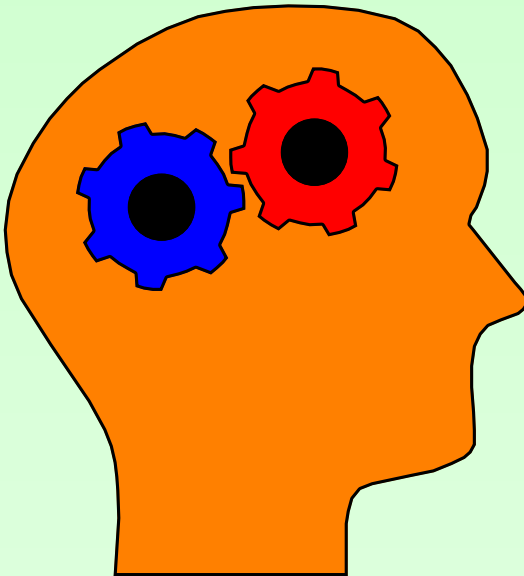




CS2104: Programming Languages Concepts

Lecture 6 : **Towards Monads**



*“Imperative Programming
in a Pure Language”*

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Can be challenging but You are Not Alone

- The midnight Monad, a journey to enlightenment.

<https://www.lambdacat.com/the-midnight-monad-a-journey-to-enlightenment/>

- Functors, Applicatives and Monads in Picture form:

http://adit.io/posts/2013-04-17-functors,_applicatives,_and_monads_in_pictures.html

- A Fistful of Monads.

<http://learnyouahaskell.com/a-fistful-of-monads>

Referential Transparency vs Opacity

Referential transparency and referential opacity are properties of parts of computer programs. An expression is called referentially transparent if it can be replaced with corresponding value without changing the program's behavior..

```
let v = e  
in ..v..v..
```



```
..v..v..
```

Pure vs Impure Code

- Imperative Programming (with side effects) - Opaque

```
print :: String -> ()
```

```
let c = print("hello")  
in c ; c
```

```
print("hello");  
print("hello")
```

different



- Pure Monadic Programming - Transparent

```
print :: String -> IO ()
```

```
let c = print("hello")  
in c >> c
```

```
print("hello") >>  
print("hello")
```

equivalent



Pictorial Introduction for Functors, Applicatives and Monads

http://adit.io/posts/2013-04-17-functors,_applicatives,_and_monads_in_pictures.html

Pure Value World

Here's a simple value:

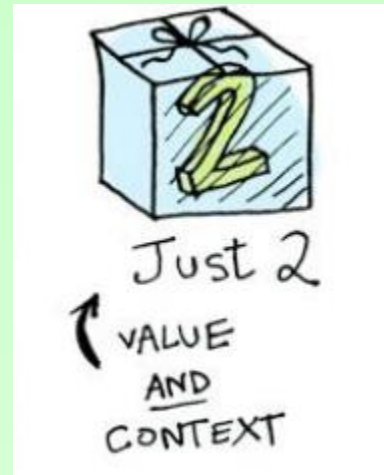


And we know how to apply a function to this value:

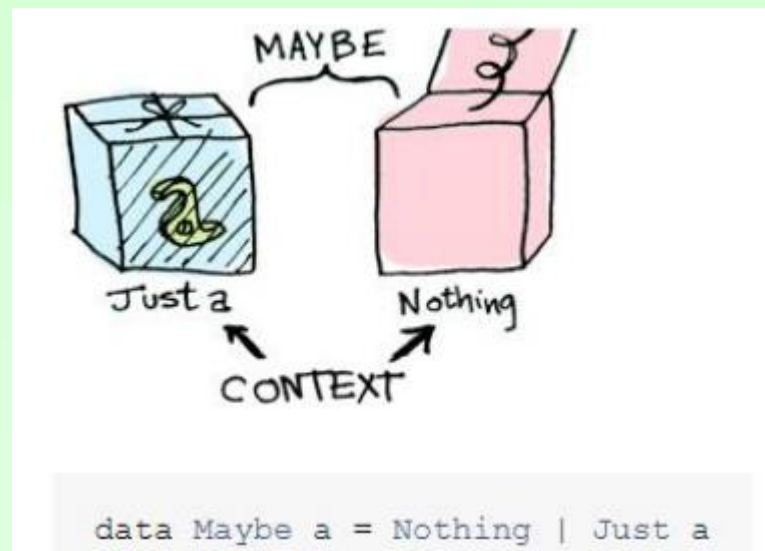


Value within Some Context

- Value and a Context.



- Maybe Type where
where
Nothing
denotes error

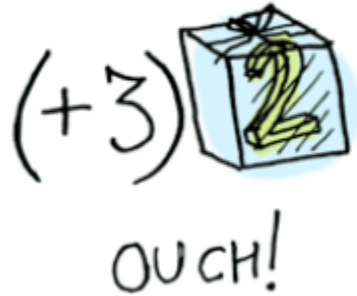


Other Examples of Context

- `[a]`
 - for non-determinism
- `state -> (state, a)`
 - for imperative state that can be updated
- `Parser a = String -> [(a, String)]`
 - For non-deterministic parsing
- `IO a`
 - for input-output interaction

Why Functor?

When a value is wrapped in a context, you can't apply a normal function to it:



- Solution : Functor.

```
> fmap (+3) (Just 2)  
Just 5
```



What is a Functor?

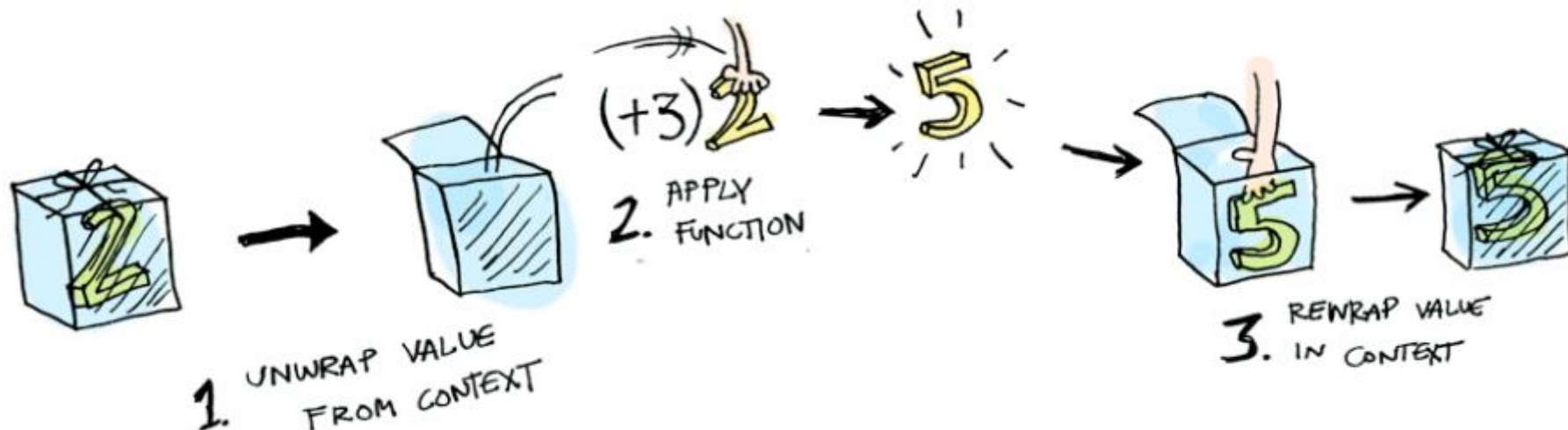
Functor is a typeclass. Here's the definition:

1. TO MAKE A DATA TYPE f
A FUNCTOR,

class Functor f where
 \rightarrow $fmap :: (a \rightarrow b) \rightarrow fa \rightarrow fb$

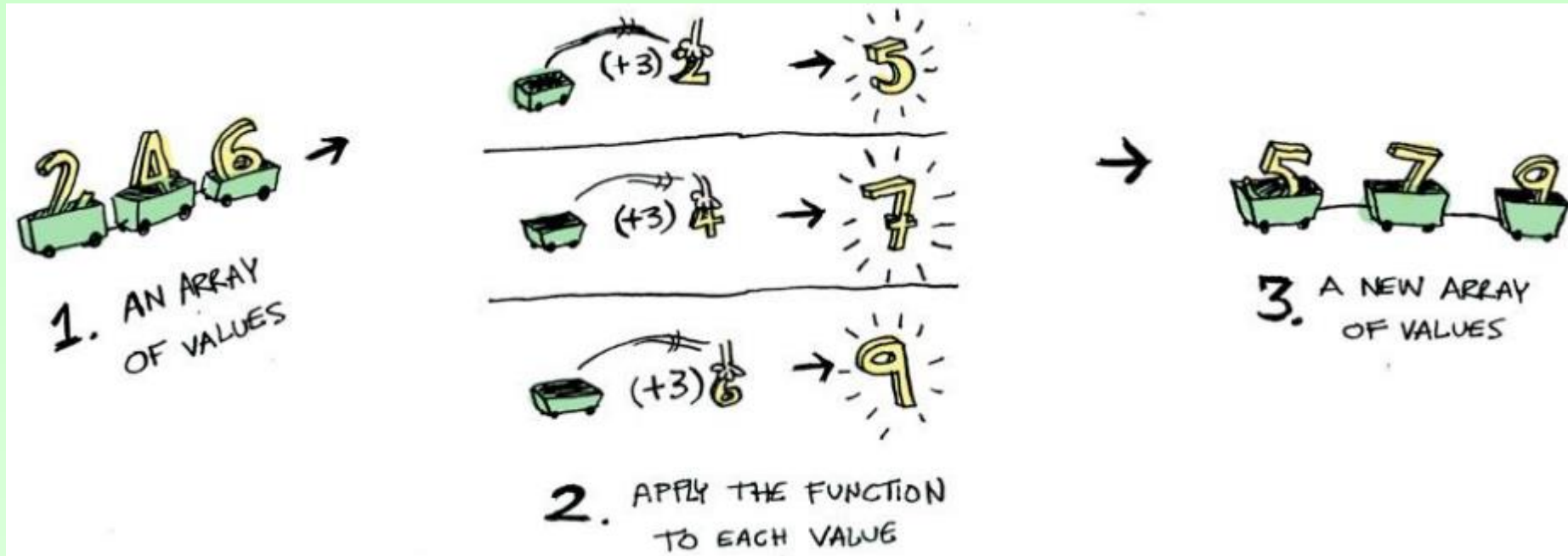
2. THAT DATA TYPE
NEEDS TO DEFINE
HOW $fmap$ WILL
WORK WITH IT.

Behind the Scene



```
instance Functor Maybe where
  fmap func (Just val) = Just (func val)
  fmap func Nothing    = Nothing
```

List/Arrays are also Functors



`fmap (+3) [2, 4, 6] → [5, 7, 9]`

`(+3) <$> [2, 4, 6] → [5, 7, 9]`

↑
infix variant

List as Functors

```
instance Functor [] where  
  fmap = map
```

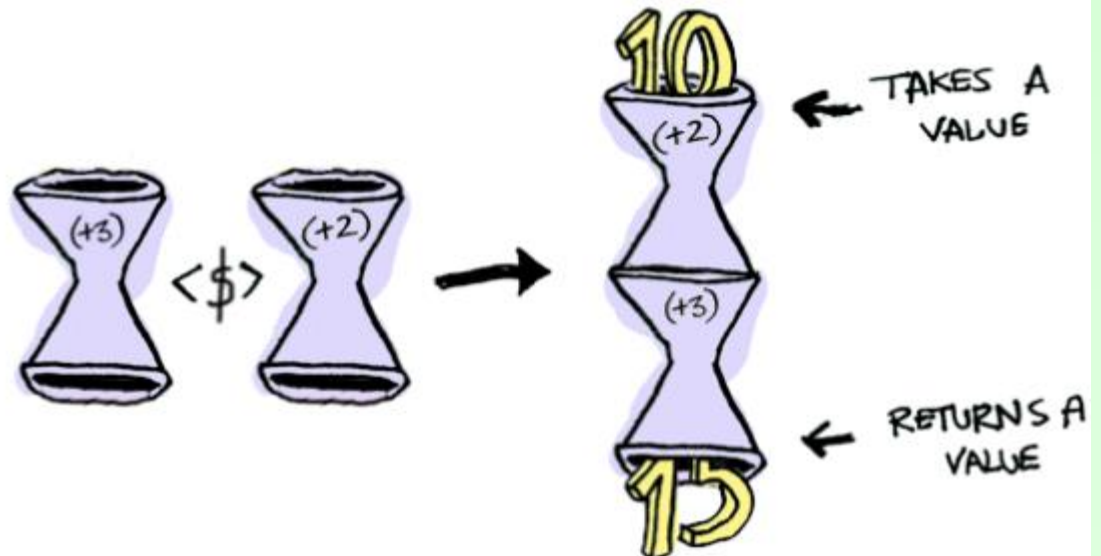
- List denotes non-determinism
- Examples:
 - `[]` means no solution
 - `[r1, r2, r3]` means three possible solutions

Functions are also Functors

Here's a function:



Here's a function applied to another function:



Functions as Functors ..

```
> let foo = fmap (+3) (+2)
> foo 10
15
```

- Implementation

```
instance Functor ((->) r) where
    fmap f g = f . g
```

What IF Functions are Wrapped in Context?



Maybe (Int -> Int)



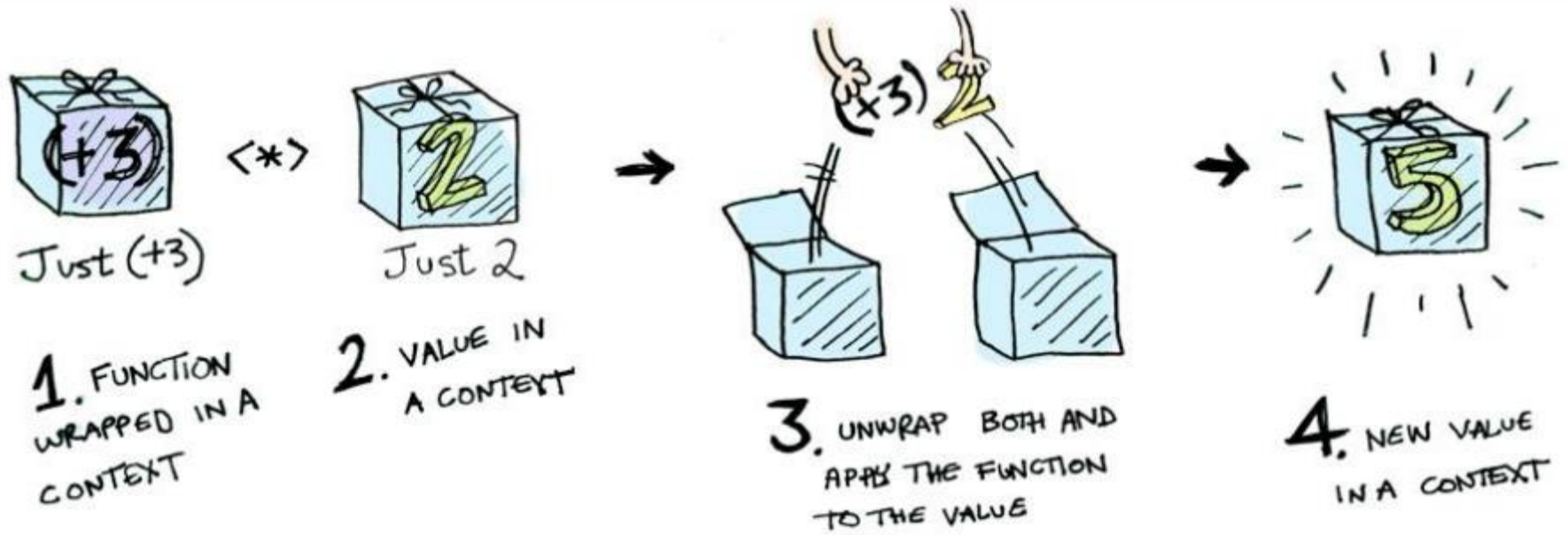
Maybe (Int)

Cannot use `fmap`

$$\text{fmap} :: (a \rightarrow b) \rightarrow f_a \rightarrow f_b$$

↑ ↑ ↑

Applicative to the Rescue

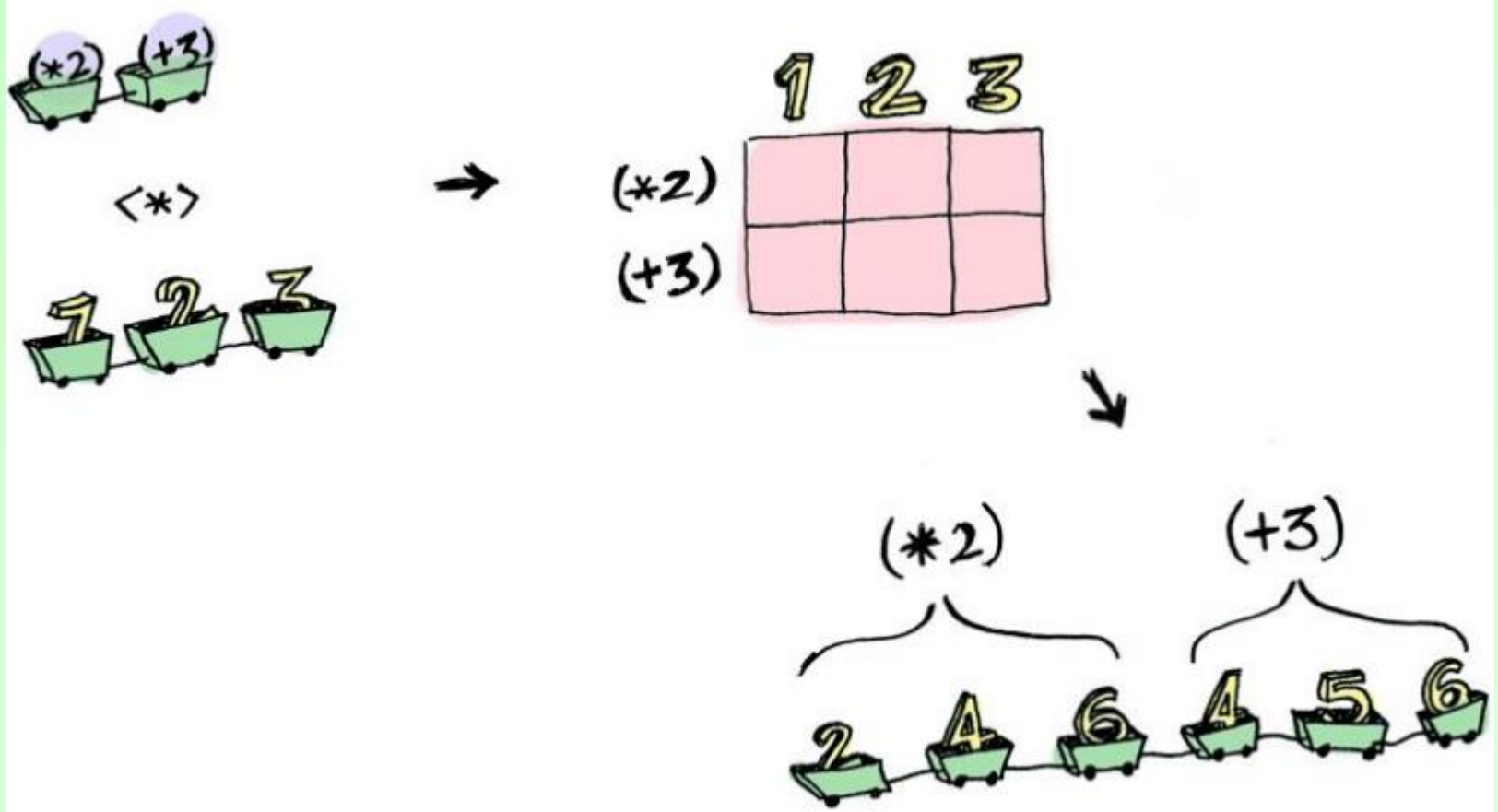


```
Just (+3) <*> Just 2 == Just 5
```

Applicative Class

```
class Functor f => Applicative (f :: * -> *) where  
    (<*>) :: f (a -> b) -> f a -> f b
```

Applicative in List Context



```
> [ (*2), (+3) ] < * > [ 1, 2, 3 ]  
[ 2, 4, 6, 4, 5, 6 ]
```

Why do we Need Applicative?

- Applicative can work with functions of any no. of arguments
 - Use fmap first

```
> let f = fmap (+) [1,2,3]
➤ :t f
➤ f :: Num a => [ (a -> a) ]
```

- Use Applicative now

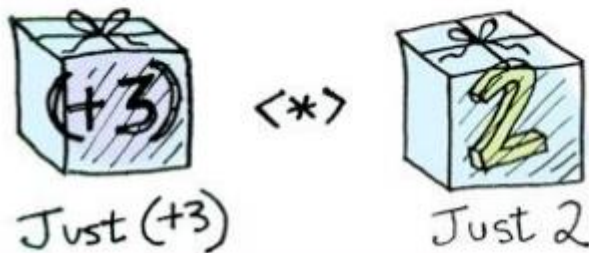
```
> f <*> [4,5]
➤ => [5,6,6,7,7,8]
```

Recap

Functors apply a function to a wrapped value:



Applicatives apply a wrapped function to a wrapped value:

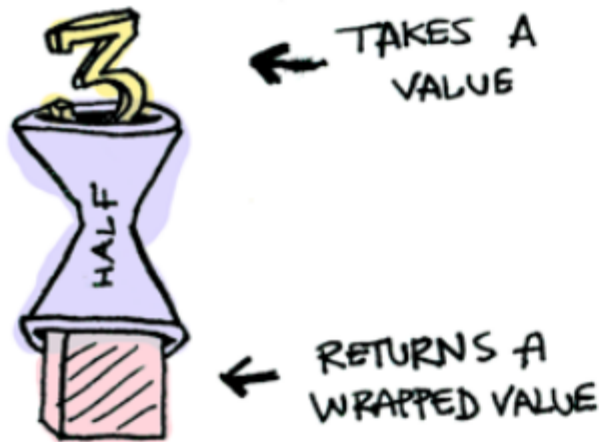


Essence of Monads

- How do we supply a wrapped value $(M\ a)$ to a function which returns a wrapped value $(a \rightarrow M\ b)$
 $half :: Int \rightarrow Maybe\ Int$

Suppose `half` is a function that only works on even numbers:

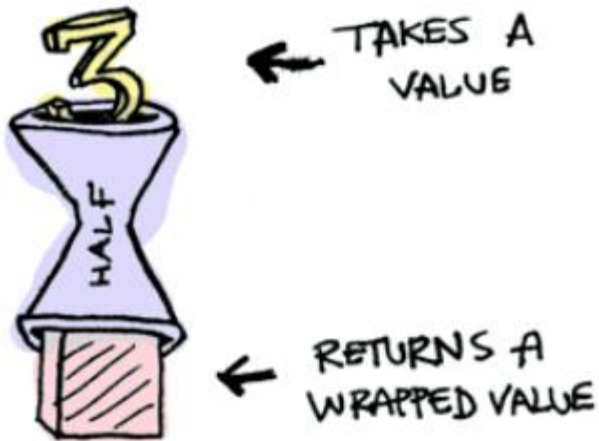
```
half x = if even x
         then Just (x `div` 2)
         else Nothing
```



What if we Apply on a Wrapped Value?

Suppose `half` is a function that only works on even numbers:

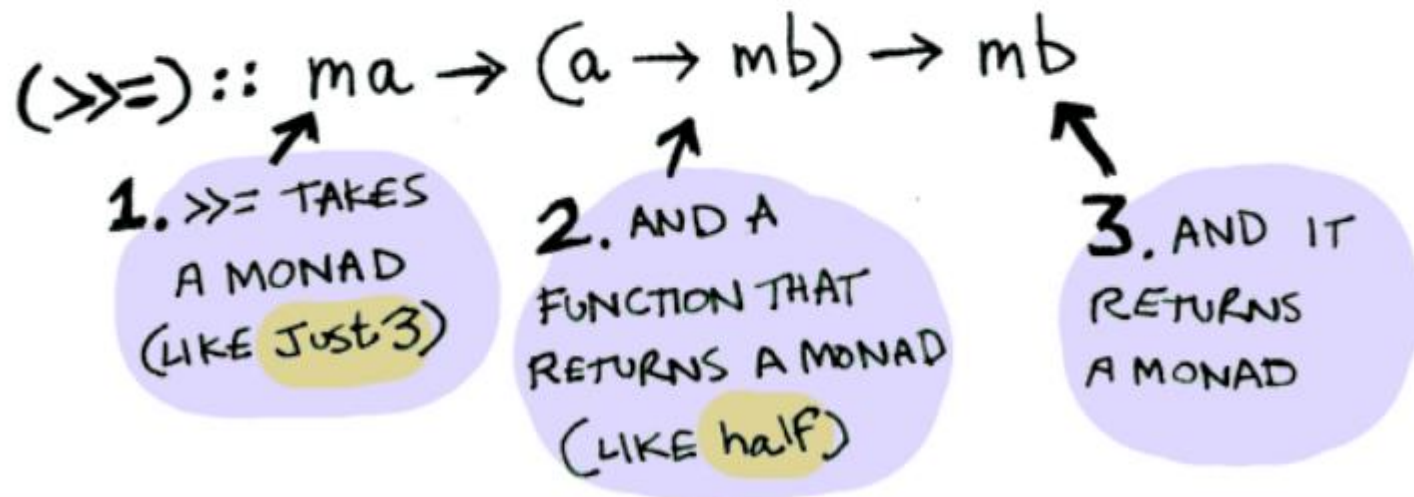
```
half x = if even x
         then Just (x `div` 2)
         else Nothing
```



Monad as a Type Class

```
class Monad m where
  (>>=) :: m a -> (a -> m b) -> m b
```

Where `>>=` is:



Chaining via Monads

```
> Just 20 >>= half >>= half >>= half  
Nothing
```

```
instance Monad Maybe where  
  Nothing >>= func = Nothing  
  Just val >>= func = func val
```



Input-Output as a Monad



getLine takes no arguments
and gets user input.



IO Monad Operation

`readFile` takes a string (a filename) and returns that file's contents



IO Monad Operation

`putStrLn` takes a string and prints it:



```
putStrLn :: String -> IO ()
```

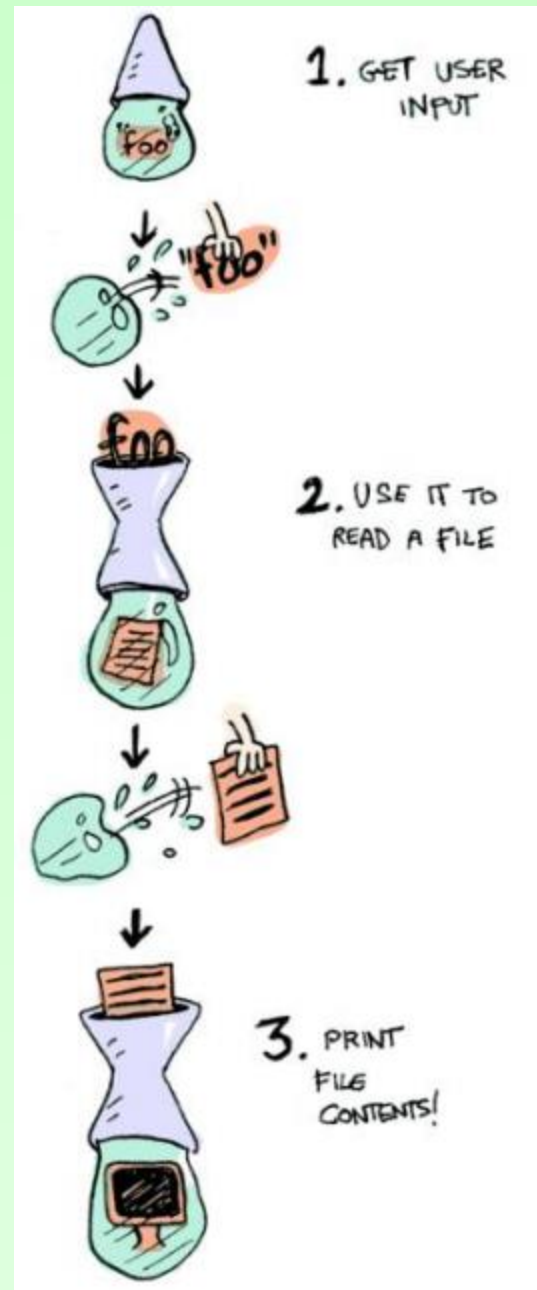
Chaining IO Operation

Chaining

```
getLine >>= readFile >>= putStrLn
```

Syntactic Sugar

```
foo = do  
  filename <- getLine  
  contents <- readFile filename  
  putStrLn contents
```

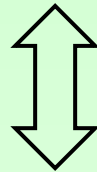


Do Comprehension

- Syntactic sugar notation for Monads.
- List is an instance of monad, and List comprehension is an instance of Do-comprehension!

```
[(x,y) | x <- xs, test x, y<-ys]
```

```
do  
  x <- xs  
  filter test  
  y <- ys  
  return (a,b)
```



```
filter test = \ x ->  
  if test x then return a else empty
```

Monads

Monads Formally

- Another example of higher-order type class is:

```
class Monad m where
  >>= :: (m a) -> (a -> m b) -> m b
  return :: a -> m a

  >> :: (m a) -> (m b) -> m b
  m1 >> m2    = m1 >>= (\ _ -> m2)
```

- Laws of `Monad` class:

```
(return a) >>= k    = k a
m >>= return        = m
(m >>= (\a -> (k a) >>= (\b -> h b)))
    = (m >>= (\a -> k a) >>= (\b -> h b))
```

- IO is an instance of `Monad` ...

Input/Output

Input/Output

- The I/O system in Haskell is purely functional but has all the expressive power of conventional imperative languages.
- Actions are *defined* rather than *invoked* in an expression-oriented style.
- These actions are modelled as *monads* of type `IO t` which is a conceptual structure with some properties that supports imperative actions.

Basic I/O Operations

- Every I/O action returns a value, e.g :

`getChar :: IO Char`

- Some IO actions also take input(s)

`putChar :: Char -> IO ()`

- `IO` is an instance of the `Monad` class.

- Actions are sequenced by `bind` operator:

`(>>=) :: Monad m => m a -> (a -> m b) -> m b`

`(>>) :: Monad m => m a -> m b -> m b`

Basic I/O Operations

- The `do` statement captures a sequence of actions, e.g.:

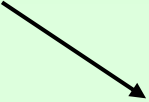
```
main    :: IO ()  
main    = do c <- getChar  
          putChar c
```

- Syntactic sugar for the following:

```
main    = getChar >>=  
          (\ c -> putChar c)
```

- How to return a value from sequence of actions?.

```
ready   = IO Bool  
ready   = do c <- getChar  
            c == 'y'
```



```
return (c == 'y')
```

Bigger I/O Operations

- Function to get a string of char may use recursion, as follows:

```
getLine  :: IO String
getLine  = do c <- getChar
            if c=='\n' then return ""
            else do l <- getLine
                    return (c:l)
```

- A pure value can be converted into an action by return, but the not the converse. Illegal to use:

```
x + print y
```

- Function `f :: Int -> Int -> Int` cannot do any IO at all, unless we make use of unsafe operations.

Building Actions

- IO operations are ordinary Haskell values that can be passed to functions, placed into data structures and returned as results etc.
- Example : we can build a list of actions.

```
todoList :: [IO ()]
todoList = [ putChar 'a' ,
              do {putChar 'b' ; putChar 'c'} ,
              do {x <- getChar; putChar x} ]
```

- Can combine them into a single action using:

```
sequence_ :: [IO ()] -> IO ()
sequence_ = foldr (>>) (return ())
```

Imperative Programming

- I/O programming in Haskell is very close to that being done for ordinary imperative programming.
- As a comparison, imperative `getLine` is simply:

```
function getLine() {  
    c := getChar();  
    if c=='\n' then return ""  
    else {l:=getLine();  
          return c:l} }
```

- Main difference is that no special semantics is needed and the entire code is still purely functional. Monad cleanly separates the pure from imperative.

Recap / Comparison

- Imperative `getLine` in C:

```
function getLine() {  
    c := getChar();  
    if c=='\n' then return ""  
    else {l:=getLine();  
         return c:l} }
```

- Monadic IO in Haskell

```
getLine  :: IO String  
getLine  = do c <- getChar  
             if c=='\n' then return ""  
             else do l <- getLine  
                   return (c:l)
```