

Engineering Mathematics III

Discrete Mathematics

Lecture 20

Theory of Inference & Predicate Calculus

This course is taught to Computer Science Engineering students in SMIT, India during Jun-Dec, 2019.

Premises & Conclusion

In logic, an argument requires a set of (at least) two declarative sentences (or "propositions") known as the **premises** or premisses along with another declarative sentence (or "proposition") known as the **conclusion**.

Validity of a formula

Let p and q be two statement formulas. We define ``q is a valid conclusion of the premise p" if and only if $p \rightarrow q$ is a tautology.

Rules of inference

Rule P:

A premise can be inserted at any point in the derivation

Rule T:

If the formula q is tautologically implied by any one or more of the previous formulas in a derivation, then q can be inserted in the derivation.

Rule CP:

Rule CP stands for rule of conditional proof. It is also known ad *deduction theorem*. It states that, if we can derive a formula q from p and a set of premises then we can derive $p \rightarrow q$ from the set of the premises alone.

Predicate Calculus

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Dog is an animal. Ent is an animal.

M: San animed -> predicate.

Dog is M = M(Dog)

Cat is M = M(Cat)

2 å an integer 3 å an integer I: an integer -> predicate 2 is an enteger: I(2)

3 is an entegere: I(3)

Quantifiers There exists a Student in the closurs who have Scored 50/00. quadrifier.

For all Students en the class I mark will be added as bornes

there exists belongs to the mose a flower, Rose is red: There Ps a flower which is hed x, xER(X)

for all x, if x is a dog then n is animal. \Rightarrow $\forall x$, $D(n) \rightarrow A(n)$

for all n, if n is a cat then n is an animal. A: animal C: Cat $\forall x \in C(n) \rightarrow A(n)$.

A(n); nis an animal C(n); nic a cat Everey rose is red.

Juantifier Juantifier Mariable medicate, All dogs are animal-als voriouxe predicate, xidog
A: Animal,

X's role,

R: Red, R(n): nis Red, tx, A(n) Hre, R(n)

For all M, if nica hove then xis ned. K(n); x is a now (3) All dogs are animals. R'(n) = n is red Alm): X is animal $\forall x, R(n) \Rightarrow R'(n)$ DIN): n is a dug nous are Red. $\forall x, D(n) \rightarrow A(n)$ R(n): x is a Rose R'(n): x is red $\exists x, R(x) \rightarrow R'(x)$