Algorithm 1: fill_runs_slice(T, t, from, to, v) Input: Tree, T, supporting WeinerLink and Parent; Query string t; Interval over t, [from, to); Starting node v. Output: Tripple (first, last, v): first index of the first failing WeinerLink call; last index last failing WeinerLink call; v the node corresponding to last; 1 currnode \leftarrow T.WeinerLink(T.root(), t[to -1]); $2 idx_set \leftarrow FALSE;$ $3 \text{ k} \leftarrow \text{to} - 2;$ 4 repeat $c \leftarrow t[k-1];$ 5 if T.IsRoot(T.WeinerLink(currnode, t[to - 1])) then 6 $\mathbf{if}\ \mathtt{idx_set} = \mathtt{FALSE}\ \mathbf{then}$ 7 $first \leftarrow k$: 8 $idx_set \leftarrow TRUE;$ 10 end $runs[k] \leftarrow 0;$ 11 repeat 12 $currnode \leftarrow T.parent(v);$ 13 until T.IsRoot(T.WeinerLink(currnode, c)); 14 $v \leftarrow T.WeinerLink(currnode, c);$ 15 $\texttt{last} \leftarrow \texttt{k};$ 16 else 17 $| runs[k] \leftarrow 1$

Algorithm 2: Build RUNS in parallel.

currnode ← T.WeinerLink(currnode, c);

18 19

20

21

end

22 until k > from;

 $k \leftarrow k - 1$;

Algorithm 3: fill_ms_slice

```
Input: Cst of s, T, supporting WeinerLink and Parent; Query string t;
                 Vector ms initialized at 0; Vector runs; Interval over t, [from, to)
     Output: Last index j of ms
 1 k \leftarrow from, i \leftarrow 0, h^* \leftarrow k + 1;
 \mathbf{z} \ \mathsf{c} \leftarrow \mathsf{t}[\mathtt{k}];
 \mathbf{3} \ \mathtt{v} \leftarrow \mathtt{T.root}();
 4 repeat
 5
         h \leftarrow h^*;
          u \leftarrow v;
 6
          repeat
 7
               c \leftarrow t[h^*];
 8
               u \leftarrow \texttt{T.WeinerLink}(v,c);
 9
               if !T.IsRoot(u) then
10
                     v \leftarrow u;
11
                    h^* \leftarrow h^* + 1;
\bf 12
               end
13
          \mathbf{until} \ !T. \mathsf{IsRoot}(\mathtt{u}) \ \mathsf{and} \ \mathtt{h}^* < |\mathtt{t}|;
14
          // resize if needed;
15
          \mathtt{i} \leftarrow \mathtt{i} - \mathtt{h}^* + \mathtt{1};
16
          \mathtt{ms}[\mathtt{i} + +] \leftarrow \mathtt{1};
17
          if h^* < |t| then
18
               repeat
19
                     v \leftarrow T.Parent(v);
20
                    u \leftarrow T.WeinerLink(v, c);
21
               until !T.IsRoot(u);
22
               h^* \leftarrow h^* + 1;
\mathbf{23}
          end
\mathbf{24}
          k' \leftarrow k + 1;
25
          repeat
26
            k' \leftarrow k' + 1;
27
          until k' < |t| and runs[k'] \neq 0;
28
29
          // resize if needed;
          \texttt{ms}[\texttt{k}+1..\texttt{k}'-1] \leftarrow 1;
30
          j \leftarrow \max(k'-1, to);
31
32 until k < to;
```

Algorithm 4: Build RUNS in parallel.

```
Input: Tree, T, supporting WeinerLink and Parent; Query string t;
              Number of threads nthreads dividing |t|.
\mathbf{1} \ \mathbf{for} \ i \leftarrow 1 \ \mathbf{to} \ \mathbf{nthreads} \ \mathbf{do}
\mathbf{2} \quad | \quad \mathtt{ms_i} \leftarrow [1..(2*|\mathtt{t}|)/\mathtt{nthreads}];
з end
4 for i \leftarrow 1 to nthreads do
5 | J_i \leftarrow fill_ms_slice(T, t, ms_i, runs, (i-1) * S, i * S);
6 end
7 for i \leftarrow 1 to nthreads -1 do
s \mid ms_i \leftarrow resize\_to(J_i);
9 end
```