

1 Introduction

Computational design and analysis has become a fundamental part of industry and academia for use in research, development, and manufacturing. In general, the accuracy of computational analysis depends heavily on the fidelity of the computational representation of a real-world object or phenomenon. However, the task of creating high fidelity models of an actual geometry can be time-consuming—sometimes consuming up to seventy-five percent of the time required to produce a solution [1]. Most mesh generation strategies focus on element quality—with the justification being that downstream applications require high quality geometries in order to achieve a desired level of accuracy. However, element quality should be secondary to accurately representing the underlying physical object or phenomenon.

The justification for the development of these methods is to explore the concept of optimal mesh generation through mesh refinement, and subsequently optimal representation, by developing a formal description of fundamental mesh operations. These mesh operations (triangle split, edge split, edge collapse, edge swap, and vertex smoothing) were formulated with respect to optimal representation of a discrete element for a continuous, parameterized surface. The proposed algorithm is a general-use method that can be applied to any surface regardless of its representation. This is due to every step being developed without the use of derivatives. A result of not using derivatives is that each step in the algorithm is robust to large or discontinuous changes in derivatives. Most other methods operate on a specific type of surface geometry, e.g. NURBS surfaces, and use the specific information available for the type of surface in use. [TRANSITION SENTENCE]

This process can only be automated if some way of judging “how well” a surface grid represents a surface is present. To this end, a method of generating surface grids through constrained, local optimization is detailed below. The concept of discrete element quality is also defined with respect to optimal representation. The above concepts are contrasted with traditional methods for grid generation through direct comparison of representation optimality and efficiency.

2 Related Work

In general, grid generation is a name for any process that creates a grid. For example, the advancing-front algorithm advances boundaries into space to generate a grid [10, 3][CITE]. Other methods generate grids from iterative refinement or enrichment from initial, coarse configurations [5, 6, 9][CITE]. Usually the benchmark for separating the two methods, generation and refinement, are the prioritization of grid quality and grid accuracy (both of these issues will be addressed later).

Traditionally, surface grid generation processes produce good quality grids from the combination of geometric growth rates and smoothing. However, the process requires input and if the input, or starting place, is not appropriate, then the geometry can be

under and/or over sampled for the intended use. That fact is not an indictment of the grid generation process, but instead implies that the final grid is heavily dependent on the inputs. In addition, if some way of controlling the point spacing in the middle of the surface (away from the boundaries) is not present, then more points could be wasted/omitted in an attempt to accurately represent the geometry.

Other efforts have gone into creating a locally or globally optimal surface grid. Many names have been assigned to the particular task, but the underlying goal is very similar — represent a surface as accurately as possible—whatever that means for each application. For example, [CITE, EXPLAIN, CITE, EXPLAIN, CITE, EXPLAIN]

3 Nomenclature and Definitions

3.1 Parametric Surface

In the field of surface mesh generation there exists an overwhelming precedent for mapping physical space to parametric space and performing the actual mesh generation in parametric space [CITE]. This strategy has been proven to be efficient and accurate—with the assumption that the parameterization is of reasonable quality. Here, mesh generation in parametric space also enables the development of unambiguous optimization constraints and objective functions. The search space for the local optimization problem is straightforward and intuitive to define—as well as efficient to search. Therefore, this method will be developed for an orientable, parameterized surface, $\vec{S}(u, v) : \mathbb{R}^2 \rightarrow \mathbb{R}^3$

3.2 Discretization

Consider a discretization, D , defined on S , comprised of n_p points, $P \in \{p_1, \dots, p_{n_p}\}$, and a non-overlapping, non-degenerate, consistently oriented, conformal, triangular connectivity $T \in \{t_1, \dots, t_{n_t}\}$. Each triangle in T , t_i , is defined by points (p_j, p_k, p_m) , and edges (e_n, e_o, e_p) . For the purposes of defining an edge in this work, the ordering of the nodes does not matter. However, the ordering of the triangle connectivity (sometimes referred to as winding) is used to maintain consistent orientation during topological operations and to define constraints in the developed optimization strategy.

Each point, p_i , in D is defined at a $(u, v)_i$ coordinate, and an edge is a straight line in both parametric space and in \mathbb{R}^3 . This is not the case for a discretization, $D : \mathbb{R}^3$, which is mapped onto a surface, $S : \mathbb{R}^3$. In the absence of a parameterization the mapping of an edge in \mathbb{R}^3 to a curve (possibly a straight line, but not guaranteed) on a surface is an ambiguous task which is outside the scope of this work. Additionally, the development of topological constraints for optimization, which will be detailed later, would not be as straightforward as is possible when using a parameterized, planar space.

4 Discretization Error, Discrete Curvature Approximation

The accuracy, or discretization error, of a piecewise, linear representation of an analytical surface in \mathbb{R}^3 can be defined in many ways depending on the intended application. The error associated with the discretization is discussed in terms of the deviation from the surface—most often quantified by calculating or approximating the distance from the surface for each linear entity (triangles and edges) in the discretization. Another way of quantifying the error associated with a discretization is to consider how well it approximates the surface area of the surface it represents. In general the surface area is not known *a priori*; but depending on the underlying representation it can be calculated exactly or can be estimated. One of the goals of the developed method is to be general, i.e. it should be independent of the underlying geometric representation. Therefore, a method requiring the surface area of the underlying geometry violates the aforementioned concept of generality and restricts the applications for which the proposed method can be applied. Some other way of determining/generating a surface grid based on surface area is needed. This process is detailed later.

The concept of deviation defined above is relatively straightforward and intuitive. However, another related way of describing how well a *discretization* represents a *surface* is the degree to which the discrete representation approximates curvature. First, however, curvature must be defined in such a way that a discrete approximation is meaningful and appropriate. In relevant literature, there are many ways to estimate surface curvature [4]. Some of it bears repeating, because it is germane to what is being discussed here: Consider a surface, S , embedded in \mathbb{R}^3 , at point P . The concept of an osculating circle [11] does not generalize well to higher dimensions and therefore cannot be used here as an approximation of curvature. Generally, a (hyper)surface does not have a (hyper)sphere that approximates curvature, but a (hyper)ellipsoid. This is because a surface can have multiple values and directions of curvature at a point. Here, an osculating ellipsoid could be constructed by considering the principal curvature directions and surface normal at a given point on a surface. These three vectors are orthogonal and could be used to define the semi-axes of an ellipse. If some way could be found to determine a scale of the semi-axis associated with the surface normal (which is a unit-vector), this ellipse would be a good approximation of the surface in the direction of the principal curvatures. However, this requires the computation or approximation of derivatives—which we have ruled out here.

Instead of determining a three-dimensional analog to an osculating circle for S , an osculating sphere could be defined for a triangle, T , contained in a discretization, D , which is approximating S . A value of curvature can be calculated for each triangle in the D by considering the corresponding osculating sphere for T . The osculating sphere here (sphere,[FIGURE]) can be approximated by considering the circumscribed sphere (circumsphere) [2] defined by the three points of the triangle, P_0 , P_1 , and P_2 , and a point, P on the surface located within the triangle in (u, v) space. — the radius of the circumsphere will be referred to as the discrete radius of curvature, R_D . This argument for curvature

approximation is a two-dimensional analog to the one in [7] for approximating curvature along a curve. In that work, the authors showed that for an edge on a curve, as the radius of curvature of the osculating circle approaches infinity the arc length of the circular segment approaches the arc length of the edge. However, here as R_D approaches infinity, the surface area of the spherical cap does not approach the area of T —it approaches that of the circumscribed circle defined by T .

Finally, what is most appropriate is to consider the surface area of the tetrahedron, \mathbb{T} , formed by P_0 , P_1 , P_2 , and P . This is a meaningful comparison: as R_D approaches infinity (as the distance between P and the plane defined by T approaches P_0) the surface area of \mathbb{T} approaches two times the area of T [ELABORATE?]. Stated another way, as D is refined the surface area of D approaches that of S .

Surface area convergence of a discretization is a sufficient condition for other schemes of surface grid generation/refinement. That is: if the difference between the surface area of the surface and the sum of the surface area of the discrete entities in the discretization approaches zero then that is sufficient to conclude that the distance between the discretization and surface is also approaching zero, also the dihedral angles between segments approaches 0 degrees [see PROOF once definition of REPRESENTATION DEFICIT is detailed in APPENDIX]. However, the converse of that statement is not true. The pathological case of a highly oscillatory, low amplitude surface approximated by two triangles [MORE DETAIL] can have a small deviation or angles between elements but be a poor estimate for surface area.

5 Representation Deficit based Mesh Operations

5.1 Representation Deficit

The concept of representation deficit is discussed in [8] in the context of scale-independent measure of curve discretizations. Here, the concept will be extended to two dimensions. It is important that any criteria for refinement maintain scale independence. This is so that the values for representation deficit can be compared between different candidate operations. Below, representation deficit is defined and discussed for four fundamental mesh operations: triangle split, edge split, edge swap, and node smoothing.

5.2 Triangle Split

A triangle is a planar object, and is representing a possibly non-planar portion of a surface. Any triangle, with area A_T , in the discretization can have at most the same area as the portion of the surface which it represents, A_S . Therefore, the representation deficit for a triangle, RD_T , can be defined as $RD_T = \frac{A_S - A_T}{A_T}$. Note that the areas are calculated in \mathbb{R}^3 . The difference in surface area, $(A_S - A_T)$, is normalized by A_T so that the result is scale independent. Also, this is a representation *deficit* since $A_S \geq A_T$ is always true.

In order to apply the above definition of representation deficit to a mesh generation, a replacement for A_S must be determined since the area of the surface represented by A_S is not always able to be determined — or, most often, the area calculation is impractical. Generally, in order to reduce the representation deficit for a triangle the triangle is split by inserting an interior point. Any point that is inserted into the interior of the triangle would decrease the representation deficit — or at worst it will remain the same. However, the determination of where to split the triangle should be done in such a way that the representation deficit is minimized. This strategy of refinement, refining each triangle in such a way that the representation deficit is locally minimized, would take advantage of the optimal substructure the discrete topology.

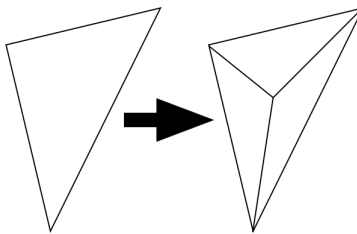


Figure 1: Triangle Split

The process of determining where to split a triangle is defined here by a locally optimizing an objective function defined for a triangle: Let $S(u, v)$ be a parameterized surface, D be a discretization of the surface, and T be a triangle (included in D) defined by an ordered tuple of nodes, (n_i, n_j, n_k) . These nodes are ordered such that the triangle normal is positive. Specifically, if $\vec{V}_0 = \{n_j - n_i\}$ and $\vec{V}_1 = \{n_k - n_i\}$ then the triangle normal, $N_T = \vec{V}_0 \times \vec{V}_1$, is positive. Note that a two-dimensional cross-product (or two-dimensional curl) is a scalar quantity. Additionally, let a node on the interior of T be defined as n_O . The four nodes, n_i, n_j, n_k , and n_O define three triangles, $T_i(n_i, n_j, n_O), T_j(n_j, n_k, n_O), T_k(n_k, n_i, n_O)$. If $A(T)$ is a function that calculates the area of a triangle in (x, y, z) space, then the optimization problem for finding the optimal position for n_O in T can be stated as:

$$\begin{aligned} \underset{n_O}{\text{minimize}} \quad O(T) &= -\frac{A(T_i) + A(T_j) + A(T_k)}{A(T)} \\ \text{subject to} \quad N_{T_i} &> 0 \\ N_{T_j} &> 0 \\ N_{T_k} &> 0 \end{aligned}$$

It should be noted that this definition of representation deficit would be difficult to derive or express for a topological entity other than a triangle. This is due to the inherent

ambiguity in the definition of not only the surface area, but also the surface representation of (possibly) non-planar elements, e.g. non-planar, or skew, quadrilateral. [MORE POSSIBLY]

5.3 Edge Split

An edge in parametric space represents a (possible) curve on the surface. Any edge, with length L_E , can have at most the same length as the portion of the surface which it represents, L_S . Therefore, the representation deficit for an edge, RD_E , can be defined as $RD_E = \frac{L_S - L_E}{L_E}$. Note that the lengths are calculated in \mathbb{R}^3 . The difference in length, $(L_S - L_E)$, is normalized by L_E so that the result is scale independent. Also, this is a representation *deficit* since $L_S \geq L_E$ is always true;

In order to apply the above definition of representation deficit to mesh generation, a replacement for L_S must be determined since the length along the surface represented by L_S is not always able to be determined – or, most often, the arc-length calculation is impractical. Generally, in order to reduce the representation deficit for an edge the edge is split by inserting an interior point. Any point that is inserted into the interior of the edge would decrease the representation deficit — or at worst it will remain the same. However, the determination of where to split the edge should be done in such a way that the representation deficit is minimized. This strategy of refinement, refining each edge in such a way that the representation deficit is locally minimized, would take advantage of the optimal substructure of the discrete topology. A method for generating a locally optimal edge split is detailed in [7, 8][OTHERS].

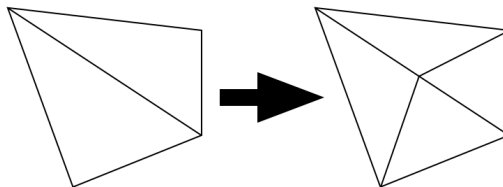


Figure 2: Edge Split

In addition, the fact that an edge split will change the surface area of the discretization should be considered. Since the overall justification of this work is reduce the *area* representation deficit of a discretization, the representation deficit will not be defined for an edge but for an edge-split. Therefore, for a given edge, only the edge-split that minimizes *area* representation deficit is considered. . The process of determining where to split a triangle is defined here by locally optimizing an objective function defined for an edge-split. Let $E(n_i, n_j)$ be an edge in D that is topologically adjacent to two triangles,

$T_i(n_i, n_j, n_k)$, and $T_j(n_i, n_l, n_j)$. Additionally, let a node on the interior of the edge be defined as n_O . The five nodes, $\{n_i, n_j, n_k, n_l, n_O\}$ define four triangles, $T_a(n_i, n_O, n_k)$, $T_b(n_O, n_j, n_k)$, $T_c(n_i, n_l, n_O)$, $T_d(n_O, n_l, n_j)$. The optimization problem for finding the optimal position for n_O on E can be stated as:

$$\begin{aligned} \underset{n_O}{\text{minimize}} \quad O(E) &= -\frac{A(T_a)+A(T_b)+A(T_c)+A(T_d)}{A(T_i)+A(T_j)} \\ \text{subject to} \quad N_{T_a} &> 0 \\ N_{T_b} &> 0 \\ N_{T_c} &> 0 \\ N_{T_d} &> 0 \end{aligned}$$

In considering the edge swap operation it became apparent that the constraint of finding a candidate point, P_c , that topologically lies on the edge, E , is an “implicit” constraint. Is it required that P_c lie on E or should P_c be allowed to be inserted where most appropriate in the hull formed by the two triangles topologically adjacent to E ?. [DO I EVEN WANT TO MENTION THIS? THEN I’LL HAVE TO GET INTO, “WELL WHY CONSTRAIN THE SEARCH TO ONLY TWO TRIANGLES? WHY NOT LET THE NODE TRAVEL UNTIL IT FINDS THE BEST PLACE?” AND THEN I’D HAVE TO COMPARE METHODS AND MORE CODING... IT’S INTERESTING, BUT I DON’T THINK THERE’S ROOM HERE. PRACTICALLY, EDGE SPLITS ARE USUALLY NEEDED SINCE THE TRIA SPLITS AND RECONNECTIONS TAKE CARE OF MOSTLY EVERYTHING.]

5.4 Edge Swap

Is it possible to have an optimal topology? If so, then it should be reachable through repeated edge flipping. Therefore, we should be able to write some sort of formulation for a pair of triangles and go from there? [SUZANNE] Comment: for the sake of not doing any more coding, if your ethics can let you come to the conclusion that this is either not possible or not practical then this would fit into the logic/justification in the next section.

5.5 Nodal Smoothing

For a node, N_i , the representation deficit can be defined only by comparing it to the node when located at another point in space. Here the comparison is bound by limiting the range of comparison within the edge-hull topologically adjacent to N_i . Formally, let N_i be a node in D that is shared topologically by n triangles, where n is the face-valence of N_i . The optimization for finding the optimal position for N_i can be stated as:

$$\begin{aligned}
\underset{N_O}{\text{minimize}} \quad & O(N) = -\sum_{j=1}^{n_t-1} A(T_j) \\
\text{subject to} \quad & N_{T_1} > 0 \\
& N_{T_2} > 0 \\
& N_{T_3} > 0 \\
& \vdots \\
& N_{T_{n-1}} > 0 \\
& N_{T_n} > 0
\end{aligned}$$

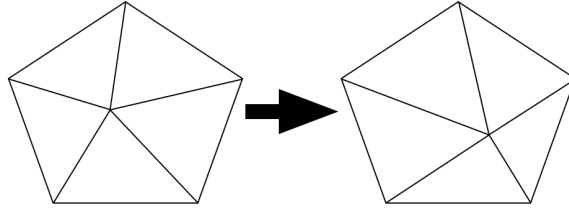


Figure 3: Nodal Smoothing

5.6 Boundary Refinement

[Is there a section needed on this? It'd be straightforward to talk about but would be pretty repetitive and I don't have it coded up yet. If I do talk about it then due diligence says I should show some results and there are already tons of variables to consider. This might get left out due to space concerns.] [This should be an argument for applying [CITE SELF] to the boundary curves first because the interior of the surface is not what the boundary curves are really trying to represent]

6 Optimal Substructure vs Global Optimization

[NEED TO COME TO THE CONCLUSION HERE THAT GLOBAL OPTIMIZATION IS NOT PRACTICAL] [COMMENT ON HOW ITERATIVE REFINEMENT IS USED HERE, NOT SETTING UP SOME GLOBAL EQUATIONS]

6.1 Local Optimization via Pattern Matching

[Suzanne] I used a simple "cross" pattern (up down left right center) and sized it based on the local geometry so that it could be moved inside of the triangles/hulls many times before encountering the boundary.

7 Element Quality

[SHOULD THIS EVEN BE A SECTION. IT NEEDS TO BE DISCUSSED BUT I'M NOT SURE THAT IT NEEDS IT'S OWN SECTION

8 Results

9 Conclusions

10 Acknowledgements

References

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