

COMP 302 Assignment 2 Question 1

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Problem 1. *Prove $\text{reflect}(\text{reflect } t) = t$.*

Proof. Proof by induction:

Base case:

```
reflect Empty = Empty
reflect (reflect Empty) = Empty
```

Induction hypothesis:

Assume $\text{reflect}(\text{reflect } \text{someNode}) = \text{someNode}$ for left and right
in $\text{someNode} = \text{Node}(x, \text{left}, \text{right})$

```
let someTree = Node(x, left, right)
reflect someTree = Node(x, reflect right, reflect left)
```

```
reflect (reflect someTree)
= Node(x, reflect (reflect left), reflect (reflect right))
= Node(x, left, right)    //from induction hypothesis
= someTree
```

□

Problem 2. *Prove that for all m : 'a tree, $\text{size } m = \text{size}' m 0$.*

Proof. Auxiliary proof by induction, then show that the statement is a special case:

Begin by proving $\text{size}' \text{someTree } a = \text{size } \text{someTree} + a$

Base case:

```
size' Empty a
= a = 0 + a
= size Empty + a
```

Induction hypothesis:

Assume $\text{size}' \text{someNode } a = \text{size } \text{someNode} + a$ for left and right
in $\text{someNode} = \text{Node}(x, \text{left}, \text{right})$

```
let someTree = Node (x, left, right)
```

```
size' someTree a
= size' left (size' right (a + 1)) //by function definition
= size left + size' right (1 + a) //from induction hypothesis
= size left + size right + 1 + a  //apply IH again
= size Node (x', left, right) + a //by function definition
```

The value of x' is irrelevant so we can set it to x to obtain:

```
= size someTree + a
```

Now that I have proven $\text{size}' \text{ someTree } a = \text{size someTree} + a$,
I can show that there is the special case when $a = 0$

```
size m = size' m 0
```

□