

MIPS Architecture and Assembly Language Overview

Adapted from: [http://edge.mcs.dre.g.el.edu/GICL/people/sevy/architecture/MIPSRef\(SPIM\).html](http://edge.mcs.dre.g.el.edu/GICL/people/sevy/architecture/MIPSRef(SPIM).html)

[\[Register Description\]](#) [\[I/O Description\]](#)

Data Types and Literals

Data types:

- Instructions are all 32 bits
- byte(8 bits), halfword (2 bytes), word (4 bytes)
- a character requires 1 byte of storage
- an integer requires 1 word (4 bytes) of storage

Literals:

- numbers entered as is. e.g. 4
- characters enclosed in single quotes. e.g. 'b'
- strings enclosed in double quotes. e.g. "A string"

Registers

- 32 general-purpose registers
- register preceded by \$ in assembly language instruction
two formats for addressing:
 - using register number e.g. \$0 through \$31
 - using equivalent names e.g. \$t1, \$sp
- special registers Lo and Hi used to store result of multiplication and division
 - not directly addressable; contents accessed with special instruction mfhi ("move from Hi") and mflo ("move from Lo")
- stack grows from high memory to low memory

This is from Figure 9.9 in the Goodman&Miller text

Register Number	Alternative Name	Description
0	zero	the value 0
1	\$at	(assembler temporary) reserved by the assembler
2-3	\$v0 - \$v1	(values) from expression evaluation and function results
4-7	\$a0 - \$a3	(arguments) First four parameters for subroutine. Not preserved across procedure calls
8-15	\$t0 - \$t7	(temporaries) Caller saved if needed. Subroutines can use w/out saving. Not preserved across procedure calls
16-23	\$s0 - \$s7	(saved values) - Callee saved. A subroutine using one of these must save original and restore it before exiting. Preserved across procedure calls
24-25	\$t8 - \$t9	(temporaries) Caller saved if needed. Subroutines can use w/out saving. These are in addition to \$t0 - \$t7 above. Not preserved across procedure calls.
26-27	\$k0 - \$k1	reserved for use by the interrupt/trap handler
28	\$gp	global pointer. Points to the middle of the 64K block of memory in the static data segment.
29	\$sp	stack pointer Points to last location on the stack.
30	\$s8/\$fp	saved value / frame pointer Preserved across procedure calls
31	\$ra	return address

See also Britton section 1.9, Sweetman section 2.21, Larus Appendix section A.6

Program Structure

- just plain text file with data declarations, program code (name of file should end in suffix .s to be used with SPIM simulator)
- data declaration section followed by program code section

Data Declarations

- placed in section of program identified with assembler directive **.data**
- declares variable names used in program; storage allocated in main memory (RAM)

Code

- placed in section of text identified with assembler directive **.text**
- contains program code (instructions)
- starting point for code e.g. execution given label **main:**
- ending point of main code should use exit system call (see below under System Calls)

Comments

- anything following # on a line
This stuff would be considered a comment
- Template for a MIPS assembly language program:

```
# Comment giving name of program and description of function
# Template.s
# Bare-bones outline of MIPS assembly language program

        .data          # variable declarations follow this line
        # ...

        .text          # instructions follow this line

main:    # indicates start of code (first instruction to execute)
        # ...

# End of program, leave a blank line afterwards to make SPIM happy
```

Data Declarations

format for declarations:

```
name:    storage_type    value(s)
```

- create storage for variable of specified type with given name and specified value
- value(s) usually gives initial value(s); for storage type `.space`, gives number of spaces to be allocated

Note: labels always followed by colon (:)

example

```
var1:    .word    3      # create a single integer variable with initial value 3
array1:   .byte    'a','b' # create a 2-element character array with elements initialized
                        # to a and b
array2:   .space   40     # allocate 40 consecutive bytes, with storage uninitialized
                        # could be used as a 40-element character array, or a
                        # 10-element integer array; a comment should indicate which!
```

Load / Store Instructions

- RAM access only allowed with load and store instructions
- all other instructions use register operands

load:

```
lw      register_destination, RAM_source

#copy word (4 bytes) at source RAM location to destination register.

lb      register_destination, RAM_source

#copy byte at source RAM location to low-order byte of destination register,
# and sign-e.g. extend to higher-order bytes
```

store word:

```
sw      register_source, RAM_destination

#store word in source register into RAM destination

sb      register_source, RAM_destination

#store byte (low-order) in source register into RAM destination
```

load immediate:

```
li      register_destination, value
```

#load immediate value into destination register

example:

```
.data
var1: .word 23          # declare storage for var1; initial value is 23

.text
__start:
    lw    $t0, var1      # load contents of RAM location into register $t0: $t0 = var1
    li    $t1, 5         # $t1 = 5 ("load immediate")
    sw    $t1, var1      # store contents of register $t1 into RAM: var1 = $t1
done
```

Indirect and Based Addressing

- Used only with load and store instructions

load address:

```
la    $t0, var1
```

- copy RAM address of var1 (presumably a label defined in the program) into register \$t0

indirect addressing:

```
lw    $t2, ($t0)
```

- load word at RAM address contained in \$t0 into \$t2

```
sw    $t2, ($t0)
```

- store word in register \$t2 into RAM at address contained in \$t0

based or indexed addressing:

```
lw    $t2, 4($t0)
```

- load word at RAM address (\$t0+4) into register \$t2
- "4" gives offset from address in register \$t0

```
sw    $t2, -12($t0)
```

- store word in register \$t2 into RAM at address (\$t0 - 12)
- negative offsets are fine

Note: based addressing is especially useful for:

- arrays; access elements as offset from base address
- stacks; easy to access elements at offset from stack pointer or frame pointer

example

```
.data
array1: .space 12      # declare 12 bytes of storage to hold array of 3 integers

.text
__start:
    la    $t0, array1    # load base address of array into register $t0
    li    $t1, 5         # $t1 = 5 ("load immediate")
    sw    $t1, ($t0)      # first array element set to 5; indirect addressing
    li    $t1, 13        # $t1 = 13
    sw    $t1, 4($t0)     # second array element set to 13
    li    $t1, -7        # $t1 = -7
    sw    $t1, 8($t0)     # third array element set to -7
done
```

Arithmetic Instructions

- most use 3 operands
- all operands are registers; no RAM or indirect addressing
- operand size is word (4 bytes)

```
add    $t0,$t1,$t2      # $t0 = $t1 + $t2; add as signed (2's complement) integers
sub    $t2,$t3,$t4      # $t2 = $t3 - $t4
addi   $t2,$t3, 5        # $t2 = $t3 + 5; "add immediate" (no sub immediate)
addu   $t1,$t6,$t7      # $t1 = $t6 + $t7; add as unsigned integers
subu   $t1,$t6,$t7      # $t1 = $t6 + $t7; subtract as unsigned integers

mult   $t3,$t4          # multiply 32-bit quantities in $t3 and $t4, and store 64-bit
```

Control Structures

- The `print_string` service expects the address to start a null-terminated character string. The directive `.asciiz` creates a null-terminated character string.
- The `read_int`, `read_float` and `read_double` services read an entire line of input up to and including the newline character.
- The `read_string` service has the same semantics as the UNIX library routine `fgets`.

- It reads up to n-1 characters into a buffer and terminates the string with a null character.
- If fewer than n-1 characters are in the current line, it reads up to and including the newline and terminates the string with a null character.
- The sbrk service returns the address to a block of memory containing n additional bytes. This would be used for dynamic memory allocation.
- The exit service stops a program from running.

e.g. Print out integer value contained in register \$t2

```
li    $v0, 1           # load appropriate system call code into register $v0;
                        # code for printing integer is 1
move  $a0, $t2         # move integer to be printed into $a0: $a0 = $t2
syscall                # call operating system to perform operation
```

e.g. Read integer value, store in RAM location with label int_value (presumably declared in data section)

```
li    $v0, 5           # load appropriate system call code into register $v0;
                        # code for reading integer is 5
syscall                # call operating system to perform operation
sw    $v0, int_value   # value read from keyboard returned in register $v0;
                        # store this in desired location
```

e.g. Print out string (useful for prompts)

```
string1    .data
            .ascii "Print this.\n"      # declaration for string variable,
                                         # .ascii directive makes string null terminated

main:      .text
            li    $v0, 4                # load appropriate system call code into register $v0;
                                         # code for printing string is 4
            la    $a0, string1          # load address of string to be printed into $a0
            syscall                    # call operating system to perform print operation
```

e.g. To indicate end of program, use **exit** system call; thus last lines of program should be:

```
li    $v0, 10          # system call code for exit = 10
syscall                # call operating sys
```