

CMPT 435

Assignment 1

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Problem 1A $\max(f(n), g(n)) = \theta(f(n) + g(n))$

Solution: True

$$\begin{aligned}\max(f(n), g(n)) &\leq f(n) + g(n) \\ &\leq 2 \\ &\leq \max(f(n), g(n)) \text{ for } n = 1\end{aligned}$$

Problem 1B $\min(f(n), g(n)) = \theta(f(n) + g(n))$

Solution: False

$$\begin{aligned}\min(f(n), g(n)) &\not\leq f(n) + g(n) \\ &\not\leq 2 \\ &\not\leq \max(f(n), g(n)) \text{ for } n = 2\end{aligned}$$

Problem 2 $T(n) = 2T(\frac{n}{2}) + n^4$

Solution:

$$\begin{aligned}&= 2[2T(\frac{n}{4}) + (\frac{n}{2})^4] + n^4 \\ &= 4[2T(\frac{n}{8}) + (\frac{n}{4})^4] + 2(\frac{n}{2})^4 + n^4 \\ &= 8[2T(\frac{n}{16}) + (\frac{n}{8})^4] + 4(\frac{n}{2})^4 + 2(\frac{n}{2})^4 + n^4 \\ &= 2^r T(\frac{n}{2^r}) + n^4(1 + \frac{1}{2^3} + \frac{1}{2^6} + \frac{1}{2^9} + \dots + \frac{1}{2^{(k-1)^3}}) \\ &\therefore \theta(n^4)\end{aligned}$$

Problem 3A $T(1) = 1, T(n) = T(n-1) + n$

Solution: We will show that $T(n) = O(n^2), T(n) \leq cn^2$

Basis:

$$T(2) = 3, 4 \leq c * 2 \text{ for } c \geq 2$$

$$T(3) = 4, 9 \leq c * 3 \text{ for } c \geq 3$$

Inductive Hypothesis: Assume for $2 \leq k < n$ that

$$T(k) \text{ is } O(n^2), T(k) \leq cn^2$$

Inductive Step: Prove $k = n$

$$\begin{aligned} T(n) &= T(n-1) + n \\ &\leq c * (n-1)^2 + n \\ &\leq c * n^2 - 2n + 1 + n \\ &\leq c(n^2 - n + 1) \\ &\leq c * n^2 - c * n + c \\ &\leq c * n^2 - n(c-1) \\ &\leq c * n^2 \text{ for } c \geq 1 \\ &O(n^2) \end{aligned}$$

Problem 3B $T(1) = 1, T(n) = 3T(\frac{n}{3}) + 7n$

Solution: We will show that $T(n) = O(n), T(n) \leq cn$

Basis:

$$T(2) = 16, 16 \leq c * 2 \text{ for } c \geq 8$$

$$T(3) = 27, 27 \leq c * 3 \text{ for } c \geq 9$$

Inductive Hypothesis: Assume for $2 \leq k < n$ that

$$T(k) \text{ is } O(n), T(k) \leq cn$$

Inductive Step: Prove $k = n$

$$\begin{aligned}
 T(n) &= 3T\left(\frac{n}{3}\right) + 7n \\
 &\leq 3\left(\frac{n}{3} + 7\left(\frac{n}{3}\right)\right) \\
 &\leq c * \frac{n}{3} + 7\frac{n}{3} \\
 &\leq \frac{n(c+7)}{3} \\
 &\leq 3n(c+7) \\
 &\leq 3cn + 21 \\
 &\leq cn \text{ for } c \geq 3 \\
 &O(n)
 \end{aligned}$$

Problem 4A $T(n) = 2T\left(\frac{n}{2}\right) + n^3$

Solution: A=2, B=2, $f(n) = n^3$, $\log_b a = \log_2 2 = 1$

$$\begin{aligned}
 f(n) &= n^3 = \Omega(n^{\log_b a + \epsilon}) \\
 &= \Omega(n^{1+\epsilon}) \quad \epsilon = 1.
 \end{aligned}$$

So, part 3 of the Master Theorem applies.

$af(n/b) \leq cf(n)$ for some constant $c < 1$

$$\begin{aligned}
 2f(n/2) &= 2\left(\frac{n}{2}\right)^3 \\
 &= \frac{2}{8}n^3 \\
 &= \frac{n^3}{4}
 \end{aligned}$$

$$\frac{n^3}{4} \leq cn^3 \text{ for } c = \frac{1}{4} \text{ So, } T(n) = \theta(n^3)$$

Problem 4B $T(n) = 16T\left(\frac{n}{2}\right) + n^2$

Solution: A=16, B=4, $f(n) = n^2$, $\log_b a = \log_4 16 = 2$

$$\begin{aligned}
 f(n) &= n^2 = \theta(n^{\log_b a}) \\
 &= \theta(n^2)
 \end{aligned}$$

So, part 2 of the Master Theorem applies.

$$\theta(n^{\log_b a} \log n) = \theta(n^2 \log n)$$

Problem 4C $T(n) = 16T(\frac{n}{3}) + n^2$

Solution: A=7, B=3, $f(n) = n^2$, $\log_b a = \log_3 7 \geq 1$ AND ≤ 2 Approximately: 1.77

$$\begin{aligned} f(n) = n^2 &= \Omega(n^{\log_b a + \epsilon}) \\ &= \Omega(n^{\log_3 7 + \epsilon}) \quad \epsilon = 1 - \log_3 7 \end{aligned}$$

So, part 3 of the Master Theorem applies.

$af(n/b) \leq cf(n)$ for some constant $c < 1$

$$\begin{aligned} 7f(n/3) &= 7\left(\frac{n}{3}\right)^2 \\ &= \frac{7}{9}n^2 \end{aligned}$$

$\frac{7}{9}n^2 \leq cn^2$ for $c = \frac{7}{9}$ So, $T(n) = \theta(n^2)$

Problem 4D $T(n) = 7T(\frac{n}{2}) + n^2$

Solution: A=7, B=2, $f(n) = n^2$, $\log_b a = \log_2 7 \geq 2$ AND ≤ 3 Approximately: 2.8

$$\begin{aligned} f(n) = n^2 &= O(n^{\log_b a - \epsilon}) \\ &= O(n^{\log_2 7 - \epsilon}) \quad \epsilon = \log_2 7 - 2 \end{aligned}$$

So, part 1 of the Master Theorem applies.

$\theta(n^{\log_b a}) = \theta(n^{\log_2 7})$