### PARTIALLY ORDERED SETS

A partially ordered set or poset is a set P and a binary relation  $\leq$  such that for all  $a, b, c \in P$ 

- $\bullet$   $a \leq a$  (reflexivity).
- 2  $a \leq b$  and  $b \leq c$  implies  $a \leq c$  (transitivity).
- 3  $a \le b$  and  $b \le a$  implies a = b. (anti-symmetry).

### **Examples**

- $P = \{1, 2, ..., \}$  and  $a \le b$  has the usual meaning.
- $P = \{1, 2, \dots, \}$  and  $a \leq b$  if a divides b.
- $P = \{A_1, A_2, \dots, A_m\}$  where the  $A_i$  are sets and  $\leq = \subseteq$ .

A pair of elements a, b are **comparable** if  $a \leq b$  or  $b \leq a$ . Otherwise they are **incomparable**.

A poset without incomparable elements (Example 1) is a linear or total order.

We write a < b if  $a \leq b$  and  $a \neq b$ .

A **chain** is a sequence  $a_1 < a_2 < \cdots < a_s$ .

A set *A* is an **anti-chain** if every pair of elements in *A* are incomparable.

Thus a Sperner family is an anti-chain in our third example.

#### Theorem

Let P be a finite poset, then  $\min\{m: \exists \text{ anti-chains } A_1, A_2, \dots, A_\mu \text{ with } P = \bigcup_{i=1}^\mu A_i\} = \max\{|C|: A \text{ is a chain}\}.$ 

The minimum number of anti-chains needed to cover *P* is at least the size of any chain, since a chain can contain at most one element from each anti-chain.

We prove the converse by induction on the maximum length  $\mu$  of a chain. We have to show that P can be partitioned into  $\mu$  anti-chains.

If  $\mu = 1$  then *P* itself is an anti-chain and this provides the basis of the induction.

So now suppose that  $C = x_1 < x_2 < \cdots < x_{\mu}$  is a maximum length chain and let A be the set of maximal elements of P.

(An element is x maximal if  $\exists y$  such that y > x.)

A is an anti-chain.

Now consider  $P' = P \setminus A$ . P' contains no chain of length  $\mu$ . If it contained  $y_1 < y_2 < \cdots < y_{\mu}$  then since  $y_{\mu} \notin A$ , there exists  $a \in A$  such that P contains the chain  $y_1 < y_2 < \cdots < y_{\mu} < a$ , contradiction.

Thus the maximum length of a chain in P' is  $\mu-1$  and so it can be partitioned into anti-chains  $A_1 \cup A_2 \cup \cdots A_{\mu-1}$ . Putting  $A_{\mu} = A$  completes the proof.

Suppose that  $C_1, C_2, \ldots, C_m$  are a collection of chains such that  $P = \bigcup_{i=1}^m C_i$ .

Suppose that A is an anti-chain. Then  $m \ge |A|$  because if m < |A| then by the pigeon-hole principle there will be two elements of A in some chain.

#### Theorem

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(Dilworth) Let P be a finite poset, then \min\{m: \exists chains C_1, C_2, \ldots, C_m \text{ with } P = \bigcup_{i=1}^m C_i\} = \max\{|A|: A \text{ is an anti-chain}\}.
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We have already argued that  $\max\{|A|\} \le \min\{m\}$ .

We will prove there is equality here by induction on |P|.

The result is trivial if |P| = 0.

Now assume that |P|>0 and that  $\mu$  is the maximum size of an anti-chain in P. We show that P can be partitioned into  $\mu$  chains.

Let  $C = x_1 < x_2 < \cdots < x_p$  be a *maximal* chain in P i.e. we cannot add elements to it and keep it a chain.

**Case 1** Every anti-chain in  $P \setminus C$  has  $\leq \mu - 1$  elements. Then by induction  $P \setminus C = \bigcup_{i=1}^{\mu-1} C_i$  and then  $P = C \cup \bigcup_{i=1}^{\mu-1} C_i$  and we are done.

## Case 2

There exists an anti-chain  $A = \{a_1, a_2, \dots, a_{\mu}\}$  in  $P \setminus C$ . Let

- $P^- = \{x \in P : x \leq a_i \text{ for some } i\}.$
- $P^+ = \{x \in P : x \succeq a_i \text{ for some } i\}.$

#### Note that

- $P = P^- \cup P^+$ . Otherwise there is an element x of P which is incomparable with every element of A and so  $\mu$  is not the maximum size of an anti-chain.
- 2  $P^- \cap P^+ = A$ . Otherwise there exists x, i, j such that  $a_i < x < a_j$  and so A is not an anti-chain.
- 3  $x_p \notin P^-$ . Otherwise  $x_p < a_i$  for some i and the chain C is not maximal.



Applying the inductive hypothesis to  $P^-$  ( $|P^-| < |P|$  follows from 3) we see that  $P^-$  can be partitioned into  $\mu$  chains  $C_1^-, C_2^-, \ldots, C_\mu^-$ .

Now the elements of A must be distributed one to a chain and so we can assume that  $a_i \in C_i^-$  for  $i = 1, 2, ..., \mu$ .

 $a_i$  must be the maximum element of chain  $C_i^-$ , else the maximum of  $C_i^-$  is in  $(P^- \cap P^+) \setminus A$ , which contradicts 2.

Applying the same argument to  $P^+$  we get chains  $C_1^+, C_2^+, \ldots, C_{\mu}^+$  with  $a_i$  as the minimum element of  $C_i^+$  for  $i = 1, 2, \ldots, \mu$ .

Then from 2 we see that  $P = C_1 \cup C_2 \cup \cdots \cup C_{\mu}$  where  $C_i = C_i^- \cup C_i^+$  is a chain for  $i = 1, 2, \dots, \mu$ .

# Three applications of Dilworth's Theorem

(i) Another proof of

#### Theorem

Erdős and Szekerés

 $a_1, a_2, \dots, a_{n^2+1}$  contains a monotone subsequence of length n+1.

Let 
$$P = \{(i, a_i) : 1 \le i \le n^2 + 1\}$$
 and let say  $(i, a_i) \le (j, a_j)$  if  $i < j$  and  $a_i \le a_j$ .

A chain in P corresponds to a monotone increasing subsequence. So, suppose that there are no monotone increasing sequences of length n+1. Then any cover of P by chains requires at least n+1 chains and so, by Dilworths theorem, there exists an anti-chain A of size n+1.

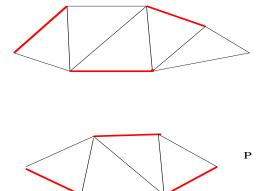
Let  $A = \{(i_t, a_{i_t}): 1 \le t \le n+1\}$  where  $i_1 < i_2 \le \cdots < i_{n+1}$ .

Observe that  $a_{i_t} > a_{i_{t+1}}$  for  $1 \le t \le n$ , for otherwise  $(i_t, a_{i_t}) \le (i_{t+1}, a_{i_{t+1}})$  and A is not an anti-chain.

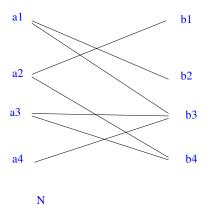
Thus A defines a monotone decreasing sequence of length n + 1.

# Matchings in bipartite graphs

Re-call that a matching is a set of vertex disjoint edges.



Let  $G = (A \cup B, E)$  be a bipartite graph with bipartition A, B. For  $S \subseteq A$  let  $N(S) = \{b \in B : \exists a \in S, (a, b) \in E\}$ .

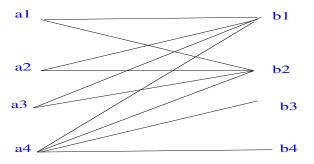


Clearly,  $|M| \leq |A|, |B|$  for any matching M of G.

#### Theorem

(Hall) G contains a matching of size |A| iff

$$|N(S)| \ge |S|$$
  $\forall S \subseteq A$ .



 $N(\{a_1, a_2, a_3\}) = \{b_1, b_2\}$  and so at most 2 of  $a_1, a_2, a_3$  can be saturated by a matching.

If G contains a matching M of size |A| then  $M = \{(a, f(a)) : a \in A\}$ , where  $f : A \to B$  is a 1-1 function.

But then,

$$|N(S)| \geq |f(S)| = S$$

for all  $S \subseteq A$ .

Let  $G = (A \cup B, E)$  be a bipartite graph which satisfies Hall's condition. Define a poset  $P = A \cup B$  and define < by a < b only if  $a \in A, b \in B$  and  $(a, b) \in E$ .

Suppose that the largest anti-chain in P is  $A = \{a_1, a_2, \dots, a_h, b_1, b_2, \dots, b_k\}$  and let s = h + k.

Now

$$N(\{a_1,a_2,\ldots,a_h\})\subseteq B\setminus\{b_1,b_2,\ldots,b_k\}$$

for otherwise A will not be an anti-chain.

From Hall's condition we see that

$$|B| - k \ge h$$
 or equivalently  $|B| \ge s$ .

Now by Dilworth's theorem, *P* is the union of *s* chains:

A matching M of size m, |A|-m members of A and |B|-m members of B.

But then

$$m + (|A| - m) + (|B| - m) = s \le |B|$$

and so  $m \ge |A|$ .

### **Marriage Theorem**

#### Theorem

Suppose  $G = (A \cup B, E)$  is k-regular.  $(k \ge 1)$  i.e.  $d_G(v) = k$  for all  $v \in A \cup B$ . Then G has a perfect matching.

#### **Proof**

$$k|A| = |E| = k|B|$$

and so |A| = |B|.

Suppose  $S \subseteq A$ . Let m be the number of edges incident with S. Then

$$k|S| = m \le k|N(S)|.$$

So Hall's condition holds and there is a matching of size |A| i.e. a perfect matching.



### König's Theorem

We will use Hall's Theorem to prove König's Theorem. Given a bipartite graph  $G = (A \cup B), E$ ) we say that  $S \subseteq V$  is a cover if  $e \cap S \neq \emptyset$  for all  $e \in E$ .

#### Theorem

 $\min\{|S|: S \text{ is a cover}\} = \max\{|M|: M \text{ is a matching}\}.$ 

**Proof** One part is easy. If S is a cover and M is a matching then  $|S| \ge |M|$ . This is because no vertex in S can belong to more than one edge in M.

To begin the main proof, we first prove a lemma that is a small generalisation of Hall's Theorem.

#### Lemma

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Assume that |A| \le |B|. Let d = \max\{(|X| - |N(X)|)^+ : X \subseteq A\} where \xi^+ = \max\{0, \xi\}. Then \mu = \max\{|M| : M \text{ is a matching }\} = |A| - d.
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**Proof** That  $\mu \leq |A| - d$  is easy. For the lower bound, add d dummy vertices D to B and add an edge between each vertex in D and each vertex in A to create the graph  $\Gamma$ . We now find that  $\Gamma$  satisfies the conditions of Hall's Theorem.

If  $M_1$  is a matching of size |A| in  $\Gamma$  then removing the edges of  $M_1$  incident with D gives us a matching of size |A| - d in G.

Continuing the proof of König's Theorem let  $S \subseteq A$  be such that |N(S)| = |S| - d.

Let  $T = A \setminus S$ . Then  $T \cup N(S)$  is a cover, since there are no edges joining S to  $B \setminus N(S)$ .

Finally observe that

$$|T \cup N(S)| = |A| - |S| + |S| - d = |A| - d = \mu.$$

#### **Intervals Problem**

 $I_1, I_2, \ldots, I_{mn+1}$  are closed intervals on the real line i.e.  $I_j = [a_j, b_j]$  where  $a_j \le b_j$  for  $1 \le j \le mn + 1$ .

#### Theorem

Either (i) there are m + 1 intervals that are pair-wise disjoint or (ii) there are n + 1 intervals with a non-empty intersection

Define a partial ordering  $\leq$  on the intervals by  $I_r \leq I_s$  iff  $b_r \leq a_s$ . Suppose that  $I_{i_1}, I_{i_2}, \ldots, I_{i_t}$  is a collection of pair-wise disjoint intervals. Assume that  $a_{i_1} < a_{i_2} \cdots < a_{i_t}$ . Then  $I_{i_1} < I_{i_2} \cdots < I_{i_t}$  form a chain and conversely a chain of length t comes from a set of t pair-wise disjoint intervals.

So if (i) does not hold, then the maximum length of a chain is m.



This means that the minimum number of chains needed to cover the poset is at least  $\lceil \frac{mn+1}{m} \rceil = n+1$ .

Dilworth's theorem implies that there must exist an anti-chain  $\{J_{j_1}, J_{j_2}, \dots, J_{j_{n+1}}\}$ .

Let 
$$a = \max\{a_{j_1}, a_{j_2}, \dots, a_{j_{n+1}}\}$$
 and  $b = \min\{b_{j_1}, b_{j_2}, \dots, b_{j_{n+1}}\}$ .

We must have  $a \le b$  else the two intervals giving a, b are disjoint.

But then every interval of the anti-chain contains [a, b].

#### **Möbius Inversion**

Suppose that |P| = n. We argue next that we can label the elements of  $P = \{p_1, p_2, \dots, p_n\}$  so that

$$p_i \leq p_j \text{ implies } i \leq j.$$
 (1)

We prove this by induction on n. The base case n = 1 is trivial.

Choose a maximal element of P and label it  $p_n$ . Assume that (1) can be achieved for posets with fewer than n elements. Let  $P' = P \setminus \{p_n\}$ .

We can, by induction, re-label  $P' = \{p_1, p_2, \dots, p_{n-1}\}$  so that (1) holds. Because  $p_n$  is maximal, we now have a labelling for all of P.



We now define  $\zeta: P^2 \to \{0, 1\}$  by

$$\zeta(x,y) = \begin{cases} 1 & x \leq y. \\ 0 & Otherwise. \end{cases}$$

Given (1) the  $n \times n$  matrix  $A_{\zeta} = [\zeta(x, y)]$  is an upper triangular matrix with an all 1's diagonal.

 $A_{\zeta}$  is invertible and its inverse is called  $A_{\mu} = [\mu(x, y)]$ . The function  $\mu$  is called the Möbius function of P. The equation  $A_{\mu}A_{\zeta} = I$  implies the following:

$$\sum_{z \in P} \mu(x, z) \zeta(z, y) = \sum_{x \le z \le y} \mu(x, z) = \begin{cases} 1 & x = y. \\ 0 & Otherwise. \end{cases}$$
 (2)

#### Theorem

For P equal to the subsets of some finite set X and ≤=⊆ we have

$$\mu(A,B) = egin{cases} (-1)^{|A|-|B|} & A \subseteq B \ 0 & \textit{Otherwise}. \end{cases}$$

**o** For P = [n] and  $a \le b$  if a divides b we have

$$\mu(a,b) = \begin{cases} (-1)^r & b/a \text{ is the product of } r \text{ distinct primes} \\ 0 & Otherwise. \end{cases}$$

# **Proof**

We just have to verify (2):

(a) We have

$$\sum_{A \subseteq C \subseteq B} x^{|C|-|A|} = (1+x)^{|B|-|A|}.$$

Putting x = -1 we get a RHS of zero, unless A = B, in which case we get  $0^0 = 1$ .

(b) Suppose that  $b/a = p_1^{k_1} p_2^{k_2} \cdots p_r^{k_r}$  where  $p_1, p_2, \dots, p_r$  are primes and  $k_1, k_2, \dots, k_r \ge 1$ .

$$\sum_{a|c|b} \mu(c,b) = \sum_{S \subseteq [r]} (-1)^{|S|} = \begin{cases} 1 & r = 0. \\ 0 & r \ge 1. \end{cases}$$



# **Möbius Inversion**

#### Theorem

Suppose that f, g, h are functions from P to R such that

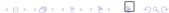
$$g(x) = \sum_{a \leq x} f(a)$$
 and  $h(x) = \sum_{b \succeq x} f(b)$ . (3)

Then,

$$f(x) = \sum_{a \leq x} \mu(a, x) g(a)$$
 and  $f(x) = \sum_{b \succeq x} \mu(x, b) h(b)$ . (4)

**Proof** Treating f, g, h as column vectors  $\mathbf{f}, \mathbf{g}, \mathbf{h}$  we see that (3) is equivalent to  $\mathbf{g} = A_{\mathcal{L}}^T \mathbf{f}$  and  $\mathbf{h} = A_{\mathcal{L}} \mathbf{f}$ . Thus

$$\mathbf{f} = A_{\zeta}^{-T} \mathbf{g} = A_{\mu}^{T} \mathbf{g}$$
 and  $\mathbf{f} = A_{\zeta}^{-1} \mathbf{h} = A_{\mu} \mathbf{h}$ .



## Inclusion-Exclusion

Let  $A_i$ ,  $i \in I$  be a family of subsets of a finite set X.

For  $J \subseteq I$  let f(J) equal the number of elements in  $\bigcap_{i \in J} A_i$  that are also in  $\bigcap_{i \notin I} (X \setminus A_i)$ .

Let h(J) be the number of elements in  $\bigcap_{i \in J} A_i$ . Then

$$h(J) = \sum_{K \supseteq J} f(K) = \sum_{K \succeq J} f(K).$$

Möbius inversion gives us

$$f(J) = \sum_{K \succeq J} \mu(K, J) h(K) = \sum_{K \supset J} (-1)^{|K| - |J|} h(K).$$

Putting  $J = \emptyset$  we get

$$\left|\bigcap_{i\in I}(X\setminus A_i)\right|=\sum_{K\subseteq I}(-1)^{|K|-|J|}\left|\bigcap_{j\in K}A_j\right|.$$

# **Divisibility Poset**

Supose now that  $f: \mathbb{N} \to \mathbb{R}$  and that g is given by

$$g(n)=\sum_{d\mid n}f(d).$$

Then Möbius inversion gives

$$f(n) = \sum_{d \mid n} \mu(d, n) g(d) = \sum_{\substack{d \mid n \ n/d \text{ square free}}} (-1)^{p(n/d)} g(d)$$

where p(m) is the number of distinct prime factors of m.

### **Totient function**

For a natural number n, let  $\phi(n)$  denote the number of integers  $m \le n$  such that m, n have n common factors (other than one) – co-prime.

#### Lemma

$$n = \sum_{d|n} \phi(d) = \sum_{d|n} \phi(n/d). \tag{5}$$

**Proof** If (m, n) = d then  $m = m_1 d$ ,  $n = n_1 d$  where  $(m_1, n_1) = 1$ . So the number of choices for m is the number of choices for  $m_1$  i.e.  $\phi(n_1) = \phi(n/d)$ .

Möbius inversion with g(n) = n and  $f(n) = \phi(n)$  applied to (5) gives

$$\phi(n) = \sum_{d|n} (-1)^{p(n/d)} d = \sum_{d|n} (-1)^{p(d)} \frac{n}{d}.$$
 (6)

$$\phi(n) = n \sum_{d|n} \frac{(-1)^{p(d)}}{d}$$

$$= n \prod_{i=1}^{k} \left(1 - \frac{1}{p_i}\right),$$
(7)

assuming that  $n = p_1^{k_1} p_2^{k_2} \dots, p_r^{k_r}$  where  $p_1, p_2, \dots, p_r$  are primes and  $k_1, k_2, \dots, k_r \ge 1$ .

# 2-colored necklace

A *necklace* is a sequence  $x_1x_2 \cdots x_n$  of n 0' and 1's arranged in circle.

Two necklaces x, y are said to *equivalent* if there exists d|n such that  $y_i = x_{i+d}$ , i = 1, 2, ..., n where we interpret  $i + d \mod n$ . In this case we say that x is *periodic* with period d.

Let  $N_n$  denote the number of distinct i.e. non-equivalent necklaces and let M(d) denote the number of aperiodic necklaces of length d.

Thus

$$N_n = \sum_{d|n} M(d)$$
 and  $\sum_{d|n} dM(d) = 2^n$ .



$$N_n = \sum_{d|n} M(d)$$
 and  $\sum_{d|n} dM(d) = 2^n$ .

For the second equation think about rotating a periodic necklace one step at a time for d steps. If we do this for all periodic necklaces then we get all  $2^n$  sequences.

Applying Möbius inversion to the second equation with  $f(d) = dM(d), g(n) = 2^n$ , we get

$$M(n) = \frac{1}{n} \sum_{d|n} \mu(d, n) 2^d.$$

So,

$$N_n = \sum_{d|n} M(d) = \sum_{d|n} \sum_{\ell|d} \frac{1}{d} \mu(\ell, d) 2^d = \sum_{d|n} \frac{1}{d} \sum_{\ell|d} \mu(\ell, d) 2^\ell.$$

Now substitute  $d = k\ell$  and tidy up to get

$$N_n = \sum_{\ell \mid n} \frac{2^{\ell}}{\ell} \sum_{k \mid \frac{n}{\ell}} \frac{\mu(1,k)}{k} = \frac{1}{n} \sum_{\ell \mid n} \phi(n/\ell) 2^{\ell}.$$

For the second equation, we use the expression (7).