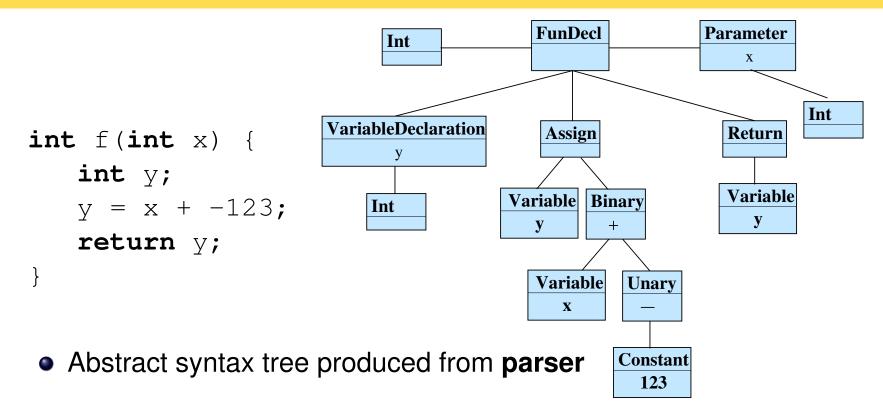
SWEN430 - Compiler Engineering

Lecture 12 - Bytecode Generation

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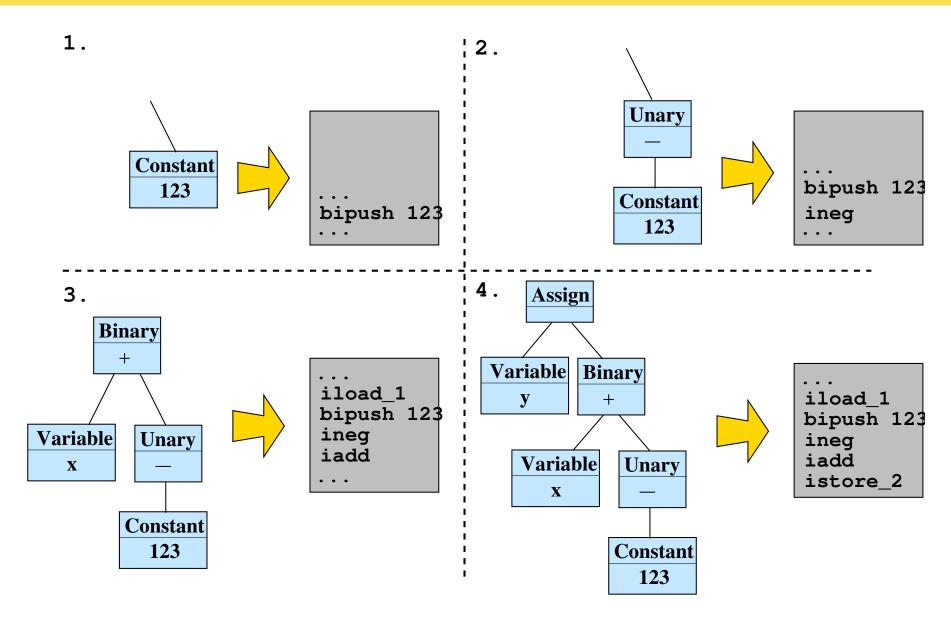
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Abstract Syntax Trees (AST)



- Abstract syntax tree is programmatic representation of source
- Abstract syntax tree used for e.g. type checking
- Abstract syntax tree turned into intermediate language or target code

Bytecode Generation Basics



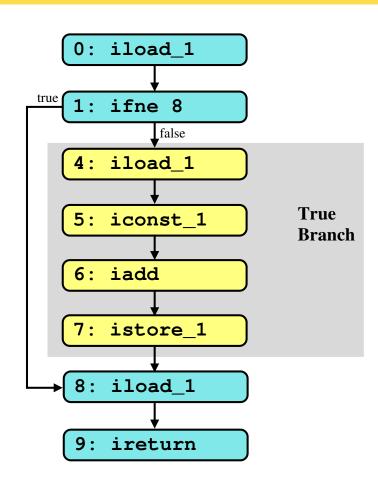
Generating Simple Bytecodes

• Often several choices of bytecode:

Bytecode	Format	Description
iload iload_0 iload_1	[1 byte op][1 byte X] [1 byte op] [1 byte op]	Push local variable X Push local variable 0 Push local variable 1
istore istore_0 istore_1	[1 byte op] [1 byte X] [1 byte op] [1 byte op]	Pop stack to local variable X Pop stack to local variable 0 Pop stack to local variable 1
bipush sipush ldc iconst_0 iconst_1	[1 byte op] [1 byte] [1 byte op] [2 bytes] [1 byte op] [1 byte idx] [1 byte op] [1 byte op]	Push int constant (-128+127) Push int constant (-32768+32767) Push int constant from constant pool Push int zero Push int one

Translating If-Statements

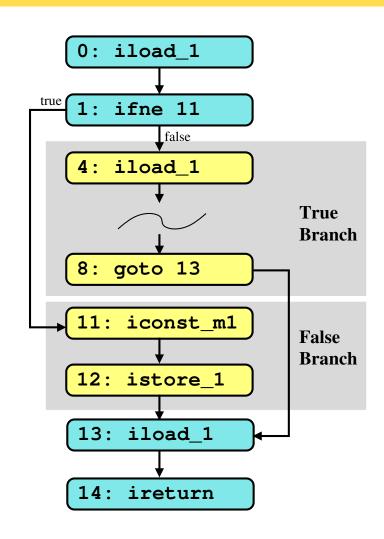
```
int f (int y) {
  if (y == 0) {
    y = y + 1;
  }
  return y;
}
```



• In this case, no **else** branch — easy!

Translating If-Else-Conditionals

```
int f (int y) {
  if (y == 0) {
    y = y + 1;
  } else {
    y = -1;
  }
  return y;
}
```

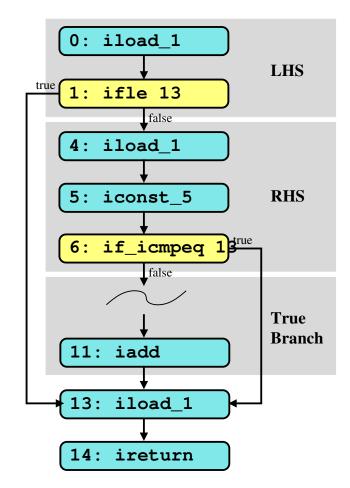


• The true branch **jumps over** the false branch!

Short Circuiting

Logical connectives are translated using short-circuiting:

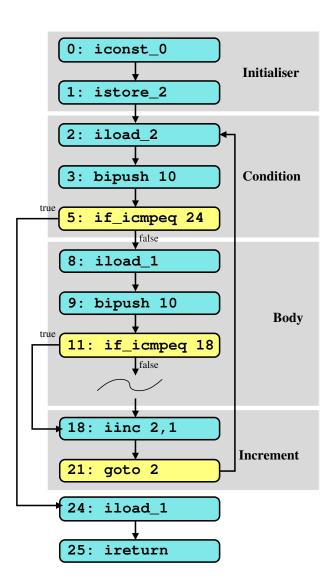
```
int f(int y) {
  if(y > 0 && y!=5) {
    y = y + 1;
  }
  return y;
}
```



Here, right-hand expression only executed if left-hand gives true

Translating Loops

```
int f(int y) {
  for(int i=0;i!=10;++i) {
    if(y==10) continue;
    y = y * 2;
  }
  return y;
}
```



Generating Branch Bytecodes

```
goto [1 byte op][2 bytes offset]

Unconditional Branch (range -32768...+32767)

goto_w [1 byte op][4 bytes offset]

Unconditional Wide Branch (range -2<sup>31</sup> - 1...2<sup>31</sup>)

ifeq [1 byte op][2 bytes offset]

Branch if top two stack locations equal (range -32768...+32767)
```

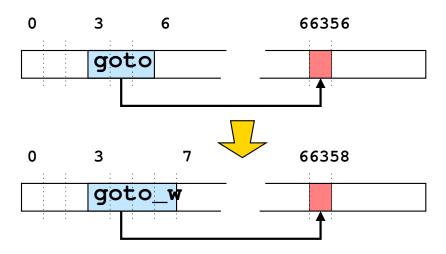
. . .

- Branch bytecodes employ relative addressing, not absolute addressing
- Target address calculated by adding branch bytecode address to offset:

```
void f(int):
    ...
24: goto +35
    ...
59: ...
```

Here, target address of goto bytecode is (24 + 35) = 59

Calculating Branch Offsets



- Algorithm for calculating branch offsets:
 - Generate all bytecodes, assuming branches take 3 bytes
 - If branch exists which cannot reach target:

Replace it with a *wide branch*:

Update offsets of all branches (since they may have changed)

- 3 Repeat step 2 until all branches can reach destination
- Does this algorithm always terminate?
 (need to consider padding of tableswitch + lookupswitch)

Generating Invoke Bytecodes

invokevirtual	[1 byte op][2 bytes index]
	Invoke method on a receiver of class type. The method and receiver types are located in the constant pool at the given index.
invokeinterface	[1 byte op][2 bytes index]
	Invoke method on a receiver of interface type. The method and receiver types are located in the constant pool at the given index.
invokestatic	[1 byte op][2 bytes index]
	Invoke static method. The method and receiver types are located in the constant pool at the given index.
invokespecial	[1 byte op][2 bytes index]
	Invoke special method (e.g. constructor). The method and receiver types are located in the constant pool at the given index.

Generating Invoke Bytecodes (Cont'd)

```
class Test {
  Test(int x) { }
  int f(String s, int i) {
    return 1;
  }

static void m(String[] s) {
  Test t = new Test(123);
  t.f(s[0],2);
  }
}
```

```
static void m(String[] s):
Code:
 0:
     new Test
 3:
     dup
 4:
     bipush 123
      invokespecial Test. <init>: (I) V
 6:
 9:
     astore 1
 10:
     aload 1
 11:
     aload 0
12:
     iconst 0
 13:
     aaload
 14:
     iconst 2
 15:
     invokevirtual Test.f:(L...;I)I
 18:
     pop
 19:
     return
```

- Receiver pushed on stack first (line 10)
- Parameters pushed on stack next in order (lines 13-14)
- Return value is popped afterwards since its not used (line 18)

Generating Switch Bytecodes

• Two bytecodes for switch statements:

tableswitch [op][padding][default][low][high][offsets]

Padding: 0-3 zeroed bytes, so next byte word-aligned.

Default: target address for default label

Low: lowest value in case range

High: Highest value in case range

Offsets: Array of (high-low+1) Case Offsets

lookupswitch [op][padding][default][npairs][pairs]

padding: 0-3 zeroed bytes, so next byte word-aligned.

default: target address for default label

npairs: *number of case value pairs*

pairs: array of pairs mapping case values to offsets

Generating Switch Bytecodes (cont'd)

```
void f(int x) {
 int y;
 switch(x) {
  case 0:
  y = 1;
  break;
  case 1:
   y = 2;
  case 2:
  v = 3;
 default:
   y = -1;
```

```
public void f(int);
  0: iload 1
  1: tableswitch
       default: 37
        low: 0
         high: 2
         offsets: +27, +32, +34
  28: iconst_1
  29: istore_2
  30: goto 39
  33: iconst_2
  34: istore 2
  35: iconst 3
  36: istore_2
  37: iconst_m1
  38: istore_2
  39: return
```

- Tableswitch is useful for contiguous case values
- How many bytes of padding required here?

Generating Switch Bytecodes (cont'd)

```
void f(int x) {
 int y;
 switch(x) {
  case 0:
   v = 1;
  break;
  case 12:
   y = 2;
  case 2046:
   y = 3;
  default:
   v = -1;
```

```
public void f(int);
   0: iload 1
   1: lookupswitch
           default: 45
          npairs: 3
           pairs: 0 \rightarrow +35, 12 \rightarrow +40, 2046 \rightarrow +42
   36: iconst 1
   37: istore 2
   38: goto 47
   41: iconst 2
   42: istore_2
   43: iconst 3
   44: istore 2
   45: iconst m1
   46: istore 2
   47: return
```

- Lookupswitch is useful for non-contiguous case values
- Notice that lookupswitch bytecode is much larger than before.