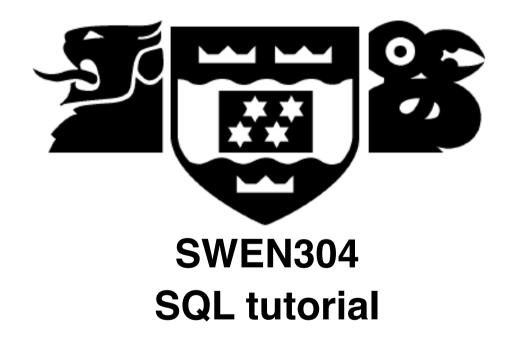
# Database System Engineering



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# Structured Query Language (SQL)

- We are concerned with three major problems
  - How to define a database schema and its relation schemas?
  - How to store and manipulate data in relations?
  - How to retrieve data from a database?
- SQL is used as
  - a Data Definition Language (DDL)
  - a Data Manipulation Language (DML)
  - a Query Language (QL)

## SQL as DDL

- Catalog
- Domains
- Tables (Relation)
- Constraints
- Modify schemas (DROP and ALTER)

# Catalog

- The catalog records essential information (metadata) about this database
  - Schema table: information on all relation schemas
  - Attribute table: information on all attributes
  - Data types table: information on all data types (domains)
  - Constraints table: information on all constraints

## **SQL** Definitions

CREATE SCHEMA UNIVERSITY AUTHORIZATION huima;

```
CREATE DOMAIN idno
```

AS INT

**DEFAULT 300001** 

**NOT NULL** 

CONSTRAINT idnoconstr

CHECK (VALUE > 300000 AND VALUE <= 399999);

### **UNIVERSITY Database**

UNIVERSITY ={STUDENT(<u>StudId</u>, Lname, Fname, Major),
COURSE(<u>CourId</u>, Cname, Points, Dept),
GRADES(<u>StudId</u>, <u>CourId</u>, Grade)}

| STUDENT |       |       |       |
|---------|-------|-------|-------|
| StudId  | Lname | Fname | Major |
| 300111  | Smith | Susan | COMP  |
| 300121  | Bond  | James | MATH  |
| 300143  | Bond  | Jenny | MATH  |
| 300132  | Smith | Susan | COMP  |

| Course  |          |        |             |
|---------|----------|--------|-------------|
| Courld  | Cname    | Points | Dept        |
| COMP302 | DB sys   | 15     | Engineering |
| COMP301 | softEng  | 20     | Engineering |
| COMP201 | Pr & Sys | 22     | Engineering |
| MATH214 | DisMat   | 15     | Mathematics |

| GRADES |         |       |  |
|--------|---------|-------|--|
| StudId | Courld  | Grade |  |
| 300111 | COMP302 | A+    |  |
| 300111 | COMP301 | Α     |  |
| 300111 | MATH214 | Α     |  |
| 300121 | COMP301 | В     |  |
| 300132 | COMP301 | С     |  |
| 300121 | COMP302 | B+    |  |
| 300143 | COMP201 | ω     |  |
| 300132 | COMP201 | ω     |  |
| 300132 | COMP302 | C+    |  |

### Relation Schema: COURSE

```
CREATE TABLE COURSE (
Courld CHAR(7) CONSTRAINT cspk PRIMARY KEY,
CName CHAR(15) NOT NULL,
Points INT NOT NULL

CONSTRAINT pointschk

CHECK (Points >= 0 AND Points <= 50),
Dept CHAR(25)
);
```

## With/Without Domain

```
CREATE TABLE STUDENT (
  StudId INT
    NOT NULL
    DEFAULT 30000
    CONSTRAINT stpk PRIMARY KEY
    CONSTRAINT StidRange CHECK
  (Studid BETWEEN 300000 AND 399999),
LName CHAR(15) NOT NULL,
  FName CHAR(15) NOT NULL,
Major CHAR(25) DEFAULT 'Comp'
CREATE TABLE STUDENT (
  Studld idno CONSTRAINT stpk PRIMARY KEY,
  LName CHAR(15) NOT NULL,
  FName CHAR(15) NOT NULL,
  Major CHAR(25) DEFAULT 'Comp'
);
```

# Second Key

Unique and Not Null

```
CREATE TABLE STAFF (
staff_id INT PRIMARY KEY,
ird_number CHAR(7) NOT NULL UNIQUE,
address VARCHAR(255)
);
```

### DROP

- CASCADE behavior means deleting the construct itself and all the other constructs related to it,
- RESTRICT behavior means that the construct will be deleted only if it is empty (schema), or not referenced by any other construct (like: table, attribute, view)
- e.g. DROP TABLE COURSE RESTRICT
   DROP TABLE STUDENTS CASCADE

### **ALTER**

ALTER TABLE STUDENT ADD NoOfPoints INT DEFAULT 320;

ALTER TABLE GRADES ALTER Grade SET DEFAULT 'C';

ALTER TABLE GRADES DROP CONSTRAINT gsri;

ALTER TABLE STUDENT DROP CONSTRAINT stpk CASCADE;

# SQL as DML

UNIVERSITY ={STUDENT(StudId, Lname, Fname, Major),

COURSE(Courld, Cname, Points, Dept),

GRADES(StudId, Courld, Grade)

 $IC = \{GRADES[Id] \subseteq STUDENT[Id],$  $GRADES[Course\_id] \subseteq COURSE[Course\_id]\}$ 

| STUDENT |       |       |       |
|---------|-------|-------|-------|
| StudId  | Lname | Fname | Major |
| 300111  | Smith | Susan | COMP  |
| 300121  | Bond  | James | MATH  |
| 300143  | Bond  | Jenny | MATH  |
| 300132  | Smith | Susan | COMP  |

| Course  |          |        |             |
|---------|----------|--------|-------------|
| Courld  | Cname    | Points | Dept        |
| COMP302 | DB sys   | 15     | Engineering |
| COMP301 | softEng  | 20     | Engineering |
| COMP201 | Pr & Sys | 22     | Engineering |
| MATH214 | DisMat   | 15     | Mathematics |

| GRADES |         |       |  |
|--------|---------|-------|--|
| StudId | Courld  | Grade |  |
| 300111 | COMP302 | A+    |  |
| 300111 | COMP301 | Α     |  |
| 300111 | MATH214 | Α     |  |
| 300121 | COMP301 | В     |  |
| 300132 | COMP301 | С     |  |
| 300121 | COMP302 | B+    |  |
| 300143 | COMP201 | ω     |  |
| 300132 | COMP201 | ω     |  |
| 300132 | COMP302 | C+    |  |

### SQL as DML

- Three commands used to modify the database:
  - INSERT, DELETE, and UPDATE
- Full insert

```
INSERT INTO STUDENT
VALUES (111111, 'Bole', 'Ann', Math);
```

Partial insert

```
INSERT INTO STUDENT (FName, LName, StudId)
VALUES('Ann', 'Bole', 111111);
```

### **UPDATE** and **DELETE**

```
UPDATE GRADES

SET Grade = 'A+'

WHERE Courld = 'C302';
```

DELETE FROM STUDENT WHERE FName = 'Susan';

```
DELETE FROM STUDENT
WHERE StudId IN
(SELECT s.StudId
FROM STUDENT s, GRADES g
WHERE s.StudId = g.StudId AND Courld = 'C302');
```

### DROP vs DELETE

- DELETE statement performs conditional based deletion, whereas DROP command deletes entire records in the table
- DELETE statement removes only the rows in the table and it preserves the table structure as same whereas DROP command removes all the data in the table and the table structure
- DELETE operation can be rolled back and it is not auto committed, while DROP operation cannot be rolled back in any way as it is an auto committed statement
- DROP is a DDL statement while DELETE is a DML statement

## SQL as QL

 SELECT is the basic SQL statement for retrieving data from a database

## SQL as QL

- SQL does not automatically eliminate duplicate tuples in query results
- Use the keyword DISTINCT in the SELECT clause
  - Only distinct tuples should remain in the result

## **The University Database**

Suppose for each student only pass grades are recorded in the database

#### **STUDENT**

| LName | FName | <u>StudId</u> | Major |
|-------|-------|---------------|-------|
| Smith | Susan | 131313        | Comp  |
| Bond  | James | 007007        | Math  |
| Smith | Susan | 555555        | Comp  |
| Cecil | John  | 010101        | Math  |

#### **COURSE**

| PName  | Courld | Points | Dept |
|--------|--------|--------|------|
| DB Sys | C302   | 15     | Comp |
| SofEng | C301   | 15     | Comp |
| DisMat | M214   | 22     | Math |
| Pr&Sys | C201   | 22     | Comp |

#### **GRADES**

| <u>StudId</u> | Courld | Grade |
|---------------|--------|-------|
| 007007        | C302   | A+    |
| 555555        | C302   | D     |
| 007007        | C301   | Α     |
| 007007        | M214   | A+    |
| 131313        | C201   | B-    |
| 555555        | C201   | С     |
| 131313        | C302   | D     |
| 007007        | C201   | Α     |
| 010101        | C201   | D     |

## **Single Table Queries**

Retrieve the first and last names of Comp students

```
SELECT FName, LName
FROM STUDENT
WHERE Major = 'Comp';
```

| FName | LNam<br>e |
|-------|-----------|
| Susan | Smith     |
| Susan | Smith     |

Find all different grades

```
SELECT DISTINCT Grade FROM GRADES;
```

| Grade      |
|------------|
| <b>A</b> + |
| Α          |
| B-         |
| С          |

## **Arithmetic Operations, Sorting**

 SQL provides capability to perform four basic arithmetic operations (+, -, \*, / ) that can be applied to numeric attributes and constants only

```
SELECT 2 + 2;
```

Sorting of the query result tuples is done using

```
ORDER BY { \( \langle attribute_name \rangle \) [ (ASC | DESC) ] , ... } clause after the WHERE clause (ASC is default)
```

```
SELECT *
FROM GRADES
ORDER BY StudId ASC, CourId DESC;
```

## **Qualification and Aliasing**

- Attributes in different relation schemas can have the same names. How do we prevent ambiguity?
- In the SELECT clause, we prefix attributes by table name: SELECT STUDENT.StudId ...
- In the FROM clause, specify a tuple variable from the table: ...FROM COURSE c, GRADES g, STUDENT s
- In the WHERE clause, prefix an attribute by the tuple variable: WHERE c.CourId = g.CourId

### Join

- To retrieve data from more than one table, we need a new operation: JOIN
- There are different joins:
  - INNER (theta join, equi-join, natural join)
  - OUTER (left, right, full)
  - Most often, we use the equi-join

## **A JOIN Example**

3

4

5

b2

b3

 $\omega$ 

b2

b3

4

b2

b3

c1

c1

If there is no join condition what is the result?

b3

b4

c1

c2

# **Explicit Join**

 Q1: Retrieve first name, course id and corresponding grades of the student with Student Id = 007007

```
SELECT FName, Courld, Grade
FROM (STUDENT NATURAL JOIN GRADES)
WHERE StudId = 007007;
```

The FROM clause contains a single joined table

# Types of Join

- Inner Joins:
  - JOIN, INNER JOIN, EQUI-JOIN, NATURAL JOIN,
- Outer Joins:
  - LEFT OUTER JOIN, RIGHT OUTER JOIN, FULL OUTER JOIN
  - The keyword OUTER may be omitted
  - CROSS JOIN is used to specify the CARTESIAN PRODUCT operation and should be used only with the utmost care

### **JOIN**

Using JOIN operator

SELECT \* FROM R1 JOIN R2 ON R1.B = R2.B;

**R1** 

| Α              | В              |
|----------------|----------------|
| a <sub>1</sub> | ω              |
| a <sub>2</sub> | b <sub>1</sub> |

R2

| С                     |
|-----------------------|
| <b>C</b> <sub>1</sub> |
| C <sub>1</sub>        |
| C <sub>1</sub>        |
|                       |

R1 JOIN R2

| Α              | В              | В              | С              |
|----------------|----------------|----------------|----------------|
| $\mathbf{a_2}$ | b <sub>1</sub> | b <sub>1</sub> | C <sub>1</sub> |

Using NATURAL JOIN operator

SELECT \* FROM R1 NATURAL JOIN R2;

| A              | В              | С                     |
|----------------|----------------|-----------------------|
| a <sub>2</sub> | b <sub>1</sub> | <b>c</b> <sub>1</sub> |

**SWEN304** 

Lect8: SQL(3)

### **LEFT OUTER JOIN**

- Every tuple in left table must appear in result
- If no matching tuple
  - Padded with NULL values for attributes of right table

SELECT \* FROM R1 LEFT JOIN R2 ON R1.B = R2.B;

| R1             |                |
|----------------|----------------|
| Α              | В              |
| a <sub>1</sub> | ω              |
| a <sub>2</sub> | b <sub>1</sub> |

| С              |
|----------------|
| C <sub>1</sub> |
| C <sub>1</sub> |
| C <sub>1</sub> |
|                |

R<sub>2</sub>

| KILEFT JOIN RZ |                |                |                       |
|----------------|----------------|----------------|-----------------------|
| Α              | В              | В              | С                     |
| a <sub>1</sub> | ω              | ω              | ω                     |
| a <sub>2</sub> | b <sub>1</sub> | b <sub>1</sub> | <b>C</b> <sub>1</sub> |
|                |                | _              | _                     |

### **RIGHT OUTER JOIN**

- Every tuple in right table must appear in result
- If no matching tuple
  - Padded with NULL values for attributes of right table

SELECT \* FROM R1 RIGHT JOIN R2 ON R1.B = R2.B;

**R1** 

| Α              | В              |
|----------------|----------------|
| a <sub>1</sub> | ω              |
| a <sub>2</sub> | b <sub>1</sub> |

**R2** 

| В                     | С                     |
|-----------------------|-----------------------|
| b <sub>1</sub>        | <b>c</b> <sub>1</sub> |
| b <sub>2</sub>        | <b>c</b> <sub>1</sub> |
| <b>b</b> <sub>3</sub> | C <sub>1</sub>        |

R1 RIGHT JOIN R2

| Α              | В              | В              | С                     |
|----------------|----------------|----------------|-----------------------|
| a <sub>2</sub> | b <sub>1</sub> | b <sub>1</sub> | C <sub>1</sub>        |
| ω              | ω              | b <sub>2</sub> | <b>c</b> <sub>1</sub> |
| ω              | ω              | b <sub>3</sub> | <b>c</b> <sub>1</sub> |

### **FULL OUTER JOIN**

All tuples from both relations must appear in result

SELECT \* FROM R1 FULL JOIN R2 ON R1.B = R2.B;

| ı |   | 1 | 1 |
|---|---|---|---|
|   | 7 |   | ı |

| Α              | В              |
|----------------|----------------|
| a <sub>1</sub> | ω              |
| a <sub>2</sub> | b <sub>1</sub> |

R2

| В              | С                     |
|----------------|-----------------------|
| b <sub>1</sub> | <b>c</b> <sub>1</sub> |
| b <sub>2</sub> | <b>c</b> <sub>1</sub> |
| b <sub>3</sub> | <b>c</b> <sub>1</sub> |

R1 RIGHT JOIN R2

| Α              | В              | В              | С              |
|----------------|----------------|----------------|----------------|
| a <sub>1</sub> | ω              | ω              | ω              |
| a <sub>2</sub> | b <sub>1</sub> | b <sub>1</sub> | C <sub>1</sub> |
| ω              | ω              | b <sub>2</sub> | C <sub>1</sub> |
| ω              | ω              | b <sub>3</sub> | C <sub>1</sub> |

### **Nested Query**

Retrieve first names of students that passed M214

```
SELECT FName

FROM STUDENT s

WHERE s.StudId IN

(SELECT StudId FROM GRADES

WHERE CourId = 'M214' AND Grade IS NOT NULL);
```

 A nested query defined by using IN (or =ANY) operator can be expressed as a single block query

```
SELECT FName

FROM STUDENTs, GRADES g

WHERE s.StudId = g.StudId AND g.CourId = 'M214'

AND g.Grade IS NOT NULL;
```

### **EXISTS and NOT EXISTS**

 Retrieve Id's and surnames of students who passed at least one course:

```
SELECT s.StudId, s.LName FROM STUDENTs
WHERE EXISTS

(SELECT * FROM GRADES
WHERE s.StudId = StudId AND Grade IS NOT NULL);
```

Retrieve Id's and surnames of students who didn't pass any course:

```
SELECT s.StudId, s.LName FROM STUDENTs

WHERE NOT EXISTS

(SELECT * FROM GRADES

WHERE s.StudId = StudId AND Grade IS NOT NULL);
```

## **Aggregate Functions**

Q2: What is the total point value of the courses in the COMP dept?

```
SELECT SUM (Points)
FROM COURSE
WHERE Dept = 'Comp';
```

```
sum_Points
52
```

Q3: How many different first names do COMP majors have?

```
SELECT COUNT(DISTINCT FName) AS NoOfNames
FROM STUDENT
WHERE Major = 'Comp';
```

NoOfNames 1

### **Grouping**

For each student, retrieve the number of courses passed

```
SELECT StudId, COUNT(*)
FROM GRADES
WHERE Grade IS NOT NULL
GROUP BY StudId;
```

| StudId | COUNT(*) |  |
|--------|----------|--|
| 007007 | 4        |  |
| 131313 | 1        |  |
| 555555 | 1        |  |

### **HAVING Clause**

Retrieve the number of courses passed for students that passed at least two courses

```
SELECT StudId, COUNT(*)

FROM STUDENTs NATURAL JOIN GRADES g

WHERE Grades IS NOT NULL

GROUP BY s.StudId

HAVING COUNT(*) > 1;
```

| StudId | COUNT(*) |
|--------|----------|
| 007007 | 4        |

### **SQL Views**

Create a view

```
CREATE VIEW StudOccupied AS

SELECT g.StudId, SUM(Hours) AS Occupied

FROM Grades g, Course p

WHERE g.CourId = p.CourId AND Grade IS NULL

GROUP BY StudId;
```

Deleting a view

DROP VIEW StudOccupied;