fc	fractional capability variable
x, g, a, b	expression variable
k	integer variable
el	array-element variable
l	location variable
M	matrix variable

m	::=			matrix expressions
	|	M		matrix variables
	|	$m + m'$		matrix addition
	|	$m\, m'$		matrix multiplication
	|	(m)	S	parentheses
f	::=			fractional capability
	|	fc		variable
	|	1		whole capability
	|	$\frac{1}{2} \cdot f$		
t	::=			linear type
	|	**unit**		unit
	|	**bool**		boolean (true/false)
	|	**int**		63-bit integers
	|	**elt**		array element
	|	f **arr**		arrays
	|	f **mat**		matrices
	|	$!t$		multiple-use type
	|	$\forall\, fc.t$	bind fc in t	frac. cap. generalisation
	|	$t \otimes t'$		pair
	|	$t \multimap t'$		linear function
	|	(t)	S	parentheses
p	::=			primitive
	|	**not**		boolean negation
	|	$(+)$		integer addition
	|	$(-)$		integer subtraction
	|	$(*)$		integer multiplication
	|	$(/)$		integer division
	|	$(=)$		integer equality
	|	$(\langle)$		integer less-than
	|	$(+.)$		element addition
	|	$(-.)$		element subtraction
	|	$(*.)$		element multiplication
	|	$(/.)$		element division
	|	$(= .)$		element equality
	|	$(< .)$		element less-than
	|	**set**		array index assignment
	|	**get**		array indexing
	|	**share**		share array
	|	**unshare**		unshare array
	|	**free**		free arrary
	|	**array**		Owl: make array
	|	**copy**		Owl: copy array
	|	**sin**		Owl: map sine over array
	|	**hypot**		Owl: $x_i := \sqrt{x_i^2 + y_i^2}$

	asum		BLAS: $\sum_i \lvert x_i \rvert$
	axpy		BLAS: $x := \alpha x + y$
	dot		BLAS: $x \cdot y$
	rotmg		BLAS: see its docs
	scal		BLAS: $x := \alpha x$
	amax		BLAS: $\operatorname{argmax} i : x_i$
	setM		matrix index assignment
	getM		matrix indexing
	shareM		share matrix
	unshareM		unshare matrix
	freeM		free matrix
	matrix		Owl: make matrix
	copyM		Owl: copy matrix
	copyM_to		Owl: copy matrix onto another
	sizeM		dimension of matrix
	trnsp		transpose matrix
	gemm		BLAS: $C := \alpha A^{T?} B^{T?} + \beta C$
	symm		BLAS: $C := \alpha AB + \beta C$
	posv		BLAS: Cholesky decomp. and solve
	potrs		BLAS: solve with given Cholesky

v	::=			values
	|	p		primitives
	|	x		variable
	|	$()$		unit introduction
	|	**true**		true
	|	**false**		false
	|	k		integer
	|	l		heap location
	|	el		array element
	|	**Many** v		!-introduction
	|	**fun** $fc \to v$		frac. cap. abstraction
	|	$v[f]$		frac. cap. specialisation
	|	(v, v')		pair introduction
	|	**fun** $x : t \to e$	bind x in e	abstraction
	|	**fix** $(g, x : t, e : t')$	bind $g \cup x$ in e	fixpoint
	|	(v)	S	parentheses

e	::=			expression
	|	p		primitives
	|	x		variable
	|	**let** $x = e$ **in** e'	bind x in e'	let binding
	|	$()$		unit introduction
	|	**let** $() = e$ **in** e'		unit elimination
	|	**true**		true
	|	**false**		false

	$\mid$ $\mathbf{if}\ e\ \mathbf{then}\ e_1\ \mathbf{else}\ e_2$		if
	$\mid$ k		integer
	$\mid$ l		heap location
	$\mid$ el		array element
	$\mid$ $\mathbf{Many}\ e$		!-introduction
	$\mid$ $\mathbf{let}\,\mathbf{Many}\ x = e\ \mathbf{in}\ e'$		!-elimination
	$\mid$ $\mathbf{fun}\ fc \to e$		frac. cap. abstraction
	$\mid$ $e[f]$		frac. cap. specialisation
	$\mid$ (e, e')		pair introduction
	$\mid$ $\mathbf{let}\,(a, b) = e\ \mathbf{in}\ e'$	$\mathsf{bind}\ a \cup b\ \mathsf{in}\ e'$	pair elimination
	$\mid$ $\mathbf{fun}\ x : t \to e$	$\mathsf{bind}\ x\ \mathsf{in}\ e$	abstraction
	$\mid$ $e\ e'$		application
	$\mid$ $\mathbf{fix}\,(g, x : t, e : t')$	$\mathsf{bind}\ g \cup x\ \mathsf{in}\ e$	fixpoint
	$\mid$ (e)	S	parentheses
erased_v, *ev*	::=		values
	$\mid$ p		primitives
	$\mid$ x		variable
	$\mid$ $()$		unit introduction
	$\mid$ $\mathbf{true}$		true
	$\mid$ $\mathbf{false}$		false
	$\mid$ k		integer
	$\mid$ l		heap location
	$\mid$ el		array element
	$\mid$ (ev, ev')		pair introduction
	$\mid$ $\mathbf{fun}\ x \to er$	$\mathsf{bind}\ x\ \mathsf{in}\ er$	abstraction
	$\mid$ $\mathbf{fix}\,(g, x, er)$	$\mathsf{bind}\ g \cup x\ \mathsf{in}\ er$	fixpoint
	$\mid$ (ev)	S	parentheses
erased, *er*	::=		expression
	$\mid$ p		primitives
	$\mid$ x		variable
	$\mid$ $\mathbf{let}\ x = er\ \mathbf{in}\ er'$	$\mathsf{bind}\ x\ \mathsf{in}\ er'$	let binding
	$\mid$ $()$		unit introduction
	$\mid$ $\mathbf{let}\,() = er\ \mathbf{in}\ er'$		unit elimination
	$\mid$ $\mathbf{true}$		true
	$\mid$ $\mathbf{false}$		false
	$\mid$ $\mathbf{if}\ er\ \mathbf{then}\ er_1\ \mathbf{else}\ er_2$		if
	$\mid$ k		integer
	$\mid$ l		heap location
	$\mid$ el		array element
	$\mid$ (er, er')		pair introduction
	$\mid$ $\mathbf{let}\,(a, b) = er\ \mathbf{in}\ er'$	$\mathsf{bind}\ a \cup b\ \mathsf{in}\ er'$	pair elimination
	$\mid$ $\mathbf{fun}\ x \to er$	$\mathsf{bind}\ x\ \mathsf{in}\ er$	abstraction
	$\mid$ $er\ er'$		application
	$\mid$ $\mathbf{fix}\,(g, x, er)$	$\mathsf{bind}\ g \cup x\ \mathsf{in}\ er$	fixpoint

	(er)	S	parentheses

C	::=			evaluation contexts
	$\mid$	$\mathbf{let}\, x = [-]\, \mathbf{in}\, er$	$\mathsf{bind}\ x\ \mathsf{in}\ er$	let binding
	$\mid$	$\mathbf{let}\, () = [-]\, \mathbf{in}\, er$		unit elimination
	$\mid$	$\mathbf{if}\, [-]\, \mathbf{then}\, er_1\, \mathbf{else}\, er_2$		if
	$\mid$	$([-], er)$		pair introduction
	$\mid$	$(ev, [-])$		pair introduction
	$\mid$	$\mathbf{let}\, (a, b) = [-]\, \mathbf{in}\, er$	$\mathsf{bind}\ a \cup b\ \mathsf{in}\ er$	pair elimination
	$\mid$	$[-]er$		application
	$\mid$	$ev[-]$		application

Θ	::=		fractional capability environment
	$\mid$	$.$	
	$\mid$	Θ, fc	

Γ	::=		linear types environment
	$\mid$	$.$	
	$\mid$	$\Gamma, x : t$	
	$\mid$	Γ, Γ'	

Δ	::=		intuitionistic types environment
	$\mid$	$.$	
	$\mid$	$\Delta, x : t$	

σ	::=		heap
	$\mid$	$\{\}$	empty heap
	$\mid$	$\sigma \uplus \{l \mapsto m_{k_1,k_2}\}$	location l points to matrix m

$\boxed{\Theta \vdash f\ \mathsf{Cap}}$ Valid fractional capabilities

$$\frac{fc \in \Theta}{\Theta \vdash fc\ \mathsf{Cap}}\ \textsc{WF\_Cap\_Var}$$

$$\frac{}{\Theta \vdash 1\ \mathsf{Cap}}\ \textsc{WF\_Cap\_Zero}$$

$$\frac{\Theta \vdash f\ \mathsf{Cap}}{\Theta \vdash \frac{1}{2} \cdot f\ \mathsf{Cap}}\ \textsc{WF\_Cap\_Succ}$$

$\boxed{\Theta \vdash t\ \mathsf{Type}}$ Valid types

$$\frac{}{\Theta \vdash \mathbf{unit}\ \mathsf{Type}}\ \textsc{WF\_Type\_Unit}$$

$$\frac{}{\Theta \vdash \mathbf{bool}\ \mathsf{Type}}\ \textsc{WF\_Type\_Bool}$$

$$\frac{}{\Theta \vdash \mathbf{int}\ \mathsf{Type}}\ \textsc{WF\_Type\_Int}$$

$$\frac{}{\Theta \vdash \mathbf{elt}\ \mathsf{Type}}\ \textsc{WF\_Type\_Elt}$$

$$\frac{\Theta \vdash f\ \mathsf{Cap}}{\Theta \vdash f\ \mathbf{arr}\ \mathsf{Type}}\ \textsc{WF\_Type\_Array}$$

$$\frac{\Theta \vdash t\ \mathsf{Type}}{\Theta \vdash\ !t\ \mathsf{Type}}\ \textsc{WF\_Type\_Bang}$$

$$\frac{\Theta, fc \vdash t\ \mathsf{Type}}{\Theta \vdash \forall\, fc.t\ \mathsf{Type}}\ \textsc{WF\_Type\_Gen}$$

$$\frac{\Theta \vdash t\ \mathsf{Type} \quad \Theta \vdash t'\ \mathsf{Type}}{\Theta \vdash t \otimes t'\ \mathsf{Type}}\ \textsc{WF\_Type\_Pair}$$

$$\frac{\Theta \vdash t\ \mathsf{Type} \quad \Theta \vdash t'\ \mathsf{Type}}{\Theta \vdash t \multimap t'\ \mathsf{Type}}\ \textsc{WF\_Type\_Lolly}$$

$\boxed{\Theta; \Delta; \Gamma \vdash e : t}$ Typing rules for expressions

$$\frac{}{\Theta; \Delta; \cdot, x : t \vdash x : t}\ \textsc{Ty\_Var\_Lin}$$

$$\frac{x : t \in \Delta}{\Theta; \Delta; \cdot \vdash x : t}\ \textsc{Ty\_Var}$$

$$\frac{\Theta; \Delta; \Gamma \vdash e : t \quad \Theta; \Delta; \Gamma', x : t \vdash e' : t'}{\Theta; \Delta; \Gamma, \Gamma' \vdash \mathbf{let}\, x = e\, \mathbf{in}\, e' : t'}\ \textsc{Ty\_Let}$$

$$\frac{}{\Theta; \Delta; \cdot \vdash () : \mathbf{unit}}\ \textsc{Ty\_Unit\_Intro}$$

$$\frac{\Theta; \Delta; \cdot \vdash e : \mathbf{unit} \quad \Theta; \Delta; \Gamma \vdash e' : t}{\Theta; \Delta; \Gamma \vdash \mathbf{let}\, () = e\, \mathbf{in}\, e' : t}\ \textsc{Ty\_Unit\_Elim}$$

$$\frac{}{\Theta; \Delta; \cdot \vdash \mathbf{true} : !\mathbf{bool}}\ \textsc{Ty\_Bool\_True}$$

$$\frac{}{\Theta; \Delta; \cdot \vdash \mathbf{false} : !\mathbf{bool}}\ \textsc{Ty\_Bool\_False}$$

$$\frac{\Theta; \Delta; \Gamma \vdash e : \mathbf{bool} \quad \Theta; \Delta; \Gamma' \vdash e_1 : t' \quad \Theta; \Delta; \Gamma' \vdash e_2 : t'}{\Theta; \Delta; \Gamma, \Gamma' \vdash \mathbf{if}\, e\, \mathbf{then}\, e_1\, \mathbf{else}\, e_2 : t}\ \textsc{Ty\_Bool\_Elim}$$

$$\frac{}{\Theta; \Delta; \cdot \vdash k : !\mathbf{int}}\ \textsc{Ty\_Int\_Intro}$$

$$\frac{}{\Theta; \Delta; \cdot \vdash el : !\mathbf{elt}}\ \textsc{Ty\_Elt\_Intro}$$

$$\frac{\Theta; \Delta; \cdot \vdash v : t \quad v \neq l}{\Theta; \Delta; \cdot \vdash \mathbf{Many}\, v : !t}\ \textsc{Ty\_Bang\_Intro}$$

$$\frac{\Theta; \Delta; \Gamma \vdash e : !t \quad \Theta; \Delta, x : t; \Gamma' \vdash e' : t'}{\Theta; \Delta; \Gamma, \Gamma' \vdash \mathbf{let}\, \mathbf{Many}\, x = e\, \mathbf{in}\, e' : t'}\ \textsc{Ty\_Bang\_Elim}$$

$$\frac{\Theta; \Delta; \Gamma \vdash e : t \quad \Theta; \Delta; \Gamma' \vdash e' : t'}{\Theta; \Delta; \Gamma, \Gamma' \vdash (e, e') : t \otimes t'}\ \textsc{Ty\_Pair\_Intro}$$

$$\frac{\Theta;\Delta;\Gamma \vdash e_{12} : t_1 \otimes t_2 \quad \Theta;\Delta;\Gamma', a : t_1, b : t_2 \vdash e : t}{\Theta;\Delta;\Gamma,\Gamma' \vdash \mathbf{let}\,(a,b) = e_{12}\,\mathbf{in}\,e : t}\ \textsc{Ty\_Pair\_Elim}$$

$$\frac{\Theta \vdash t' \,\mathsf{Type} \quad \Theta;\Delta;\Gamma, x : t' \vdash e : t}{\Theta;\Delta;\Gamma \vdash \mathbf{fun}\,x : t' \to e : t' \multimap t}\ \textsc{Ty\_Lambda}$$

$$\frac{\Theta;\Delta;\Gamma \vdash e : t' \multimap t \quad \Theta;\Delta;\Gamma' \vdash e' : t'}{\Theta;\Delta;\Gamma,\Gamma' \vdash e\ e' : t}\ \textsc{Ty\_App}$$

$$\frac{\Theta, fc;\Delta;\Gamma \vdash e : t}{\Theta;\Delta;\Gamma \vdash \mathbf{fun}\,fc \to e : \forall\, fc.t}\ \textsc{Ty\_Gen}$$

$$\frac{\Theta \vdash f \,\mathsf{Cap} \quad \Theta;\Delta;\Gamma \vdash e : \forall\, fc.t}{\Theta;\Delta;\Gamma \vdash e[f] : t[f/fc]}\ \textsc{Ty\_Spc}$$

$$\frac{\Theta;\Delta, g : t \multimap t';\cdot, x : t \vdash e : t'}{\Theta;\Delta;\cdot \vdash \mathbf{fix}\,(g, x : t, e : t') : !(t \multimap t')}\ \textsc{Ty\_Fix}$$

$\boxed{\langle \sigma, er\rangle \to \langle \sigma', er'\rangle}$ operational semantics

$$\frac{}{\langle \sigma, \mathbf{let}\,() = ()\,\mathbf{in}\,er\rangle \to \langle \sigma, er\rangle}\ \textsc{Op\_Let\_Unit}$$

$$\frac{}{\langle \sigma, \mathbf{let}\,x = ev\,\mathbf{in}\,er\rangle \to \langle \sigma, er[x/ev]\rangle}\ \textsc{Op\_Let\_Var}$$

$$\frac{}{\langle \sigma, \mathbf{if\,true\,then}\,er_1\,\mathbf{else}\,er_2\rangle \to \langle \sigma, er_1\rangle}\ \textsc{Op\_If\_True}$$

$$\frac{}{\langle \sigma, \mathbf{if\,false\,then}\,er_1\,\mathbf{else}\,er_2\rangle \to \langle \sigma, er_2\rangle}\ \textsc{Op\_If\_False}$$

$$\frac{}{\langle \sigma, \mathbf{fix}\,(g, x, er)\ ev\rangle \to \langle \sigma, er[x/ev][g/\mathbf{fix}\,(g, x, er)]\rangle}\ \textsc{Op\_Let\_Fix}$$

$$\frac{}{\langle \sigma, (\mathbf{fun}\,x \to er)\ ev\rangle \to \langle \sigma, er[x/ev]\rangle}\ \textsc{Op\_App}$$

$$\frac{\langle \sigma, er\rangle \to \langle \sigma, er'\rangle}{\langle \sigma, C[er]\rangle \to \langle \sigma, C[er']\rangle}\ \textsc{Op\_Context}$$

$$\frac{0 \le k_1, k_2}{\langle \sigma, \mathbf{matrix}\,k_1\,k_2\rangle \to \langle \sigma \uplus \{l \mapsto M_{k_1,k_2}\}, l\rangle}\ \textsc{Op\_Matrix}$$

$$\frac{}{\langle \sigma \uplus \{l \mapsto m_{k_1,k_2}\}, \mathbf{free}\,l\rangle \to \langle \sigma, ()\rangle}\ \textsc{Op\_Free}$$

$$\frac{}{\langle \sigma, \mathbf{share}\,l\rangle \to \langle \sigma, (l, l)\rangle}\ \textsc{Op\_Share}$$

$$\frac{}{\langle \sigma, \mathbf{unshare}\,l\,l\rangle \to \langle \sigma, l\rangle}\ \textsc{Op\_Unshare}$$

$$\frac{\sigma(l_1) = M1_{k_1,k_2} \quad \sigma(l_2) = M2_{k_2,k_3}}{\langle \sigma \uplus \{l_3 \mapsto M1_{k_1,k_3}\}, \mathbf{gemm}\,l_1\,l_2\,l_3\rangle \to \langle \sigma \uplus \{l_3 \mapsto (M_1\,M_2 + M_3)_{k_1,k_3}\}, ((l_1, l_2), l_3)\rangle}\ \textsc{Op\_Gemm}$$