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Module: IT3031 - Database Systems and Data Driven Applications

Practical: 07 (Indexes)

--Check what are the tables already you have

SELECT table_name FROM user_tables;

--Tables created in SQL Plus.

--Client table

CLN NAME	ADDRESS
	3 East Av, Bentley, WA 6102 6 West Av, Westray, NY 6103

--Stock Table

COMPANY	PRICE	DIVIDEND	EPS
BHP IBM INTEL FORD GM INFOSYS	10.5 70 76.5 40 60 45	1.5 4.25 5 2 2.5	3.2 10 12.4 8.5 9.2 7.8
6 rows sel	ected.		

--Trading Table

COMPANY	EXCHANGE
ВНР	Sydney
ВНР	New York
IBM	New York
IBM	London
IBM	Tokyo
INTEL	New York
INTEL	London
FORD	New York
GM	New York
INFOSYS	New York
10 rows	selected.

--Purchase Table

CLN	COMPANY	PDATE	QTY	PRICE	
c01	BHP	02-OCT-01	1000	12	
c01	BHP	08-JUN-02	2000	10.5	
c01	IBM	12-FEB-00	500	58	
c01	IBM	10-APR-01	1200	65	
c01	INFOSYS	11-AUG-01	1000	64	
c02	INTEL	30-JAN-00	300	35	
c02	INTEL	30-JAN-01	400	54	
c02	INTEL	02-0CT-01	200	60	
c02	FORD	05-0CT-99	300	40	
c02	GM	12-DEC-00	500	55.5	
10 r	rows sele	ected.			

--Plan Table created

```
SQL> create table PLAN_TABLE (
2 statement_id varchar2(30),
3 timestamp date,
4 remarks varchar2(30),
5 operation varchar2(20),
6 options varchar2(255),
7 object_node varchar2(128),
8 object_owner varchar2(30),
9 object_name varchar2(30),
10 object_tinstance numeric,
11 object_type varchar2(255),
12 optimizer varchar2(255),
13 search_columns number,
14 id numeric,
15 parent_id numeric,
16 position numeric,
17 cost numeric,
18 cardinality numeric,
19 bytes numeric,
20 other_tag varchar2(255),
21 partition_start varchar2(255),
22 partition_start varchar2(255),
23 partition_id numeric,
24 other long,
25 distribution varchar2(30),
26 cpu_cost numeric,
27 io_cost numeric,
28 temp_space numeric);

Table created.
```

Display the client number and name of clients who have made large purchases. A large purchase occurs when a client purchases more than 1000 shares from a company at any given time.

SELECT DISTINCT CLIENT.CLNO, CLIENT.NAME

FROM CLIENT

INNER JOIN PURCH ON CLIENT.CLNO = PURCH.CLNO

WHERE PURCH.QTY > 1000;

OUTPUT:

```
CLN NAME
--- -----
c01 John Smith
```

- --Session altered commands
- --Try to get all the rows in the most efficient way, even if the first few rows take a little longer.

```
SQL> ALTER SESSION SET OPTIMIZER_MODE = ALL_ROWS;
Session altered.
```

When choosing how to run queries, give more importance to CPU cost, not just how much disk I/O is used.

```
SQL> ALTER SESSION SET "_optimizer_cost_model"=CPU; Session altered.
```

(a) Use EXPLAIN PLAN to find Oracle's generated query plan for the above query by executing the following statement.

```
SQL> explain plan for
2 select c.clno, c.name
3 from client c, purch p
4 where c.clno = p.clno and p.qty > 1000;
Explained.
```

(b) You can execute the script utlxpls.sql to view the query plan and associated costs. Describe each step of the query plan.

Operation and options	 Obiect	 cost	cpu_cost	io cost
SELECT STATEMENT		3	18063	3
HASH JOIN		3	18063	3
NESTED LOOPS		3	18063	3
NESTED LOOPS		3	18063	3
STATISTICS COLLECTOR	j	i II I		
TABLE ACCESS FULL	PURCH	2	9721	2
INDEX UNIQUE SCAN	SYS_C0082	0	1050	0
TABLE ACCESS BY INDEX ROWID	CLIENT	1	8341	1
TABLE ACCESS FULL	CLIENT	1	8341	1

This query plan shows how Oracle will execute the SQL query, step by step, and provides insights into the operations, costs, and access methods used. Here's a breakdown of the key components in this plan:

1. SELECT STATEMENT

- Operation and options: The main operation for the query, indicating that it's a SELECT query.
- Cost: 3. The overall cost of executing this query.
- CPU Cost: 18063. The estimated CPU resources needed for this query.
- I/O Cost: 3. The estimated I/O resources needed for this query.

2. HASH JOIN

- This is the method used for joining two tables. A Hash Join is used when no indexes are available, or when the optimizer believes it's the most efficient way to join large sets of data.
- The Hash Join indicates that the database will perform a join between the two tables (client and purch) based on the common clno column.

3. Nested Loops (Repeated Twice)

- The Nested Loop Join is used when one of the tables is small and can be scanned for every row in the larger table.
- This operation is nested, meaning that the database will loop through the records of one table and, for each record, it will perform a lookup on the other table.
- This is commonly used for equi-joins (like c.clno = p.clno), where rows from both tables are compared.

4. Statistics Collector

• This step collects execution statistics for the query (e.g., how many rows were processed).

5. TABLE ACCESS FULL for PURCH

- Object: PURCH
- Cost: 2. The cost of a full table scan on the purch table.
- I/O Cost: 9721. The I/O cost of accessing this table.
- This indicates that Oracle is performing a full table scan on the purch table to find all matching records.

6. INDEX UNIQUE SCAN on SYS C0082

- Object: SYS_C0082 (likely an index on client.clno).
- Cost: 0. The cost for the index scan is negligible.
- CPU Cost: 1050. The CPU cost associated with scanning the index.
- This means Oracle is using an index to efficiently access the client table based on the clno column.

7. TABLE ACCESS BY INDEX ROWID for CLIENT

- Object: CLIENT
- Cost: 1. The cost of retrieving data from the client table using the index rowid.
- I/O Cost: 8341. The I/O cost for accessing the client table.
- This indicates that Oracle is using an index to fetch rows from the client table by accessing the rows identified by the index.

8. TABLE ACCESS FULL for CLIENT

- Object: CLIENT
- Cost: 1. The cost of performing a full table scan on the client table.
- I/O Cost: 8341. The I/O cost for scanning the client table.

(C). Next, create indexes for client and purchase tables. In Oracle you can use the CREATE INDEX statement to create indexes.

```
SQL> CREATE INDEX index_1 ON client(CLNO, NAME);
Index created.
```

```
SQL> CREATE INDEX index_2 ON purch(QTY, CLNO);
Index created.
```

--Re-executed plan table for the same query

```
SQL> explain plan for
2 select c.clno, c.name
3 from client c, purch p
4 where c.clno = p.clno and p.qty > 1000;
Explained.
```

(c) Re-execute the explain plan in part (a). Use utlxpls.sql to view the query plan and associated costs. Explain the query plan and compare it with the query plan in part (b.)

```
Plan Table
                                                           cpu_cost | io_cost
 Operation and options
                                       Object
                                               cost
  SELECT STATEMENT
                                                           14843
                                                 2
                                                                        2
                                                 2
   HASH JOIN
                                                           14843
                                                                        2
                                                           14843
    NESTED LOOPS
                                                 | | |
     STATISTICS COLLECTOR
      INDEX RANGE SCAN
                                     INDEX_2
                                                 1
                                                           7521
                                                                        1
     INDEX RANGE SCAN
                                     INDEX_1
                                                 1
                                                           7321
                                                                        1
    INDEX FULL SCAN
                                    INDEX_1
                                                | 1
                                                           7321
                                                                       1
10 rows selected.
```

Explanation of the query plan

SELECT STATEMENT

- This is the root operation for the SQL query execution.
- Cost: 2
- CPU Cost: 14843
- I/O Cost: 2

HASH JOIN

- The query uses a hash join to combine data from the involved tables.
- Cost: 2
- CPU Cost: 14843
- I/O Cost: 2

Nested Loops

• The join method is nested loops, indicating a row-by-row comparison between tables.

STATISTICS COLLECTOR

• Collecting statistics for the query execution, such as the number of rows processed.

INDEX RANGE SCAN on INDEX 2

- An index range scan is used on INDEX_2, which is likely a range-based lookup on one of the columns.
- Cost: 1
- CPU Cost: 7521
- I/O Cost: 1

INDEX RANGE SCAN on INDEX_1

- Another index range scan is performed on INDEX_1, similar to the previous step.
- Cost: 1
- CPU Cost: 7321
- I/O Cost: 1

INDEX FULL SCAN on INDEX_1

- Index Full Scan indicates that the query is scanning all the rows in INDEX_1 without filtering by a range.
- Cost: 1
- CPU Cost: 7321
- I/O Cost: 1

Comparison of two plans and their features

Hash Join Method:

- Plan 1: Uses a Hash Join for combining tables.
- Plan 2: Also uses a Hash Join, but there's a slight difference in cost and underlying operations.

Nested Loops:

- Plan 1: Uses Nested Loops for the join operation.
- Plan 2: Similarly uses Nested Loops, but the operations involve different indexes.

Statistics Collection:

- Plan 1: The statistics collection operation occurs at a deeper level within the nested loops.
- Plan 2: Statistics collection happens within the nested loops, indicating a similar approach.

Table Access for PURCH:

- Plan 1: Accesses the PURCH table using Full Table Scan.
- Plan 2: Plan 2 doesn't reference the PURCH table at all, suggesting a different underlying data structure or query design.

Table Access for CLIENT:

- Plan 1: Uses Full Table Scan for the CLIENT table.
- Plan 2: Uses Index Scans (both range scan and full scan) for the CLIENT table.

Index Scan Method:

- Plan 1: Uses Index Unique Scan on SYS_C0082.
- Plan 2: Uses Index Range Scan on INDEX_1 and INDEX_2, showcasing a different indexing strategy.

Cost Estimates:

- Plan 1: Has a higher CPU Cost and I/O Cost overall, reflecting a more expensive execution.
- Plan 2: Shows a slightly lower Cost and a more balanced distribution between CPU and I/O Cost.

Object Access:

- Plan 1: Accesses CLIENT using both Index RowID and Full Table Scan.
- Plan 2: Accesses CLIENT using Index Range Scan and Index Full Scan, optimizing the index access.

In summary, Plan 2 is likely the better plan because:

- It uses index range scans and index full scans, which are typically more efficient than full table scans.
- It relies on index-based access for the CLIENT table, making the query more optimized for data retrieval, especially if the CLIENT table is large.
- Plan 1, on the other hand, relies on full table scans for both PURCH and CLIENT tables, which can be slower, particularly for large tables.