## CSCI 5535: Homework Assignment 2: Language Design and Implementation

David Baines \*\*\*
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## 1 Language Design: IMP

1.1. The two judgment forms (one for expressions, the other for functions / commands, respectively e and c in **IMP**, are:

For *e*:

For *c*:

$$\frac{\Gamma \vdash e : \mathtt{bool}}{\Gamma \vdash \mathtt{not}(e) : \mathtt{bool}}$$

1.2. (a) For the evals:

$$addr[a] val$$
  $num[n] val$   $bool[b] val$ 

<sup>\*</sup>https://courses.cs.cornell.edu/cs412/2004sp/lectures/lec13.pdf

<sup>†</sup>https://csci3155.cs.colorado.edu/csci3155-notes.pdf

<sup>&</sup>lt;sup>‡</sup>https://www.hedonisticlearning.com/posts/understanding-typing-judgments.html

For "plus()":

$$\frac{n_1 + n_2 = n}{\mathtt{plus}(\mathtt{num}[n_1]; \mathtt{num}[n_2]) \longmapsto \mathtt{num}[n]} \qquad \frac{e_1 \longmapsto e_1'}{\mathtt{plus}(e_1; e_2) \longmapsto \mathtt{plus}(e_1'; e_2)}$$

$$\frac{e \ \mathtt{val} \qquad e_2 \longmapsto e_2'}{\mathtt{plus}(e_1; e_2) \longmapsto \mathtt{plus}(e_1; e_2')}$$

For "times()":

$$\frac{n_1*n_2=n}{\texttt{times}(\texttt{num}[n_1];\texttt{num}[n_2]) \longmapsto \texttt{num}[n]} \frac{e_1 \longmapsto e_1'}{\texttt{times}(e_1;e_2) \longmapsto \texttt{times}(e_1';e_2)}$$

$$\frac{e \ \texttt{val} \quad e_2 \longmapsto e_2'}{\texttt{times}(e_1;e_2) \longmapsto \texttt{times}(e_1;e_2')}$$

For "eq()":

$$\frac{n_1 == n_2 \vdash b}{\operatorname{eq}(\operatorname{num}[n_1]; \operatorname{num}[n_2]) \longmapsto \operatorname{bool}[b]} \quad \frac{e_1 \longmapsto e_1'}{\operatorname{eq}(e_1; e_2) \longmapsto \operatorname{eq}(e_1'; e_2)} \quad \frac{e \text{ val} \quad e_2 \longmapsto e_2'}{\operatorname{eq}(e_1; e_2) \longmapsto \operatorname{eq}(e_1; e_2')}$$

For "le()":

$$\frac{n_1 <= n_2 \vdash b}{\mathtt{le}(\mathtt{num}[n_1]; \mathtt{num}[n_2]) \longmapsto \mathtt{bool}[b]} \quad \frac{e_1 \longmapsto e_1'}{\mathtt{le}(e_1; e_2) \longmapsto \mathtt{le}(e_1'; e_2)} \quad \frac{e \ \mathtt{val} \quad e_2 \longmapsto e_2'}{\mathtt{le}(e_1; e_2) \longmapsto \mathtt{le}(e_1; e_2')}$$

For "not()":

$$\frac{n_1! \, n_2 \mid\mid b_1! \, b_2 \vdash b}{ \operatorname{not}(\operatorname{num}[n_1]) \operatorname{num}[n_2] \mid\mid \operatorname{not}(\operatorname{bool}[b_1]) \operatorname{bool}[b_2] \longmapsto \operatorname{bool}[b]} \qquad \frac{e_1 \longmapsto e_1'}{ \operatorname{not}(e_1) e_2 \longmapsto \operatorname{not}(e_1') e_2}$$

$$\frac{e \, \operatorname{val} \quad e_2 \longmapsto e_2'}{ \operatorname{not}(e_1) e_2 \longmapsto \operatorname{not}(e_1) e_2'}$$

For "and()":

$$\frac{b_1 \&\& b_2 \vdash b}{\texttt{and}(\texttt{bool}[b_1]; \texttt{bool}[b_2]) \longmapsto \texttt{bool}[b]} \qquad \frac{e_1 \longmapsto e_1'}{\texttt{not}(e_1)e_2 \longmapsto \texttt{not}(e_1')e_2}$$

$$\frac{e \: \texttt{val} \quad e_2 \longmapsto e_2'}{\texttt{not}(e_1)e_2 \longmapsto \texttt{not}(e_1)e_2'}$$

## 2 Language Implementation: ETPS

See hw02.ml and test\_hw02.ml.

**3 Final Project Preparation: Pre-Proposal** 

3.1.

## A Syntax of IMP

```
Тур
       	au ::= num
                                 num
                                                          numbers
                                                          booleans
                 bool
                                  bool
                                                          addresses (or "assignables")
Exp
        e ::=
                 addr[a]
                                  a
                 num[n]
                                                          numeral
                                 n
                 bool[b]
                                                          boolean
                                  b
                 plus(e_1;e_2)
                                                          addition
                                 e_1 + e_2
                                                          multiplication
                 times(e_1;e_2)
                                 e_1 * e_2
                 eq(e_1;e_2)
                                  e_1 == e_2
                                                          equal
                 le(e_1;e_2)
                                                          less-than-or-equal
                                  e_1 <= e_2
                 not(e_1)
                                                          negation
                                  !e_1
                 and(e_1;e_2)
                                                          conjunction
                                 e_1 \&\& e_2
                 or(e_1;e_2)
                                                          disjunction
                                  e_1 || e_2
\mathsf{Cmd} \quad c ::=
                 set[a](e)
                                                          assignment
                                 a := e
                 skip
                                  skip
                                                          skip
                 seq(c_1;c_2)
                                                          sequencing
                                 c_1; c_2
                 if(e;c_1;c_2)
                                  if e then c_1 else c_2
                                                          conditional
                                                          looping
                 while(e; c_1)
                                 while e 	ext{ do } c_1
Addr a
```