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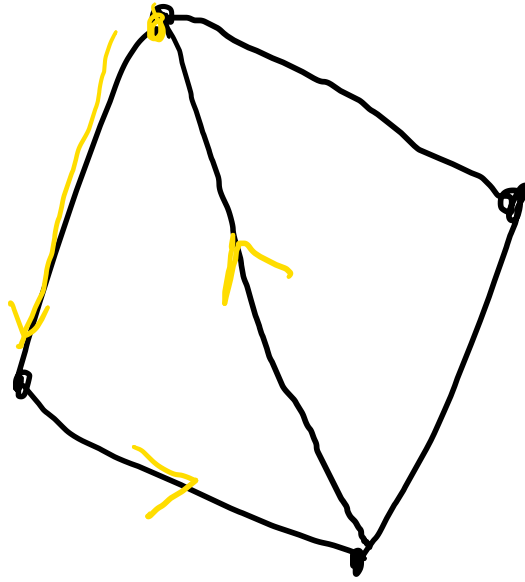
Engineering Mathematics III

Discrete Mathematics

Lecture 16

Connected Graphs, Trees & Pigeonhole Principle

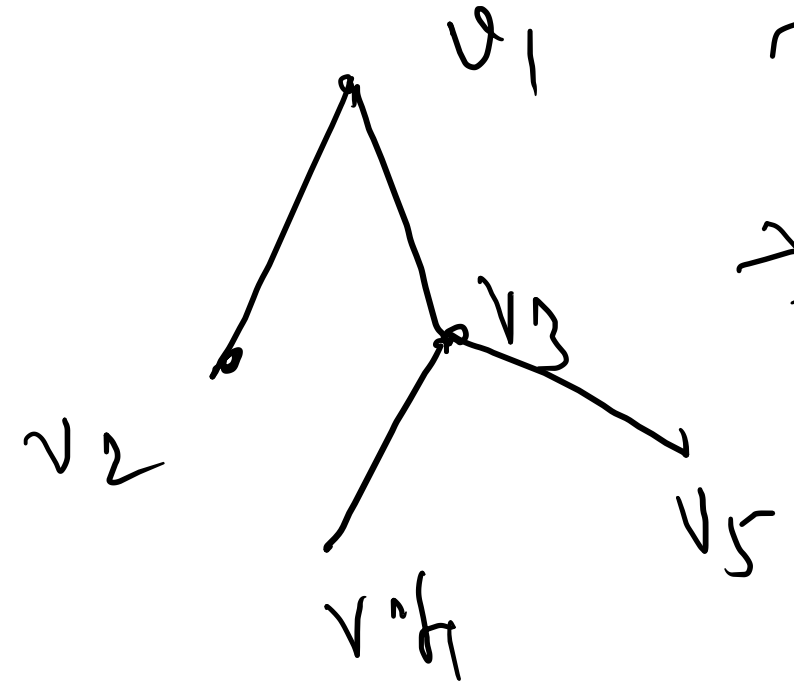
This course is taught to Computer Science Engineering students in SMIT, India during Jun-Dec, 2019.



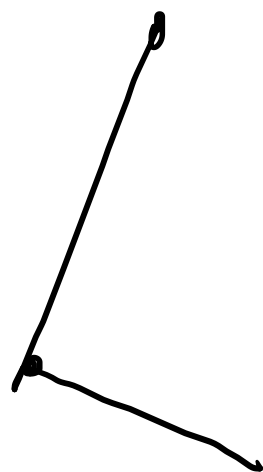
there is a
 \Rightarrow cycle.

A graph without any
 cycles is called as
 acyclic graph.

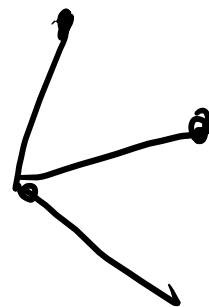
Acyclic graph



Tree
 \rightarrow Tree

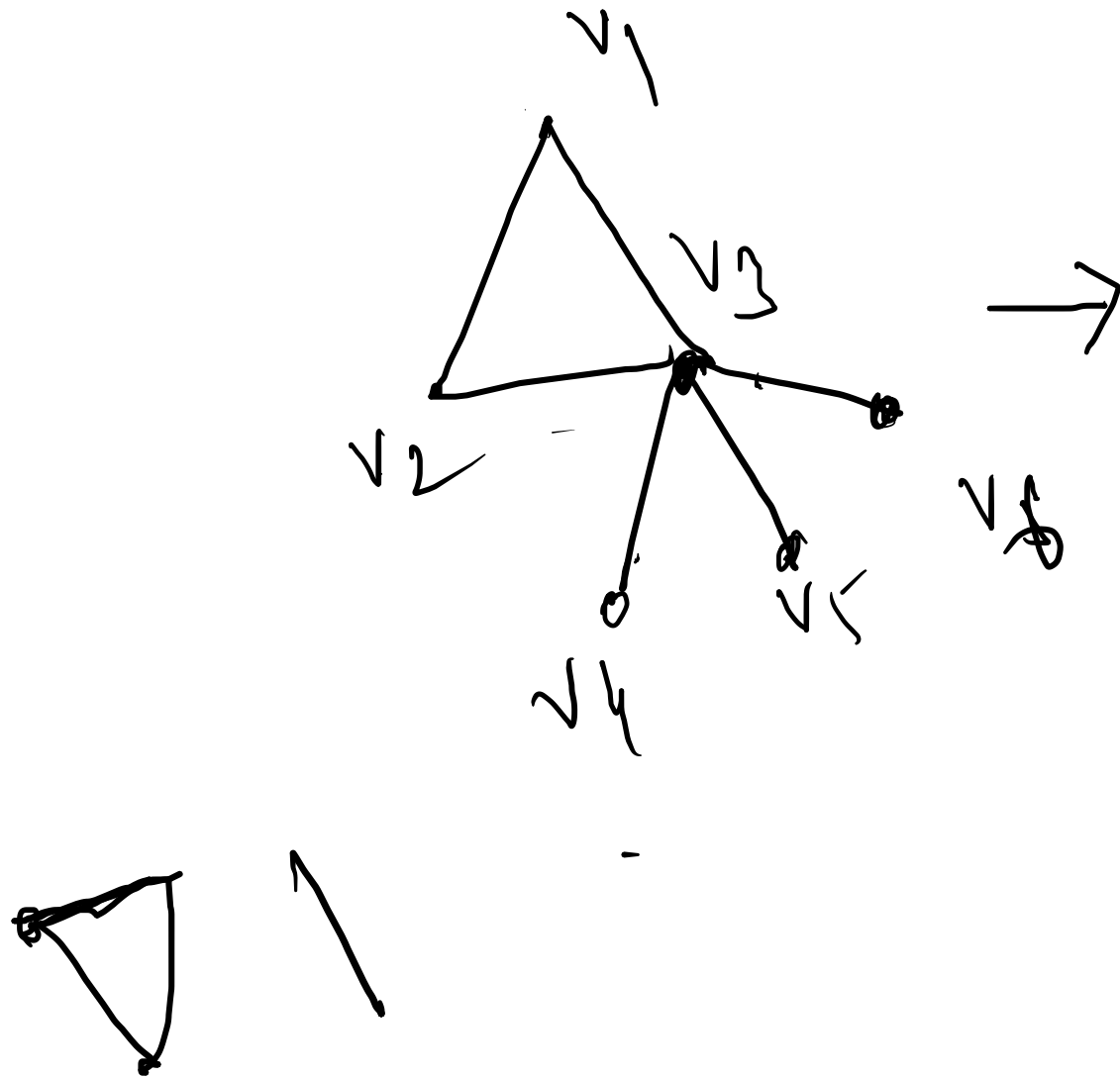


G



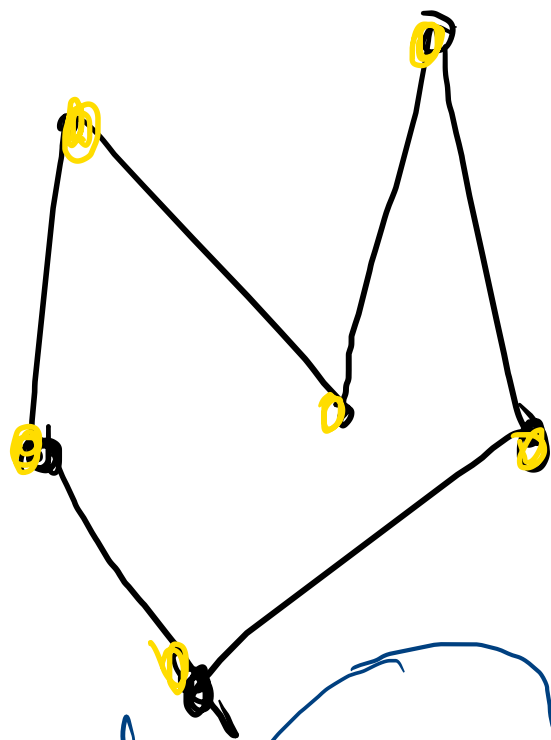
Acyclic (not free)

Connected graph:-



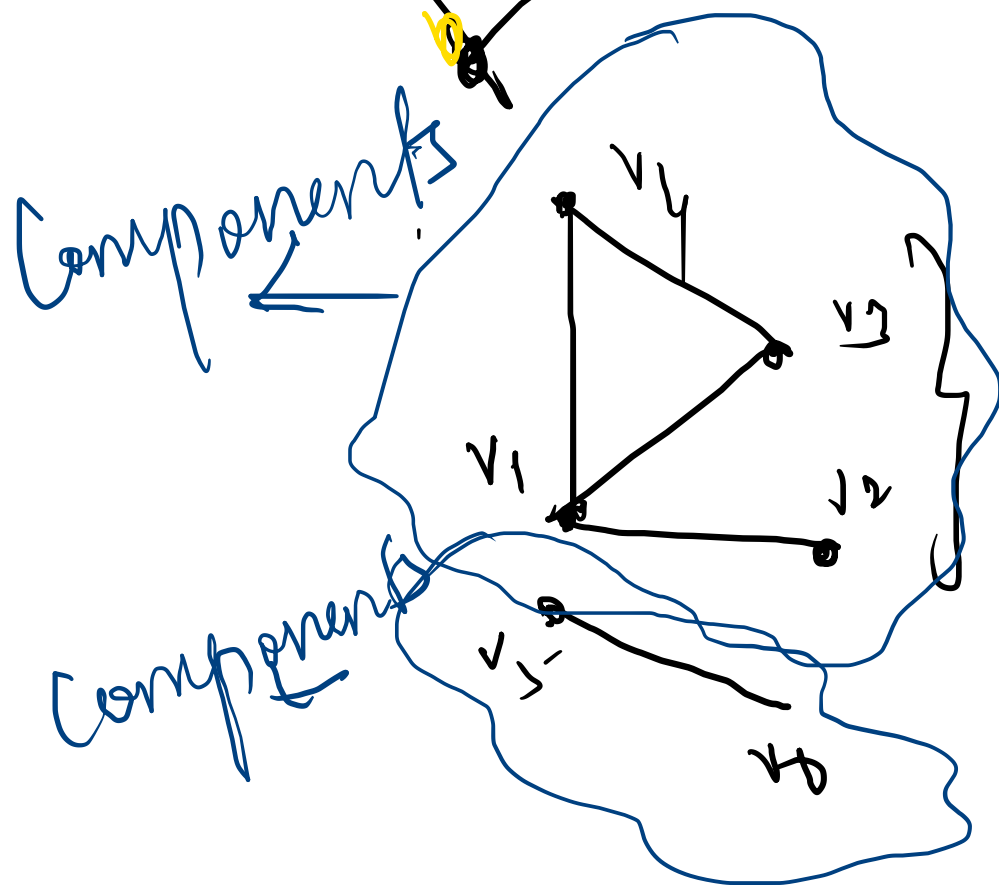
A graph G such that
any two vertices
can be connected

If there is a path b/w
any two vertices of a graph
then that graph is called
connected.

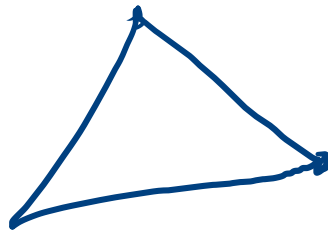
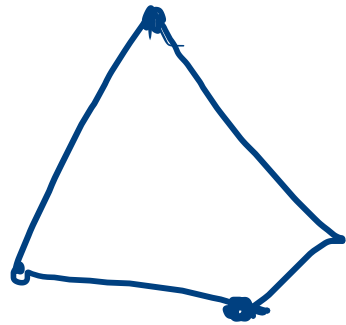


→ Cycle

→ Connected graph

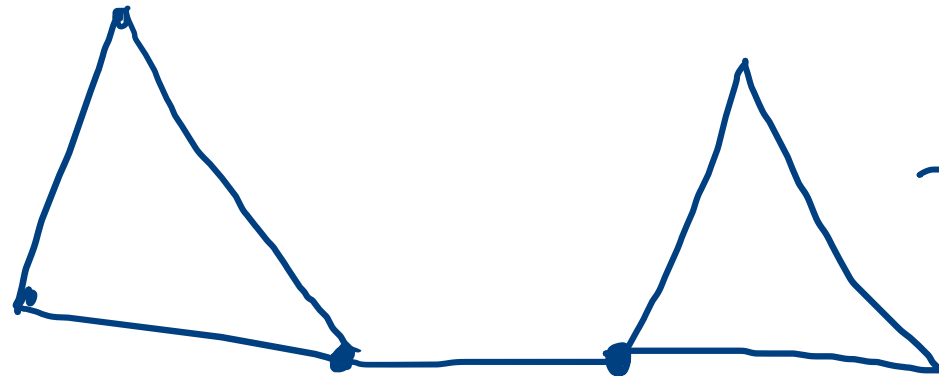


→ not Connected (disconnected)
 why? (there is no path
 b/w v_4 and v_5)



How many components does
this graph has?

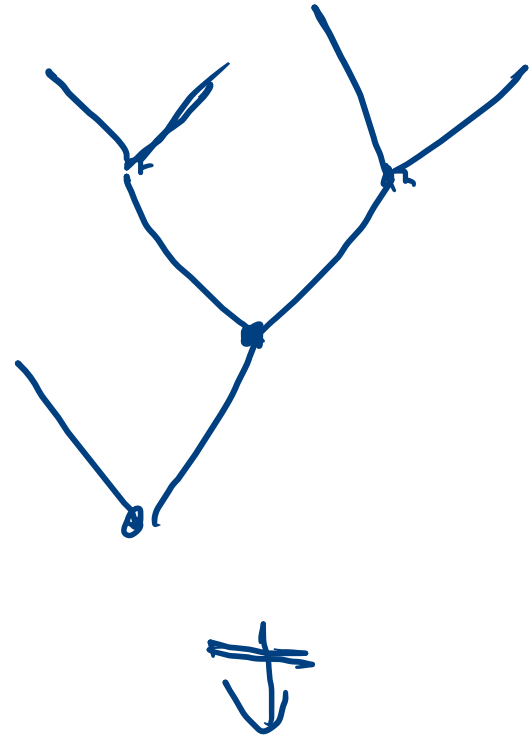
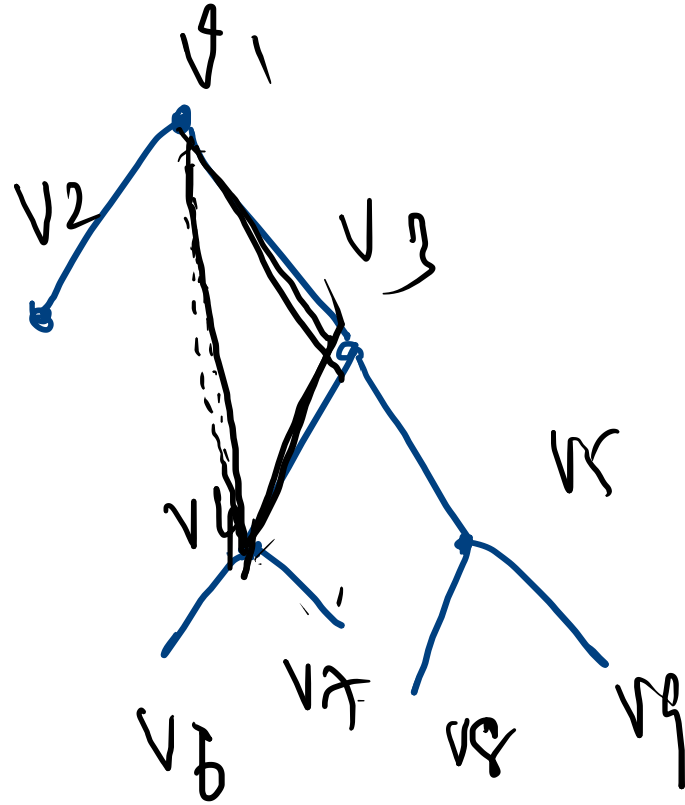
— three.



→ not a
tree.

Tree

A Connected graph
is a tree if
there is always
only one path
b/w any two vertices.



The following are equivalent

(i) G is a tree

(ii) G is a connected and acyclic graph.

Proof:

(i) \Rightarrow (ii) Let G be a tree.

\Rightarrow G is connected and any two vertices
have unique path b/w them.

TP: G does not have a cycle

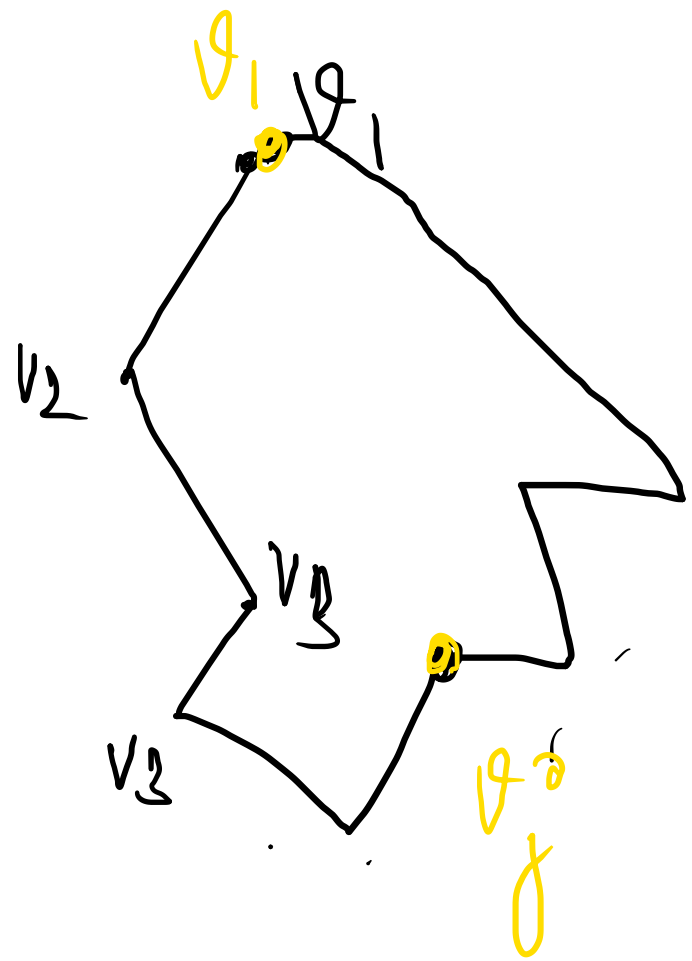
Suppose G has a cycle

Let $v_1, v_2, \dots, v_k = v_1$ forms a cycle

Let v_1 be a vertex in the cycle

and v_j^* be another vertex in G

\Rightarrow there exist two paths b/w v_1 and v_j^*
which is a contradiction to G is a tree
 $\Rightarrow G$ cannot have a cycle $\Rightarrow G$ is acyclic

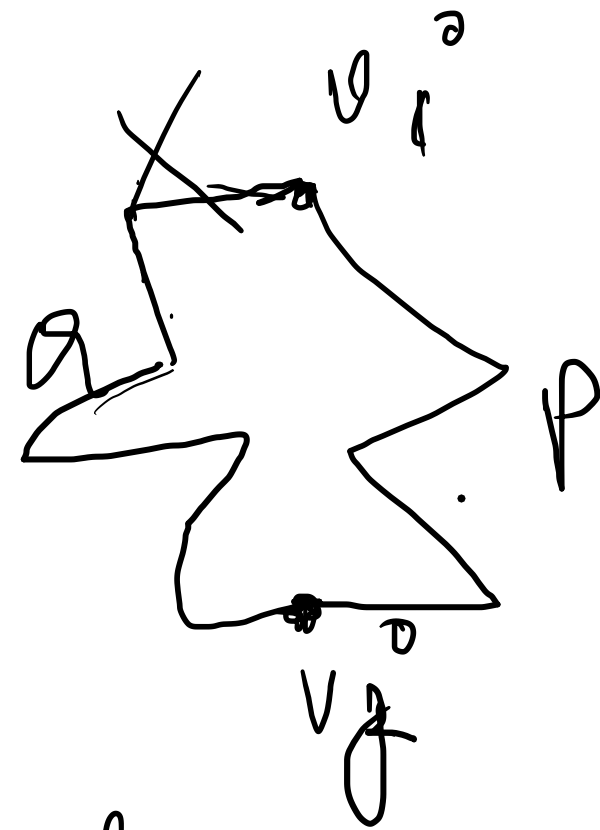


(ii) \Rightarrow (i) Let G be a connected graph
and G be acyclic graph

Let v_i and v_j have a path P

IP: There is only one path b/w v_i & v_j

Suppose there is another path b/w v_i & v_j



Then the two paths makes a cycle

which is a contradiction to G -acyclic

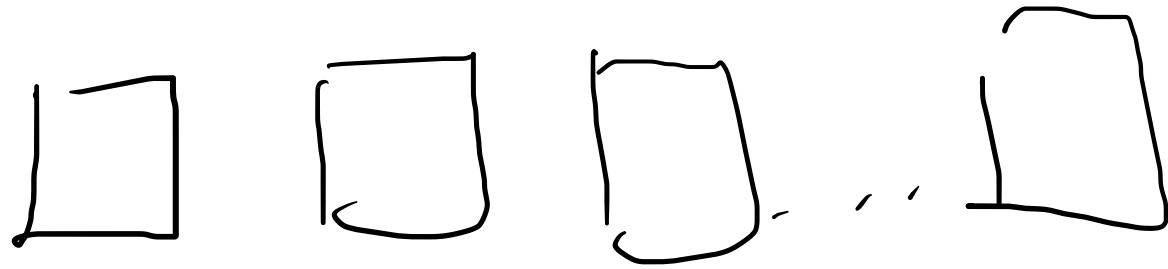
\Rightarrow there is only one path b/w v_i and v_j .

This is true for any i , and j , $1 \leq i, j \leq$ no. of vertices of the graph

$\Rightarrow G$ is a tree.

Hence proved -

Pigeonhole principle



If there are n pigeons
and $n-1$ pigeonholes

then at least one pigeonhole
contains two or more pigeons.



Question 1.9. In a party with atleast two people, there are atleast two people who have same number of friends.

Proof: Let $0, 1, 2, \dots, N$ are the number of friends
and N - no. of people in the party.

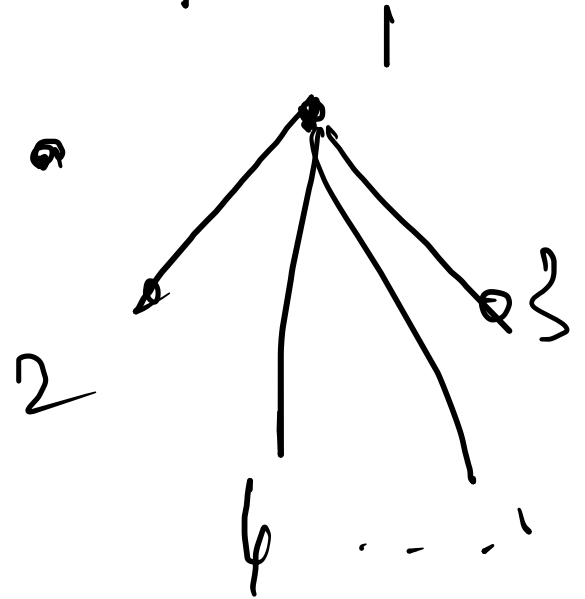
Suppose if a person does not have anybody as friend.

There is always one person who is

Friend of everybody. (?)

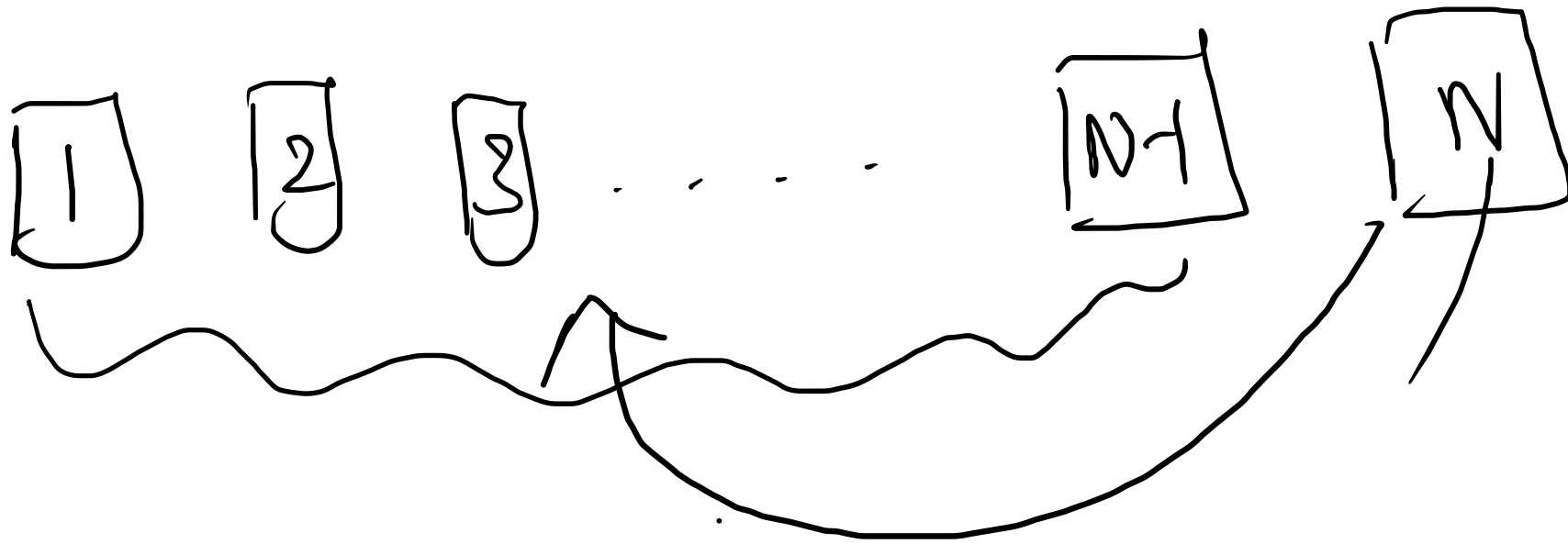
\Rightarrow the seq of number of friends

has to be $1, 2, \dots, N-1$



Pigeons = people

pigeonholes = no of friends they have



\Rightarrow The n^{th} person should occupy one of the pigeonholes $1, 2, \dots, N-1$.

\Rightarrow There are at least two people have same no of friends.