

# Combination Frequency Differencing

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## Abstract

This paper introduces a new method related to combinatorial testing and measurement, *combination frequency differencing* (CFD), and illustrates the use of CFD in machine learning applications. Combinatorial coverage measures have been defined and applied to a wide range of problems, including fault location and for evaluating the adequacy of test inputs and input space models. More recently, methods applying coverage measures have been used in applications of artificial intelligence and machine learning, for explainability and for analyzing aspects of transfer learning. These methods have been developed using measures that depend on the inclusion or absence of  $t$ -tuples of values in inputs, training data, and test cases. In this paper, we extend these combinatorial coverage measures to include the frequency of occurrence of combinations. Combination frequency differencing is particularly suited to AI/ML applications, where training data sets used in learning systems are dependent on the prevalence of various attributes of elements of class and non-class sets. We illustrate the use of this method by applying it to analyzing physically unclonable functions (PUFs) for bit combinations that disproportionately influences PUF response values, and in turn provides indication of the PUF potentially being more vulnerable to model-building attacks. Additionally, it is shown that combination frequency differences provide a simple but effective algorithm for classification problems.

40

## Keywords

combinatorial coverage; combination frequency difference; combinatorial testing; physical unclonable function (PUF); unclonable.

43

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61

## 1 Introduction

62 Methods and tools for measuring combinatorial coverage were initially developed to analyze the degree to which  
63 test sets included  $t$ -way combinations of values (for some specified level of  $t$ ) [1][2][4] and have since been  
64 studied extensively in the realm of system and software testing [7][8][9][10][11]. Combinatorial coverage  
65 measures have been defined and applied to a wide range of problems, specifically for fault location and for  
66 evaluating the adequacy of test inputs and input space models. More recently, coverage measures have been used  
67 for explainability in artificial intelligence and machine learning [24][28] and for analyzing aspects of transfer  
68 learning [27]. These methods have been developed using measures that depend on the inclusion or absence of  $t$ -  
69 tuples of values in inputs and test cases. For software testing, primarily for deterministic systems where the  
70 presence of a particular combination always triggers a specified error, it is relevant whether a  $t$ -tuple of values  
71 is present in test inputs, but the number of occurrences of a particular  $t$ -tuple of values is generally not relevant  
72 to testing. Multiple occurrences are only redundant and do not add value. These measures can also be applied in  
73 artificial intelligence and machine learning (AI/ML) systems.

74 For many aspects of assurance of autonomous systems and machine learning, this type combinatorial coverage  
75 measure is valuable and possibly essential, since the correct and safe behavior of many AI systems is dependent  
76 on the training inputs. Conventional structural coverage measures are not applicable to such black box behavior.  
77 Consequently, it is essential to evaluate the degree to which possible combinations of input attribute values have  
78 been included in training and test sets for AI and autonomy. (Attributes in a machine learning setting correspond  
79 to parameters in a test effort; they are the inputs to the system.) If the system has not been shown to function  
80 correctly for an input combination that may be encountered in use, then assurance is inadequate. However, for  
81 some questions in machine learning, consider the frequency (or rate) of occurrence of  $t$ -tuples of values in input  
82 and how two different sets may compare or differ in combinatorial coverage.

83 This paper applies combinatorial coverage measures from [13], which include the frequency of occurrence of  
84 combinations, in an approach referred to as *combination frequency differencing* (CFD). This method is  
85 particularly suited to AI/ML applications, where training data sets used in learning systems are dependent on the  
86 prevalence of various attributes of elements of class and non-class sets. This paper illustrates the use of this  
87 method by applying it to analyzing physical unclonable functions (PUFs) for potential weaknesses in design and  
88 showing how it can be extended to develop a simple but effective classification algorithm.

89

## 2 Combinatorial Coverage and Combination Frequency Differences

90 This section reviews the basic measures of combinatorial coverage and applications of these measures  
91 in Section 2.1. This idea is extended to measures that include the frequency of occurrence of  
92 combinations in Section 2.2. These measures can then be applied to the analysis of PUFs.

93

### 2.1 Basic Combinatorial Coverage and Coverage Difference Measures

94 Combinatorial methods offer an approach to coverage measurement that provides a measure directly related to  
95 fault detection. A series of studies have shown that most software bugs and failures are caused by one or two  
96 parameters and progressively fewer by three or more [19][20][21][22][5][6]. This finding means that testing  
97 parameter combinations can provide more efficient fault detection than conventional methods. This section,  
98 derived from [13], reviews the concept of measuring the combinatorial coverage of an input space [1][2][4] for  
99 use in testing or in other applications where it is important to ensure the inclusion of combinations of input  
100 parameter values.

	<b>a</b>	<b>b</b>	<b>c</b>	<b>d</b>
1	0	0	0	0
2	0	1	1	0
3	1	0	0	1
4	0	1	1	1

101

**Figure 1. Example test array for a system with four binary components**

102 Combinatorial coverage measurement concepts can be illustrated using the example in Figure 1, which shows a  
 103 test array that contains 19 of a possible set of 24 2-tuples of values. To facilitate discussion, it is helpful to  
 104 establish terminology for two related but distinct concepts:

- 105     • *t-way combination*: a set of *t* parameters or variables. For example, using the parameters in Figure 1,  
 106       (b,d) is a 2-way combination, (a,c) is a different 2-way combination, and (a,c,d) is a 3-way combination.  
 107     • *t-tuple of values*: a combination for which the parameters have specific values. (Note: in the original  
 108       definition from [1], this is referred to as a variable-value combination.) For example, (b=0, d=0) is one  
 109       *t*-tuple of values, and (b=1, d=0) is a different *t*-tuple of values for the same 2-way parameter  
 110       combination.

111 A simple combinatorial coverage of *t*-way combinations,  $S_t$ , is the fraction of possible *t*-tuples of values included  
 112 in a test array from a domain  $D_t$  that may include constraints. With no constraints, where  $v$  is the number of  
 113 values and  $k$  is the number of parameters, the size of the domain is  $v^t \binom{k}{t}$  but may be smaller with constraints.

114 For a set of *t*-tuples of values  $A_t$  in a test array,

$$115 \quad S_t = \frac{|A_t|}{|D_t|}$$

116 **Example:** Figure 1 contains 19 different 2-way combinations out of a possible domain of  $2^2 \binom{4}{2} = 24$  *t*-tuples  
 117 of values, so  $S_t = 19/24 = 0.79$ .

118 Combinatorial coverage differences have been applied to several problem domains. Initially, this approach was  
 119 used in fault identification, specifically to determine the particular combination(s) of parameter values that would  
 120 trigger a fault. Another example problem where there is a need to distinguish one class of elements from another  
 121 is anomaly-based intrusion detection, which seeks to determine if a particular exchange of packets represents an  
 122 attempted network intrusion. Thus, it is useful to generalize the approach to find combinations that are present  
 123 in one class or set and absent or rare in another, as well as to distinguish one set from another.

124 For fault location, if  $A_t$  = the set of *t*-tuples of values from passing tests and  $B_t$  = the set of *t*-tuples of values  
 125 from failing tests, then the set difference  $B_t \setminus A_t$  is of interest. These are the combinations in failing tests but not  
 126 in passing tests, and thus, those that triggered a failure are contained in this set difference [26].

127 **Example:** If test #2 from Figure 1 is a failing test, then  $B_t \setminus A_t = \{bc = 10, cd = 10\}$  is to be investigated to  
 128 identify failing combinations because the four other 2-way *t*-tuples of values in test #2 are also contained in the  
 129 passing tests #1, #3, #4, which are set  $A_t$ .

130 For transfer learning, if  $A_t$  = the set of *t*-way *t*-tuples of values from a source set of class instances and  $B_t$  = the

131 set of  $t$ -tuples of values from a target set of instances, then the size of the set difference  $B_t \setminus A_t$  as a fraction of  
132 the target set size is of interest as a metric of how similar the source is to the target set [27]. This set difference  
133 of  $t$ -tuples of values is: 
$$\frac{|B_t \setminus A_t|}{|B_t|}$$

134 **2.2 Distinguishing Combinations**

135 For many machine learning applications, the goal is to develop a model that distinguishes members of one class  
136 from another using attributes that identify them, such as distinguishing dogs from cats using attributes like size,  
137 ear shape, or hair texture. This publication will refer to sets being distinguished as either *Class* or *Non-class* sets.  
138 The terms Class and Non-class are used as generic terms for sets of objects that can be distinguished based on  
139 some attributes or properties. In a machine learning context, these sets may refer to concepts that are to be  
140 learned, such as distinguishing one animal species from others. In earlier applications, set differences of  $t$ -tuples  
141 of values have been used to identify the causes of failures [4][5]. In both cases, the process is the same – set  
142 differencing is used to identify combinations that occur in the class set that do not occur, or are rare, in the non-  
143 class set. If this difference is computed on  $t$ -tuples of values in failed tests versus passed tests, then the difference  
144 contains  $t$ -tuples of values that have triggered the failure (in a deterministic system). In machine learning, the  
145 difference represents properties or attributes that occur in the class (e.g., a particular animal species) that do not  
146 occur, or are rare, in the non-class examples (other species). Note that this is simply a generalized version of the  
147 original fault location problem, where the class whose distinguishing features are to be identified is the set of  
148 failing tests, and the features to be found are the combinations that lead to a test resulting in a failure.

149 The combinatorial coverage measures described in the previous section – as applied in fault location,  
150 explainability, and transfer learning – are based on the presence or absence of  $t$ -tuples of values in input files for  
151 testing or machine learning training. That is, a combination is counted as covered if it occurs once or multiple  
152 times in the input file, and this measure is appropriate in the applications discussed. For these applications, it is  
153 important to determine if a  $t$ -tuple of values has been included, but the number of times it occurs is less important.  
154 For testing, multiple occurrences of a combination mean some duplication of effort but do not affect the  
155 requirement for ensuring that all  $t$ -way combinations have been covered. In transfer learning evaluation, the  
156 same type of requirement holds – assurance that states and environments, as represented by  $t$ -tuples of values of  
157 the input model, are handled correctly. If it can be shown that the ML model produces the right prediction or  
158 classification for a  $t$ -tuple of values, multiple occurrences of the combination are not needed. (This does not  
159 consider the effect of input sequences; other measures are appropriate for sequence coverage.)

160 In other types of evaluations related to machine learning, it will be important to consider the number or frequency  
161 of occurrence of  $t$ -way  $t$ -tuples of values to determine the degree to which an attribute is associated with a  
162 particular class. If a particular combination of attribute values is seen in a high proportion of class members but  
163 not in non-class members, then it may be a reasonable indicator for distinguishing instances or at least for  
164 narrowing the range of possibilities for class identification. For example, many dog breeds may have a long tail,  
165 and many may have a curled tail, but a much smaller number of breeds have both attributes. Thus, it is important  
166 to have a measure that considers the quantity of instances with  $t$ -tuples of values in class and non-class instances.

167 This paper will abbreviate  $C_t$  and  $N_t$  as  $C$  and  $N$ , where interaction level  $t$  is clear or is not needed for discussion.  
168 The following discussion defines a  $t$ -way combination  $c_t$  as a distinguishing combination for the class  $C$  if it is  
169 present in a class instance of class set  $C$  and absent in non-class instances  $N$ , or if it is more common in  $C$  than  
170  $N$  as determined by a threshold value. Two ways to identify distinguishing combinations are suggested below,  
171 and others are clearly possible. The key point is to use combinations of attribute values that are *strongly*  
172 *associated* with one class but not with others based on the frequency or rate of occurrence in one class as  
173 compared with others.

At least two possible ways to define the strength of association of a  $t$ -tuple of values with a class can be considered. These are defined and presented below as CFD1 and CFD2. (In a previous publication, only CFD1 was given as the definition of this strength of association [13].) The threshold  $T$  in definition CFD1 determines if a  $t$ -tuple of values  $c_t$  is common in set  $C_t$  and rare in set  $N_t$  and, thus, distinguishes one set from the other. Specifically, the definition below identifies  $t$ -tuples of values for which one can say “ $x$  is  $T$  times more common in  $C$  than it is in  $N$ ” – an intuitive way to identify  $t$ -tuples of values that are associated closely with the class  $C$ . Note that the phrase “ $T$  times more common” suggests that  $T$  will normally be 1 or greater. For definition CFD2,  $U$  designates the threshold value.  $T$  may be any positive number, and  $U$  ranges from 0.0 to 1.0. Notice that these definitions produce the same result for inclusion or exclusion in the set of distinguishing combinations when  $T = \frac{1}{1-U}$ , or  $U = \frac{T-1}{T}$ . For example, if  $T = 4$  or  $U = 0.75$ , then for pairs  $[f(x_t, C_t); f(x_t, N_t)]$ , [.81; .2], and [.79; .2], the first will be found to be distinguishing, and the second will not.

CFD1 Definition: A combination  $x_t$  is *distinguishing* for a class  $C \Leftrightarrow f(x_t, C_t) > T \times f(x_t, N_t)$ , where  $f(x_t, Y_t)$  = frequency of  $t$ -tuple of values  $x$  in set of  $t$ -tuples of values  $Y$ . The frequency  $f$  is the number of times a  $t$ -tuple of values appears in rows of the class over the number of rows for the class.

CFD2 Definition: A combination  $x_t$  is *distinguishing* for a class  $C \Leftrightarrow \frac{f(x_t, C_t) - f(x_t, N_t)}{f(x_t, C_t)} > U$ , where  $f(x_t, Y_t)$  = frequency of  $t$ -tuple of values  $x$  in set of  $t$ -tuples of values  $Y$ . Note that, in this case, the threshold  $U$  ranges from 0.0 to 1.0. The frequency  $f$  is the number of times a  $t$ -tuple of values appears in rows of the class over the number of rows for the class.

The choice of CFD1 or CFD2 as a definition may depend on which is more intuitive for the application. Specifying  $T = 1$  or  $U = 0$  means that a combination is selected as distinguishing whenever it occurs at a higher frequency in  $C$  than  $N$ , no matter how small the difference in frequency.

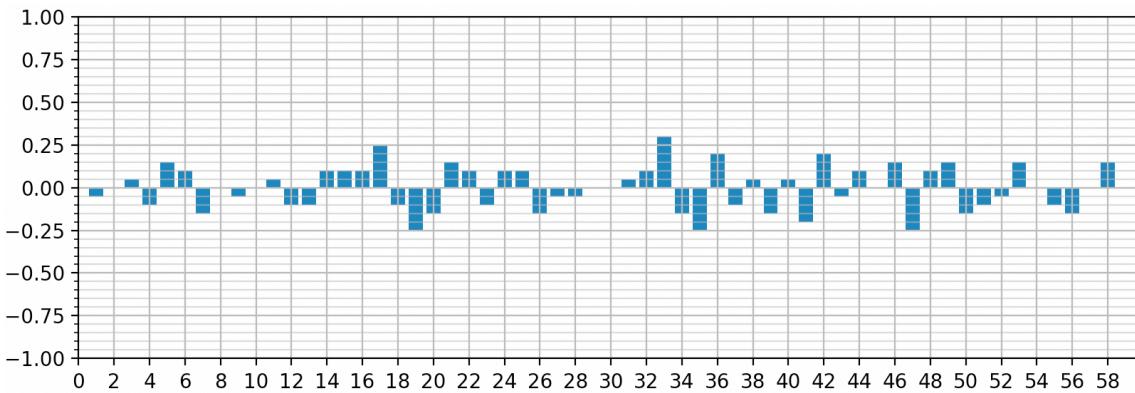
### 2.3 Combination Frequency Difference Measures

The frequency (or rate) of occurrence refers to the number of times a  $t$ -tuple of values is present per number of rows in the file or array. Therefore, the combination frequency difference, for a  $t$ -tuple of values  $x$  in two arrays of instances of two different classes can be defined as the difference between the fraction of occurrences in one array and the second. That is, using the symbols defined below,  $CFD = F_{Cx} - F_{Nx}$ , where

$R$  = number of rows of challenge-response file  
 $R_C$  = rows of class instances; for PUFs,  $R_C = R_1$  (i.e., where challenges produce a 1 response)  
 $R_N$  = rows of non-class instances; for PUFs,  $R_N = R_0$   
 $k$  = number of columns or attributes, excluding class or response variable; for PUFs,  $k = 64$   
 $v$  = number of values for attributes; for PUFs,  $v = 2$  as the attributes correspond to bits  
 $M_{Cx}$  = number of occurrences of a particular  $t$ -tuple of values  $x$  in  $C$   
 $M_{Nx}$  = number of occurrences of a particular  $t$ -tuple of values  $x$  in  $N$   
 $F_{Cx} = M_{Cx}/R_C$  = fraction of occurrences of a  $t$ -tuple of values in  $C$   
 $F_{Nx} = M_{Nx}/R_N$  = fraction of occurrences of a  $t$ -tuple of values in  $N$

The frequency difference values can be graphed, where the height on the Y axis shows the difference  $F_{Cx} - F_{Nx}$  for every  $t$ -tuple of values  $x$ . The X axis is indexed by  $v^t \binom{k}{t}$ , points for  $t$ -way combinations. Thus, for each  $t$ -way combination, there are  $v^t$  possible values or settings of the  $t$  attributes or variables in the combination. For example, 2-way  $t$ -tuples of values are displayed in the order given by:  $i, j$  for  $i$  in  $0 \leq i < k-1$  for  $j$  in  $i+1 \leq j < k$ . Thus, there are  $k-1$  iterations of the inner loop on  $j$  for each attribute  $i$ , and for each 2-way combination, the

graph displays the fraction of occurrences of each set of  $v^2$   $t$ -tuples of values on the X axis at  $v^2((k-1)i+j-1)$  through  $v^2((k-1)i+j-1) + v^2 - 1$ . For each of these 2-way combinations  $x$ ,  $F_{Cx} - F_{Nx}$  for four  $t$ -tuples of values are displayed for the four possible value settings 00, 01, 10, 11. Thus, in Figure 2, the difference in coverage for C and N for  $i=1, j=4$  will be found on the horizontal axis at  $x = 32..35$ .



**Figure 2. Example combinatorial frequency difference for two classes of 6 binary variables**

For example, with  $n = 6$  numbered 0..5, 2-way combinations will be indexed on the Y axis as (0,1,00), (0,1,01),..., (4,5,11), for a total of  $2^2 \binom{6}{2} = 60$  X-axis points, numbered 0..59. For each of these, the Y-axis shows the difference in frequency of occurrence between C and N, normalized for the size of sets C and N. For example, if the value 01 for attribute combination  $i=1, j=4$  occurs 40 times in a C file of 100 rows and 60 times in an N file of 120 rows, then the Y axis value for  $i,j = 1,4$  for value 01 is  $(40/100) - (60/120) = -0.1$ . The analysis of PUFs described in this paper can use these quantities to identify bits related to internal structure.

### 3 Application to Physical Unclonable Functions

A physical unclonable function, or PUF, may be regarded as a physical implementation of a black box function that produces a response  $r$  for a given challenge string of bits  $c$ , that is,  $r = f(c)$ . The unit response is binary and can be represented as 0 or 1. A series of PUFs can be put together to produce a larger response sequence. As the name suggests, PUFs are designed using physical hardware devices. These functions utilize unique properties of the physical elements within the hardware, such as the small variation in propagation delays between identical circuit gates or small threshold mismatches in a transistor feedback loop due to process variation. These physical characteristics are difficult to reproduce in the hardware, which is what makes them physically unclonable. Using such physical characteristics, PUFs can be utilized to combat insecure storage, hardware counterfeiting, and other security problems.

An ideal PUF should be stable over time, unique in its existence, easy to evaluate, and difficult or impossible to predict. Thus, it should not be possible to generate a function that has the same behavior or produces the same output as the PUF for challenge inputs. In this sense, the PUF function is “unclonable.” It should also be infeasible to determine components of the PUF that influence the output of the PUF, such that a 0 or 1 value in some positions of the input string makes a 0 or 1 output more likely for the output  $r$ .

The primary use of PUFs is related to authentication. In a simple use case, the physical system is subjected to

243 one or more challenges during manufacturing, and the responses to these challenges are recorded. Later, if one  
 244 of those recorded challenges is repeated and if the expected response is received, then the device is authenticated.

245 Depending on the strength of their implementation and consequent scalability, PUFs are categorized into two  
 246 levels – weak and strong. Weak PUFs have a limited number of challenge-response pairs (CRPs) that can be  
 247 generated from a single device, while strong PUFs can generate a much larger set of CRPs. One of the key  
 248 requirements for a strong PUF design is that it should not be possible to infer information about the internal  
 249 structure by observing inputs and outputs [16]. Many authors have shown that machine learning models can be  
 250 constructed to predict the output of PUFs for a given input string (i.e., “breaking” the PUF by defeating its  
 251 authentication function). Vulnerability to breaking through machine learning attacks can vary significantly with  
 252 PUF design, and one of the challenges in developing PUFs is to identify potential weaknesses before constructing  
 253 the PUF.

254 Table 1 shows ML prediction results for the five PUF designs discussed in this paper and for 10 ML algorithms  
 255 available through the Weka machine learning tool package [17]. Note that ZeroR is a baseline, where predictions  
 256 are simply the proportion of 0 or 1 results for the challenge/response pairs in the training set. The other algorithms  
 257 were selected to provide a representative sample of popular ML algorithms of different types. AdaBoost is an  
 258 adaptive ensemble algorithm that uses a phased sequence of basic decision tree algorithms, improving on  
 259 prediction results with each phase. Bayes Net and Naïve Bayes are based on Bayesian statistical concepts.  
 260 Decision Table is a majority classifier based on a nearest neighbor algorithm. J48 and Random Forest are based  
 261 on decision trees. Stochastic gradient descent minimizes a loss function that is a weighted linear combination of  
 262 the attributes, and logistic regression uses weighted attributes in a regression function. JRip is a propositional  
 263 logic-based rule learning algorithm. Although there is a wide range of results for different algorithms, it is clear  
 264 that DB1 – the arbiter design – is much more vulnerable to ML attacks, where two algorithms are able to predict  
 265 the response to challenges with near perfect accuracy. Even the best two PUF implementations (DB3 and DB4)  
 266 are not fully resistant to revealing some bias in their responses. Note that their averages are all considerably  
 267 above the baseline ZeroR, which simply guesses in proportion to 0 or 1 responses in challenge-response pairs.

268 **Table 1. ML Prediction results for five PUF designs**

	Ada Boost	BayesNet	Decision Table	J48/C45	JRip	Logistic	Naïve Bayes	Random Forest	Stoc Grad Descent	ZeroR	Average accuracy	combined diff 2-way
DB1	77.1	96.2	75.6	72.1	77.2	99.7	96.2	87.2	99.3	55.0	86.7	0.489
DB2	54.8	54.9	76.7	68.1	75.2	54.9	54.9	71.9	52.4	55.6	62.6	0.309
DB3	50.7	50.1	71.0	63.9	67.2	50.3	50.1	62.6	50.2	50.1	57.3	0.248
DB4	57.5	56.5	58.8	54.6	60.7	56.4	56.5	55.3	54.6	50.6	56.8	0.216
NN00	64.1	64.8	62.1	59.1	64.8	64.8	64.8	65.4	62.6	50.5	63.6	0.383

269 This section shows how combination frequency differences of PUF input data can be used to determine a good  
 270 deal of information about the design and internal structure of a PUF. This is achieved by measuring the difference  
 271 between occurrences of  $t$ -way combinations associated with a 0 response as compared with a 1 response. Ideally,  
 272 there should be little difference, except for random variances. As shown below, however, these differences vary  
 273 considerably and align with the differences in predictability using machine learning. Although this work is only  
 274 preliminary, this information may be useful in identifying design deficiencies and making PUFs more resistant  
 275 to breakthrough machine learning.

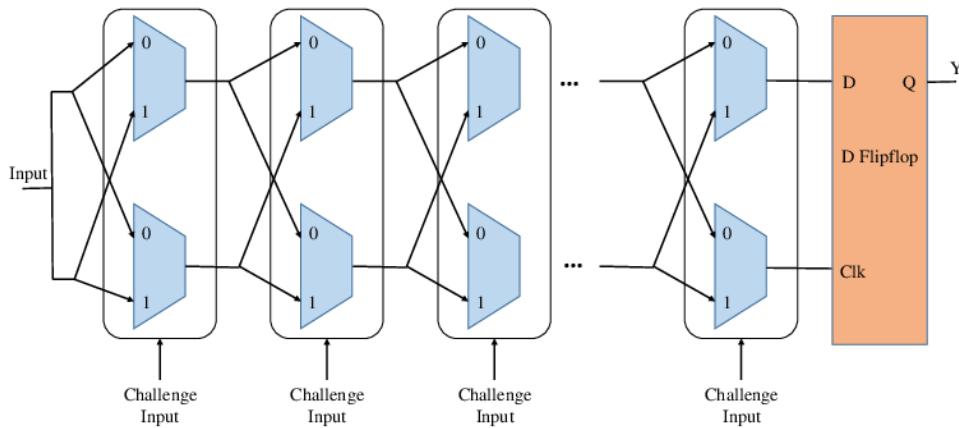
276 Comparing the accuracy of ML predictions in Table 1 with the graphs in Figures 3 through 7, it is immediately  
 277 apparent that there is a relationship between the “noisiness” of the graphs and the success of ML algorithms in  
 278 predicting or breaking the PUF. The arbiter PUF, DB1 (Figure 3), response has a very noisy graph with

279 differences for nearly every 2-way combination of bits ranging from about 0.10 to 0.25. For this PUF, ML  
 280 algorithms predict the response with up to 99.7 % accuracy. For the PUF most resistant to ML predictions, DB4  
 281 (Figure 6), the graph shows small frequency differences with nearly all under 0.05 and up to a few around 0.10.  
 282 The others fall within the range between DB1 and DB4 for both frequency differences and prediction accuracy,  
 283 which is a metric for the potential of breaking the PUF. Maximum frequency differences for DB3 are around  
 284 0.12, for DB2 about 0.15, and for the neural net PUF around 0.19 – roughly consistent with the rankings of best,  
 285 worst, and average for prediction accuracy and, hence, vulnerability to ML attacks. See the last column of Table  
 286 1, which shows the range for 2-way frequency differences above and below the center line, or  $\max(|f(x_i, C_t) - f(x_i, N_t)|) + \max(|f(x_i, N_t) - f(x_i, C_t)|)$ .  
 287

288 There are two major types of hardware implementation of PUFs: memory-based and delay-based. A typical  
 289 memory-based PUF is the SRAM PUF. Delay-based PUFs include arbiter PUFs, the pseudo-linear feedback  
 290 shift register PUF, and the ring oscillator (RO) PUF.

### 291 3.1 Arbiter PUF (DB1)

292 The main idea of an arbiter PUF is to create a digital race for signals through two paths within a chip and to have  
 293 an *arbiter* circuit that decides which signal has won the race. The two paths are designed identically. However,  
 294 the manufacturing process usually introduces a very slight longer delay in one of the paths from the other. Given  
 295 a particular challenge, the arbiter PUF will therefore produce an output dictated by the physical characteristics  
 296 of that unique hardware implementation. During an arbiter PUF design, one has to make sure that the delays  
 297 between the two paths are not too close to each other. Otherwise, the output will be dictated by noise in the signal  
 298 rather than the delay uniquely introduced through the manufacturing variation.



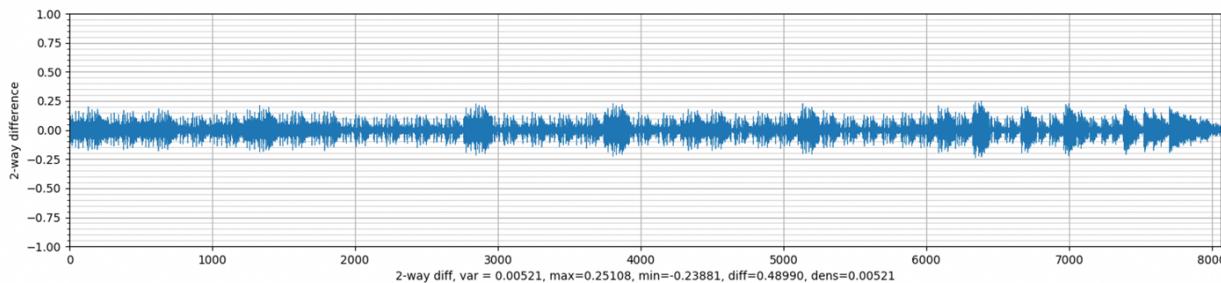
299  
 300 **Figure 3. Basic operations of an arbiter PUF**

301 As Figure 3 shows, each gate or switch-block introduces a delay for one of the outputs, which accumulates over  
 302 the blocks. This gives rise to the opportunity of building what is typically known as *model-based attacks* (also  
 303 known as *model building attacks* or *model learning attacks*). The idea is that one can build a mathematical model  
 304 of the PUF which, after observing several CRP queries, will be able to predict the response for a given challenge  
 305 with a high level of accuracy. With the proliferation of machine learning algorithms, this type of model building

306 or model learning has become easier to implement. To make model building attacks more difficult on basic  
307 arbiter PUFs, non-linearity is introduced into the delay lines of the designed circuit. For example, in case of feed-  
308 forward arbiter PUFs, some challenge bits are not set by the user. Rather, they are connected to the outputs of  
309 the intermediate arbiters evaluating the race at some intermediate point the circuit. This technique, however,  
310 increases the noise in the output of the arbiter PUF. Although initial results with feed-forward arbiter PUFs were  
311 shown to be resistant to model-building or model-learning attacks, more sophisticated learning models were able  
312 to break them [17].

313 By simply analyzing combination frequency differences (CFD) within a subset of the challenge-response pairs  
314 (CRPs) and without knowing anything about the type or design of the circuitry, one can predict which arbiter  
315 PUF design is likely to be more vulnerable to model-building attacks.

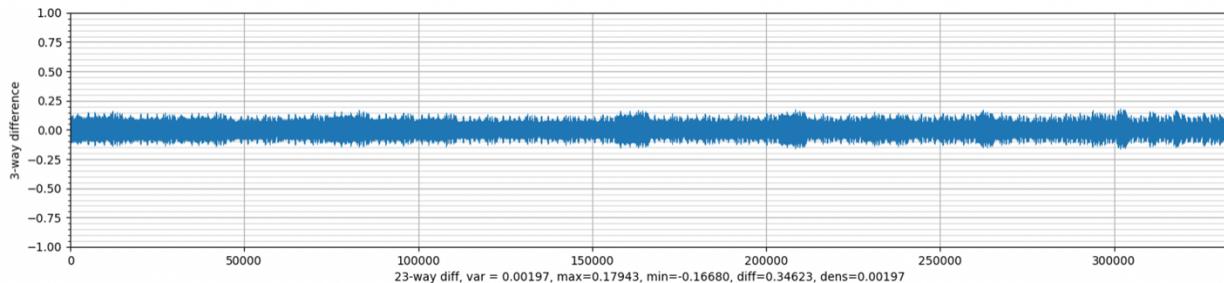
316 Figure 3(a) shows 2-way frequency differences for a 64-bit PUF, DB1, an early arbiter design with delays placed  
317 randomly in the hardware. With 64 bits, there are  $2^2 \binom{64}{2} = 8064$  2-way differences indexed. Differences range  
318 from a low of -0.23881 to a high of 0.25108 for a range of 0.48990. Note that differences are given as difference  
319  $F_{Cx} - F_{Nx}$ , so negative values are cases where non-class  $t$ -tuples of values exceed class  $t$ -tuples of values.



320

321 **Figure 3(a). 2-way frequency differences for a 64 bit arbiter PUF**

322 Figure 3(b) shows 3-way frequency differences for the same PUF. Note that variance, minimum, and  
323 maximum differences are smaller than those for 2-way combinations. The X axis indexes  $2^3 \binom{64}{3} = 333,312$   
324 combinations.



325

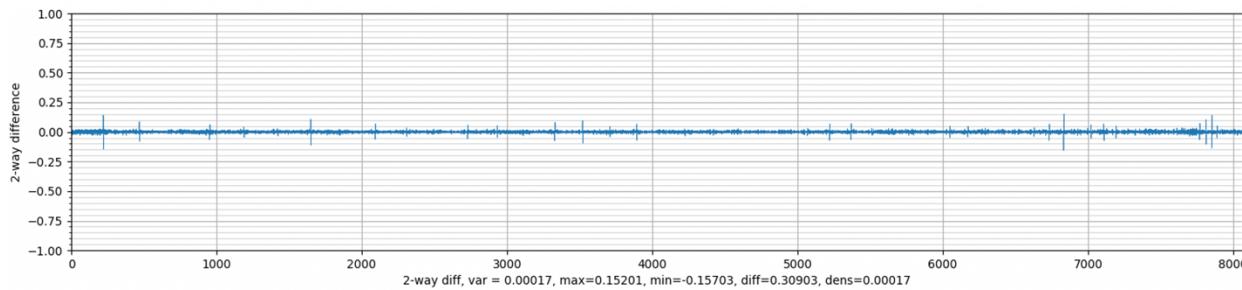
326 **Figure 3(b). 3-way frequency differences for a 64 bit arbiter PUF**

### 327 **3.2 8-bit Shift Register PUF (DB2)**

328 Shift register PUF is another delay-based PUF implementation, where a series of linear feedback shift registers

329 (LFSR) are put together to capture the unique delays associated with a physical implementation. Researchers  
330 have proposed pseudo-LFSR-based physically unclonable functions, known as PL-PUF, which are usually small  
331 in size, efficient in producing authentication ID for devices, and easy to modify to adjust the challenge-response  
332 pairs when needed [29].

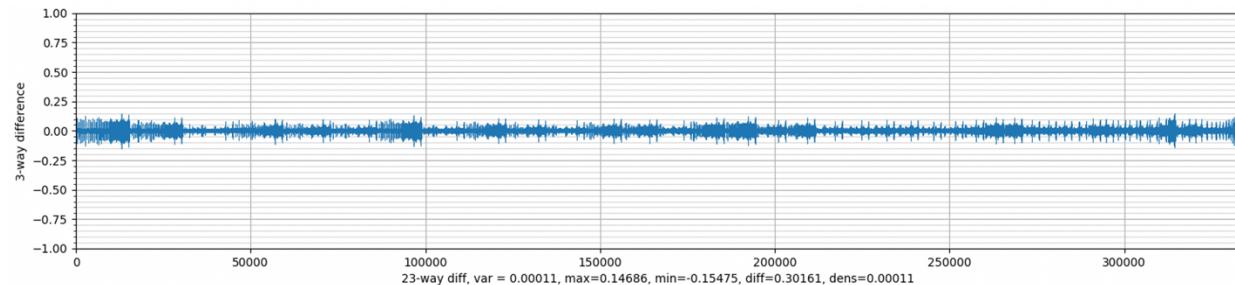
333 This section examines the security of a shift register-based PUF against a model-based attack using combinatorial  
334 frequency difference analysis. Frequency differences for an 8-bit shift register type of PUF are shown in Figure  
335 4(a) (2-way) and Figure 4(b) (3-way). Note that the variance is much smaller – 0.00017 compared to 0.0521 for  
336 2-way combinations of DB1 inputs. There is much more uniformity in the response of DB2 to 2-way and 3-way  
337 combinations of input bits, and as expected, this makes it much more difficult for ML to derive a model for the  
338 PUF that can successfully reproduce its response to inputs.



339

340

**Figure 4(a). 2-way frequency differences for an 8-bit shift register PUF**



341

342

**Figure 4(b). 3-way frequency differences for an 8-bit shift register PUF**

343 However, Figure 3(a) also shows a small number of spikes in the combination frequency chart. Combinations  
344 producing these spikes are shown in Table 2, which shows 2-way bit combinations where the frequency  
345 difference exceeds  $3\sigma$ . Combinations of almost all bits with bit 56 result in a spike that exceeds  $3\sigma$  (others have  
346 spikes that are slightly below this value but still clearly different from the other combinations). The appearance  
347 of spikes compresses towards the right end of the graph because combinations are indexed in a loop computation:  
348  $i,j,b: \text{for } i \text{ in } 0 \leq i < 63 \text{ for } j \text{ in } i+1 \leq j < 64 \text{ for } b \text{ in } \{00,01,10,11\}$ , similarly for 3-way combination indexes.

349

350

351

**Table 2. 2-way combinations with greatest frequency differences in Figure 4(a)**

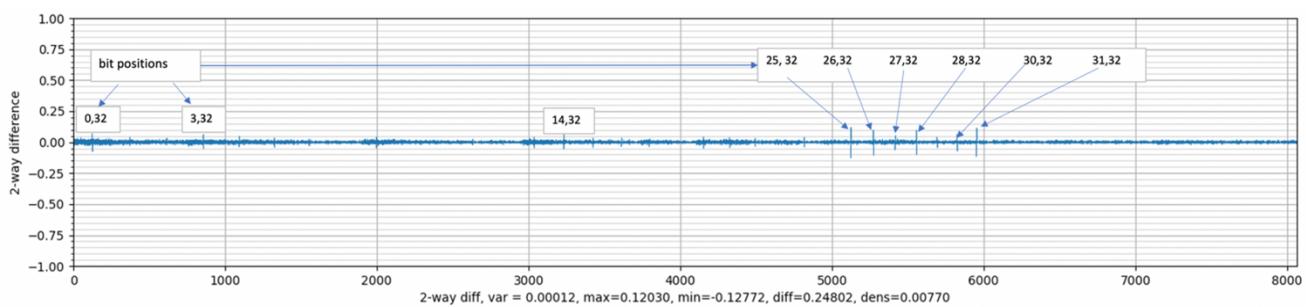
<b>bits = values</b>	<b>bits = values</b>	<b>bits = values</b>	<b>bits = values</b>
( 0,56) = (1,0)	(11,56) = (0,0)	(26,56) = (1,1)	(41,56) = (0,1)
( 0,56) = (0,1)	(12,56) = (1,0)	(26,56) = (0,0)	(42,56) = (1,1)
( 1,56) = (1,0)	(12,56) = (0,1)	(31,56) = (1,1)	(42,56) = (0,0)
( 1,56) = (0,1)	(14,56) = (1,1)	(31,56) = (0,0)	(51,56) = (1,1)
( 3,56) = (1,1)	(14,56) = (0,0)	(32,56) = (1,1)	(51,56) = (0,0)
( 3,56) = (0,0)	(15,56) = (1,0)	(37,56) = (1,1)	(52,56) = (1,0)
( 4,56) = (1,0)	(15,56) = (0,1)	(37,56) = (0,0)	(52,56) = (0,1)
( 6,56) = (1,0)	(16,56) = (0,1)	(38,56) = (1,1)	(53,56) = (1,1)
( 6,56) = (0,1)	(17,56) = (1,0)	(38,56) = (0,0)	(53,56) = (0,0)
( 8,56) = (1,1)	(17,56) = (0,1)	(40,56) = (1,0)	(54,56) = (1,1)
( 8,56) = (0,0)	(25,56) = (1,1)	(40,56) = (0,1)	(54,56) = (0,0)
(11,56) = (1,1)	(25,56) = (0,0)	(41,56) = (1,0)	

352 A potential explanation can be developed for the pattern of spikes in combinations that include bit 56 by noting  
 353 that 8 is an even divisor of 56. PUFs accumulate differences as steps progress, so bit 56 occurs at the final stage  
 354 before the last 8-bit shift register. In a design situation, the next step would be to analyze the hardware  
 355 components to determine why this irregularity was occurring.

### 356 3.3 32-bit Shift Register PUF (DB3)

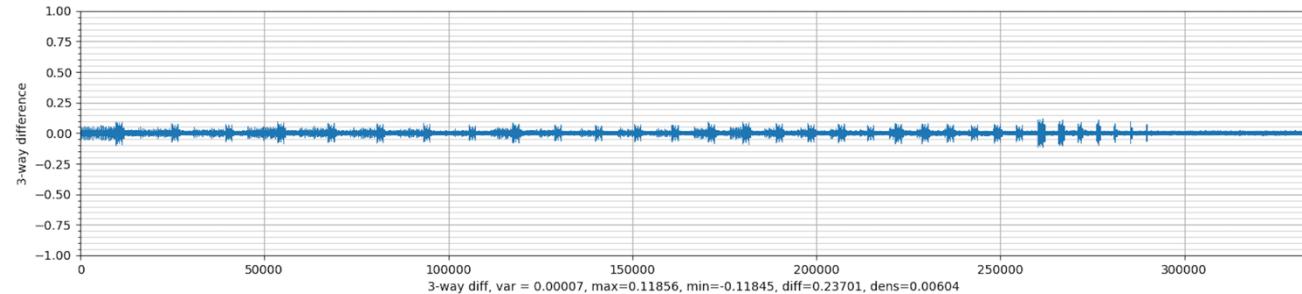
357 This section shows the results of the analysis performed on a 32-bit shift register PUF. As the name suggests, a  
 358 32-bit shift register PUF is designed the same as an 8-bit shift register, where the circuitry is four times longer.  
 359 The added circuitry increases the complexity of the PUF and, thus, likely makes it a little less susceptible to  
 360 model-building attacks.

361 The results of applying the analyses are shown in Figures 5(a) and 5(b).



362

**Figure 5(a). 2-way frequency differences for a 32-bit shift register PUF**



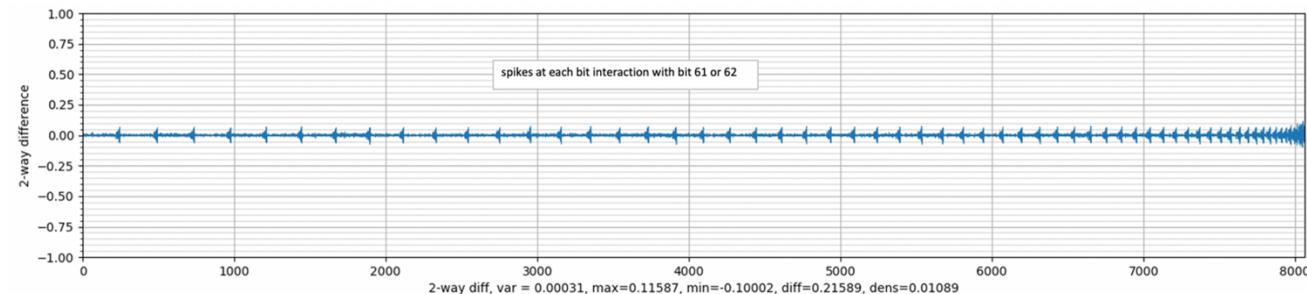
364

365

**Figure 5(b). 3-way frequency differences for a 32-bit shift register PUF**

### 366 3.4 Uniform distribution PUF (DB4)

367 Figure 6 shows results for a PUF with the most uniform distribution of all studied here. This PUF has the greatest  
 368 resistance to machine learning attacks, which are able to predict responses only somewhat better than chance  
 369 (see Table 1). In this case, the variations used in producing PUF responses accumulate uniformly with slight  
 370 frequency differences for  $t$ -tuples of bits that include either bit 61 or 62. (Compression of the spikes towards the  
 371 right side of the graph occurs because of the loop computation, as explained in Section 3.2.)

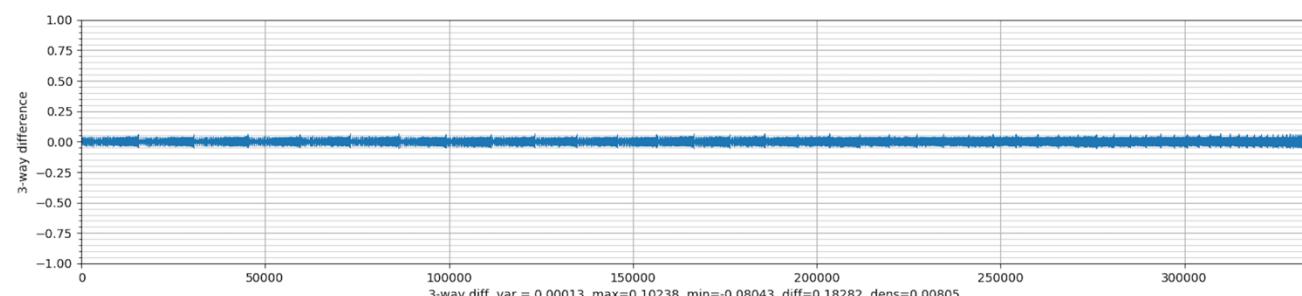


372

373

**Figure 6(a). 2-way frequency differences for a uniform distribution PUF**

374



375

376

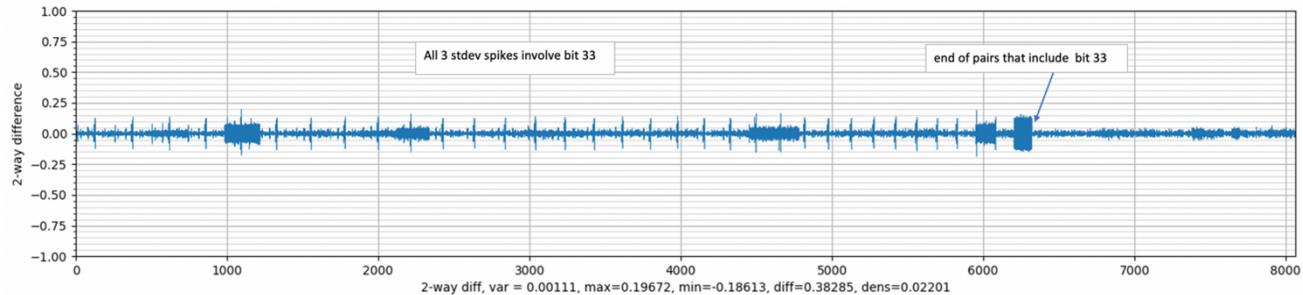
**Figure 6(b). 3-way frequency differences for a uniform distribution PUF**

377

378

### 3.5 Neural Net PUF

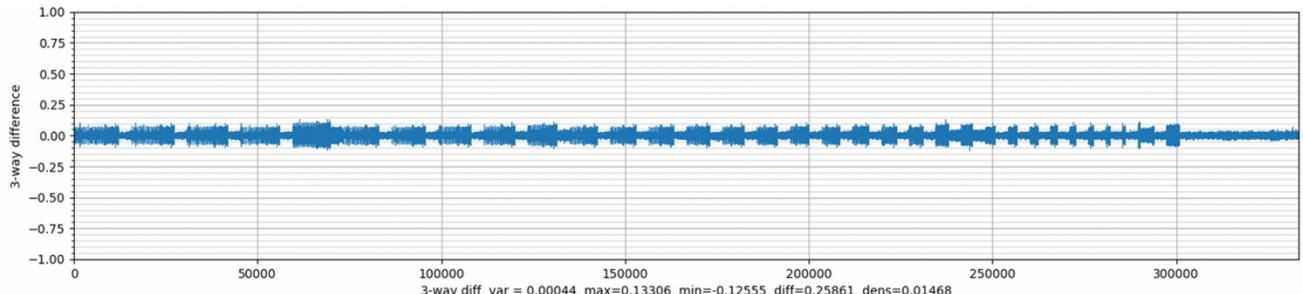
379 Researchers have pointed out the vulnerabilities of arbiter and other types of PUFs, especially against model-  
 380 building attacks [30]. To thwart the model-learning attack, researchers proposed both a simple neural network  
 381 (NN) [31] as well as recurrent neural network (RNN)-based PUFs [32]. These new models are specifically  
 382 designed for high resistance to model-building attacks achieved by introducing non-linearity between the  
 383 challenge-response pairs. The physical implementation uses current-mirrors to construct the PUF. The basic idea  
 384 is to propagate a current through two identical chains of non-linear current mirrors. In the case of RNN-based  
 385 PUF, the circuitry feeds back the challenge bits into the PUF. [32]



386

387 **Figure 7(a). 2-way frequency differences for a neural net PUF**

388



389

390 **Figure 7(b). 3-way frequency differences for a neural net PUF**

391

## 4 Extension to Machine Learning

392 A *distinguishing combination* has been defined as one present in a class instance of class set  $C$  and absent in  
 393 non-class instances  $N$ , or if it is more strongly associated with  $C$  than  $N$ , as determined by a threshold value. As  
 394 the name suggests, a distinguishing combination is one that differentiates one type or class of instance from  
 395 others. Thus, it is natural to consider if these combinations can be used directly in machine learning problems  
 396 for predicting class membership from instance attributes. If an instance contains many  $t$ -tuples of values that are  
 397 associated with a particular class but not with other classes, then it is likely to be a member of the class with  
 398 which the  $t$ -tuples of values are strongly associated. This section shows that initial results suggest this approach  
 399 works quite well in many cases. No ML algorithm is best for all problems, and the CFD approach to classification  
 400 performs better than other ML algorithms for some problems and less well for others. This section reviews some  
 401 of these empirical results and suggests future work to characterize the conditions under which CFD machine

402 learning will be advantageous.

403 Given a set of distinguishing combinations, a simple algorithm for classification seems natural: if an instance  
 404 has more attribute combinations that are associated with a class  $C$  than another class, then assign it to  $C$ , and if  
 405 there are fewer combinations associated with  $C$  than another class, then assign it to the other class. (For  
 406 simplicity, only two classes are considered here, but the method can be extended to multiple classes by  
 407 considering each one as “ $C$ ” in turn). If the  $C$  and  $N$  combinations are equally present, then the result is  
 408 undermined. As the saying goes, “if it looks like a duck and walks like a duck and quacks like a duck (a 3-way  
 409 combination), it’s probably a duck!”

410 CFD algorithm:

```
411     dist_c = {distinguishing combinations for instances in class C}
412     dist_n = {distinguishing combinations for instances not in class C}
413
414     dc = sum(1 for t-way combinations x_i in row if x_i in dist_c)
415     dn = sum(1 for t-way combinations x_i in row if x_i in dist_n)
416     if dc > dn: predict C
417     if dn > dc: predict N
418     if dc == dn: indeterminate
```

419 A number of possible alternatives to the basic algorithm can be conceived. Perhaps the most obvious is to weigh  
 420 the presence of distinguishing combinations in instances, shown below as CFDw. Using a weight of  $|F_{Cx} - F_{Nx}|$ ,  
 421 the CFDw algorithm has been compared with the basic CFD for several examples. Comparisons of the weighted  
 422 method with the basic method are shown in the following sections along with frequency difference graphs.  
 423 Accuracy scores for CFD and CFDw are relatively close, and there is no clear winner between these two  
 424 variations.

425 CFDw algorithm:

```
426     dc = sum(weight(x_i) for t-way combinations x_i in row if x_i in dist_c)
427     dn = sum(weight(x_i) for t-way combinations x_i in row if x_i in dist_n)
428     if dc > dn: predict C
429     if dn > dc: predict N
430     if dc == dn: indeterminate
```

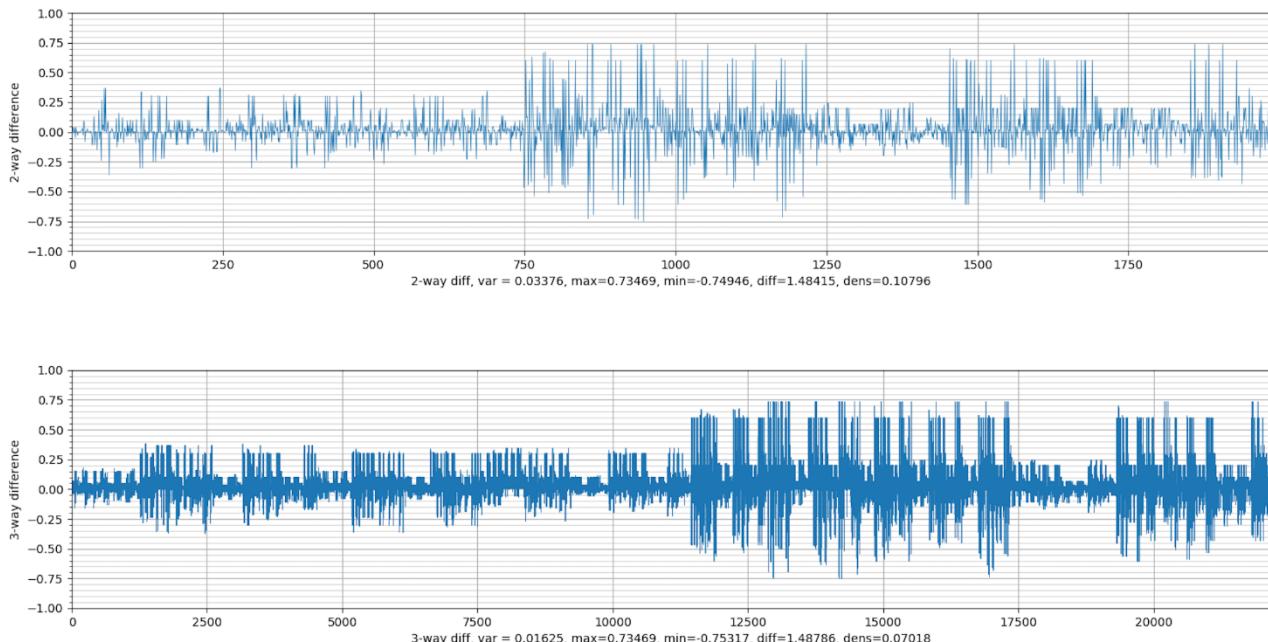
431 Using this approach on the PUF data presented in the previous section produces results that are relatively  
 432 comparable to the ML algorithms shown in Table 1 for 10,000 rows using 4-way combinations shown in Table  
 433 3.

434 **Table 3. Comparison of CFD accuracy with average, best, worst from Table 1**

	CFD	Avg, Table 1	Best, Table 1	Worst, Table 1
DB1	.953	86.7	.997 (logistic)	.721 (J48)
DB2	.547	62.6	.767 (dec tbl)	.524 (SGD)
DB3	.520	57.3	.710 (dec tbl)	.501 (Bayesnet)
DB4	.546	56.8	.607 (JRip)	.546 (SGD)
NN00	.621	63.6	.654 (Rand Forest)	.591 (J48)

435 As previously discussed, PUFs are designed to be “unclonable” (i.e., difficult to replicate, including through  
 436 strategies such as machine learning). In most ML applications, the classes of interest are in nature or may be  
 437 industrial products not designed to be resistant to modeling. This difference is also immediately apparent in the  
 438 graphs in Appendix A, which show much wider variation for these “natural” or practical datasets. An example  
 439 is shown in Figure 8 below (mushroom data set from Appendix):

440



441

442 **Figure 8. Frequency difference graph for 2-way and 3-way differences, mushroom example**

443 As shown in this graph and others in Appendix A, there is a much wider variation in frequency differences – up  
 444 to roughly 75 % or more. The much smaller variation for PUFs is likely due to the fact that they are designed to  
 445 be difficult to clone or replicate. The wider range of frequency differences in these natural examples make the  
 446 CFD approach more effective, using the differences to distinguish between classes. On these applications, CFD  
 447 class prediction does quite well, as shown in Table 4. Accuracy scores in the column labeled “CFD4 @ T” are  
 448 the average of 10 random assignments of the total number of rows given by “*n* rows” split into 66 % training  
 449 and 34 % test for the threshold of T shown using 4-way combinations.

450

**Table 4. Comparison of CFD accuracy with other ML algorithms**

Dataset	n row	n col	n class	n non	CFD4 @ T	Ada	Baye	DecTbl	J48	JRip	Log	NB	Rand	SGD	ZR
Bcanc	286	9	68	218	.970@1.0	.759	.766	.745	.769	.720	.752	.752	.745	.766	.762
Coupon	12684	25	7210	5474	.730@5.0	.644	.663	.688	.718	.725	.693	.663	.757	.684	.569
Credit	1000	20	37	963	.991 @ 5.1	.963	.950	.962	.963	.957	.958	.949	.963	.963	.963
Diab	768	8	367	401	.992@1.0	.698	.723	.709	.694	.692	.728	.723	.674	.715	.522
Heart2	47786	21	23893	23893	.755@5.0	.745	.741	.745	.757	.754	.767	.741	.753	.762	.500
Mush	5644	22	2156	3488	1.00 @ 1.0	.963	.985	1.00	1.00	1.00	1.00	.974	1.00	1.00	.618
Soyb	684	31	133	551	.986 @ 15.0	.991	.968	.988	.981	.972	.975	.929	.983		.845

451 It is important to note that a small number of threshold values have been tried. Further experimentation with  
452 threshold values and characterization of their applicability will be the subject of future research. An additional  
453 issue to be investigated is the possibility of overfitting. Two of the sample machine learning data sets have less  
454 than 10 attributes. Using 4-way combinations to test for membership in class or non-class sets may have a  
455 potential for overfitting because a 4-way combination could include roughly half of the attributes available for  
456 classifying an instance. The other data sets were chosen with more than 20 attributes to reduce the possibility of  
457 overfitting. A detailed investigation of this issue will be the subject of future research.

458 **5 Conclusions**

459 This paper presents a method for measuring and visualizing differences in the frequency or rate of occurrence of  
460 t-way combinations for two data sets. This measure, combination frequency differencing (CFD), has potential  
461 use in a variety of applications. Initially applied to challenge-response pairs for physical unclonable functions of  
462 PUFs, CFD was shown to provide the ability to identify combinations of bits in the challenge that are more or  
463 less strongly associated with particular output values of 0 or 1. The level of difference appears to correlate with  
464 the effectiveness of machine learning attacks on PUFs. In future research, the authors hope to develop ways to  
465 trace these strongly non-uniform bit combination associations to the hardware components that produce them.  
466 This ability might be useful in the design and development of PUFs to identify design weaknesses and correct  
467 them before production.

468 It was also shown that the basic idea behind CFD can be extended to produce a new type of machine learning  
469 algorithm. CFD identifies and measures differences between two data sets using attribute value combinations,  
470 and this approach lends itself naturally to identifying instances in classification problems. An instance that is  
471 very similar to others of a particular class is likely to be a member of that class. This paper shows that the  
472 accuracy of this CFD approach to classification problems is comparable to the accuracy of well-known  
473 algorithms across a variety of problem types. Further research is planned to investigate developing this method  
474 into a practical approach for classification problems. In previous work, the authors have used the concept of  
475 unique or distinguishing combinations for explainability in AI/ML systems [23][28], so there may be effective  
476 methods for combining the CFD method for classification with explainability.

477

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555 [32] Shah, Nimesh, et al. "A 0.16 pj/bit recurrent neural network based PUF for enhanced machine learning attack resistance."  
556 *Proceedings of the 24th Asia and South Pacific Design Automation Conference*. 2019.  
557

558

559

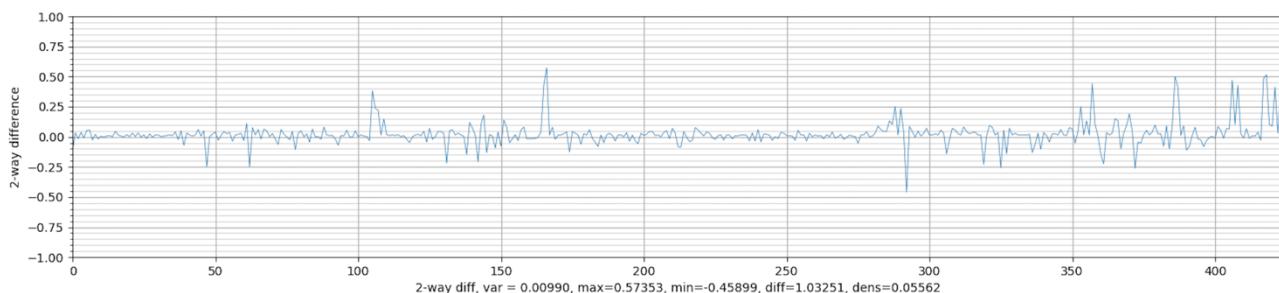
560

## 561 Appendix A—Difference Graphs of Classification Problems

562 This section presents examples of typical machine learning classification problems taken from the UCI  
 563 Machine Learning Repository (<https://archive.ics.uci.edu/ml>) or from Kaggle (<https://www.kaggle.com>). Each  
 564 example includes the data source, associated publication, and results from the tools described in this paper.

565 **Bcanc** - <https://archive.ics.uci.edu/ml/datasets/Breast+Cancer>

566 Michalski,R.S., Mozetic,I., Hong,J., & Lavrac,N. (1986). The Multi-Purpose Incremental Learning System  
 567 AQ15 and its Testing Application to Three Medical Domains. *Proceedings of the Fifth National Conference*  
 568 *on Artificial Intelligence*, 1041-1045, Philadelphia, PA: Morgan Kaufmann.



569

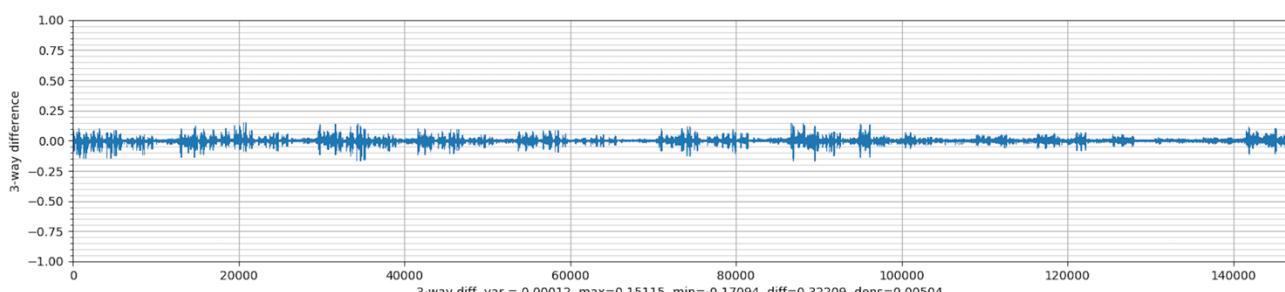
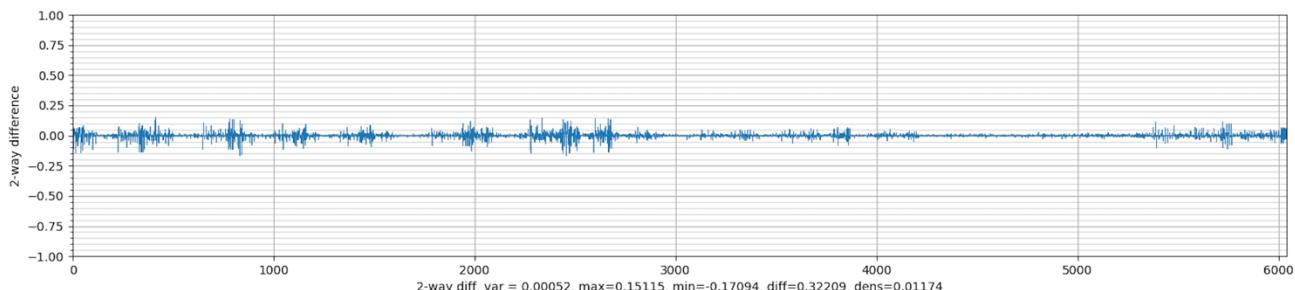
570 **Figure 9. Breast cancer data frequency differences.**

571 CFD results:  
 572 == confusion matrix 4-way ==  
 573 | C | N <- predicted  
 574 C | 75 | 0  
 575 N | 3 | 21  
 576 ======  
 577 Accuracy: 0.970  
 578

579 CFDw results:  
 580 == confusion matrix 4-way ==  
 581 | C | N <- predicted  
 582 C | 24 | 0  
 583 N | 2 | 73  
 584 ======  
 585 Accuracy: 0.980  
 586 ======  
 587

**588 Coupon** - <https://www.kaggle.com/mathurinache/invehicle-coupon-recommendation>

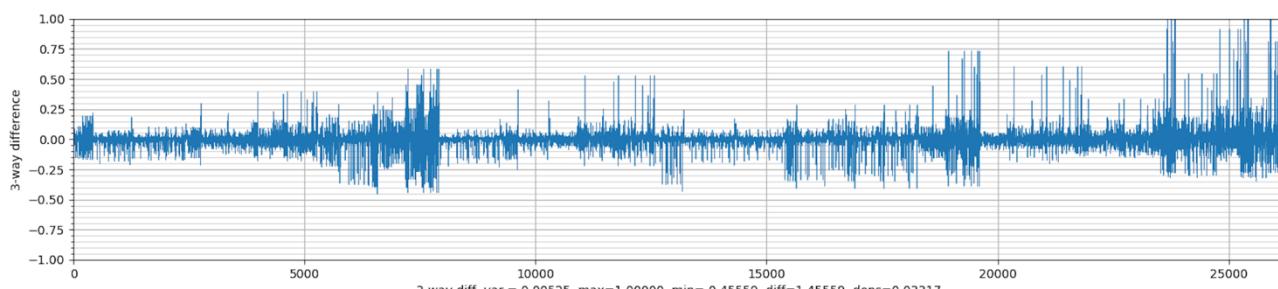
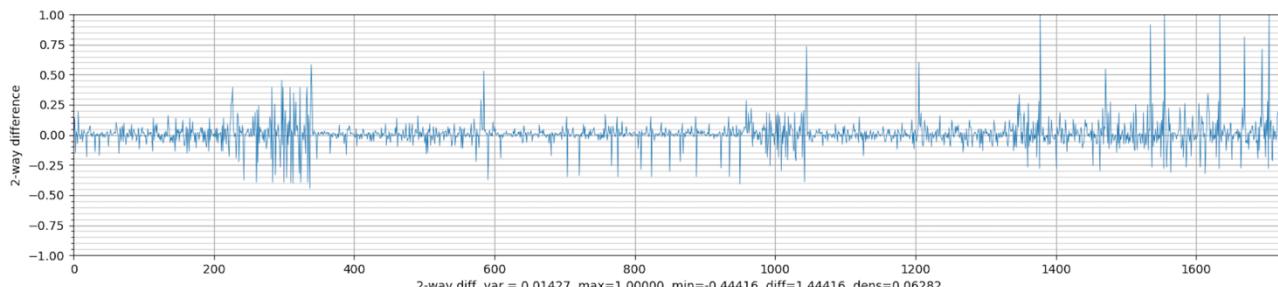
589 Wang, Tong, Cynthia Rudin, Finale Doshi-Velez, Yimin Liu, Erica Klampfl, and Perry MacNeille. 'A  
590 Bayesian framework for learning rule sets for interpretable classification.' *The Journal of Machine Learning*  
591 Research 18, no. 1 (2017): 2357-2393.



**Figure 10. Coupon data frequency differences.**

```
594      == confusion matrix 4-way ==
595      |       C       |       N      <- predicted
596      C |     1931    |      521
597      N |      661    |    1201
598      =====
599      Accuracy:  0.726
600      =====
601
```

603 Credit - [https://archive.ics.uci.edu/ml/citation\\_policy.html](https://archive.ics.uci.edu/ml/citation_policy.html)  
 604 Dua, D. and Graff, C. (2019). UCI Machine Learning Repository [http://archive.ics.uci.edu/ml]. Irvine, CA: University of  
 605 California, School of Information and Computer Science.  
 606  
 607



608

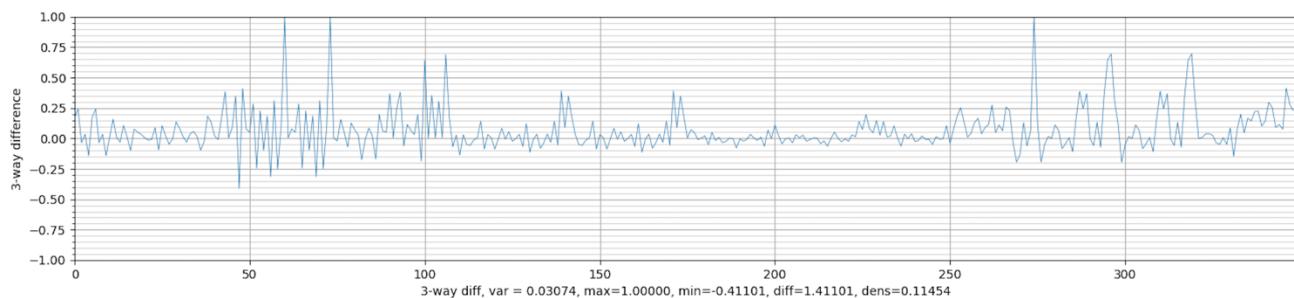
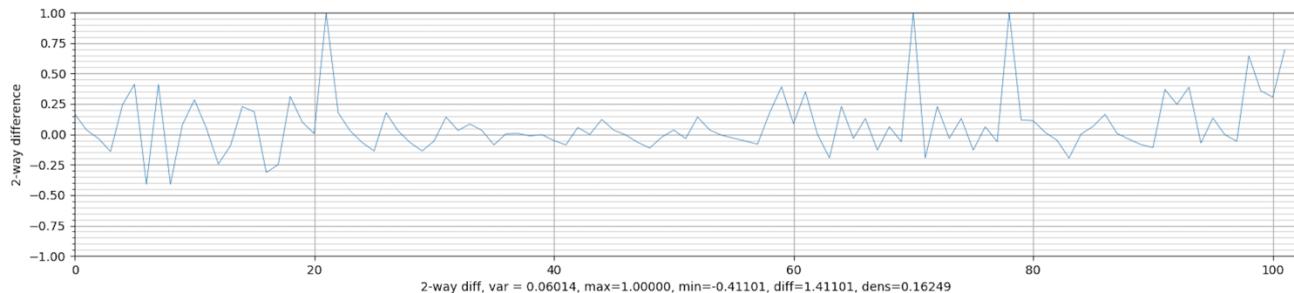
609 **Figure 11. German credit check data frequency differences.**

610 == confusion matrix 4-way ==  
 611 | C | N <- predicted  
 612 C | 10 | 3  
 613 N | 0 | 328  
 614 ======  
 615 Accuracy: 0.991  
 616 ======  
 617  
 618  
 619 CFDw results:  
 620 == confusion matrix 4-way ==  
 621 | C | N <- predicted  
 622 C | 293 | 35  
 623 N | 0 | 13  
 624 ======  
 625 Accuracy: 0.897  
 626 ======

627

628 Diab - <https://archive.ics.uci.edu/ml/datasets/diabetes>

629 Smith, J.W., Everhart, J.E., Dickson, W.C., Knowler, W.C., Johannes, R.S. (1988). Using the ADAP learning  
630 algorithm to forecast the onset of diabetes mellitus. *Proceedings of the Symposium on Computer Applications*  
631 and Medical Care (pp. 261--265). IEEE.



632

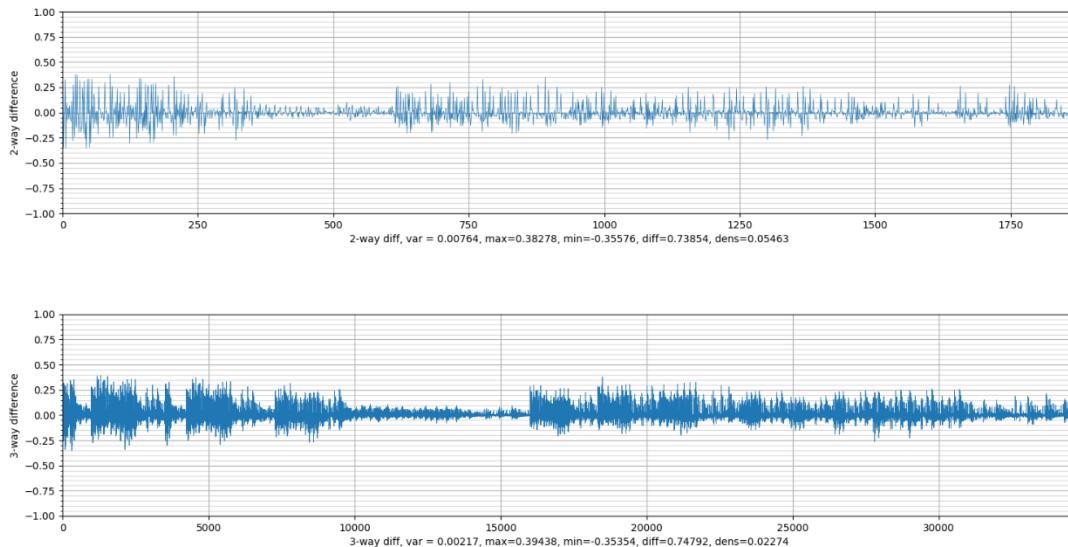
633 **Figure 12. Diabetes data frequency differences.**

634 CFD results:  
635 == confusion matrix 4-way ==  
636 | C | N <- predicted  
637 C | 124 | 1  
638 N | 1 | 136  
639 =====  
640 Accuracy: 0.992  
641  
642 CFDw results:  
643 == confusion matrix 4-way ==  
644 | C | N <- predicted  
645 C | 125 | 0  
646 N | 1 | 136  
647 =====  
648 Accuracy: 0.996  
649 =====

650

651 **Heart2 - <https://github.com/doguilmak/Heart-Diseaseor-Attack-Classification>**

652 Large set of data containing 253,681 instances, with 23,893 heart disease or attack, and the rest healthy. To  
653 make instance sets equal size, a random set of 23,893 disease/attack instances were extracted.



654

655 **Figure 13. Heart disease data frequency differences.**

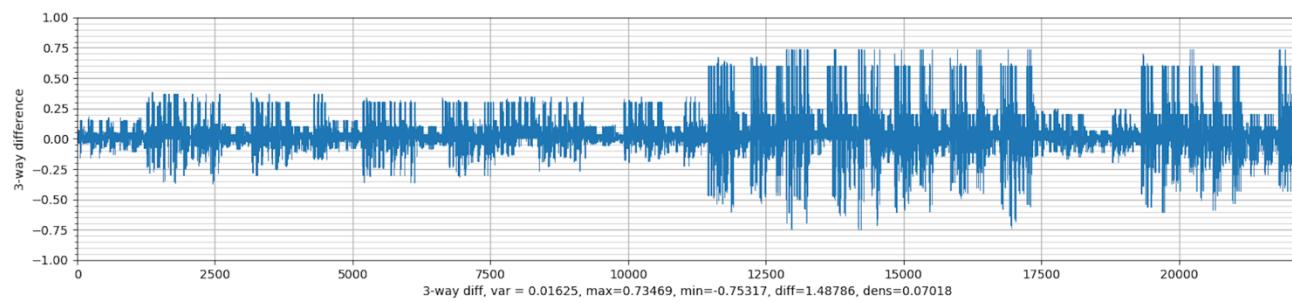
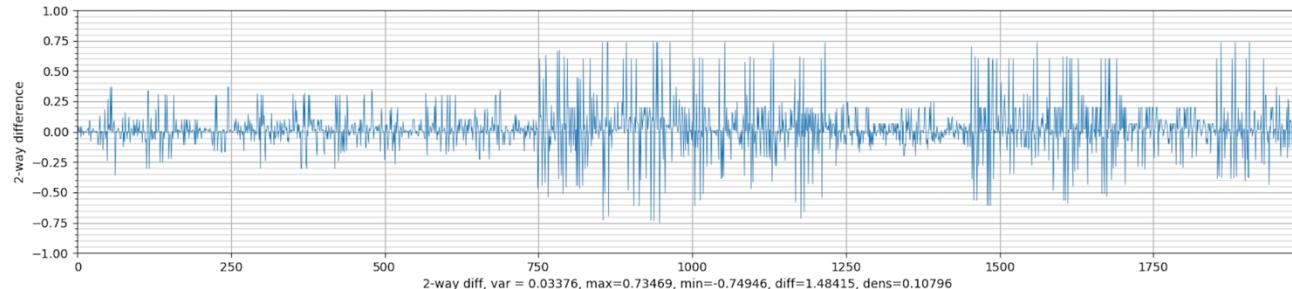
```

656 == confusion matrix 2-way ==
657 | C | N <- predicted
658 C | 6254 | 1870
659 N | 2472 | 5652
660 =====
661 Accuracy 2-way = 0.733, SD_2 = 0.08709, dt = 1.046
662
663 == confusion matrix 3-way ==
664 | C | N <- predicted
665 C | 6017 | 2107
666 N | 1950 | 6174
667 =====
668 Accuracy 3-way = 0.750, SD_3 = 0.04643, dt = 1.046
669 == confusion matrix 4-way ==
670 | C | N <- predicted
671 C | 5905 | 2219
672 N | 1764 | 6360
673 =====
674 Accuracy 4-way = 0.755, dt = 1.046
675
676 == confusion matrix 5-way ==
677 | C | N <- predicted
678 C | 5859 | 2265
679 N | 1991 | 6133
680 =====
681 Accuracy 5-way = 0.738, dt = 1.046
682 =====
683

```

684 **Mush - <https://archive.ics.uci.edu/ml/datasets/Mushroom>**

685 Schlimmer, J.S. (1987). Concept Acquisition Through Representational Adjustment (Technical Report 87-19).  
686 Doctoral dissertation, Department of Information and Computer Science, University of California, Irvine.



687

688 **Figure 14. Edible mushroom data frequency differences.**

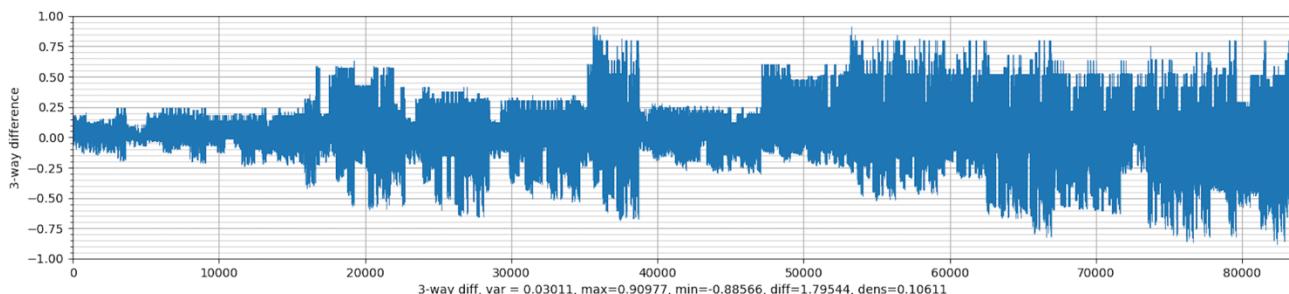
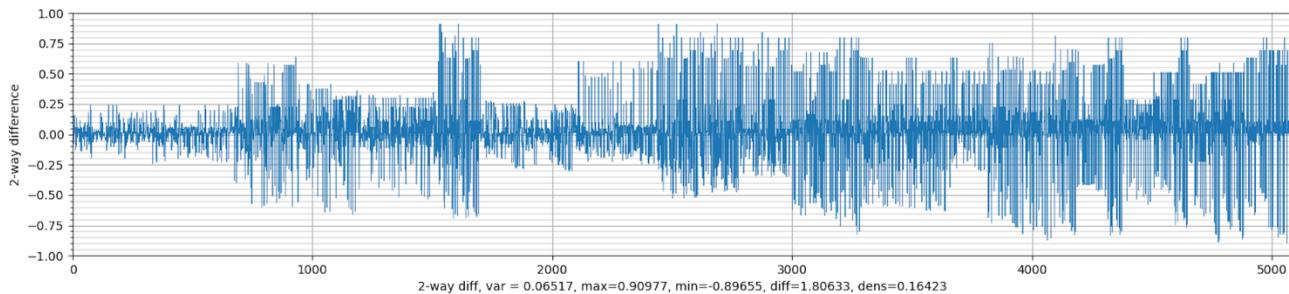
689 CFD results:  
690 == confusion matrix 4-way ==  
691 | C | N <- predicted  
692 C | 725 | 9  
693 N | 58 | 1128  
694 =====  
695 Accuracy: 0.965  
696

697 CFDw results:  
698 == confusion matrix 4-way ==  
699 | C | N <- predicted  
700 C | 553 | 181  
701 N | 0 | 1186  
702 =====  
703 Accuracy: 0.906  
704 =====

706

707 707 **Soyb - <https://archive.ics.uci.edu/ml/datasets/Soybean+Large%29>**

708 R.S. Michalski and R.L. Chilausky. "Learning by Being Told and Learning from Examples: An Experimental  
709 Comparison of the Two Methods of Knowledge Acquisition in the Context of Developing an Expert System  
710 for Soybean Disease Diagnosis", *International Journal of Policy Analysis and Information Systems*, Vol. 4,  
711 No. 2, 1980.



712

713 **Figure 15. Soybean disease data frequency differences.**

714 CFD results:  
715 == confusion matrix 4-way ==  
716 | C | N <- predicted  
717 C | 42 | 4  
718 N | 0 | 188  
719 =====  
720 Accuracy: 0.983  
721

722 CFDw results:  
723 == confusion matrix 4-way ==  
724 | C | N <- predicted  
725 C | 41 | 5  
726 N | 0 | 188  
727 =====  
728 Accuracy: 0.979  
729 =====

730