

Elementary Number Theory

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Modular Arithmetic

- $a \equiv b \pmod{m}$ iff $a - b$ is divisible by m
- If $a \equiv c$ and $b \equiv d \pmod{m}$, then
$$a \pm b \equiv c \pm d \pmod{m}$$
$$a \times b \equiv c \times d \pmod{m}$$
$$a^k \equiv c^k \pmod{m} \text{ for all integers } k$$
$$\frac{a}{b} \equiv ? \pmod{m}$$

Prime Number

- Positive integers which have exactly 2 positive divisors
- 2, 3, 5, 7, ...

Primality Testing

- Obvious $O(\sqrt{n})$ loop
- Sieve Eratosthenes ($O(n \log \log n)$)

```
24 -     for (int i = 1; i <= MAXN; i++) {  
25         p[i] = i;  
26     }  
27 -     for (int i = 2; i * i <= MAXN; i++) {  
28         if (p[i] == i) {  
29             for (int j = i * i; j <= MAXN; j += i) {  
30                 p[j] = i;  
31             }  
32         }  
33     }
```

– It is fast; 186ms for $n = 10^7$

Theorems Involving Primes

- Fermat Little Theorem

$$a^p \equiv a \pmod{p}$$

$$a^{p-1} \equiv 1 \pmod{p} \text{ for } a \not\equiv 0 \pmod{p}$$

- Wilson's Theorem

$$(n - 1)! \equiv -1 \pmod{n} \text{ if and only if } n \text{ is prime}$$

Multiplicative Functions

- $f: N \rightarrow N$ such that $f(xy) = f(x)f(y)$ for all relatively prime numbers x and y (i.e. $\gcd(x, y) = 1$)
- Obvious example: $f(x) = x$

Sigma, Tau, Phi

- $\sigma(x)$ denotes the *sum* of positive divisors of x
- $\tau(x)$ denotes the *number* of positive divisors of x
- $\phi(x)$ (also called *Euler totient function*) denotes the number of i such that $i \leq x$ and $\gcd(i, x) = 1$
- All of them are multiplicative functions

Sigma, Tau, Phi

Let $x = p_1^{e_1} p_2^{e_2} \cdots p_k^{e_k}$

- $\tau(x) = \prod_{i=1}^k (e_i + 1)$
- $\sigma(x) = \prod_{i=1}^k \left(\frac{p_i^{e_i+1} - 1}{p_i - 1} \right)$
- $\phi(x) = x \prod_{i=1}^k \left(1 - \frac{1}{p_i} \right)$

Euler's Theorem

- $a^{\phi(n)} \equiv 1 \pmod{n}$ if $\gcd(a, n) = 1$

Euclidean Algorithm

- $\text{gcd}(a, b) = \text{gcd}(b - a, a)$

```
4 + long long gcd(long long a, long long b) {  
5 +     if (a == 0) {  
6         return b;  
7     }  
8 +     else {  
9         return gcd(b % a, a);  
10    }  
11 }
```

Extended Euclidean Algorithm

- Find x and y such that $ax + by = \gcd(a, b)$
- Computes sequence of remainders and quotients:
 - $r_0 = a$
 - $r_1 = b$
 - ...
 - $r_{i+1} = r_{i-1} - q_i r_i$ and $0 \leq r_{i+1} < |r_i|$

Extended Euclidean Algorithm

- Also maintain two other sequences s_i and t_i

$$s_0 = 1 \qquad s_1 = 0$$

$$t_0 = 0 \qquad t_1 = 1$$

...

$$s_{i+1} = s_{i-1} - q_i s_i$$

$$t_{i+1} = t_{i-1} - q_i t_i$$

- We stop when $r_{k+1} = 0$
 $\gcd(a, b) = r_k = as_k + bt_k$

Extended Euclidean Algorithm

```
4 + pair<long long, long long> extendedEuclid(long long a, long long b){
5     long long s = 0, old_s = 1;
6     long long t = 1, old_t = 0;
7     long long r = b, old_r = a;
8 +     while (r != 0){
9         long long quotient = old_r / r, temp;
10        // (old_r, r) = (r, old_r - quotient * r)
11        temp = old_r;
12        old_r = r;
13        r = temp - quotient * r;
14        // (old_s, s) = (s, old_s - quotient * s)
15        temp = old_s;
16        old_s = s;
17        s = temp - quotient * s;
18        // (old_t, t) = (t, old_t - quotient * t)
19        temp = old_t;
20        old_t = t;
21        t = temp - quotient * t;
22    }
23    return make_pair(old_s, old_t);
24 }
```

Multiplicative Inverse

- $\frac{a}{b} = a * ? \pmod{p}$

- Fermat Little Theorem
 $? = b^{p-2}$

- Extended Euclid

Find x and y such that $xp + yb = 1$

$? = y$

Linear Diophantine Equation

- Find x and y satisfying $ax + by = c$
- Solution exists iff c is divisible by $\gcd(a, b)$
- If (x_0, y_0) is a solution, then
$$\left(x_0 + k \frac{y}{\gcd(x, y)}, y_0 - k \frac{x}{\gcd(x, y)}\right)$$
 is

Fast Exponentiation

$$x^{yz} = (x^y)^z$$

$$x^{2y} = (x^y)^2$$

Matrix Multiplication

$M \times N$

- $M \rightarrow a \times b$
- $N \rightarrow b \times c$
- $MN \rightarrow a \times c$

$$MN(i,j) = \sum_{k=1}^b M_{ik} N_{kj}$$

$MN \neq NM$!!!!

Matrix Exponentiation

Misal $M_i \times N = M_{i+1}$

- $M_2 = M_1 \times N$
- $M_3 = M_2 \times N = M_1 \times N \times N = M_1 \times N^2$
- $M_4 = M_3 \times N = M_1 \times N^2 \times N = M_1 \times N^3$

Fast Exponentiation $^{\wedge} \cdot ^{\wedge} 9$

$M_i \rightarrow a \times b$ $N \rightarrow b \times b$
 $N \rightarrow b \times b$ $M_i \rightarrow b \times a$

$M \times I = M$

$I = \text{Matriks identitas} = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$

Problem: Benefit (UVa11889)

- Given a and c , find the smallest positive integer b such that $\text{lcm}(a, b) = c$

Problem: Permasalahan DP (P1 2014)

- If $f(a, b)$ is known, then so is $f(a \times k, b \times k)$
- How many values of f should be calculated to find all $f(a, b)$, $1 \leq a, b \leq n$?

Problem: Fungsi Haha

- Diketahui $\text{haha}(0) = x$ dan $\text{haha}(1) = y$
- Diketahui $\text{haha}(i) = \text{haha}(i-1) - \text{haha}(i-2)$
- Cari nilai $t(n)$, $0 \leq n \leq 10^9$