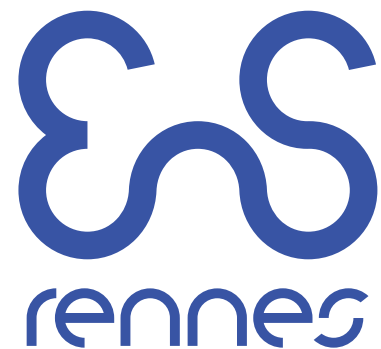


# Compilation vérifiée

David Pichardie - David Cachera



école   
normale  
supérieure

# Les logiciels critiques

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Certains logiciels demandent un très haut niveau de confiance

# Les logiciels critiques

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- est-ce que ce programme est sûr ?

```
.const 2
.align 2
__stringlit_1:
.asciz "integr(square, 0.0, 1.0,
.text
.align 4
_square:
subl $12, %esp
leal 16(%esp), %edx
movl %edx, 0(%esp)
movl 0(%esp), %edx
movsd 0(%edx), %xmm0
mulsd %xmm0, %xmm0
movsd %xmm0, 0(%esp)
fldl 0(%esp)
addl $8, %esp
addl $12, %esp
ret
.text
.align 4
_integr:
subl $60, %esp
leal 64(%esp), %edx
movl %edx, 8(%esp)
movl %edx, 20(%esp)
movl %esi, 24(%esp)
movl 8(%esp), %edx
movl 0(%edx), %esi
movl 8(%esp), %edx
movsd 4(%edx), %xmm6
movsd %xmm6, 48(%esp)
movl 8(%esp), %edx
movsd 12(%edx), %xmm6
movsd %xmm6, 32(%esp)
movl 8(%esp), %edx
movl 20(%edx), %ebx
movsd 32(%esp), %xmm6
movsd 48(%esp), %xmm7
subsd %xmm7, %xmm6
movsd %xmm6, 32(%esp)
cvtsizsd %ebx, %xmm1
movsd 32(%esp), %xmm6
divsd %xmm1, %xmm6
movsd %xmm6, 32(%esp)
xorpd %xmm6, %xmm6 # +0.0
movsd %xmm6, 40(%esp)
L100:
cmpl $0, %ebx
jle L101
movsd 48(%esp), %xmm6
movsd %xmm6, 0(%esp)
call *%esi
subl $8, %esp
fstpl 0(%esp)
movsd 0(%esp), %xmm0
addl $8, %esp
movsd 40(%esp), %xmm6
addsd %xmm0, %xmm6
movsd %xmm6, 40(%esp)
leal -1(%ebx), %ebx
movsd 48(%esp), %xmm6
movsd 32(%esp), %xmm7
addsd %xmm7, %xmm6
movsd %xmm6, 48(%esp)
jmp L100
L101:
movsd 40(%esp), %xmm6
movsd 32(%esp), %xmm7
mulsd %xmm7, %xmm6
movsd %xmm6, 40(%esp)
fldl 40(%esp)
movl 20(%esp), %ebx
movl 24(%esp), %esi
addl $60, %esp
ret
.text
.align 4
.globl _test
_test:
subl $44, %esp
leal 48(%esp), %edx
movl %edx, 24(%esp)
movl %ebx, 36(%esp)
movl 24(%esp), %edx
movl 0(%edx), %eax
leal _square, %ebx
xorpd %xmm2, %xmm2 # +0.0
movsd L102, %xmm1 # 1
movl %ebx, 0(%esp)
movsd %xmm2, 4(%esp)
movsd %eax, 12(%esp)
call _integr
subl $8, %esp
fstpl 0(%esp)
movsd 0(%esp), %xmm0
addl $8, %esp
subl $8, %esp
movsd %xmm0, 0(%esp)
fldl 0(%esp)
addl $8, %esp
movl 36(%esp), %ebx
addl $44, %esp
ret
.const_data
.align 3
L102:
.quad 0x3ff0000000000000
.text
.align 4
.globl _main
_main:
subl $44, %esp
leal 48(%esp), %edx
movl %edx, 16(%esp)
movl %ebx, 28(%esp)
movl 16(%esp), %edx
movl 0(%edx), %ebx
movl 16(%esp), %edx
movl 4(%edx), %eax
cmpl $2, %ebx
jge L103
movl $100000000, %ebx
jmp L104
L103:
movl 4(%eax), %eax
movl %eax, 0(%esp)
call _atoi
movl %eax, %ebx
L104:
movl %ebx, 0(%esp)
call _test
subl $8, %esp
fstpl 0(%esp)
movsd 0(%esp), %xmm0
addl $8, %esp
leal __stringlit_1, %eax
movl %eax, 0(%esp)
movl %ebx, 4(%esp)
movsd %xmm0, 8(%esp)
call _printf
xorl %eax, %eax
movl 28(%esp), %ebx
addl $44, %esp
ret
```

critical\_prog.ppc

# Les logiciels critiques

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- est-ce que ce programme est sûr ?

Deux options:

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.const 2
.align 2
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leal 16(%esp), %edx
movl %edx, 0(%esp)
movl 0(%esp), %edx
movsd 0(%edx), %xmm0
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fldl 0(%esp)
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_integr:
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movl %edx, 8(%esp)
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movl 0(%edx), %esi
movl 8(%esp), %edx
movsd 4(%edx), %xmm6
movsd %xmm6, 48(%esp)
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subsd %xmm7, %xmm6
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cvtsizsd %ebx, %xmm1
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divsd %xmm1, %xmm6
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fstpl 0(%esp)
movsd 0(%esp), %xmm0
addl $8, %esp
movsd 40(%esp), %xmm6
addsd %xmm0, %xmm6
movsd %xmm6, 40(%esp)
leal -1(%ebx), %ebx
movsd 48(%esp), %xmm6
movsd 32(%esp), %xmm7
addsd %xmm7, %xmm6
movsd %xmm6, 48(%esp)
jmp L100
L101:
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movsd 32(%esp), %xmm7
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movl 20(%esp), %ebx
movl 24(%esp), %esi
addl $60, %esp
ret
.text
.align 4
.globl _test
_test:
subl $44, %esp
leal 48(%esp), %edx
movl %edx, 24(%esp)
movl %ebx, 36(%esp)
movl 24(%esp), %edx
movl 0(%edx), %eax
leal _square, %ebx
xorpd %xmm2, %xmm2 # +0.0
movsd L102, %xmm1 # 1
movl %ebx, 0(%esp)
movsd %xmm2, 4(%esp)
movsd %eax, 12(%esp)
call _integr
subl $8, %esp
fstpl 0(%esp)
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movl 16(%esp), %edx
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cmpl $2, %ebx
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call _test
subl $8, %esp
fstpl 0(%esp)
movsd 0(%esp), %xmm0
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leal __stringlit_1, %eax
movl %eax, 0(%esp)
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movsd %xmm0, 8(%esp)
call _printf
xorl %eax, %eax
movl 28(%esp), %ebx
addl $44, %esp
ret
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critical\_prog.ppc

# Les logiciels critiques

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- est-ce que ce programme est sûr ?

Deux options:

- *qualifier* le programme PPC comme s'il avait été écrit à la main  
(tests intensive, revue manuel de code...)

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movl %edx, 0(%esp)
movl 0(%esp), %edx
movsd 0(%edx), %xmm0
mulsd %xmm0, %xmm0
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fldl 0(%esp)
addl $8, %esp
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_integr:
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movl %edx, 20(%esp)
movl %esi, 24(%esp)
movl 8(%esp), %edx
movl 0(%edx), %esi
movl 8(%esp), %edx
movsd 4(%edx), %xmm6
movsd %xmm6, 48(%esp)
movl 8(%esp), %edx
movsd 12(%edx), %xmm6
movsd %xmm6, 32(%esp)
movl 8(%esp), %edx
movl 20(%edx), %ebx
movsd 32(%esp), %xmm6
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cvtsizsd %ebx, %xmm1
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xorpd %xmm6, %xmm6 # +0.0
movsd %xmm6, 40(%esp)
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movsd %xmm6, 0(%esp)
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movsd 0(%esp), %xmm0
addl $8, %esp
movsd 40(%esp), %xmm6
addsd %xmm0, %xmm6
movsd %xmm6, 40(%esp)
leal -1(%ebx), %ebx
movsd 48(%esp), %xmm6
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movsd %xmm6, 48(%esp)
jmp L100
L101:
movsd 40(%esp), %xmm6
movsd 32(%esp), %xmm7
mulsd %xmm7, %xmm6
movsd %xmm6, 40(%esp)
fldl 40(%esp)
movl 20(%esp), %ebx
movl 24(%esp), %esi
addl $60, %esp
ret
.text
.align 4
.globl _test
_test:
subl $44, %esp
leal 48(%esp), %edx
movl %edx, 24(%esp)
movl %ebx, 36(%esp)
movl 24(%esp), %edx
movl 0(%edx), %eax
leal _square, %ebx
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call _integr
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movsd 0(%esp), %xmm0
addl $8, %esp
subl $8, %esp
movsd %xmm0, 0(%esp)
fldl 0(%esp)
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movsd 0(%esp), %xmm0
addl $8, %esp
leal __stringlit_1, %eax
movl %eax, 0(%esp)
movl %ebx, 4(%esp)
movsd %xmm0, 8(%esp)
call _printf
xorl %eax, %eax
movl 28(%esp), %ebx
addl $44, %esp
ret
```

critical\_prog.ppc

# Les logiciels critiques

## Certains logiciels demandent un très haut niveau de confiance

- est-ce que ce programme est sûr ?

## Deux options:

- *qualifier* le programme PPC comme s'il avait été écrit à la main (tests intensive, revue manuel de code...)
- *qualifier* le programme au niveau source (analyse statique, model checking, preuve de programme)

```
#include <math.h>
#include <stdio.h>

static double square(double x)
{
    return x * x;
}

static double integr(double (*f)(double), double low, double high, int n)
{
    double h, x, s;
    int i;

    h = (high - low) / n;
    s = 0;
    for (i = n, x = low; i > 0; i--, x += h) s += f(x);
    return s * h;
}

double test(int n)
{
    return integr(square, 0.0, 1.0, n);
}

int main(int argc, char ** argv)
{
    int n; double r;
    if (argc >= 2) n = atoi(argv[1]); else n = 10000000;
    r = test(n);
    printf("integr(square, 0.0, 1.0, %d) = %g\n", n, r);
    return 0;
}
```

critical\_prog.c

# compilateur

[illegible]

critical prog.ppc



# Les logiciels critiques

## Certains logiciels demandent un très haut niveau de confiance

- est-ce que ce programme est sûr ?

## Deux options:

- *qualifier* le programme PPC comme s'il avait été écrit à la main (tests intensive, revue manuel de code...)
  - *qualifier* le programme au niveau source (analyse statique, model checking, preuve de programme)
- La seconde option est préférée en pratique
- comment faire confiance au compilateur ?
  - cet exposé : comment vérifier le compilateur !

## La seconde option est préférée en pratique

```
#include <stdlib.h>
#include <stdio.h>

static double square(double x)
{
    return x * x;
}

static double integr(double (*f)(double), double low, double high, int n)
{
    double h, x, s;
    int i;

    h = (high - low) / n;
    s = 0;
    for (i = n, x = low; i > 0; i--, x += h) s += f(x);
    return s * h;
}

double test(int n)
{
    return integr(square, 0.0, 1.0, n);
}

int main(int argc, char ** argv)
{
    int n; double r;
    if (argc >= 2) n = atoi(argv[1]); else n = 10000000;
    r = test(n);
    printf("integr(square, 0.0, 1.0, %d) = %g\n", n, r);
    return 0;
}
```

critical\_prog.c

# compilateur

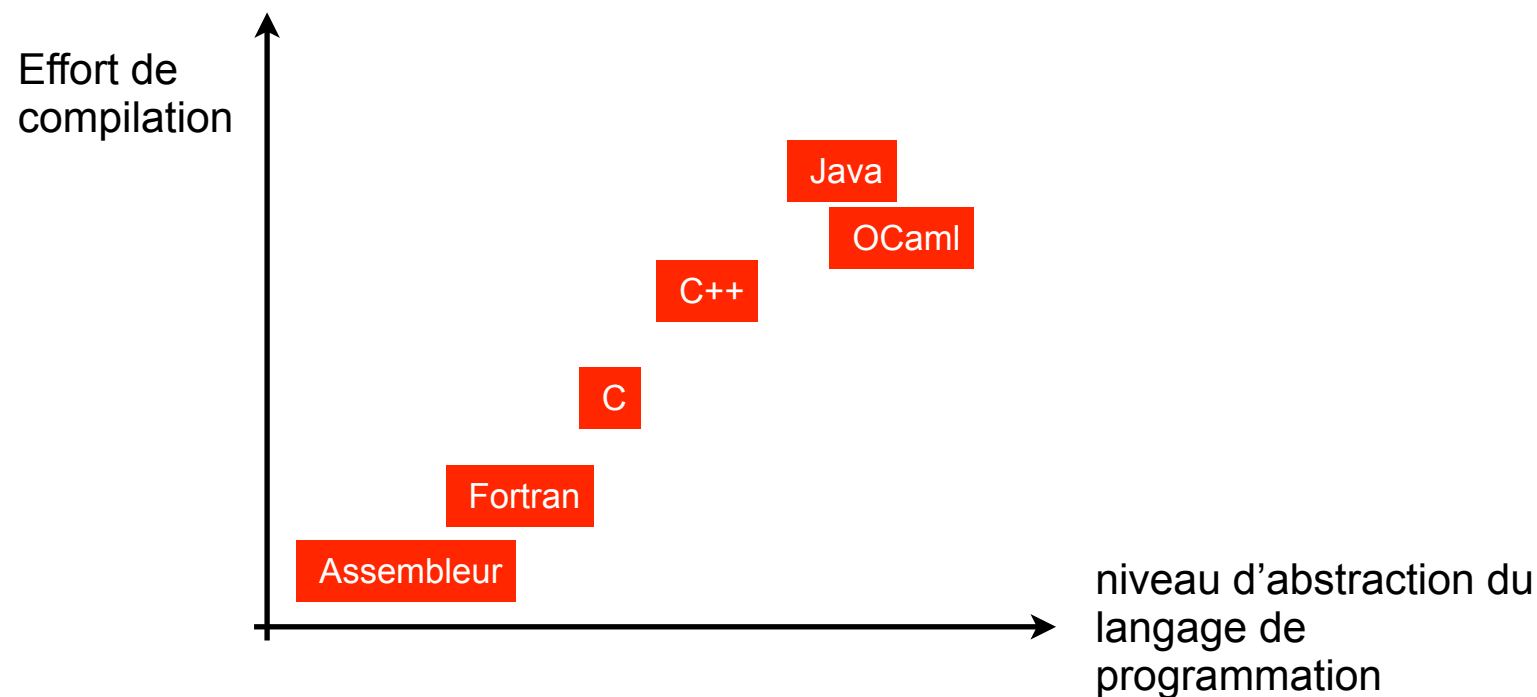
[illegible]

critical prog.ppc



# Compiler : est-ce difficile ?

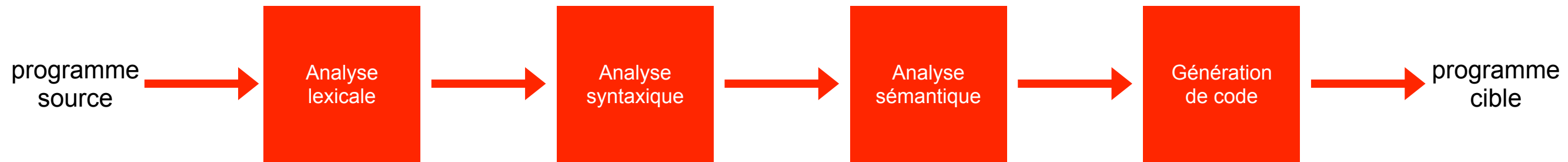
Les langages de haut-niveau d'abstraction diminuent le travail du programmeur mais augmentent celui du compilateur.



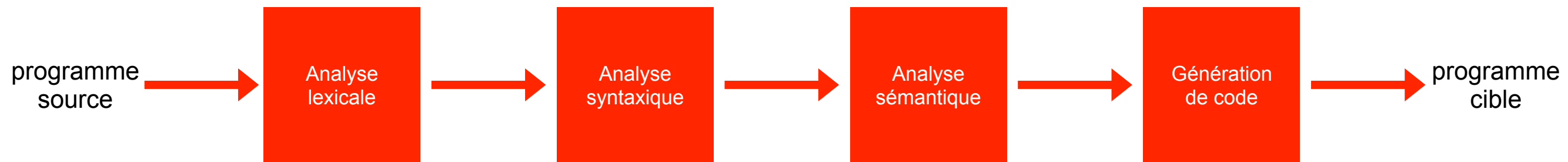
## Autres difficultés

- il faut passer à l'échelle des programmes gigantesques
- les langages dynamiques (Java, Javascript) doivent être compilés pendant l'exécution : *Just-in-time compiler*
- il faut tenir compte des spécificités de chaque architecture cible

# Compilation : les grandes étapes



# Compilation : les grandes étapes



Théorie des  
automates finis et  
des langages  
rationnels



Théorie des  
automates à pile et  
des langages  
algébriques



Sémantique des  
langages de  
programmation

Interprétation  
abstraite



Algorithmique des  
problèmes NP-complets

Algorithmique parallèle

Architecture des  
ordinateurs

# Compilation optimisante

Un exemple : le produit scalaire

```
double p = 0.0;
for (int i = 0; i < N, i++)
    p = p + A[i] * B[i];
return p;
```

# Compilation optimisante

Passage en code intermédiaire *3-adresses*

```
double p = 0.0;
for (int i = 0; i < N, i++)
    p = p + A[i] * B[i];
return p;
```

```
p = 0.0;
for (int i = 0; i < N, i++)
    a = load(A + i * 8);
    b = load(B + i * 8);
    c = a * b;
    p = p + c;
}
...
```

# Compilation optimisante

Les multiplications sont coûteuses...

```
p = 0.0;
for (int i = 0; i < N, i++)
    a = load(A + i * 8);
    b = load(B + i * 8);
    c = a * b;
    p = p + c;
}
...
```

```
p = 0.0;
for (int i = 0; i < N, i++)
    a = load(A);
    b = load(B);
    c = a * b;
    p = p + c;
    A = A + 8;
    B = B + 8;
}
...
```



# Compilation optimisante

On regroupe l'itération  $i$  et  $i+1$  dans un même corps de boucle

```
p = 0.0;
for (int i = 0; i < N, i++)
    a = load(A);
    b = load(B);
    c = a * b;
    p = p + c;
    A = A + 8;
    B = B + 8;
}
...
```

```
p = 0.0;
for (int i = 0; i < N-1, i += 2)
    a1 = load(A);
    b1 = load(B);
    c1 = a1 * b1;
    p = p + c1;
    a2 = load(A + 8);
    b2 = load(B + 8);
    c2 = a2 * b2;
    p = p + c2;
    A = A + 16;
    B = B + 16;
}
if (i < N) { ... }
```

# Compilation optimisante

On réordonne certaines instructions pour profiter du pipeline d'instructions

```
p = 0.0;
for (int i = 0; i < N-1, i += 2)
    a1 = load(A);
    b1 = load(B);
    c1 = a1 * b1;
    p = p + c1;
    a2 = load(A + 8);
    b2 = load(B + 8);
    c2 = a2 * b2;
    p = p + c2;
    A = A + 16;
    B = B + 16;
}
if (i < N) { ... }
```

```
p = 0.0;
for (int i = 0; i < N-1, i += 2)
    a1 = load(A);
    b1 = load(B);
    a2 = load(A + 8);
    b2 = load(B + 8);

    c1 = a1 * b1;
    c2 = a2 * b2;
    A = A + 16;

    p = p + c1;
    B = B + 16;

    p = p + c2;
}
if (i < N) { ... }
```

# Compilation optimisante

On effectue les lectures mémoires une itération en avance

```
p = 0.0;
for (int i = 0; i < N-1, i += 2)
    a1 = load(A);
    b1 = load(B);
    a2 = load(A + 8);
    b2 = load(B + 8);

    c1 = a1 * b1;
    c2 = a2 * b2;
    A = A + 16;

    p = p + c1;
    B = B + 16;

    p = p + c2;
}
if (i < N) { ... }
```

```
...
a3 = load(A); a4 = load(A + 8);
b3 = load(B); b4 = load(B + 8);
for (int i = 0; i < N-3, i += 2)
    a1 = a3; a3 = load(A + 16);
    b1 = b3; b3 = load(B + 16);
    c1 = a1 * b1;
    a2 = a4; a4 = load(A + 24);
    b2 = b4; b4 = load(B + 24);
    c2 = a2 * b2;
    A = A + 16;
    p = p + c1;
    B = B + 16;
    p = p + c2;
}
c1 = a3 * b3;
c2 = a4 * b4;
p = p + c1;
p = p + c2;
```

# Compilation optimisante

Version complète

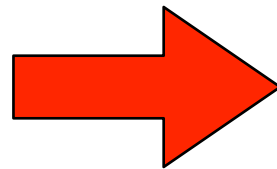
```
...
a3 = load(A); a4 = load(A + 8);
b3 = load(B); b4 = load(B + 8);
for (int i = 0; i < N-3, i += 2)
    a1 = a3; a3 = load(A + 16);
    b1 = b3; b3 = load(B + 16);
    c1 = a1 * b1;
    a2 = a4; a4 = load(A + 24);
    b2 = b4; b4 = load(B + 24);
    c2 = a2 * b2;
    A = A + 16;
    p = p + c1;
    B = B + 16;
    p = p + c2;
}
c1 = a3 * b3;
c2 = a4 * b4;
p = p + c1;
p = p + c2;
```

```
p = 0.0;
i = 0;
if (N >=4) {
    a3 = load(A); a4 = load(A + 8);
    b3 = load(B); b4 = load(B + 8);
    for (; i < N-3, i += 2)
        a1 = a3; a3 = load(A + 16);
        b1 = b3; b3 = load(B + 16);
        c1 = a1 * b1;
        a2 = a4; a4 = load(A + 24);
        b2 = b4; b4 = load(B + 24);
        c2 = a2 * b2;
        A = A + 16; p = p + c1;
        B = B + 16; p = p + c2;
    }
    c1 = a3 * b3;
    c2 = a4 * b4;
    p = p + c1; p = p + c2;
    A = A + 16; B = B + 16;
}
for (; i < N; i++) {
    a = load(A); A = A + 8;
    b = load(B); B = B + 8;
    c = a * b;
    p = p + c;
}
```

# Compilation optimisante

Version complète

```
double p = 0.0;
for (int i = 0; i < N, i++)
    p = p + A[i] * B[i];
return p;
```



programme  
2 à 3 fois plus  
efficace

```
p = 0.0;
i = 0;
if (N >=4) {
    a3 = load(A); a4 = load(A + 8);
    b3 = load(B); b4 = load(B + 8);
    for (; i < N-3, i += 2)
        a1 = a3; a3 = load(A + 16);
        b1 = b3; b3 = load(B + 16);
        c1 = a1 * b1;
        a2 = a4; a4 = load(A + 24);
        b2 = b4; b4 = load(B + 24);
        c2 = a2 * b2;
        A = A + 16; p = p + c1;
        B = B + 16; p = p + c2;
    }
    c1 = a3 * b3;
    c2 = a4 * b4;
    p = p + c1; p = p + c2;
    A = A + 16; B = B + 16;
}
for (; i < N; i++) {
    a = load(A); A = A + 8;
    b = load(B); B = B + 8;
    c = a * b;
    p = p + c;
}
```

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*We created a tool that generates random C programs, and then spent two and a half years using it to find compiler bugs. So far, we have reported more than 290 previously unknown bugs to compiler developers. Moreover, every compiler that we tested has been found to crash and also to silently generate wrong code when presented with valid inputs.*

J. Regehr et al, 2010

# Le projet CompCert

Xavier Leroy et al.

Un compilateur développé et prouvé correct, dans l'assistant de preuve Coq

- langage source: un très grand sous-ensemble de C
- langage cible : assembleur PowerPC/ARM/x86
- génère du code *raisonnablement compacte et efficace*  
⇒ génération de code minutieuse; quelques optimisations
- en cours de transfert pour l'industrie aéronautique...

# Comment vérifier un compilateur ?

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# Vérifier un compilateur en Coq

## La recette

Assistant de preuve

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Programmer le compilateur en Coq

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Definition compiler  
  (p:source) : option target := ...
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Écrire son théorème de correction, en s'appuyant sur une description formelles des sémantiques des langages

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Theorem compiler_is_correct :  
   $\forall p, \text{compiler } p = \text{Some } p' \rightarrow$   
    behaviors(p')  $\subseteq$  behaviors(p)
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Proof. ... (* few days later *) ... Qed.
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Nous utilisons l'*extraction* de Coq pour générer un compilateur écrit en OCaml

```
Extraction compiler.
```

Compilateur

Sémantiques  
sources & cibles

Preuve de correction

Assistant de preuve



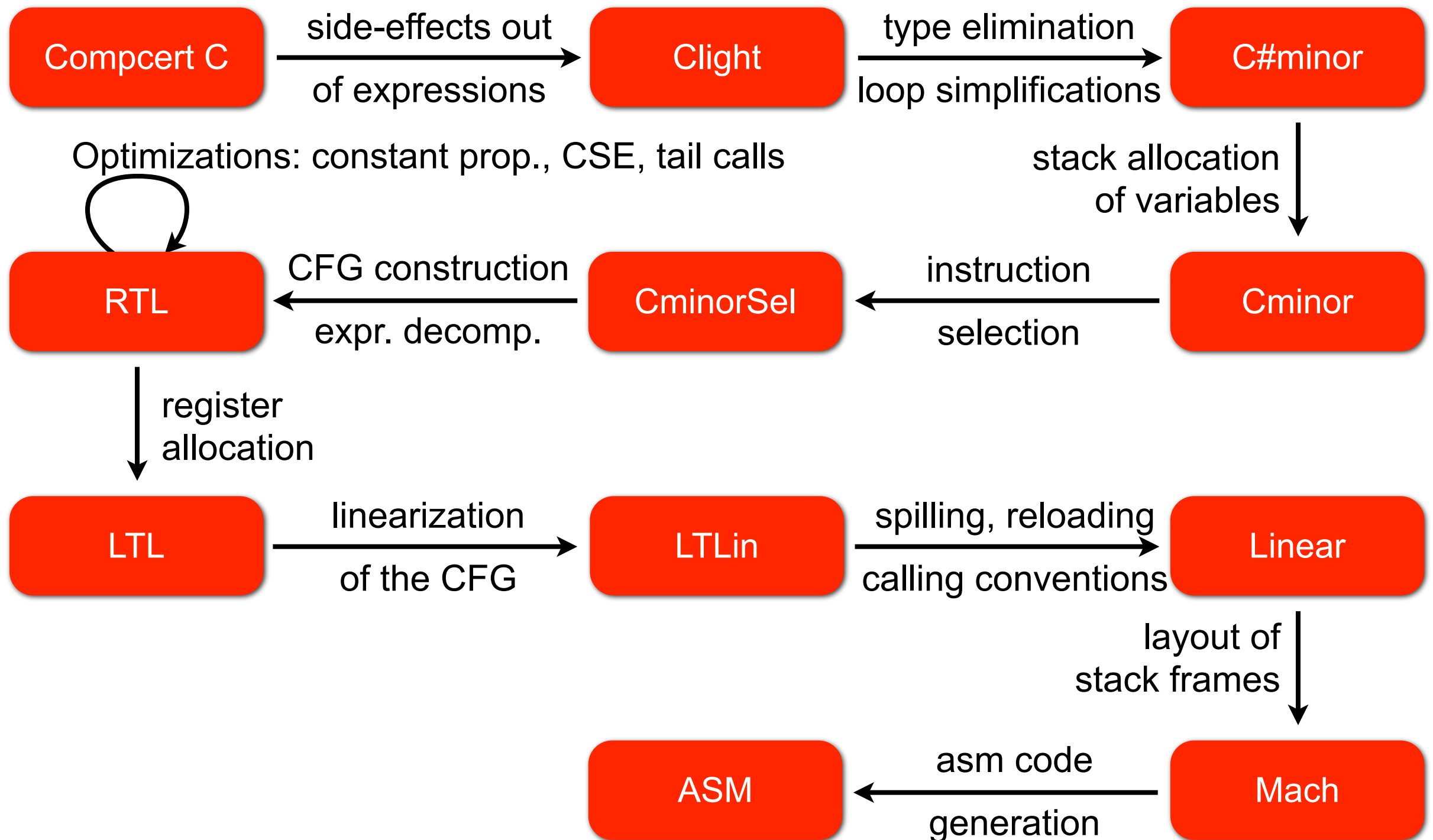
parser.ml

compiler.ml

linker.ml



# Les *entrailles* d'un compilateur C vérifié



# Le théorème principal

Après 50 000 lignes de Coq et 4 hommes/ans d'effort

**Theorem** `transf_c_program_is_refinement :`

$\forall p \text{ } tp,$   
 $\text{transf\_c\_program } p = \text{OK } tp \rightarrow$   
 $(\forall \text{ beh, exec\_C\_program } p \text{ beh} \rightarrow \text{not\_wrong beh}) \rightarrow$   
 $(\forall \text{ beh, exec\_asm\_program } tp \text{ beh} \rightarrow \text{exec\_C\_program } p \text{ beh}).$

Le comportement `beh` modélise la termination ou la divergence ou la fin abrupte (erreur) d'une exécution.

# Compiler verification patterns

(for each pass)

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## Verified transformation

transformation



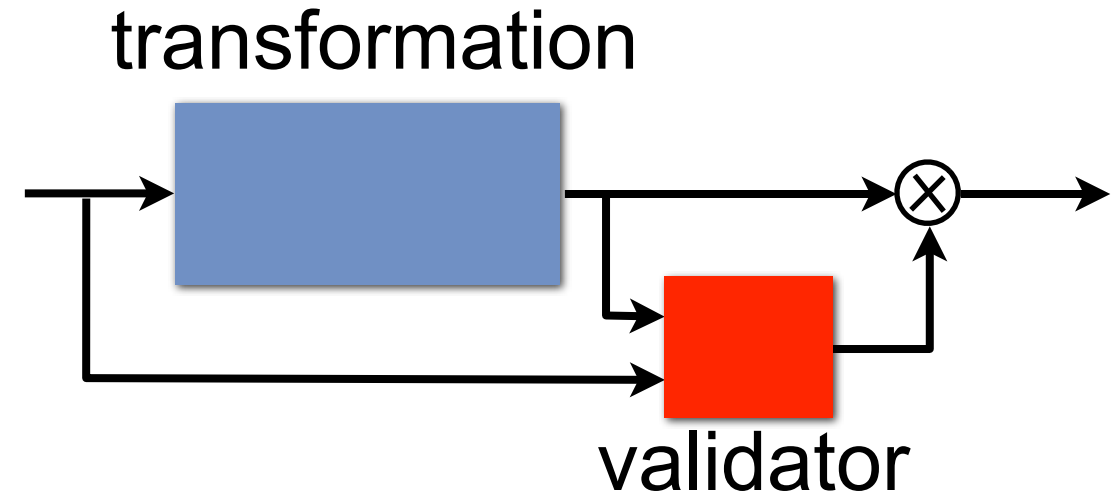
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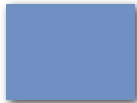
(for each pass)

## Verified transformation



## Verified translation validation



 = formally verified  
 = not verified

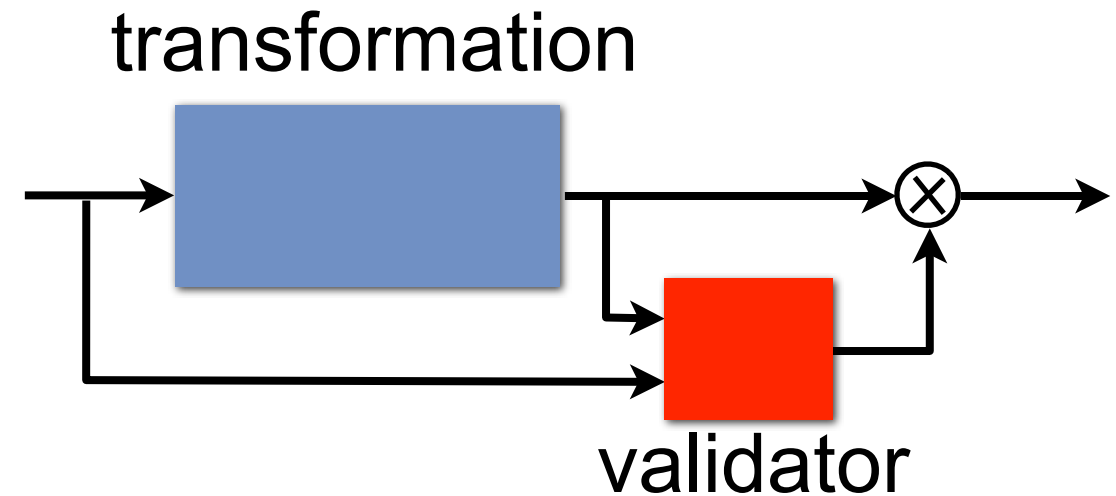
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(for each pass)

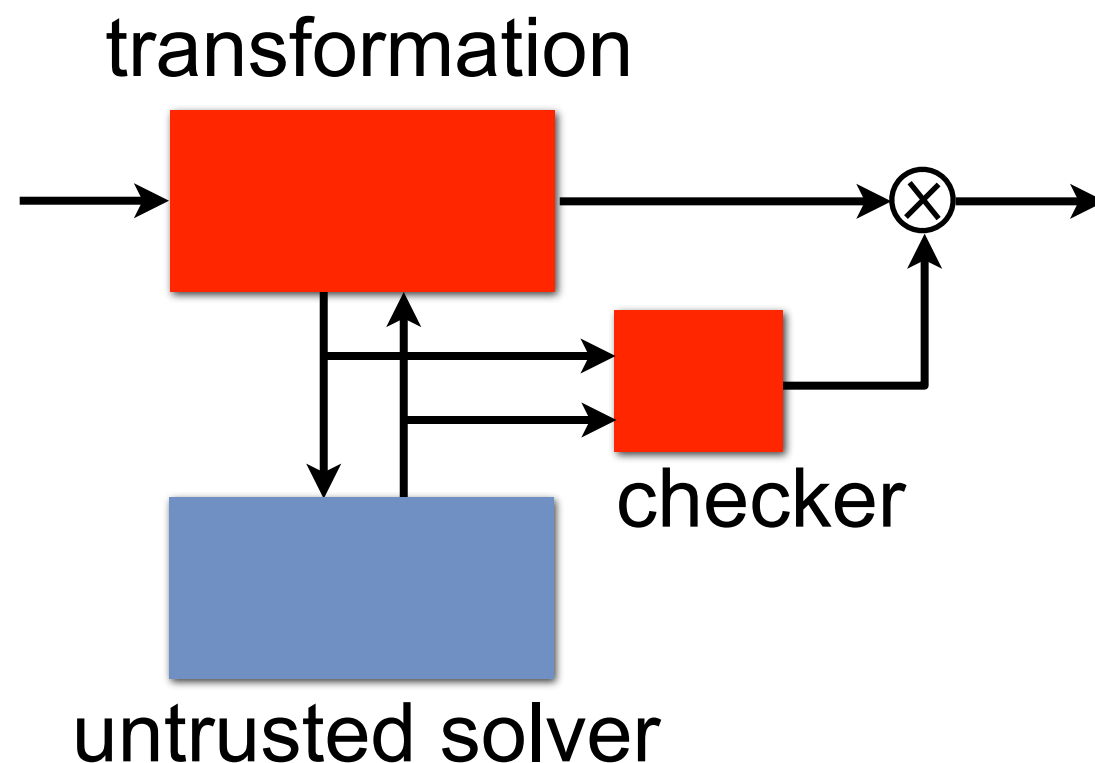
## Verified transformation





## Verified translation validation



## External solver with verified validation

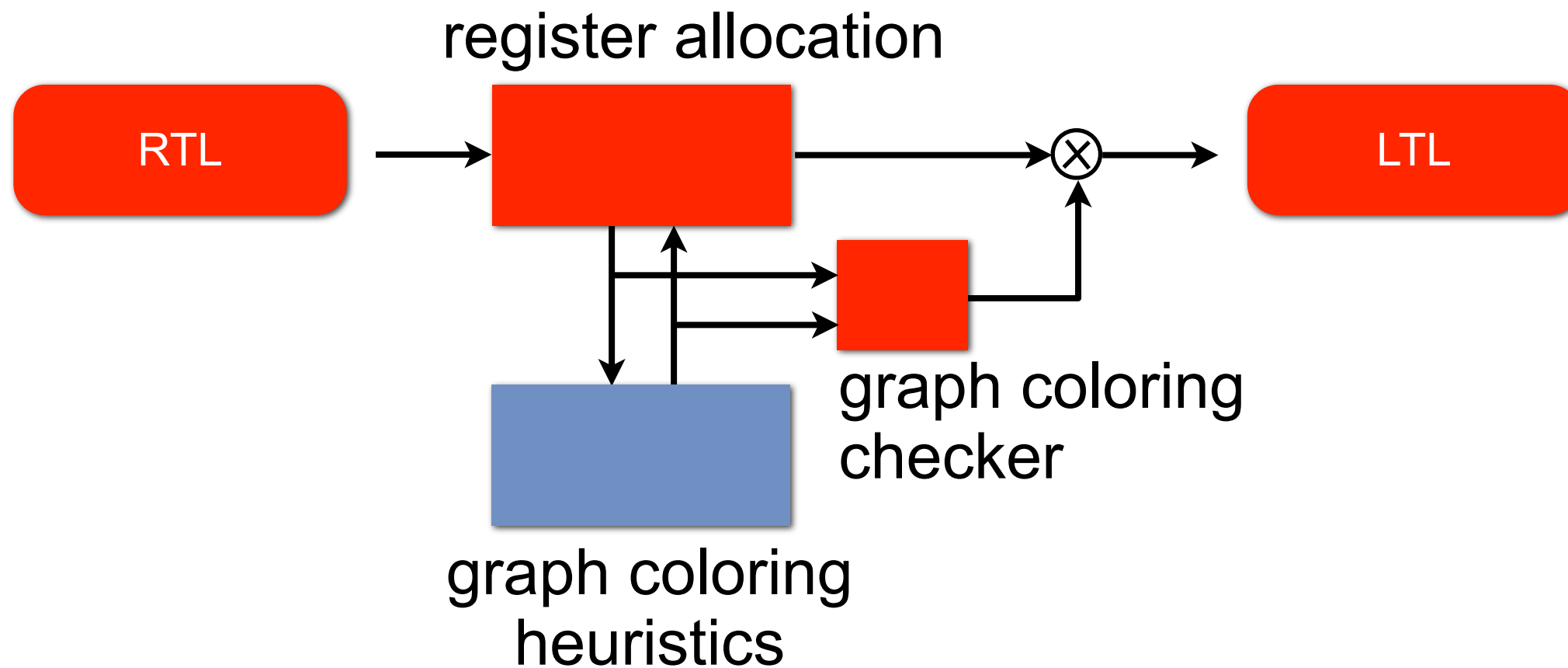


 = formally verified  
 = not verified



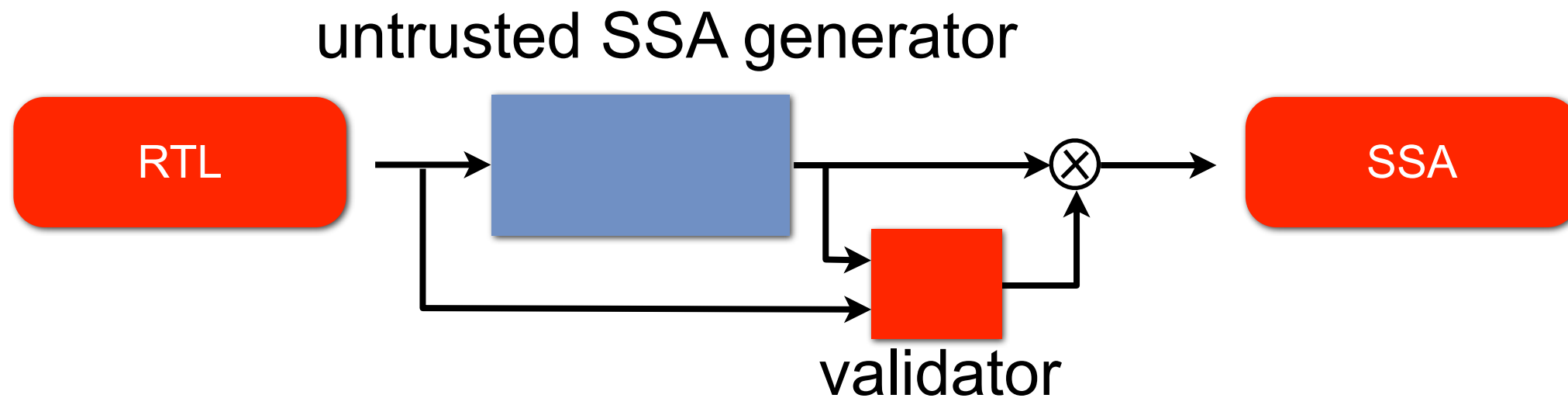
# External solver with verified validation

Example: register allocation



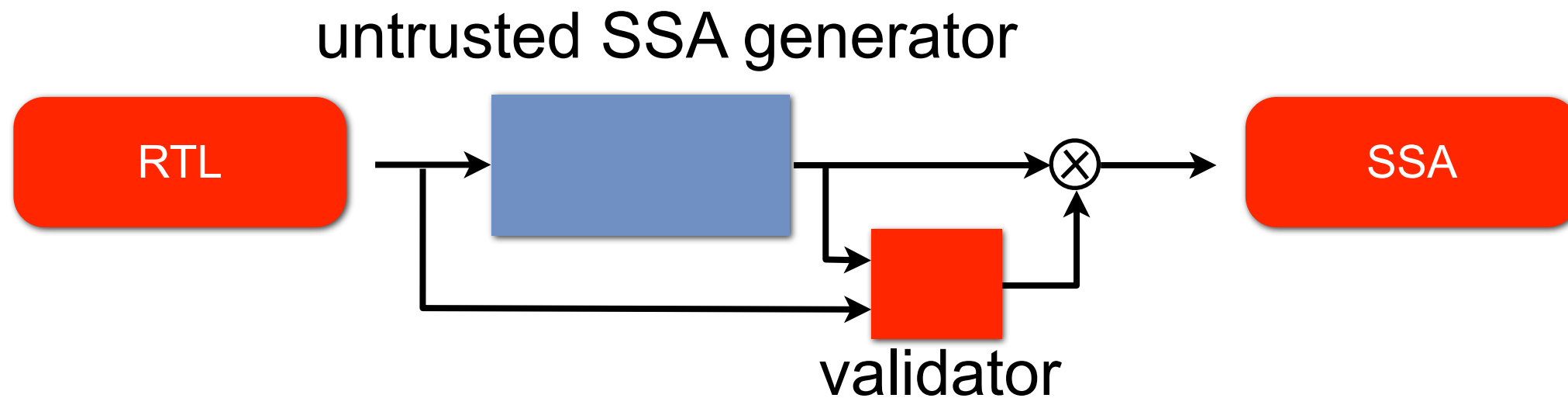
# Verified translation validation

Example: SSA generation (in CompCert SSA extension)



# Verified translation validation

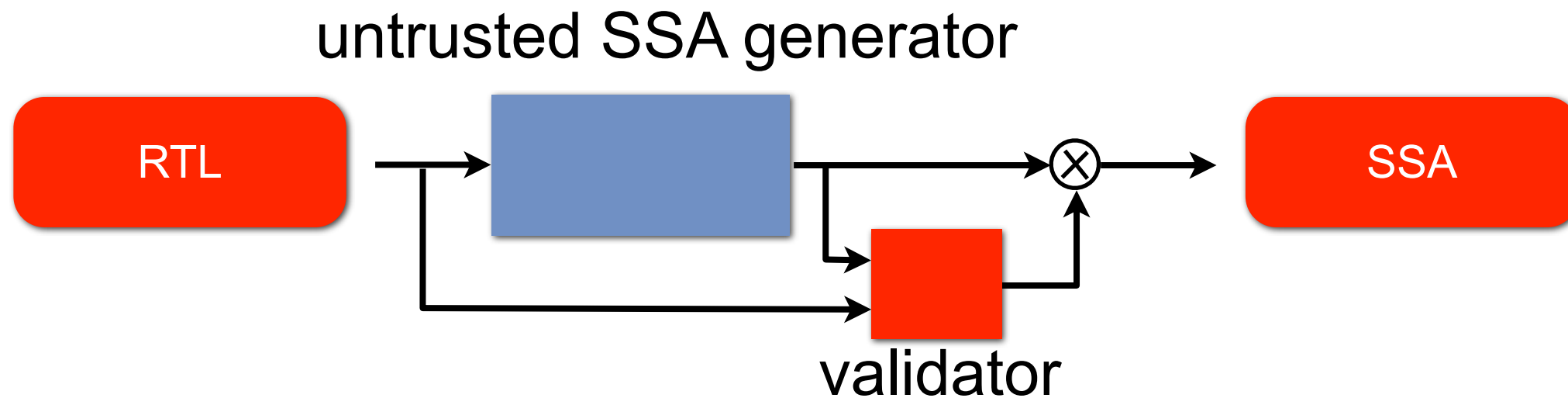
Example: SSA generation (in CompCert SSA extension)



The untrusted generator can rely on advanced graph algorithms as Lengauer and Tarjan's dominator tree construction and frontier dominance computation.

# Verified translation validation

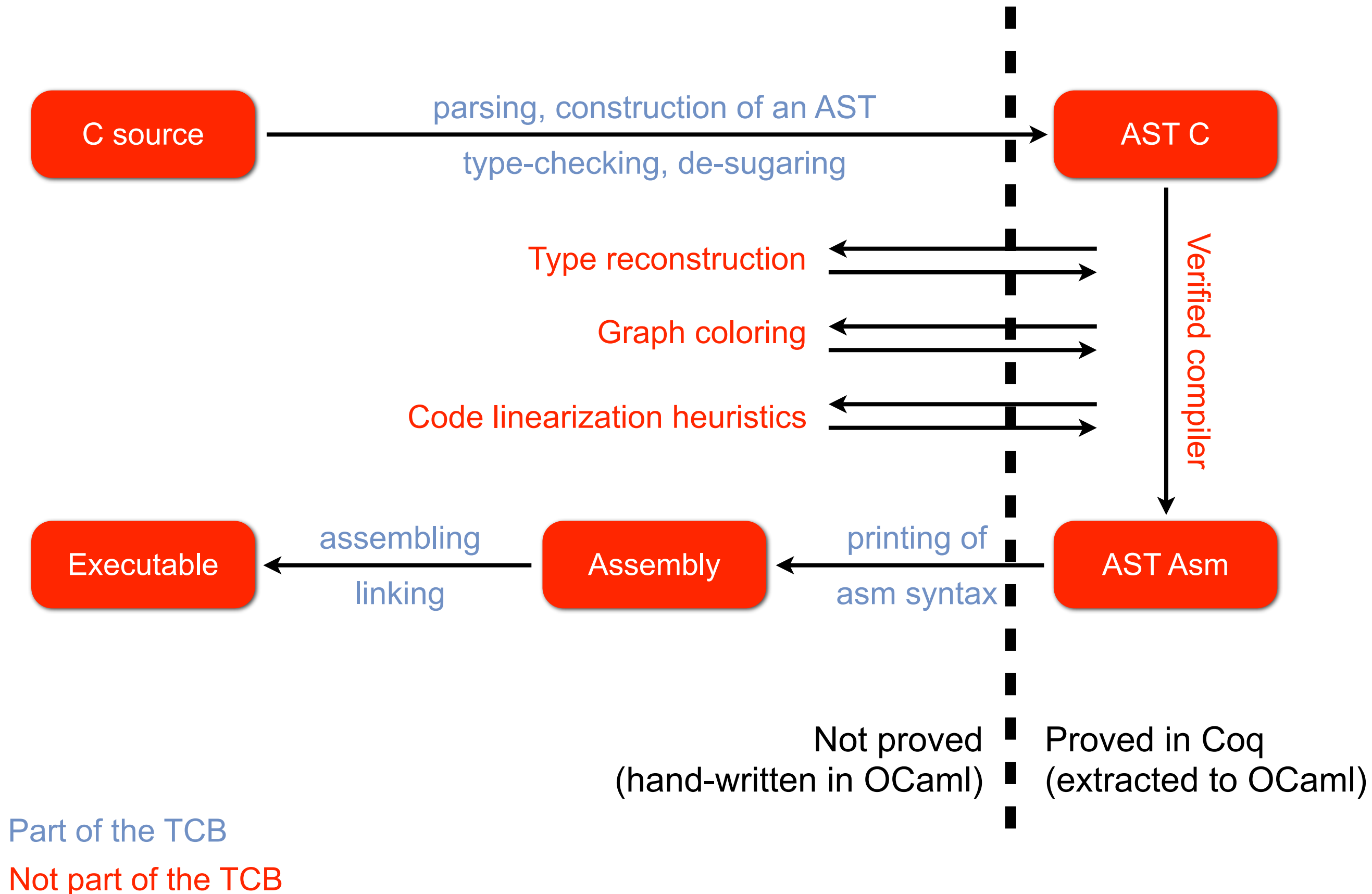
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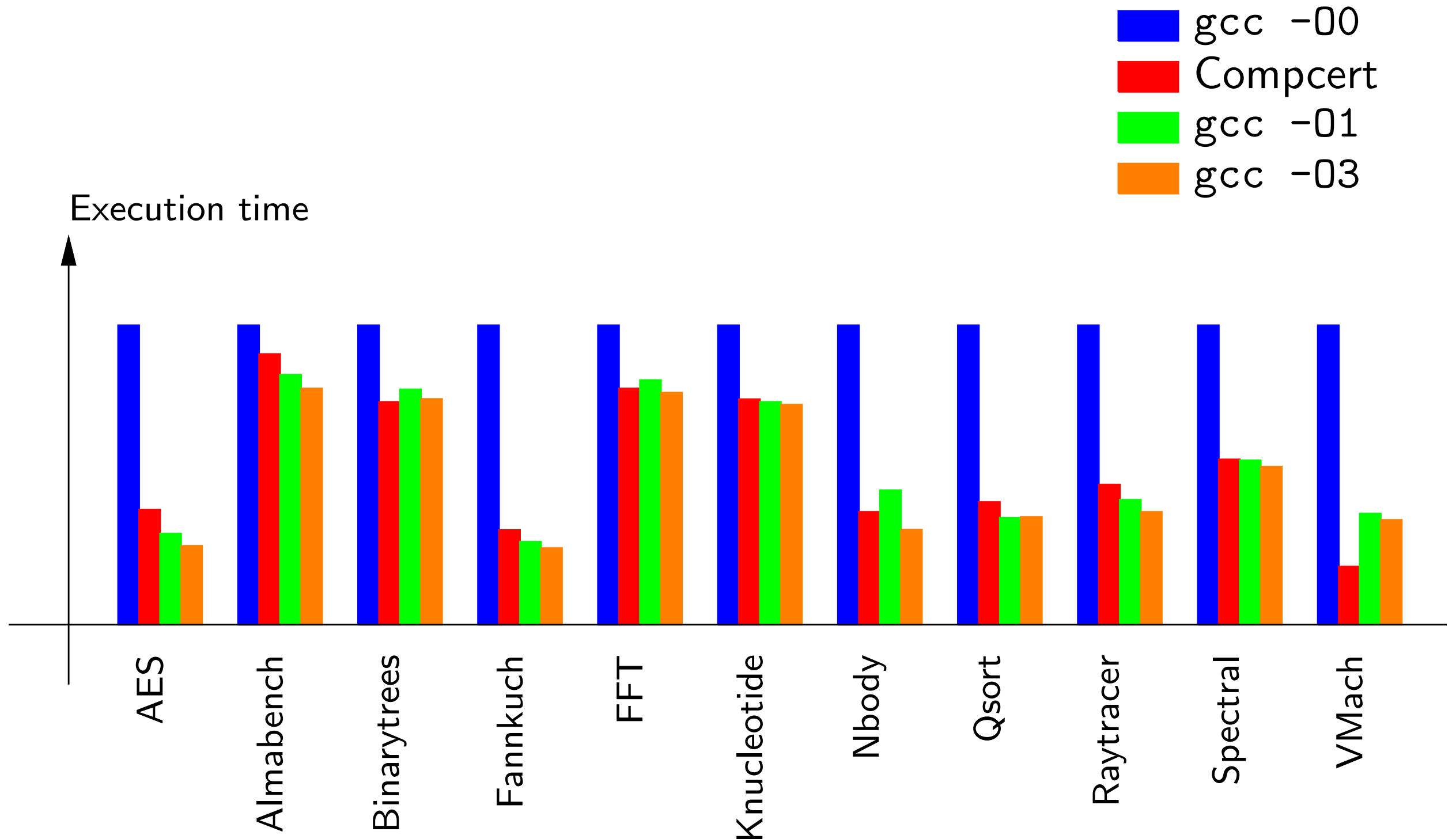
We prove the validator is *complete* with respect to this family of algorithms.

# The whole CompCert compiler



# Performance of generated code

(On a PowerPC G5 processor)



# Conclusions

La vérification formelle d'un compilateur réaliste est *faisable*.

## Encore de nombreux défis

- Réduire la base de confiance  
(e.g. analyse lexicale/syntaxique, assemblage, édition de liens).
- Plus d'optimisations  
(par exemple CompCert SSA).
- Cible pour des langages sources de plus haut niveau (Caml, Java)
- Parallélisme des langages sources/cibles
- Connexion avec vérificateur source *vérifié*  
(projet ANR Verasco sur un analyseur statique C vérifié)

# Pour aller plus loin

## Compilation

- Du langage à l'action: compilation et typage, Xavier Leroy, Collège de France 2008

[http://pauillac.inria.fr/~xleroy/talks/compilation\\_typage\\_College\\_de\\_France.pdf](http://pauillac.inria.fr/~xleroy/talks/compilation_typage_College_de_France.pdf)

## Découvrir Coq

- Preuves de programme en Coq, Yves Bertot, Fuscia.info

<http://fuscia.inrialpes.fr/cours/coq/>