

# Theoretical Computer Science: Exercise 3 - Reductions

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Show that the language:

$X = \{\langle M, w \rangle \mid M \text{ is a 1-band TM and } M \text{ does not modify the part of the band where the input } w \text{ is}\}$  is not decidable.

Let us assume that  $X$  is decidable and let  $R$  be the Turing machine that decides  $X$ . Now, we can create a Turing machine  $S$  that decides  $A_{TM}$ <sup>1</sup>. The construction of  $S$  is the following:

- The input is  $\langle M, w \rangle$ ,  $M$  being a Turing Machine and  $w$  being a string.
- Now, we create a new Turing Machine  $M'$  using the following steps:
  - The input is code  $y$
  - Run  $M$  with input  $w$ , without editing the input  $y$
  - If  $M$  accepts,  $M'$  accepts
  - Otherwise delete the input string  $y$  and write 1 and reject
- Now run  $R$  with  $M'$  as input
- $S$  accepts if  $R$  accepts,  $S$  rejects if  $R$  rejects

If  $M$  accepts  $w$ ,  $M'$  will not modify the input (which means it will be accepted by  $R$ ). If  $M$  does not accept  $w$ ,  $M'$  will modify the input (it will be rejected by  $R$ ). This means that  $R$  will accept  $M'$  if and only if  $M$  accepts  $w$ . Consequently,  $S$  decides  $A_{TM}$ .

However, since we know that  $A_{TM}$  is undecidable,  $S$  cannot exist. Due to this contradiction  $R$  cannot exist either, so  $X$  must be undecidable.

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<sup>1</sup> $A_{TM} = \{\langle T, w \rangle \mid T \text{ is a TM and } T \text{ accepts } w\}$