31}$$
$$= \delta_x(B)\kappa\pi_S(x; A) \tag{32}$$

Where 32 follows from Lemma 2.4.

$$R(x; A \times B) = \int \delta_{(x,x)}(dy \times dy')\kappa\pi_S \otimes \mathrm{Id}_X(y, y'; A \times B) \tag{33}$$
$$= \kappa\pi_S(x; A)\delta_x(B) \qquad\qquad = L \tag{34}$$

□

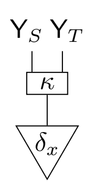
**Theorem 2.6.** *Given countable $X$ and standard measurable $Y$, $n, m \in \mathbb{N}$, $S, T \subset [m]$, $\kappa$ with label signature $\mathsf{X}_{[n]} \dashrightarrow \mathsf{Y}_{[m]}$ a g-disintegration exists from $S$ to $T$ if $\kappa\pi_S$ is right-invertible*

*via a Markov kernel*

*Proof.* In addition, as $R$ is a composition of Markov kernels, and hence a Markov kernel itself, $L$ must also be a Markov kernel even if $\kappa^\dagger$ is not.

For all $x \in X$ we have a (regular) disintegration $c_x : Y_S \to \Delta(Y_T)$ of $\kappa(x; \cdot)$ by standard measurability of $Y$. Define $c : X \otimes Y_S \to \Delta(Y_T)$ by $c : (x, y_S) \mapsto c_x(y_S)$. Clearly, $c(x, y_S)$ is a probability distribution on $Y_T$ for all $(x, y_S) \in X \otimes Y_S$. It remains to show $c(\cdot)^{-1}(B)$ is measurable for all $B \in \mathcal{B}([0, 1])$. But $c(\cdot)^{-1}(B) = \cap_{x \in X} c_y(\cdot)^{-1}(B)$. The right hand side is measurable by measurability of $c_y(\cdot)^{-1}(B)$ countability of $X$, so $c$ is a Markov kernel.

By the definition of $c_x$, we have for all $x \in X$

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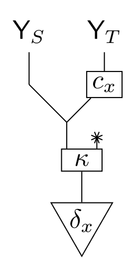


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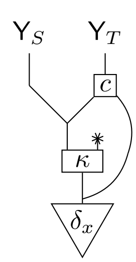
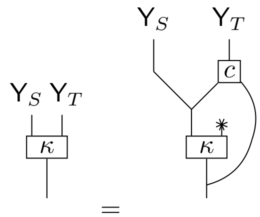


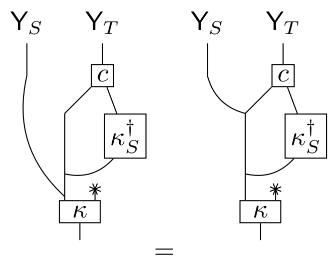

(36)

Which implies

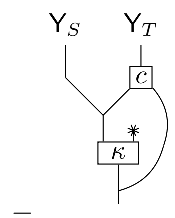



(37)

Finally, we have




(38)




(39)

Where the first line follows from 20 and the second line from 27. If $\kappa_S^\dagger$ is a Markov kernel, then $\curlyvee(\mathrm{Id}_{Y_S} \otimes \kappa_S^\dagger)c$ is a g-disintegration. □

In the reverse direction, suppose $\kappa$ is such that $\kappa\pi_T = \mathrm{Id}_X$; that is, $\pi_T$ is a right inverse of $\kappa$. If $\kappa\pi_S$ is not right invertible then, by definition, there is no $d$ such that $\kappa\pi_S d\pi_T = \mathrm{Id}_X$. However, if a g-disintegration of $\kappa$ exists then there is a $d$ such that $\kappa\pi_S d = \kappa$, a contradiction. Thus if $\kappa\pi_S$ is not right invertible then there is *in general* no g-disintegration from $S$ to $T$.

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