# Phil/LPS 31 Introduction to Inductive Logic Lecture 5

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April 12th 2023

#### **Topics**

- ▶ Inference in Quantified Relational Logic
- ► Models and Counterexamples

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  - ∴2. Someone is wise.  $(\exists x(Wx))$  (From 1 and By Existential Generalization)

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- Here is an instance of an inference that uses Universal Instantiation.
  - 1. All logicians are wise.  $(\forall x(Lx \rightarrow Wx))$
  - $\therefore$  2. If Ruth Barcan Marcus is a logician, then she is wise. ( $La \rightarrow Wa$ ) (From 1 and By Universal Instantiation)

- 4. From  $\exists x \mathcal{F}$ , you may suppose that  $\mathcal{F}(a)$  and use  $\mathcal{F}(a)$  together with other premises to infer a sentence G using other truth preserving rules of inference. Here  $\mathcal{F}(a)$  is the result of uniformly replacing every free occurrence of x in  $\mathcal{F}$  with an arbitrary constant a, which has not occurred before. This rule is known as Existential Supposition.
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- ► These are the only truth-preserving rules of inference that we will avail for ourselves in quantified relational logic but other rules of inference can be derived from these.

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- ► This will require a tool which we will call a model or intended interpretation of some sentences in quantified relational logic.