

Thm. CG algorithm:

$$x^{(0)} = \underline{0}$$

$$r^{(0)} = p^{(0)} = b$$

$\Rightarrow$

$$i) \quad x^{(k+1)} = x^{(k)} + \alpha_k p^{(k)}$$

$$\alpha_k = \frac{p^{(k)T} r^{(k)}}{p^{(k)T} A p^{(k)}}$$

$$ii) \quad r^{(k+1)} = b - A x^{(k+1)}$$

$$iii) \quad p^{(k+1)} = r^{(k+1)} + \beta_k p^{(k)}$$

$$\beta_k = \frac{-r^{(k+1)T} A p^{(k)}}{p^{(k)T} A p^{(k)}}$$

$\Rightarrow$  CG converges to the true soln  $x^*$   
in at most "n" steps

$$(a) \quad r^{(k)T} r^{(i)} = 0, \quad \forall i < k$$

$$(b) \quad p^{(k)T} A p^{(i)} = 0, \quad \forall i < k$$

$$\mathcal{K}^{(k)} = \text{Span} \{ x^{(1)}, \dots, x^{(k)} \}$$

$\nearrow$

Krylov  
Space

$$= \text{Span} \{ p^{(0)}, \dots, p^{(k-1)} \}$$

$$= \text{Span} \{ r^{(0)}, \dots, r^{(k-1)} \}$$

$$= \text{span} \{ b, Ab, \dots, A^k b \}$$

$$(c) \quad e^{(k)} = x^* - x^{(k)}$$

$$\|e^{(k)}\|_A = \inf_{u \in K^{(k)}} \|x^* - u\|_A$$

Rank.  $\|e^{(k+1)}\|_A \leq \|e^{(k)}\|_A$

Ⓟ  $K^{(k)} \subseteq K^{(k+1)}$

Note

$$a > b > 0$$

$$\frac{1}{a+b} < \frac{1}{a} < \frac{1}{b}$$

Rank. Rate of convergence of CG

$$\|e^{(n)}\|_A \leq \frac{2}{(\oplus)^n + (\otimes)^{-n}} \|e^{(0)}\|_A$$

$$\left( \oplus = \frac{\sqrt{k} + 1}{\sqrt{k} - 1} (> 1), \quad k = \text{cond}(A) \right)$$

$$\leq \frac{2}{(\otimes)^n} \|e^{(0)}\|_A$$

i)  $\otimes \sim 1$   $\Rightarrow$  slow convergence  
(if  $k$  is too large)

ii)  $\otimes \gg 1$   $\Rightarrow$  fast convergence  
(if  $k$  is not too large)

Alg. Basic CG

$$x = x_0 \quad (= 0)$$

$$r = b - Ax \quad (= b)$$

$$p = r \quad (= b)$$

do while  $\|r\| > \text{large}$

$$y = Ap$$

$$\alpha = \frac{p^T r}{p^T y}$$

$$x = x + \alpha p$$

$$\begin{aligned} r &= r - \alpha y &= (b - Ax) - \alpha Ap \\ & &= b - A(x + \alpha p) \\ & &= b - Ax \end{aligned}$$

calculate  $\|r\|$

$$\beta = - \frac{r^T y}{p^T y}$$

$$p = r + \beta p$$

end do

[3] Preconditioning for CG

$$Ax = b$$

Goal: Find  $M$  (invertible) s.t

i)  $M \sim A$ ,

ii)  $M^{-1}$  is easy.

iii)  $M^{-1} A x = M^{-1} b$

vi) we hope:  $k(M^{-1} A) \ll k(A)$

Rank:  $M$  is called a preconditioner

(Ex)

i)  $M = A$

$\Downarrow$

$$M^T A x = M^T b$$

$\Downarrow$

$$k(M^T A) = k(I) = 1, \quad Ix = A^{-1}b$$

$\rightarrow$  useless!

ii)  $M = I, \quad M^T = I$  (easy!)

$$M^T A x = M^T b$$

$$\Leftrightarrow I A x = I b$$

(useless!)

iii)  $M \sim A$  &  $M^{-1}$ : easy

What we want

$$\rightarrow M = \text{diag}(A) \text{ if}$$

$$M^{-1} = \begin{bmatrix} \frac{1}{a_{11}} & & \\ & \ddots & \\ & & \frac{1}{a_{nn}} \end{bmatrix}$$

$A$  is diagonally dominant.

$$|a_{ii}| \geq \sum_{j \neq i} |a_{ij}|$$

$$\begin{bmatrix} 0 & \frac{1}{a_{max}} \end{bmatrix}$$

(Jacobi preconditioner)

This works well for  $A$ : diagonally dominant

Rule. Not all  $M = \text{diag}(A)$  works all the time

(Ex) Consider  $A : \begin{cases} a_{ii} = D, \forall i. \\ a_{ij} = 1, \forall i \neq j \end{cases}$

For example  $A_{3 \times 3} = \begin{bmatrix} D & 1 & 1 \\ 1 & D & 1 \\ 1 & 1 & D \end{bmatrix}$

$$\rightarrow M = \begin{bmatrix} \frac{1}{D} & & \\ & \frac{1}{D} & \\ & & \frac{1}{D} \end{bmatrix}$$

$$\rightarrow M^{-1}A = \begin{bmatrix} 1 & \frac{1}{D} & \frac{1}{D} \\ \frac{1}{D} & 1 & \frac{1}{D} \\ \frac{1}{D} & \frac{1}{D} & 1 \end{bmatrix} = \left(\frac{1}{D}\right)A$$

$$\rightarrow \text{cond}(M^{-1}A) = \text{cond}\left(\frac{1}{D}A\right) = \text{cond}(A)!$$

Not good!

$\rightarrow$  If  $d_{ii} \neq \text{const}$ , then

$M = \text{diag}(A)$  would work!

Ⓚ  $M^{-1}A$  is Not spd even if  $A$  is spd!

Ⓐ Find  $C$  s.t.  $M = CC^T$   
 $\nwarrow$  (invertible)

$$\rightarrow M^{-1} = C^{-T}C^{-1}$$

$$\rightarrow M^{-1}Ax = \cancel{C^{-T}}C^{-1}Ax$$

||

$$M^+ b = \cancel{C}^T C^{-1} b$$

$$\rightarrow \boxed{C^{-1} A x = C^{-1} b}$$

$$\rightarrow \underbrace{(C^{-1} A C^{-T})}_{=\tilde{A}} \underbrace{(C^T x)}_{=\tilde{x}} = \underbrace{C^{-1} b}_{=\tilde{b}}$$

$$\rightarrow \boxed{\tilde{A} \tilde{x} = \tilde{b}}, \text{ when } \tilde{A} \text{ is spd}$$

(pf) i)  $\tilde{A}$  is symm

$$\tilde{A}^T = (C^{-1} A C^{-T})^T$$

$$= C^{-1} A C^{-T} = \tilde{A}$$

ii)  $\tilde{A}$  is pd. ;

$$u^T \tilde{A} u = (u^T C^{-1}) A \underbrace{(C^{-T} u)}_{=v}$$

$$= \underbrace{v^T A v}_{> 0}, \text{ plus}$$

(  $v \neq 0$ , since  $C^{-T}$  is )



invertible

$$\textcircled{ii} \quad \underline{\underline{\tilde{A} : \text{spd}}}$$

$$\rightarrow \Phi \subset G$$

$$(\text{preconditioned} - CG)$$

Rank. "Exit condition"

$$i) \quad \|\tilde{r}^{(k)}\| > \underline{\underline{\text{large}}} \quad \leftarrow \text{one good way}$$

or

$$ii) \quad \alpha_k \sim 0 \Rightarrow \text{converged!}$$

## ④ GMRES (Generalized Minimal Residual Method)

$$A \in \mathbb{R}^{m \times m}$$

$(n < m)$

i) Search space =  $K^{(n)}$

ii) works for non-spd.

iii) create  $K^{(n)} \subset \mathbb{R}^m$ ,  $n$ : fixed

$$\text{find } x_n \text{ s.t. } \|b - Ax_n\| = \inf_{u \in K^{(n)}} \|b - Au\|$$

### §4.1 Krylov subspace

Def.  $\mathcal{K}_n(A, v)$ : Krylov subspace of  $A$  w.r.t.  $v$

$$\left( \begin{array}{l} \text{def} \\ \text{span} \{ v, Av, A^2v, \dots, A^{n-1}v \} \\ (\subset \mathbb{R}^m) \end{array} \right)$$

Def.  $K_n$ :  $m \times n$  Krylov matrix

$$\stackrel{(\text{def})}{=} \begin{bmatrix} v & Av & \dots & A^{n-1}v \end{bmatrix}$$

$$= \begin{bmatrix} k_1 & k_2 & \dots & k_n \end{bmatrix}$$

Def.  $\mathcal{K}_n(A, v) = \text{span} \{k_1, \dots, k_n\}$

$$= \text{span} \{v, Av, \dots, A^{n-1}v\}$$

Def. Goal:

$$\min \|b - Ax_n\|_2 \quad \text{where } x_n \in \mathcal{K}_n(A, v)$$

$$\rightarrow \exists c \in \mathbb{R}^n \text{ s.t.}$$

$$x_n = c_1 k_1 + \dots + c_n k_n$$

$$= K_n c$$

$$c = \begin{bmatrix} c_1 \\ \vdots \\ c_n \end{bmatrix}$$

$$\rightarrow \min_{x_n \in \mathcal{K}_n(A, v)} \|b - Ax_n\| \Leftrightarrow \min_{c \in \mathbb{R}^n} \|b - \overbrace{A K_n c}^{m \times n \cdot n \times 1}\|$$

$$\rightarrow \underbrace{A}_{m \times m} \underbrace{K_n}_{m \times n} : m \times n \quad (m > n)$$

$$\rightarrow \hat{Q} \hat{R} \approx A K_n \quad \text{(reduced QR)} \quad \leftarrow \text{over-determined system}$$

$$\rightarrow P_{AK_n} = \hat{Q} \hat{Q}^T$$

$$\rightarrow \underbrace{(A K_n)}_{P_{AK_n}} c = \hat{Q} \hat{Q}^T b \quad \leftarrow$$

$$\rightarrow \underbrace{\hat{Q}^T \hat{Q}}_{= I} \hat{R} c = \underbrace{\hat{Q}^T \hat{Q}}_{\leq I} \hat{Q}^T b$$

$$\rightarrow \boxed{\hat{R} c = \hat{Q}^T b}$$

Def.  $A: m \times n$

$\rightarrow A$  minimal poly of  $A$  is "p"

uniquely defined as a monic  
poly with the least <sup>possible</sup> degree ( $\alpha_n = 1$ )

satisfying  $p(A)v = \underline{0}$ ,  $\forall v \in \mathbb{R}^m$

$$\text{i.e., } p(A)v = \left( \sum_{i=0}^n \alpha_i A^i \right) v$$

$$= (\alpha_0 I + \alpha_1 A + \dots + A^n) v$$

$$= \underline{0}$$

Rank. Why  $K_n(A, b)$  is a good  
choice?

$$\textcircled{A} 0 = p(A)b = (\alpha_0 I + \dots + A^n) b$$

$$\rightarrow I = -\frac{1}{\alpha_0} \left( \sum_{i=1}^n \alpha_i A^i \right) b$$

$$\rightarrow \underbrace{A^{-1}b}_{//} = -\frac{1}{\alpha_0} \left( \sum_{i=0}^{n-1} \alpha_{i+1} A^i \right) b$$

$$\rightarrow \underline{\underline{x}} = A^{-1}b = -\frac{1}{\alpha_0} \sum_{i=0}^{n-1} \alpha_{i+1} A^i b$$