


1 NumLin: Linear Types for Linear Algebra

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8 — Abstract —

9 We present NUMLIN, a functional programming language whose type system is designed to enforce
10 the safe usage of the APIs of low-level linear algebra libraries (such as BLAS/LAPACK). We do
11 so through a brief description of its key features and several illustrative examples. We show that
12 NUMLIN's type system is sound and that its implementation improves upon naïve implementations
13 of linear algebra programs, almost towards C-levels of performance. By doing so, we demonstrate (a)
14 that linear types are well-suited to expressing the APIs of low-level linear algebra libraries accurately
15 and concisely and (b) that, despite the complexity of prior work on it, fractional permissions
16 can actually be implemented using simple, well-known techniques and be used practically in real
17 programs.

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21 **Supplement Material** www.github.com/dc-mak/NumLin

22 1 Introduction

23 Programmers writing numerical software often find themselves caught on the horns of a
24 dilemma. The foundational, low-level linear algebra libraries such as BLAS and LAPACK
25 offer programmers very precise control over the memory lifetime and usage of vector and
26 matrix values. However, this power comes paired with the responsibility to manually manage
27 the memory associated with each array object, and in addition to bringing in the familiar
28 difficulties of reasoning about lifetimes, aliasing and sharing that plague low-level systems
29 programming; this also moves the APIs away from the linear-algebraic, mathematical style
30 of thinking that numerical programmers want to use.

31 As a result, programmers often turn to higher-level languages such as Matlab, R and
32 NumPy, which offer very high-level array abstractions that can be viewed as ordinary
33 mathematical values. This makes programming safer, as well as making prototyping and
34 verification much easier, since it lets programmers write programs which bear a closer
35 resemblance to the formulas that the mathematicians and statisticians designing these
36 algorithms prefer to work with, and ensures that program bugs will reflect incorrectly-
37 computed values rather than heap corruption.

38 The intention is that these languages can use libraries BLAS and LAPACK, without
39 having to expose programmers to explicit memory management. However, this benefit comes
40 at a price: because user programs do not worry about aliasing, the language implementations
41 cannot in general exploit the underlying features of the low-level libraries that let them
42 explicitly manage and reuse memory. As a result, programs written in high-level statistical
43 languages can be much less memory-efficient than programs that make full use of the powers
44 the low-level APIs offer.



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So in practice, programmers face a tradeoff: they can eschew safety and exploit the full power of the underlying linear algebra libraries, or they can obtain safety at the price of unneeded copies and worse memory efficiency. In this work, we show that this tradeoff is not a fundamental one.

NUMLIN is a functional programming language whose type system is designed to enforce the safe usage of the APIs of low-level linear algebra libraries (such as BLAS/LAPACK). It does so by combining linear types, fractional permissions, runtime errors and recursion into a small, easily understandable, yet expressive set of core constructs.

NUMLIN allows a novice to understand and work with complicated linear algebra library APIs, as well as point out subtle aliasing bugs and reduce memory usage in existing programs. In fact, we were able to use NUMLIN to find linearity and aliasing bugs in a linear algebra algorithm that was *generated* by another program *specifically designed to translate matrix expressions into an efficient sequence of calls to linear algebra routines*. We were also able to reduce the number of temporaries used by the same algorithm, using NUMLIN’s type system to guide us.

NUMLIN’s implementation supports several syntactic conveniences as well as a *usable* integration with real OCaml libraries.

1.1 Contributions

In this paper

- we describe NUMLIN, a linearly typed language for linear algebra programs
- we illustrate that NUMLIN’s design and features are well-suited to its intended domain with progressively sophisticated examples
- we prove NUMLIN’s soundness, using a step-indexed logical relation
- we describe a very simple, unification based type-inference algorithm for polymorphic fractional permissions (similar to ones used for parametric polymorphism), demonstrating an alternative approach to dataflow analysis [8]
- we describe an implementation that is both compatible with and usable from existing code
- we show an example of how using NUMLIN helped highlight linearity and aliasing bugs, and reduce the memory usage of a *generated* linear algebra program
- we show that using NUMLIN, we can achieve parity with C for linear algebra routines, whilst having much better static guarantees about the linearity and aliasing behaviour of our programs.

2 NumLin Overview and Examples

2.1 Syntax

NUMLIN’s concrete syntax is inspired by that of OCaml. It desugars (Figure 2) into a smaller, core type and expression grammar (Figure 1).

2.2 Type System and Other Features

The core type theory of NUMLIN is a nearly off-the-shelf linear type theory [3], supporting familiar features such as linear function spaces $A \multimap B$ and tensor products $A \otimes B$. We adopt linearity – the restriction that each program variable be used at most once – since it allows us to express purely functional APIs for numerical library routines that mutate arrays

$$\begin{aligned}
f &::= 'fc \mid \mathbf{z} \mid f \ \mathbf{s} \\
t &::= \mathbf{unit} \mid \mathbf{bool} \mid \mathbf{int} \mid \mathbf{elt} \mid f \ \mathbf{arr} \mid f \ \mathbf{mat} \mid !t \mid 'fc. t \mid t \otimes t' \mid t \multimap t' \\
e &::= p \text{ (primitives)} \mid x \text{ (variable)} \mid \mathbf{let} \ x = e \ \mathbf{in} \ e' \mid () \mid \mathbf{let} \ () = e \ \mathbf{in} \ e' \mid \mathbf{true} \mid \mathbf{false} \\
&\quad \mathbf{if} \ e \ \mathbf{then} \ e_1 \ \mathbf{else} \ e_2 \mid k \text{ (integer)} \mid l \text{ (heap location)} \mid el \text{ (array element)} \\
&\quad \mathbf{Many} \ v \mid \mathbf{let} \ \mathbf{Many} \ x = e \ \mathbf{in} \ e' \mid \mathbf{fun} \ 'fc \rightarrow e \text{ (frac. perm. abstraction)} \\
&\quad e[f] \text{ (frac. perm. specialisation)} \mid (e, e') \mid \mathbf{let} \ (a, b) = e \ \mathbf{in} \ e' \\
&\quad \mathbf{fun} \ x : t \rightarrow e \mid e \ e' \mid \mathbf{fix} \ (g, x : t, e : t')
\end{aligned}$$

■ **Figure 1** Core fraction f , type t and expression e grammar of NUMLIN. Values v are a subset of the expressions, their full definition and a list of all primitives p is in Appendix E.

$$\begin{aligned}
x[e] &\Rightarrow \mathbf{get} \ _ \ x \ (e) && \text{(similarly for matrices)} \\
x[e_1] := e_2 &\Rightarrow \mathbf{set} \ x \ (e_1) \ (e_2) && \text{(similarly for matrices)} \\
pat &::= () \mid x \mid !x \mid \mathbf{Many} \ pat \mid (pat, pat) \\
\mathbf{let} \ !x = e_1 \ \mathbf{in} \ e_2 &\Rightarrow \mathbf{let} \ \mathbf{Many} \ x = e_1 \ \mathbf{in} \\
&\quad \mathbf{let} \ \mathbf{Many} \ x = \mathbf{Many} \ (\mathbf{Many} \ x) \ \mathbf{in} \ e_2 \\
\mathbf{let} \ \mathbf{Many} \ \langle pat_x \rangle = e_1 \ \mathbf{in} \ e_2 &\Rightarrow \mathbf{let} \ \mathbf{Many} \ x = x \ \mathbf{in} \\
&\quad \mathbf{let} \ \langle pat_x \rangle = x \ \mathbf{in} \ e_2 \\
\mathbf{let} \ (\langle pat_a \rangle, \langle pat_b \rangle) = e_1 \ \mathbf{in} \ e_2 &\Rightarrow \mathbf{let} \ (a, b) = a_b \ \mathbf{in} \ \mathbf{let} \ \langle pat_a \rangle = a \ \mathbf{in} \\
&\quad \mathbf{let} \ \langle pat_b \rangle = b \ \mathbf{in} \ e_2 \\
\mathbf{fun} \ (\langle pat_x \rangle : t) \rightarrow e &\Rightarrow \mathbf{fun} \ (x : t) \rightarrow \mathbf{let} \ \langle pat_x \rangle = x \ \mathbf{in} \ e \\
arg &::= \langle pat \rangle : t \mid 'x \text{ (fractional permission variable)} \\
\mathbf{fun} \ \langle arg_{1..n} \rangle \rightarrow e &\Rightarrow \mathbf{fun} \ \langle arg_1 \rangle \rightarrow .. \mathbf{fun} \ \langle arg_n \rangle \rightarrow e \\
\mathbf{let} \ f \ \langle arg_{1..n} \rangle = e_1 \ \mathbf{in} \ e_2 &\Rightarrow \mathbf{let} \ f = \mathbf{fun} \ \langle arg_{1..n} \rangle \rightarrow e_1 \ \mathbf{in} \ e_2 \\
\mathbf{let} \ !f \ \langle arg_{1..n} \rangle = e_1 \ \mathbf{in} \ e_2 &\Rightarrow \mathbf{let} \ \mathbf{Many} \ f = \mathbf{Many} \ (\mathbf{fun} \ \langle arg_{1..n} \rangle \rightarrow e_1) \ \mathbf{in} \ e_2 \\
\mathbf{fixpoint} &\equiv \mathbf{fix} \ (f, x : t, \mathbf{fun} \ \langle arg_{0..n} \rangle \rightarrow e_1 : t') \\
\mathbf{let} \ \mathbf{rec} \ f \ (x : t) \ \langle arg_{0..n} \rangle : t' = e_1 \ \mathbf{in} \ e_2 &\Rightarrow \mathbf{let} \ f = \mathbf{fixpoint} \ \mathbf{in} \ e_2 \\
\mathbf{let} \ \mathbf{rec} \ !f \ (x : t) \ \langle arg_{0..n} \rangle : t' = e_1 \ \mathbf{in} \ e_2 &\Rightarrow \mathbf{let} \ \mathbf{Many} \ f = \mathbf{Many} \ \mathbf{fixpoint} \ \mathbf{in} \ e_2
\end{aligned}$$

■ **Figure 2** Desugaring NUMLIN.

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87 and matrices [19]. Due to linearity, values cannot alias and are only used once, which means
88 that linearly-typed updates result in no *observable* mutation.

89 As a result, programmers can reason about NUMLIN expressions as if they were ordinary
90 mathematical expressions – as indeed they are! We are merely adopting a stricter type
91 discipline than usual to make managing memory safe.

92 2.2.1 Intuitionism: ! and Many

93 However, linearity by itself is not sufficient to produce an expressive enough programming
94 language. For values such as booleans, integers, floating-point numbers as well as pure
95 functions, we need to be able to use them *intuitionistically*, that is, more than once or not at
96 all. For this reason, we have the ! constructor at the type level and its corresponding Many
97 constructor and let Many <id> = .. in .. eliminator at the term level. Because we want
98 to restrict how a programmer can alias pointers and prevent a programmer from ignoring
99 them (a memory leak), NUMLIN enforces simple syntactic restrictions on which values can
100 be wrapped up in a Many constructor (details in Section 3).

101 2.2.2 Fractional Permissions

102 There are also valid cases in which we would want to alias pointers to a matrix. The most
103 common is exemplified by the BLAS routine `gemm`, which (rather tersely) stands for *GEneric*
104 *Matrix Multiplication*. A *simplified* definition of `gemm`(α , A, B, β , C) is $C := \alpha AB + \beta C$.
105 In this case, A and B may alias each other but neither may alias C, because it is being written
106 to. Related to *mutating* arrays and matrices is *freeing* them. Here, we would also wish to
107 restrict aliasing so that we do not free one alias and then attempt to use another. Although
108 linearity on its own suffices to prevent use-after-free errors when values are *not* aliased (a
109 freed value is *out of scope* for the rest of the expression), we still need another simple, yet
110 powerful concept to provide us with the extra expressivity of aliasing *without* losing any of
111 the benefits of linearity.

112 Fractional permissions provide exactly this. Concretely, types of (pointers to) arrays
113 and matrices are *parameterised* by a *fraction*. A fraction is either 1 (2^0) or exactly *half* of
114 another fraction (2^{-k} , for natural k). The former represents complete ownership of that
115 value: the programmer may mutate or free that value as they choose; the latter represents
116 read-only access or a *borrow*: the programmer may read from the value but not write to or
117 free it. Creating an array/matrix gives you ownership of it, so too does having one (with a
118 fractional permission of 2^0) passed in as an argument.

119 In NUMLIN, we can produce two aliases of a single array/matrix, by *sharing* it. If the
120 original alias had a fractional permission of 2^{-k} then the two new aliases of it will have a
121 fractional permission of $2^{-(k+1)}$ each. Thanks to linearity, the original array/matrix with a
122 fractional permission of 2^{-k} will be out of scope after the sharing. When an array/matrix is
123 shared as such, we can prevent the programmer from freeing or mutating it by making the
124 types of `free` and `set` (for mutation) require a *whole* (2^0) permission.

125 If we have two aliases *to the same matrix* with *identical* fractional permissions ($2^{-(k+1)}$),
126 we can recombine or *unshare* them back into a single one, with a larger 2^{-k} permission. As
127 before, thanks to linearity, the original two aliases will be out of scope after unsharing.

128 2.2.3 Runtime Errors

129 Aside from out-of-bounds indexing, matrix unsharing is one of only *two* operations that
130 can fail at runtime (the other being dimension checks, such as for `gemm`). The check being

```

let rec factorial ( !x : !int ) : !int =
  if x < 0 || x = 0 then
    1
  else
    x * factorial (x - 1) in
  factorial ;;

```

■ **Figure 3** Factorial function in NUMLIN.

131 performed is a simple sanity check that the two aliasing pointers passed to `unshare` point to
 132 the same array/matrix. Section 5 contains an overview of how we could remove the need for
 133 this by tracking pointer identities statically by augmenting the type system further.

134 2.2.4 Recursion

135 The final feature of NUMLIN which makes it sufficiently expressive is recursion (and of
 136 course, conditional branches to ensure termination). Conditional branches are implemented
 137 by ensuring that both branches use the same set of linear values. A function can be recursive
 138 if it captures no linear values from its environment. Like with `Many`, this is enforced via
 139 simple syntactic restrictions on the definition of recursive functions.

140 2.3 Examples

141 2.3.1 Factorial

142 Although a factorial function (Figure 3) may seem like an aggressively pedestrian first example,
 143 in a linearly typed language such as NUMLIN it represents the culmination of many features.

144 To simplify the design and implementation of NUMLIN’s type system, recursive functions
 145 must have full type annotations (non-recursive functions need only their argument types
 146 annotated). Its body is a closed expression (with respect to the function’s arguments), so it
 147 type-checks (since it does not capture any linear values from its environment).

148 The only argument is `!x : !int`. The `!` annotation on `x` is a syntactic convenience for
 149 declaring the value to be used intuitionistically, its full and precise meaning is described in
 150 Section 3.1.

151 The condition for an `if` may or may not use linear values (here, with `x < 0 || x = 0`, it
 152 does not). Any linear values used by the condition would not be in scope in either branch of
 153 the `if`-expression. Both branches use `x` differently: one ignores it completely and the other
 154 uses it twice.

155 All numeric and boolean literals are implicitly wrapped in a `Many` and all primitives
 156 involving them return a `!int`, `!bool` or `!elt` (types of elements of arrays/matrices, typically
 157 64-bit floating-point numbers). The short-circuiting `||` behaves in exactly the same way as a
 158 boolean-valued `if`-expression.

159 2.3.2 Summing over an Array

160 Now we can add fractional permissions to the mix: Figure 4 shows a simple, tail-recursive
 161 implementation of summing all the elements in an array. There are many new features; first
 162 among them is `!x0 : !elt`, the type of array/matrix elements (64-bit floating point).

163 Second is `(’x) (row: ’x arr)` which is an array with a universally-quantified fractional
 164 permission. In particular, this means the body of the function cannot mutate or free the

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```

let rec sum_array (!i : !int) (!n : !int) (!x0 : !elt)
  ('x) (row : 'x arr) : 'x arr * !elt =
  if i = n then
    (row, x0)
  else
    let (row, !x1) = row[i] in
    sum_array (i + 1) n (x0 +. x1) 'x row in
  sum_array ;;

```

■ **Figure 4** Summing over an array in NUMLIN.

```

let row = row[i] := x1 in (* or *) let () = free row in
(* Could not show equality: *)
(*      z arr *)
(* with *)
(*      'x arr *)
(* *)
(* Var 'x is universally quantified *)
(* Are you trying to write to/free/unshare an array you don't own? *)
(* In examples/sum_array.lt, at line: 7 and column: 19 *)

```

■ **Figure 5** Attempting to write to or free a read only array in NUMLIN.

165 input array, only read from it. If the programmer did try to mutate or free `row`, then they
 166 would get a helpful error message (Figure 5).

167 Alongside taking a `row : 'x arr`, the function also returns an array with exactly the
 168 same fractional permission as the `row` (which can only be `row`). This is necessary because of
 169 linearity: for the caller, the original array passed in as an argument would be out of scope
 170 for the rest of the expression, so it needs to be returned and then rebound to be used for the
 171 rest of the function.

172 An example of this consuming and re-binding is in `let (row, !x1) = row[i]`. Indexing
 173 is implemented as a primitive `get : 'x. 'x arr \multimap !int \multimap 'x arr \otimes !elt`. Although fractional
 174 permissions can be passed around explicitly (as done in the recursive call), they can also
 175 be *automatically inferred at call sites*: `row[i] == get _ row i` takes advantage of this
 176 convenience.

177 2.3.3 One-dimensional Convolution

178 Figure 6 extends the set of features demonstrated by the previous examples by mutating one
 179 of the input arrays. A one-dimensional convolution involves two arrays: a read-only kernel
 180 (array of weights) and an input vector. It modifies the input vector *in-place* by replacing
 181 each `write[i]` with a weighted (as per the values in the kernel) sum of it and its neighbours;
 182 intuitively, sliding a dot-product with the kernel across the vector.

183 What's implemented in Figure 6 is a *simplified* version of this idea, so as to not distract
 184 from the features of NUMLIN. The simplifications are:

- 185 ■ the kernel has a length 3, so only the value of `write[i-1]` (prior to modification in the
 186 previous iteration) needs to be carried forward using `x0`
- 187 ■ `write` is assumed to have length `n+1`
- 188 ■ `i`'s initial value is assumed to be 1
- 189 ■ `x0`'s initial value is assumed to be `write[0]`
- 190 ■ the first and last values of `write` are ignored.

```

let rec simp_oned_conv
  (!i : !int) (!n : !int) (!x0 : !elt)
  (write : z arr) ('x) (weights : 'x arr)
  : 'x arr * z arr =
  if n = i then (weights, write) else
  let !w0 <- weights[0] in
  let !w1 <- weights[1] in
  let !w2 <- weights[2] in
  let !x1 <- write[i] in
  let !x2 <- write[i + 1] in
  let written = write[i] := w0 *. x0 +. (w1 *. x1 +. w2 *. x2) in
  simp_oned_conv (i + 1) n x1 written _ weights in
simp_oned_conv ;;

```

■ **Figure 6** *Simplified* one-dimensional convolution.

```

symm  : !bool → !elt → 'x. 'x mat → 'y. 'y mat → !elt → z mat →
        ('x mat ⊗ 'y mat) ⊗ z mat
gemm  : !elt → 'x. 'x mat ⊗ !bool → 'y. 'y mat ⊗ !bool → !elt →
        z mat → ('x mat ⊗ 'y mat) ⊗ z mat
gesv  : z mat → z mat → z mat ⊗ z mat
posv  : z mat → z mat → z mat ⊗ z mat
potrs : 'x. 'x mat → z mat → 'x mat ⊗ z mat
syrrk : !bool → !elt → 'x. 'x mat → !elt → z mat → 'x mat ⊗ z mat

```

■ **Figure 7** Types of some NUMLIN primitives.

191 Mutating an array is implemented similarly to indexing one – a primitive `set : z arr →`
 192 `!int → !elt → z arr`. It consumes the original array and returns a new array with the
 193 updated value. `let written = write[i] := <exp>` is just syntactic sugar for `let written`
 194 `= set write i <exp>`.

195 Since `write : z arr` (where `z` stands for $k = 0$, representing a fractional permission of
 196 $2^{-k} = 2^{-0} = 1$), we may mutate it, but since we only need to read from `weights`, its
 197 fractional permission index can be universally-quantified. In the recursive call, we see `_` being
 198 used explicitly to tell the compiler to *infer* the correct fractional permission based on the
 199 given arguments.

200 2.3.4 Digression: Types of Primitives

201 *The most pertinent aspect of NUMLIN is the types of its primitives* (Figure 7). While the types
 202 of operations such as `get` and `set` might be borderline obvious, the types of BLAS/LAPACK
 203 routines become an *incredibly useful, automated check for using the API correctly*.

204 We determine the types for these routines by consulting their documentation. Each
 205 routine has a record of the expected aliasing behaviour and whether or not it modifies or
 206 consumes its argument in any way. We use that to derive the types in Figure 7. Since
 207 most of these low-level routines are very careful not to do any allocation themselves, it is
 208 generally very easy to give each a NUMLIN type – every argument that can modify/consume

```

let !square ('x) (x : 'x mat) =
  let (x, (!m, !n)) = sizeM _ x in
  let (x1, x2) = shareM _ x in
  let answer <- new (m, n) [| x1 * x2 |] in
  let x = unshareM _ x1 x2 in
  (x, answer) in
square
;;

```

■ **Figure 8** Linear regression (OLS): $\hat{\beta} = (\mathbf{X}^T \mathbf{X})^{-1} \mathbf{X}^T \mathbf{y}$

its argument needs a full permission, and all others can be fraction-polymorphic. Taking Fortran as an example, it has a notion of `in`, `out` and `inout` parameters. The latter two would need full `z`permissions; the first would be fraction-polymorphic.

2.3.5 Squaring a Matrix

Figure 8 shows how a linearly-typed matrix squaring function may be written in NUMLIN. It is a *non-recursive* function declaration (the return type is inferred). Since we would like to be able to use a function like `square` more than once, it is marked with a `!` annotation (which also ensures it captures no linear values from the surrounding environment).

To square a matrix, first, we extract the dimensions of the argument `x`. Then, because we need to use `x` twice (so that we can multiply it by itself) but linearity only allows one use, we use `shareM`: $'x. 'x \text{ mat} \multimap 'x \text{ s mat} \otimes 'x \text{ s mat}$ to split the permission `'x` (which represents 2^{-x}) into two halves (`'x s`, which represents $2^{-(x+1)}$).

Even if `x` had type `z mat`, sharing it now enforces the assumption of all BLAS/LAPACK routines that any matrix which is written to (which, in NUMLIN, is always of type `z mat`) does not alias any other matrix in scope. So if we did try to use one of the aliases in mutating way, the expression would not type check, and we would get an error similar to the one in Figure 5.

The line `let answer <- new (m,n) [| x1 * x2 |]` is syntactic sugar for first creating a new $m \times n$ matrix (`let answer = matrix m n`) and then storing the result of the multiplication in it (`let ((x1, x2), answer) = gemm 1. _ (x1, false) _ (x2, false) 0. answer`). `false` means the matrix should not be accessed with indices transposed.

By using some simple pattern-matching and syntactic sugar (Figure 9), we can:

- write normal-looking, apparently non-linear code
- use matrix expressions directly and have a call to an efficient call to a BLAS/LAPACK routine inserted with appropriate re-bindings
- retain the safety of linear types with fractional permissions by having the compiler statically enforce the aliasing and read/write rules implicitly assumed by BLAS/LAPACK routines.

2.3.6 Linear Regression

In Figure 10, we wish to compute $\hat{\beta} = (\mathbf{X}^T \mathbf{X})^{-1} \mathbf{X}^T \mathbf{y}$. To do that, first, we extract the dimensions of matrix `x`. Then, we say we would like `xy` to be a new matrix, of dimension $m \times 1$, which contains the result of $\mathbf{X}^T \mathbf{y}$ (using syntactic sugar for `matrix` and `gemm` calls similar to that used in Figure 8, with a `^T` annotation on `x` to set `x`'s 'transpose indices'-flag to `true`).

$$\begin{aligned} \text{let } v \leftarrow x[e] \text{ in } e &\Rightarrow \text{let } (x, !v) = x[e] \text{ in } e && \text{(similarly for matrices)} \\ \text{let } x_2 \leftarrow \text{new } [| x_1 |] \text{ in } e &\Rightarrow \text{let } (x_1, x_2) = \text{copyM_} x_1 \text{ in } e \\ \text{let } x_2 \leftarrow [| x_1 |] \text{ in } e &\Rightarrow \text{let } (x_1, x_2) = \text{copyM_to_} x_1 \ x_2 \text{ in } e \\ M ::= X \mid X^T \mid \text{sym}(X) \\ \\ \text{let } Y \leftarrow \text{new } (n, k) [| \alpha M_1 M_2 |] \text{ in } e &\Rightarrow \\ \text{let } Y = \text{matrix } n \ k \text{ in let } Y \leftarrow [| \alpha M_1 M_2 + 0Y |] \text{ in } e \\ \text{let } Y \leftarrow [| \alpha X X^T + \beta Y |] \text{ in } e &\Rightarrow \\ \text{let } (X, Y) = \text{syrk false } \alpha _ X \ \beta \ Y \text{ in } e \\ \text{let } Y \leftarrow [| \alpha X^T X + \beta Y |] \text{ in } e &\Rightarrow \\ \text{let } (X, Y) = \text{syrk true } \alpha _ X \ \beta \ Y \text{ in } e \\ \text{let } Y \leftarrow [| \alpha \text{sym}(X_1) X_2 + \beta Y |] \text{ in } e &\Rightarrow \\ \text{let } ((X_1, X_2), Y) = \text{symm false } \alpha _ X_1 _ X_2 \ \beta \ Y \text{ in } e \\ \text{let } Y \leftarrow [| \alpha X_2 \text{sym}(X_1) + \beta Y |] \text{ in } e &\Rightarrow \\ \text{let } ((X_1, X_2), Y) = \text{symm true } \alpha _ X_1 _ X_2 \ \beta \ Y \text{ in } e \\ \text{let } Y \leftarrow [| \alpha X_1^{T?} X_2^{T?} + \beta Y |] \text{ in } e &\Rightarrow \\ \text{let } ((X_1, X_2), Y) = \text{gemm } \alpha _ (X_1, \text{true}_{\text{false}}) _ (X_2, \text{true}_{\text{false}}) \ \beta \ Y \text{ in } e \end{aligned}$$

■ **Figure 9** Purely syntactic pattern-matching translations of matrix expressions.

243 However, the line `let x_T_x <- new (m,m) [| x^T * x |]`, works for a slightly differ-
 244 ent reason: that pattern is matched to a BLAS call to `(syrk true 1. x 0. x_T_x)`, which
 245 only uses `x` once. Hence `x` can appear *twice* in the *pattern* without any calls to `share`.

246 After computing `x_T_x`, we need to invert it and then multiply it by `xy`. The BLAS
 247 routine `posv : z mat \rightarrow z mat \rightarrow z mat \otimes z mat` does exactly that: assuming the first
 248 argument is symmetric, `posv` mutates its second argument to contain the desired value. Its
 249 first argument is also mutated to contain the (upper triangular) Cholesky decomposition
 250 factor of the original matrix. Since we do not need that matrix (or its memory) again, we
 251 `free` it. If we forgot to, we would get a `Variable to_del not used` error. Lastly, we return
 252 the `answer` alongside the untouched input matrices `(x,y)`.

253 2.3.7 L1-norm Minimisation on Manifolds

254 L1-norm minimisation is often used in optimisation problems, as a *regularisation* term for
 255 reducing the influence of outliers. Although the below formulation [11] is intended to be used
 256 with *sparse* computations, NUMLIN's current implementation only implements dense ones.
 257 However, it still serves as a useful example of explaining NUMLIN's features.

258 Figure 11 shows even more pattern-matching. Patterns of the form `let <id> <- [|`
 259 `beta * c + alpha * a * b |]` are also desugared to `gemm` calls. Primitives like `transpose :`
 260 `'x. 'x mat \rightarrow 'x mat \otimes z mat` and `eye : 'int \rightarrow z mat` allocate new matrices; `transpose`
 261 returns the transpose of a given matrix and `eye k` evaluates to a $k \times k$ identity matrix.

262 We also see our first example of re-using memory for different matrices: like with `to_del`

```

let !lin_reg ('x) (x : 'x mat)
    ('y) (y : 'y mat) =
  let (x, (!_n, !m)) = sizeM _ x in
  let xy <- new (m, 1) [| xT * y |] in
  let x_T_x <- new (m, m) [| xT * x |] in
  let (to_del, answer) = posv x_T_x xy in
  let () = freeM to_del in
  ((x, y), answer) in
lin_reg ;;

```

■ Figure 10 Linear regression (OLS): $\hat{\beta} = (X^T X)^{-1} X^T y$

```

let !l1_norm_min (q : z mat) (u : z mat) =
  let (u, (!_n, !k)) = sizeM _ u in
  let (u, u_T) = transpose _ u in
  let (tmp_n_n, q_inv_u) = gesv q u in
  let i = eye k in
  let to_inv <- [| i + u_T * q_inv_u |] in
  let (tmp_k_k, inv_u_T) = gesv to_inv u_T in
  let () = freeM tmp_k_k in
  let answer <- [| 0. * tmp_n_n + q_inv_u * inv_u_T |] in
  let () = freeM q_inv_u in
  let () = freeM inv_u_T in
  answer in
l1_norm_min ;;

```

■ Figure 11 L1-norm minimisation on manifolds: $Q^{-1}U(I + U^T Q^{-1}U)^{-1}U^T$

263 and `posv` in the previous example, we do not need the value stored in `tmp_5_5` after the call
 264 to `gesv` (a primitive similar to `posv` but for a non-symmetric first argument). However, we
 265 can re-use its memory much later to store `answer` with `let answer <- [| 0. * tmp_5_5`
 266 `+ q_inv_u * inv_u_T |]`. Again, thanks to linearity, the identifiers `q` and `tmp_5_5` are out
 267 of scope by the time `answer` is bound. Although during execution, all three refer to the same
 268 piece of memory, logically they represent different values throughout the computation.

269 2.3.8 Kalman Filter

270 A *Kalman Filter* [15] is an algorithm for combining prior knowledge of a state, a statistical
 271 model and measurements from (noisy) sensors to produce an estimate a more reliable
 272 estimated of the current state. It has various applications (navigation, signal-processing,
 273 econometrics) and is relevant here because it is usually presented as a series of complex
 274 matrix equations.

275 Figure 12 shows a NUMLIN implementation of a Kalman filter (equations in Figure 13).
 276 A few new features and techniques are used in this implementation:

- 277 ■ `sym` annotations in matrix expressions: when this is used, a call to `symm` (the equivalent
 278 of `gemm` but for symmetric matrices so that only half the operations are performed) is
 279 inserted
- 280 ■ `copyM_to` is used to re-use memory by *overwriting* the contents of its second argument
 281 to that of its first (erroring if dimensions do not match)
- 282 ■ `let new_r <- new [| r_2 |]` creates a copy of `r_2`
- 283 ■ `posvFlip` is like `posv` except for solving $XA = B$
- 284 ■ a lot of memory re-use; the following sets of identifiers alias each other:

```

let !kalman
  ('s) (sigma : 's mat) (* n,n *)
  ('h) (h : 'h mat)      (* k,n *)
  (mu : z mat)            (* n,1 *)
  (r_1 : z mat)          (* k,k *)
  (data_1 : z mat)       (* k,1 *) =
  let (h, (!k, !n)) = sizeM _ h in
  (* could use [| sym(sigma) * hT |] but would
     need a (n,k) temporary hT = transpose _ h *)
  let sigma_hT <- new (n, k) [| sigma * h^T |] in
  let r_2 <- [| r_1 + h * sigma_hT |] in
  let (k_by_k, x) = posvFlip r_2 sigma_hT in
  let data_2 <- [| h * mu - data_1 |] in
  let new_mu <- [| mu + x * data_2 |] in
  let x_h <- new (n,n) [| x * h |] in
  let () = freeM (* n,k *) x in
  let sigma2 <- new [| sigma |] in
  let new_sigma <- [| sigma2 - x_h * sym(sigma) |] in
  let () = freeM (* n,n *) x_h in
  ((sigma, h), (new_sigma, (new_mu, (k_by_k, data_2)))) in
kalman ;;

```

■ **Figure 12** Kalman filter: see Figure 13 for the equations this code implements and Figure 20 for an equivalent CBLAS/LAPACKE implementation.

$$\mu' = \mu + \Sigma H^T (R + H \Sigma H^T)^{-1} (H \mu - \text{data})$$

$$\Sigma' = \Sigma (I - H^T (R + H \Sigma H^T)^{-1} H \Sigma)$$

■ **Figure 13** Kalman filter equations (credit: matthewrocklin.com).

285 — r_1, r_2 and k_by_k
 286 — data_1 and data_2
 287 — mu and new_mu
 288 — sigma_hT and x

289 The NUMLIN implementation is much longer than the mathematical equations for two
 290 reasons. First, the NUMLIN implementation is a let-normalised form of the Kalman equations:
 291 since there a large number of unary/binary (and occasionally ternary) sub-expressions in
 292 the equations, naming each one line at a time makes the implementation much longer.
 293 Second, NUMLIN has the additional task of handling explicit allocations, aliasing and frees
 294 of matrices. However, it is exactly this which makes it possible (and often, easy) to spot
 295 additional opportunities for memory re-use. Furthermore, a programmer can explore those
 296 opportunities easily because NUMLIN's type system statically enforces correct memory
 297 management and the aliasing assumptions of BLAS/LAPACK routines.

298 3 Formal System

299 3.1 Core Type Theory

300 The full typing rules are in Appendix A.1, but the key ideas are as follow.

301 ■ A typing judgement consists of $\Theta; \Delta; \Gamma \vdash e : t$.

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- 302 ■ Θ is the environment that tracks which fractional permission variables in scope. Fractional
- 303 permissions (the `Perm` judgement) and types (the `Type` judgement) are *well-formed* if all
- 304 of their free fractional variables are in Θ .
- 305 ■ Δ is the environment storing non-linearly or *intuitionistically* typed variables.
- 306 ■ Γ is the environment storing linearly typed variables.

307 Note that rules for typing `()`, booleans, integers and elements are typed with respect to
 308 an *empty* linear environment: this means no linear values are needed to produce a value of
 309 those types.

$$310 \quad \frac{}{\Theta; \Delta; \cdot \vdash () : \mathbf{unit}} \text{TY_UNIT_INTRO}$$

311 Conversely, whenever two or more subexpressions need to be typed, they must consume
 312 a disjoint set of linear values (pairs, let-expressions). In the case of if-expressions, both
 313 branches must consume the same set of linear values (disjoint to the ones used to evaluate
 314 the condition).

$$315 \quad \frac{\begin{array}{l} \Theta; \Delta; \Gamma \vdash e : !\mathbf{bool} \\ \Theta; \Delta; \Gamma' \vdash e_1 : t' \\ \Theta; \Delta; \Gamma' \vdash e_2 : t' \end{array}}{\Theta; \Delta; \Gamma, \Gamma' \vdash \mathbf{if } e \mathbf{ then } e_1 \mathbf{ else } e_2 : t} \text{TY_BOOL_ELIM}$$

316 The **Many** introduction and elimination rules are very important. Producing `!`-type values
 317 may only be done if the expression inside is a syntactic value which is not a location. This
 318 allows all safely duplicable resources, including functions which capture non-linear resources
 319 from their environments, but prevents producing aliases of (pointers to) arrays and matrices.
 320 This is exactly the same as value-restriction from the world of parametric polymorphism;
 321 without it, the expression `let Many x = Many (array 5) in let () = free x in x[0]` would
 322 type-check but error at runtime.

$$323 \quad \frac{\begin{array}{l} \Theta; \Delta; \cdot \vdash v : t \\ v \neq l \end{array}}{\Theta; \Delta; \cdot \vdash \mathbf{Many } v : !t} \text{TY_BANG_INTRO}$$

324 Consuming a variable that refers to a `!`-type value *moves it* from the linear environment
 325 Γ and *into* the intuitionistic environment Δ .

$$326 \quad \frac{\begin{array}{l} \Theta; \Delta; \Gamma \vdash e : !t \\ \Theta; \Delta, x : t; \Gamma' \vdash e' : t' \end{array}}{\Theta; \Delta; \Gamma, \Gamma' \vdash \mathbf{let Many } x = e \mathbf{ in } e' : t'} \text{TY_BANG_ELIM}$$

327 Using this, we can explain how the `!` annotation on variables – first introduced in the
 328 factorial example in 2.3.1 – works. That is, we can explain why the meaning of `let !x = e in e'`
 329 can be expressed using only the rules presented thus far, as `let Many x = e in let Many x =`
 330 `Many (Many x) in e'`.¹ The reader is invited to quickly convince themselves that the

¹ Why we have this at all is for the sake of ergonomics when using binary arithmetic operations (e.g. of type `!int ⊗ !int → !int`): a programmer should be able to write `let x = 5 + 5 in x - x`, which, although non-linear in `x`, is morally right because integers and operations on them rarely need to be linear. Though it should be possible to handle this using a LNL-style presentation of linear types [5] (using adjoint modalities to distinguish between intrinsically linear and intrinsically intuitionistic types) that is a pretty big digression from the stated goals of this paper.

331 following meta-rule is provable using `TY_BANG_INTRO` (twice), `TY_BANG_ELIM` (twice)
 332 and weakening the intuitionistic environment Δ (once).

$$333 \quad \frac{\begin{array}{c} \Theta; \Delta; \Gamma \vdash e : !t \\ \Theta; \Delta, x : !t; \Gamma' \vdash e' : t' \end{array}}{\Theta; \Delta; \Gamma, \Gamma' \vdash \text{let } !x = e \text{ in } e' : t'} \quad \text{META_TY_LET_BANG}$$

334 Rules `TY_GEN` and `TY_SPC` are for fractional permission generalisation and specialisa-
 335 tion respectively. They allow the definition and use of functions that are polymorphic in the
 336 fractional permission index of their results and one or more of their arguments.

$$337 \quad \frac{\Theta, fc; \Delta; \Gamma \vdash e : t}{\Theta; \Delta; \Gamma \vdash \text{fun } 'fc \rightarrow e : 'fc.t} \quad \text{TY_GEN} \qquad \frac{\begin{array}{c} \Theta \vdash f \text{ Perm} \\ \Theta; \Delta; \Gamma \vdash e : 'fc.t \end{array}}{\Theta; \Delta; \Gamma \vdash e[f] : t[f/fc]} \quad \text{TY_SPC}$$

338 Rule `TY_FIX` shows how recursive functions are typed. Even though recursive functions
 339 are fully annotated, type checking them is interesting for two reasons: to type check the
 340 body of the fixpoint, the type of the recursive function is in the *intuitionistic* environment Δ
 341 (without this, you would not be able to write a base case) whilst the argument and its type
 342 are the *only things in the linear environment* Γ . The latter means that recursive functions
 343 can be type checked in an empty environment (thus be wrapped in `Many` and used zero or
 344 multiple times).

$$345 \quad \frac{\Theta; \Delta, g : t \multimap t'; \cdot, x : t \vdash e : t'}{\Theta; \Delta; \cdot \vdash \text{fix } (g, x : t, e : t') : t \multimap t'} \quad \text{TY_FIX}$$

346 Lastly, types of almost all `NUMLIN` primitives, as embedded in OCaml's type system, are
 347 shown in Appendix F, with some similar ones (like those for binary arithmetic operators)
 348 omitted for brevity. The main difference between the OCaml type of a primitive like `gemm`
 349 and its `NUMLIN` counterpart (Figure 7) is the inclusion of explicit universal-quantification of
 350 fractional permission variables in the latter.

351 3.2 Dynamic Semantics

352 The full, small-step transition relation is in Appendix A.2, but the key ideas are as follow.

353 Heaps (σ) are multisets containing triples of an abstract location l , a fractional permission
 354 f and sized matrices $m_{n,k}$. The notation $l \mapsto_f m_{k_1, k_2}$ should be read as “location l represents
 355 f ownership over matrix m (of size $k_1 \times k_2$)”. Each heap-and-expression either steps to
 356 another heap-and-expression or a runtime error `err`. In the full grammar definition we see a
 357 definition of values and contexts in the language.

358 We draw the reader's attention to the definitions relating to fractional permissions.
 359 Specifically, unlike a lambda, the body of a `fun` $fc \rightarrow v$ must be a syntactic value. The
 360 context `fun` $fc \rightarrow [-]$ means expressions can be reduced inside a fractional permission
 361 generalisation. This is to emphasize that fractions are merely *compile-time constructs* and
 362 do not affect runtime behaviour. Correct usage of fractions is enforced by the type system,
 363 so programs do not get stuck. Fractional permissions are specialised using substitution over
 364 both the heap and an expression (`OP_FRAC_PERM`).

$$365 \quad \overline{\langle \sigma, (\text{fun } 'fc \rightarrow v)[f] \rangle \rightarrow \langle \sigma[f/fc], v[f/fc] \rangle} \quad \text{OP_FRAC_PERM}$$

366 Like with the static semantics, the interesting rules in the dynamic semantics are those
 367 relating to primitives. Creating a matrix (`matrix` $k_1 \ k_2$) successfully (`OP_MATRIX`) requires

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non-negative dimensions and returns a (fresh) location of a matrix of those dimensions, extending the heap to reflect that l represents a complete ownership over the new matrix.

$$\frac{0 \leq k_1, k_2 \quad l \text{ fresh}}{\langle \sigma, \mathbf{matrix} \ k_1 \ k_2 \rangle \rightarrow \langle \sigma + \{l \mapsto_1 M_{k_1, k_2}\}, l \rangle} \quad \text{OP_MATRIX}$$

Dually, OP_FREE, requires a location represent complete ownership before removing it and the matrix it points to from the heap.

$$\frac{}{\langle \sigma + \{l \mapsto_1 m_{k_1, k_2}\}, \mathbf{free} \ l \rangle \rightarrow \langle \sigma, () \rangle} \quad \text{OP_FREE}$$

Choosing a multiset representation as opposed to a set allows for two convenient invariants: multiplicity of a triple $l \mapsto_f m_{k_1, k_2}$ in the heap corresponds to the number of aliases of l in the expression with type f **mat** and the sum of all the fractions for l will always be 1 (for a closed, well-typed expression). With this in mind, the rules OP_SHARE and OP_UNSHARE_EQ are fairly natural.

$$\frac{}{\langle \sigma + \{l \mapsto_f m_{k_1, k_2}\}, \mathbf{share}[f] \ l \rangle \rightarrow \langle \sigma + \{l \mapsto_{\frac{1}{2}f} m_{k_1, k_2}\} + \{l \mapsto_{\frac{1}{2}f} m_{k_1, k_2}\}, (l, l) \rangle} \quad \text{OP_SHARE}$$

$$\frac{\sigma' \equiv \sigma + \{l \mapsto_{\frac{1}{2}f} m_{k_1, k_2}\} + \{l \mapsto_{\frac{1}{2}f} m_{k_1, k_2}\}}{\langle \sigma', \mathbf{unshare}[f] \ l \ l \rangle \rightarrow \langle \sigma + \{l \mapsto_f m_{k_1, k_2}\}, l \rangle} \quad \text{OP_UNSHARE_EQ}$$

Combining all of these features, we see that OP_GEMM_MATCH requires that the location being updated (l_3) has complete ownership of over matrix m_3 and can thus change what value it stores to $m_1 m_2 + m_3$. In particular, this places no restriction on l_2 and l_3 : they could be **shared** aliases of the same matrix. Transition rules for other primitives (omitted) follow the same structure: \mapsto_1 for any locations that are written to and \mapsto_{fc} for anything else.

$$\frac{\begin{aligned} \sigma' &\equiv \sigma + \{l_1 \mapsto_{fc_1} m_{1 \ k_1, k_2}\} + \{l_2 \mapsto_{fc_2} m_{2 \ k_2, k_3}\} \\ \sigma_1 &\equiv \sigma' + \{l_3 \mapsto_1 m_{3 \ k_1, k_3}\} \\ \sigma_2 &\equiv \sigma' + \{l_3 \mapsto_1 (m_1 \ m_2 + m_3)_{k_1, k_3}\} \end{aligned}}{\langle \sigma_1, \mathbf{gemm}[fc_1] \ l_1 [fc_2] \ l_2 \ l_3 \rangle \rightarrow \langle \sigma_2, ((l_1, l_2), l_3) \rangle} \quad \text{OP_GEMM_MATCH}$$

3.3 Logical Relation

First, we define an interpretation of heaps with fractional permissions in the style of Bornat et. al [9] (interpreting the multiset as a partial map from locations to the sum of all its associated fractions and a matrix) as well as the n-fold iteration of \rightarrow .

$$\mathcal{H}[\sigma] = \star_{(l, f, m) \in \sigma} [l \mapsto_f m]$$

where

$$(\varsigma_1 \star \varsigma_2)(l) \equiv \begin{cases} \varsigma_1(l) & \text{if } l \in \text{dom}(\varsigma_1) \wedge l \notin \text{dom}(\varsigma_2) \\ \varsigma_2(l) & \text{if } l \in \text{dom}(\varsigma_2) \wedge l \notin \text{dom}(\varsigma_1) \\ (f_1 + f_2, m) & \text{if } (f_1, m) = \varsigma_1(l) \wedge (f_2, m) = \varsigma_2(l) \wedge f_1 + f_2 \leq 1 \\ \text{undefined} & \text{otherwise} \end{cases}$$

We then define a step-indexed logical relation in the style of Morrisett et. al [2]. $(\varsigma, v) \in \mathcal{V}_k[t]$ means it takes a heap with exactly ς resources to produce a value v of type t in at

most k steps. So, something like a **unit** or a $!t$ need no resources, whereas a f **mat** needs exactly f ownership of a some matrix and a pair needs a \star combination of the heaps required for each component.

$$\begin{aligned} \mathcal{V}_k[\mathbf{unit}] &= \{(\emptyset, *)\} \\ \mathcal{V}_k[f \mathbf{mat}] &= \{(\{l \mapsto_{2-f} _ \}, l)\} \\ \mathcal{V}_k[!t] &= \{(\emptyset, \mathbf{Many} \ v) \mid (\emptyset, v) \in \mathcal{V}_k[t]\} \\ \mathcal{V}_k[t_1 \otimes t_2] &= \{(\varsigma_1 \star \varsigma_2, (v_1, v_2)) \mid (\varsigma_1, v_1) \in \mathcal{V}_k[t_1] \wedge (\varsigma_2, v_2) \in \mathcal{V}_k[t_2]\} \end{aligned}$$

The definition of $\mathcal{V}_k[fc. t]$ says a value and heap must be the same regardless of what fraction is substituted into both; the $k - 1$ is to take into account fraction specialisation takes ones step (OP_SPC).

$$\mathcal{V}_k[fc. t] = \{(\varsigma, \mathbf{fun} \ fc \rightarrow v) \mid \forall f. (\varsigma[f/fc], v[f/fc]) \in \mathcal{V}_{k-1}[t[f/fc]]\}$$

To understand the definition of $\mathcal{V}_k[t' \multimap t]$, we must first look at $\mathcal{C}_k[t]$, the computational interpretation of types. Intuitively, it is a combination of a frame rule on heaps (no interference), type-preservation and termination (in $j < k$ steps) to either an error or a heap-and-expression, with the further condition that if the expression is a syntactic value then it is also one semantically.

$$\begin{aligned} \mathcal{C}_k[t] &= \{(\varsigma_s, e_s) \mid \forall j < k, \sigma_r. \varsigma_s \star \sigma_r \text{ defined} \Rightarrow \langle \sigma_s + \sigma_r, e_s \rangle \rightarrow^j \mathbf{err} \vee \exists \sigma_f, e_f. \\ &\quad \langle \sigma_s + \sigma_r, e_s \rangle \rightarrow^j \langle \sigma_f + \sigma_r, e_f \rangle \wedge (e_f \text{ is a value} \Rightarrow (\varsigma_f \star \sigma_r, e_f) \in \mathcal{V}_{k-j}[t])\} \end{aligned}$$

In this light, $\mathcal{V}_k[t' \multimap t]$ simply says that v is a function and that evaluating the application of it to any argument (of the correct type, requiring its own set of resources, bounded by k steps) satisfies all the aforementioned properties.

$$\begin{aligned} \mathcal{V}_k[t' \multimap t] &= \{(\varsigma_v, v) \mid (v \equiv \mathbf{fun} \ x : t' \rightarrow e \vee v \equiv \mathbf{fix}(g, x : t', e : t)) \wedge \\ &\quad \forall j \leq k, (\varsigma_{v'}, v') \in \mathcal{V}_j[t']. \varsigma_v \star \varsigma_{v'}' \text{ defined} \Rightarrow (\varsigma_v \star \varsigma_{v'}', v v') \in \mathcal{C}_j[t]\} \end{aligned}$$

The interpretation of typing environments Δ and Γ are with respect to an arbitrary substitution of fractional permissions θ . Note that only the interpretation of Γ involves a (potentially) non-empty heap.

$$\begin{aligned} \mathcal{I}_k[\Delta, x : t]\theta &= \{\delta[x \mapsto v_x] \mid \delta \in \mathcal{I}_k[\Delta]\theta \wedge (\emptyset, v_x) \in \mathcal{V}_k[\theta(t)]\} \\ \mathcal{L}_k[\Gamma, x : t]\theta &= \{(\varsigma \star \varsigma_x, \gamma[x \mapsto v_x]) \mid (\varsigma, \gamma) \in \mathcal{L}_k[\Gamma]\theta \wedge (\varsigma_x, v_x) \in \mathcal{V}_k[\theta(t)]\} \end{aligned}$$

And so the final semantic interpretation of a typing judgement simply quantifies over all possible fractional permission substitutions θ , linear value substitutions γ , intuitionistic value substitutions δ and heaps σ . Note that, $\varsigma \equiv \mathcal{H}[\theta(\sigma)]$.

$$\begin{aligned} {}_k[\Theta; \Delta; \Gamma \vdash e : t] &= \forall \theta, \delta, \gamma, \sigma. \Theta = \text{dom}(\theta) \wedge (\varsigma, \gamma) \in \mathcal{L}_k[\Gamma]\theta \wedge \delta \in \mathcal{I}_k[\Delta]\theta \Rightarrow \\ &\quad (\varsigma, \theta(\delta(\gamma(e)))) \in \mathcal{C}_k[\theta(t)] \end{aligned}$$

3.4 Soundness Theorem

► **Theorem 1.** (*The Fundamental Lemma of Logical Relations*)

$$\forall \Theta, \Delta, \Gamma, e, t. \Theta; \Delta; \Gamma \vdash e : t \Rightarrow \forall k. {}_k[\Theta; \Delta; \Gamma \vdash e : t]$$

438 3.4.1 Explanation

439 If an expression is syntactically type-checked, then for an arbitrary number of steps k , under
 440 any substitution θ of free fractional-permission variables, linear variables γ (with a suitable
 441 heap ς) and intuitionistic variables δ , that suitable heap and expression $(\varsigma, \theta(\delta(\gamma(e))))$ are in
 442 the computational interpretation of the type of that expression $\mathcal{C}_k[\![\theta(t)]\!]$.

443 The *computational interpretation* is as defined before: a combination of a frame rule on
 444 heaps (no interference), type-preservation and termination (in $j < k$ steps) to either an error
 445 or a heap-and-expression, with the further condition that if the expression is a syntactic
 446 value then it is also one semantically.

447 It identifies executions that do no un- or ill-defined behaviours (e.g. adding a boolean and
 448 an integer). Since our operational semantics explicitly models deallocation, we now know no
 449 well-typed program will ever try to access deallocated memory, establishing the correctness
 450 of our memory management checking.

451 3.4.2 Proof Sketch

452 To prove the above theorem, we need several lemmas; the interesting ones are: the moral
 453 equivalent of the frame rule (C.1), monotonicity for the step-index (C.5), splitting up
 454 environments corresponds to splitting up heaps (C.7) and heap-and-expressions take the
 455 same steps of evaluation under any substitution of their free fractional permissions (C.8).

456 The proof proceeds by induction on the typing judgement. The case for `TY_FIX` is the
 457 reason we quantify over the step-index k in the *conclusion* of the soundness theorem. It
 458 allows us to then induct over the step-index and assume exactly the thing we need to prove
 459 at a smaller index.

460 The case for `TY_GEN` follows a similar pattern, but has the extra complication of reducing
 461 an expression with an arbitrary fractional permission variable in it, and then instantiating it
 462 at the last moment to conclude, which is where C.8 (heap-and-expressions take the same
 463 steps of evaluation under any substitution of their free fractional permissions) is used.

464 The rest of the cases are either very simple base cases (variables, unit, boolean, integer
 465 or element literals) or follow very similar patterns; for these, only `TY_LET` is presented in
 466 full and other similar cases simply highlight exactly what would be different. The general
 467 idea is to split up the linear substitution and heap along the same split of Γ/Γ' , then (by
 468 induction) use $\mathcal{C}_k[\!-\!]$ and one ‘half’ of the linear substitution and heap to conclude the ‘first’
 469 sub-expression either takes $j < k$ steps to **err** or another heap-and-expression.

470 In the first case, you use `OP_CONTEXT_ERR` to conclude the whole let-expression does
 471 the same. Similarly we use `OP_CONTEXT` j times in the second case. However, a small
 472 book-keeping wrinkle needs to be taken care of in the case that the heap-and-expression
 473 turns into a value in $i \leq j$ steps: `OP_CONTEXT` is not functorial for the n -fold iteration of
 474 \rightarrow . Basically, the following is not quite true:

$$475 \frac{\langle \sigma, e \rangle \rightarrow^j \langle \sigma', e' \rangle}{\langle \sigma, C[e] \rangle \rightarrow^j \langle \sigma', C[e'] \rangle} \quad \text{OP_CONTEXT}$$

476 because after the i steps, we need to invoke `OP_LET_VAR` to proceed evaluation for any
 477 remaining $j - i$ steps. After that, it suffices to use the induction hypothesis on the second
 478 sub-expression to finish the proof. To do so, we need to construct a valid linear substitution
 479 and heap (one in $\mathcal{L}_k[\![\Gamma', x : t]\!]\theta$). We take the other ‘half’ of the linear substitution and heap
 480 (from the initial split at the start) and extend it with $[x \mapsto v]$, (where x is the variable bound
 481 in the let-expression and v is the value we assume the first sub-expression evaluated to in i
 482 steps).

4 Implementation

4.1 Implementation Strategy

NUMLIN transpiles to OCaml and its implementation follows the structure of a typical domain-specific language (DSL) compiler. Although NUMLIN's current implementation is not as an embedded DSL, its the general design is simple enough to adapt to being so and also to target other languages.

Alongside the transpiler, a 'Read-Check-Translate' loop, benchmarking program and a test suite are included in the artifacts accompanying this paper.

1. **Parsing.** A generated, LR(1) parser parses a text file into a syntax tree. In general, this part will vary for different languages and can also be dealt with using combinators or syntax-extensions (the EDSL approach) if the host language offers such support.
2. **Desugaring.** The syntax tree is then desugared into a smaller, more concise, abstract syntax tree. This allows for the type checker to be simpler to specify and easier to implement.
3. **Matrix Expressions** are also desugared into the abstract syntax tree through pattern-matching.
4. **Type checking.** The abstract syntax tree is explicitly typed, with some inference to make writing typical programs more convenient.
5. **Code Generation.** The abstract syntax tree is translated into OCaml, with a few 'optimisations' to produce more readable code. This process is type-preserving: NUMLIN's type system is embedded into OCaml's (Figure 14), so the OCaml type checker acts as a sanity check on the generated code.

A very pleasant way to use NUMLIN is to have the build system generate code at *compile-time* and then have the generated code be used by other modules like normal OCaml functions. This makes it possible and even easy to use NUMLIN alongside existing OCaml libraries; in fact, this is exactly how the benchmarking program and test-suite use code written in NUMLIN.

4.1.1 Desugaring, Matrix Expressions and Type Checking

As seen earlier (Figure 2), desugaring is conventional. Matrix expressions are translated into BLAS/LAPACK calls via purely syntactic pattern-matching (also seen earlier in Figure 9).

4.1.2 Type checking

Type checking is mostly standard for a linearly typed language, with the exception of fractional permission inference. By restricting fractions to be non-positive integer powers of two, we only need to keep track of the logarithm of the fractions used. Explicit sharing and unsharing removes the need for performing dataflow analysis. As a result, all fractional arithmetic can be solved with unification, and in doing so, fractions become directly usable in NUMLIN's type-system as opposed to a convenient theoretical tool.

Because all functions must have their argument types explicitly annotated, inferring the correct fraction at a call-site is simply a matter of unification. We believe *full-inference of fractional permissions is similarly just matter of unification* (thanks to an experimental implementation of just this feature), even though the formal system we present here is for an explicitly-typed language.

$f ::=$	<code>module Arr =</code>	$\llbracket 'fc \rrbracket = 'fc$
$'fc$	<code>Owl.Dense.Ndarray.D</code>	$\llbracket \mathbf{z} \rrbracket = \mathbf{z}$
\mathbf{z}		$\llbracket f \mathbf{s} \rrbracket = \llbracket f \rrbracket \mathbf{s}$
$f \mathbf{s}$	<code>type z = Z</code>	$\llbracket \mathbf{unit} \rrbracket = \mathbf{unit}$
$t ::=$	<code>type 'a s = Succ</code>	$\llbracket \mathbf{bool} \rrbracket = \mathbf{bool}$
<code>unit</code>	<code>type 'a arr =</code>	$\llbracket \mathbf{int} \rrbracket = \mathbf{int}$
<code>bool</code>	<code>A of Arr.arr</code>	$\llbracket \mathbf{elt} \rrbracket = \mathbf{float}$
<code>int</code>	<code>[@@unboxed]</code>	$\llbracket f \mathbf{arr} \rrbracket = \llbracket f \rrbracket \mathbf{arr}$
<code>elt</code>	<code>type 'a mat =</code>	$\llbracket f \mathbf{mat} \rrbracket = \llbracket f \rrbracket \mathbf{mat}$
$f \mathbf{arr}$	<code>M of Arr.arr</code>	$\llbracket ! t \rrbracket = \llbracket t \rrbracket \mathbf{bang}$
$f \mathbf{mat}$	<code>[@@unboxed]</code>	$\llbracket 'fc. t \rrbracket = \llbracket t \rrbracket$
$! t$	<code>type 'a bang =</code>	$\llbracket t \otimes t' \rrbracket = \llbracket t \rrbracket * \llbracket t' \rrbracket$
$'fc. t$	<code>Many of 'a</code>	$\llbracket t \multimap t' \rrbracket = \llbracket t \rrbracket \rightarrow \llbracket t' \rrbracket$
$t \otimes t'$	<code>[@@unboxed]</code>	
$t \multimap t'$		

■ **Figure 14** NUMLIN's type grammar (left) and its embedding into OCaml (right).

There are a few differences between the type system as presented in 3.2 and how we implemented it: the environment *changes* as a result of type checking an expression (the standard transformation to avoid a non-deterministic split of the environment for checking pairs); variables are *marked as used* rather than removed for better error messages; variables are *tagged* as linear or intuitionistic in *one* environment as opposed to being stored in *two* separate ones (this allows scoping/variable look-up to be handled uniformly).

4.1.3 Code Generation

Code generation is a straightforward mapping from NUMLIN's core constructs to high-level OCaml ones. We embed NUMLIN's type- and term- constructors into OCaml as a sanity check on the output (Figure 14).

This is also useful when using NUMLIN from within OCaml; for example, we can use existing tools to inspect the type of the function we are using (Figure 15). It is worth reiterating that only the type- and term- constructors are translated into OCaml, NUMLIN's precise control over linearity and aliasing are not brought over.

We actually use this fact to our advantage to clean up the output OCaml by removing what would otherwise be redundant re-bindings (Figure 16). Combined with a code-formatter, the resulting code is not obviously correct and exactly what an expert would intend to write by hand, but now with the guarantees and safety of NUMLIN behind it. A small example is shown in Figure 17, a larger one in Figure 19.

4.2 Performance Metrics

We think that using NUMLIN has two primary benefits: safety and performance. We discuss safety in 5.1, where we describe how we used NUMLIN to find linearity and aliasing bugs in a linear algebra algorithm that was *generated* by another program.

```

1 let lt4la_kalman ~sigma ~h ~mu ~r ~data =
0   Examples.Kalman.it (M sigma) (M h) (M mu) (M r) (M data)
NORMAL test/examples_test.ml
'a mat ->
'b mat ->
'c mat ->
z mat ->
z mat -> ('a mat * ('b mat * ('c mat * (z mat * z mat)))) * (z mat * z mat)
:merlin-type-history:
0   let fact = Examples.Factorial.it in
NORMAL test/examples_test.ml
int bang -> int bang

```

■ **Figure 15** Using NUMLIN functions from OCaml.

```

let Many x = x in
let Many x = Many (Many x) in <exp>  ⇒  <exp>

(* fixp = fix (f, x:t, <exp> : t') *)
(*1*) let Many f = Many fixp in <body>  ⇒  let rec f x = <exp> in <body>
(*2*) let f = fixp in <body>

(*1*) let Many x = <exp> in
(*-*) let Many x = Many (Many x) in <body>  ⇒  let x = <exp> in <body>
(*2*) let Many x = Many <exp> in <body>
(*3*) (fun x : t -> <body>) <exp>

```

■ **Figure 16** Removing redundant re-bindings during translation to OCaml.

548 4.2.1 Setup

549 For performance, we measured the execution times of four equivalent implementations of a
550 Kalman filter: in C (using CBLAS), NUMLIN (using OWL’s low-level CBLAS bindings), OCaml
551 (using OWL’s intended, safe/copying-by-default interface), and Python (using NUMPY, with
552 the interpreter started and functions interpreted). We measured execution time in micro-
553 seconds, against an exponentially (powers of 5) increasing scaling factor for matrix size
554 parameters $n = 5$ and $k = 3$.

555 For large scaling factors ($n = 5^4, 5^5$), we triggered a full garbage-collection before
556 measuring the execution time of a single call of a function. However, due to the limitations
557 of the micro-benchmarking library we used, for smaller scaling factors ($n = 5^1, 5^2, 5^3$), we
558 measured the execution time of *multiple* calls to a function in a loop, thus including potential
559 garbage-collection effects.

560 We also measured the execution times of L1-norm minimisation and the “linear-regression”
561 $((\mathbf{X}^T \mathbf{X})^{-1} \mathbf{X}^T \mathbf{y})$ similarly, but without a C implementation.

562 4.2.2 Hypothesis

563 We expected the C implementation to be faster than the NUMLIN one because the latter has
564 the additional (but relatively low) overhead of dimension checks and crossing the OCaml/C
565 FFI for each call to a CBLAS routine, even though the calls and their order are exactly the
566 same. We expected the OCaml and Python implementations to be slower because they

```

let rec f i n x0 row =
  if Prim.extract @@ Prim.eqI i n then (row, x0)
  else
    let row, x1 = Prim.get row i in
    f (Prim.addI i (Many 1)) n (Prim.addE x0 x1) row
in
f

```

■ **Figure 17** Recursive OCaml function for a summing over an array, generated (at *compile time*) from the code in Figure 4, passed through `ocamlformat` for presentation.

567 allocate more temporaries (so possibly less cache-friendly) and carry out more floating-point
 568 operations – the CBLAS and NUMLIN implementations use ternary kernels (coalescing steps),
 569 a Cholesky decomposition (of a symmetric matrix, which is more efficient than the LU
 570 decomposition used for inverting a matrix in OWL and NUMPY) and `symm` (symmetric matrix
 571 multiplication, halving the number of floating-point multiplications required).

572 4.2.3 Results

573 The results in Figures 18 are as we expected: C is the fastest, followed by NUMLIN, with
 574 OCaml and Python last. Differences in timings are quite pronounced at small matrix sizes,
 575 but are still significant at larger ones. Specifically for the Kalman filter, for $n = 625$, CBLAS
 576 took $112 \pm 35\text{ ms}$, NUMLIN took $105 \pm 25\text{ ms}$, OWL took $124 \pm 38\text{ ms}$ and NUMPY took
 577 $112 \pm 12\text{ ms}$; for $n = 3125$, CBLAS took $10.8 \pm 0.7\text{ s}$, NUMLIN took $12.0 \pm 1.2\text{ s}$, OWL took
 578 $13.3 \pm 0.2\text{ s}$ and NUMPY took $12.7 \pm 0.6\text{ s}$.

579 Worth highlighting here is the other major advantage of using NUMLIN is reduced
 580 memory usage. Whilst the OWL and NUMPY use 11 temporary matrices for the Kalman
 581 filter, (*excluding* the 2 matrices which store the results), using $n + n^2 + 4nk + 3k^2 + 2k \approx 4n^2$
 582 (for $k = 3n/5$) words of memory, CBLAS and NUMLIN use only 2 temporary matrices
 583 (*excluding* the *one* matrix which stores one of the results), using only $n^2 + nk \leq 2n^2$ words
 584 of memory.

585 4.2.4 Analysis

586 As matrix sizes increase, assuming sufficient memory, the difference in the number of floating-
 587 point operations ($O(n^3)$) dominates execution times. However for small matrix sizes, since
 588 n is small and the measurements were over multiple calls to a function in a loop, the large
 589 number of temporaries show the adverse effect of not re-using memory at even quite small
 590 matrix sizes: creating pressure on the garbage collector.

591 5 Discussion and Related Work

592 5.1 Finding Bugs in SymPy's Output

593 Prior to this project, we had little experience with linear algebra libraries or the problem
 594 of matrix expression compilation. As such, we based our initial NUMLIN implementation
 595 of a Kalman filter using BLAS and LAPACK, on a popular GitHub gist of a Fortran
 596 implementation, one that was *automatically generated* from SymPy's matrix expression
 597 compiler [17].

598 Once we translated the implementation from Fortran to NUMLIN, we attempted to compile
 599 it and found that (to our surprise) it did not type-check. This was because the original

600 implementation contained incorrect aliasing, unused variables and unnecessary temporaries,
601 and did not adhere to Fortran's read/write permissions (with respect to `intent` annotations
602 `in`, `out` and `inout`) all of which were now highlighted by NUMLIN's type system.

603 The original implementation used 6 temporaries, one of which was immediately spotted
604 as never being used due to linearity. It also contained two variables which were marked as
605 `intent(in)` but would have been written over by calls to 'gemm', spotted by the fractional
606 capabilities feature. Furthermore, it used a matrix *twice* in a call to 'symm', once with a read
607 permission but once with a *write* permission. Fortran assumes that any parameter being
608 written to is not aliased and so this call was not only incorrect, but illegal according to the
609 standard, both aspects of which were captured by linearity and fractional capabilities.

610 Lastly, it contained another unnecessary temporary, however one that was not obvious
611 without linear types. To spot it, we first performed live-range splitting (checked by linearity)
612 by hoisting calls to `freeM` and then annotated the freed matrices with their dimensions.
613 After doing so and spotting two disjoint live-ranges of the same size, we replaced a call to
614 `freeM` followed by allocating call to `copy` with one, in-place call to `copyM_to`. We believe
615 the ability to boldly refactor code which manages memory is good evidence of the usefulness
616 of linearity as a tool for programming.

617 5.2 Related Work

618 Using linear types for BLAS routines is a particularly good domain fit (given the implicit
619 restrictions on aliasing arguments), and as a result the idea of using substructural types
620 to express array computations is not a particularly new one [18, 13, 7]. However, many
621 of these designs have been focused on building languages to *implement* the kernel linear
622 algebra functions, and as a result, they tend to add additional limitations on the language
623 design. Both Futhark [13] and Single Assignment C [18] omit higher-order functions to
624 facilitate compilation to GPUs. The work of [7] forbids term-level recursion, in order to
625 ensure that all higher-order computations can be statically normalized away and thereby
626 maximize opportunities for array fusion.

627 In contrast, our approach is to begin with the assumption that we can take existing
628 efficient BLAS-like libraries, and then enforce their correct *usage* using a linear type discipline
629 with fractional permissions.

630 This approach is similar to the one taken in linear algebra libraries for Rust – these libraries
631 typically take advantage of the distinction that Rust's type system offers between mutable
632 views/references to arrays. The work of [20] and [14] suggest that Rust's borrow-checker
633 *can be expressed in simpler terms* using fractional permissions, though to our knowledge the
634 programmer-visible lifetime analysis in Rust has never been formalized.

635 Working explicitly with fractional permissions has two main benefits. First, our type
636 system demonstrates that type systems for fractional permissions can be dramatically simpler
637 than existing state-of-the-art approaches, including both industrial languages like Rust, as
638 well as academic (such as those developed by [8]). Bierhoff *et al*'s type system, much like
639 Rust's, builds a complex dataflow analysis into the typing rules to infer when variables can
640 be shared or not. This allows for more natural-looking user programs, but can create the
641 impression that using fractional permissions requires a heavy theoretical and engineering
642 effort going well beyond that needed for supporting basic linear types.

643 Instead, our approach, of requiring sharing to be made explicit, lets us demonstrate that
644 the existing unification machinery already in place for ordinary ML-style type inference can
645 be reused to support fractions. Basically, we can view sharing a value as dividing a fraction
646 by two, and after taking logarithms all fractions are Peano numbers, whose equality can be

established with ordinary unification.

This fact is important because there are major upcoming implementations of linear types such as Linear Haskell [6], which do not have built-in support for fractional permissions. Instead, Linear Haskell takes a slightly different definition of linearity, one based on *arrows* as opposed to *kinds*: for $f : a \multimap b$, if fu is used exactly once then u is used exactly once. Whilst this has the advantage of being backwards-compatible, it also means that the type system has no built-in support for the concurrent reader, exclusive writer pattern that fractional permissions enable.

However, since our type system shows demonstrates unification is “all one needs” for fractions, it should be possible to *encode* NUMLIN’s approach to fractional permissions in Linear Haskell by adding a GADT-style natural number index to array types tracking the fraction, which should enable supporting high-performance BLAS bindings in Linear Haskell. Actually implementing this is something we leave for future work, as there remains one issue which we do not see a good encoding for. Namely, only having support for linear functions makes it a bit inconvenient to manipulate linear values directly – programs end up taking on a CPS-like structure. This seems to remain an advantage of a direct implementation of linear types over the Linear Haskell style.

5.3 Simplicity and Further Work

We are pleasantly surprised at how simple the overall design and implementation of NUMLIN is, given its expressive power and usability. So simple in fact, that fractions, a convenient theoretical abstraction until this point, could be implemented by restricting division and multiplication to be by 2 only [10], thus turning any required arithmetic into unification.

Indeed, the focus on getting a working prototype early on (so that we could test it with real BLAS/LAPACK routines as soon as possible) meant that we only added features to the type system when it was clear that they were absolutely necessary: these features were !-types and value-restriction for the `Many` constructor.

Going forwards, one may wish to eliminate even more runtime errors from NUMLIN, by extending its type system. For example, we could have used existential types to statically track pointer identities [2], or parametric polymorphism.

We could also attempt to catch mismatched dimensions at compile time as well. While this could be done with generative phantom types [1], using dependent types may offer more flexibility in *partitioning* regions [16] or statically enforcing dimensions related constraints of the arguments at compile-time. ATS [12] is an example of a language which combines linear types with a sophisticated proof layer. But although it provides BLAS bindings, it does not aim to provide aliasing restrictions as demonstrated in this paper.

Taking this idea one step even further, since matrix dimensions are typically fixed at runtime, we could *stage* NUMLIN programs and compile matrix expressions using more sophisticated algorithms [4]. However, it is worth noting that without care, such algorithms [17], usually based on graph-based, ad-hoc dataflow analysis, can produce erroneous output which would not get past a linear type system with fractions.

We also think that this concept (and the general design of its implementation) need not be limited to linear algebra: we could conceivably ‘backport’ this idea to other contexts that need linearity (concurrency, single-use continuations, zero-copy buffer, streaming I/O) or combine it with dependent types to achieve even more expressive power to split up a single block of memory into multiple regions in an arbitrary manner [16].

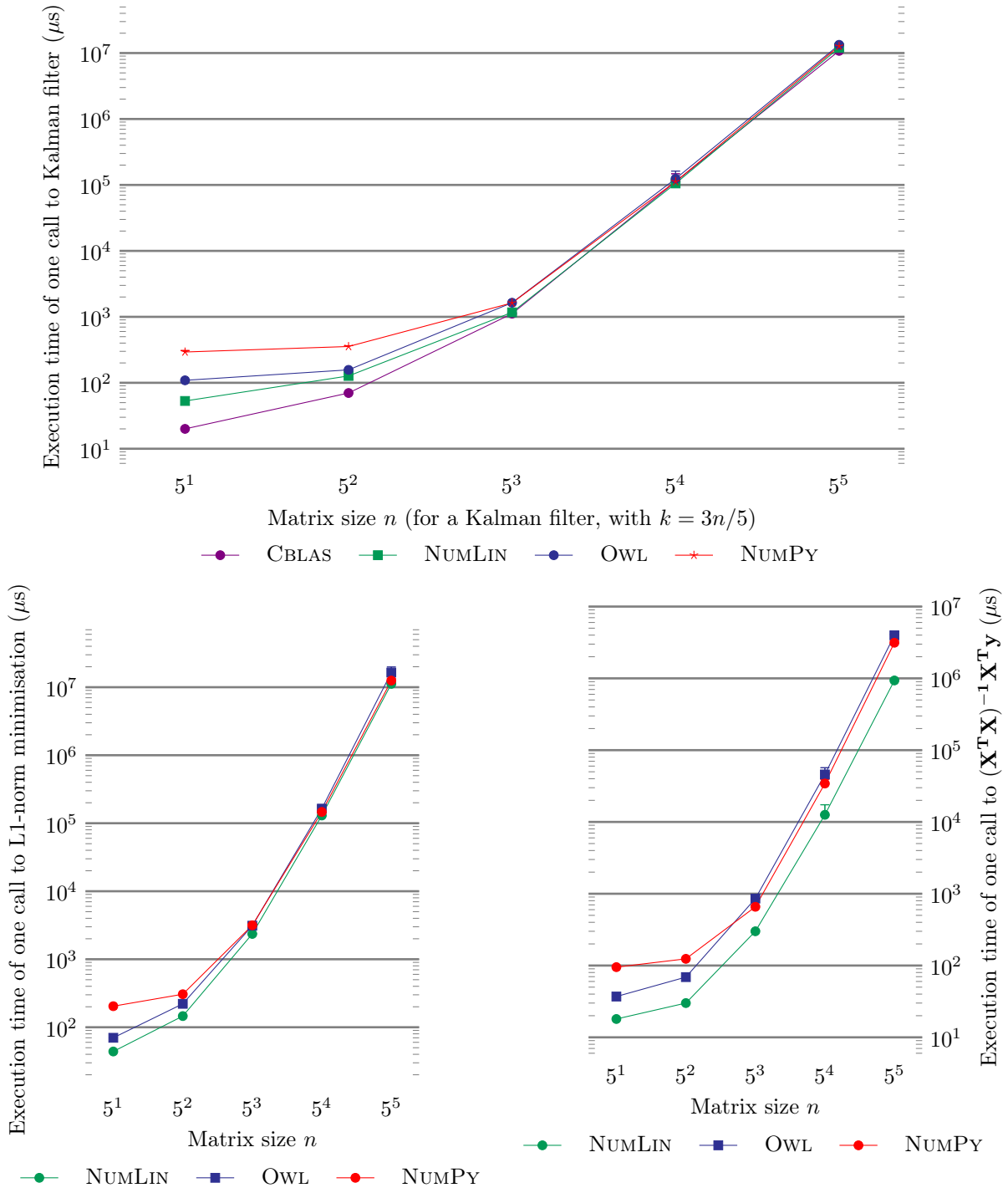


Figure 18 Comparison of execution times (error bars are present but quite small). Small matrices and timings $n \leq 5^3$ were micro-benchmarked with the `Core_bench` library. Larger ones used Unix's `getrusage` functionality, sandwiched between calls to `Gc.full_major` for the OCaml implementations.

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740 **A** NumLin Specification

741 **A.1** Static Semantics

742 $\boxed{\Theta; \Delta; \Gamma \vdash e : t}$ Typing rules for expressions

743 $\frac{}{\Theta; \Delta; \cdot, x : t \vdash x : t} \text{TY_VAR_LIN}$

744 $\frac{x : t \in \Delta}{\Theta; \Delta; \cdot \vdash x : t} \text{TY_VAR}$

745 $\frac{\begin{array}{l} \Theta; \Delta; \Gamma \vdash e : t \\ \Theta; \Delta; \Gamma', x : t \vdash e' : t' \end{array}}{\Theta; \Delta; \Gamma, \Gamma' \vdash \text{let } x = e \text{ in } e' : t'} \text{TY_LET}$

746 $\frac{}{\Theta; \Delta; \cdot \vdash () : \text{unit}} \text{TY_UNIT_INTRO}$

747 $\frac{\begin{array}{l} \Theta; \Delta; \Gamma \vdash e : \text{unit} \\ \Theta; \Delta; \Gamma' \vdash e' : t \end{array}}{\Theta; \Delta; \Gamma, \Gamma' \vdash \text{let } () = e \text{ in } e' : t} \text{TY_UNIT_ELIM}$

748 $\frac{}{\Theta; \Delta; \cdot \vdash \text{true} : \text{bool}} \text{TY_BOOL_TRUE}$

749 $\frac{}{\Theta; \Delta; \cdot \vdash \text{false} : \text{bool}} \text{TY_BOOL_FALSE}$

750 $\frac{\begin{array}{l} \Theta; \Delta; \Gamma \vdash e : !\text{bool} \\ \Theta; \Delta; \Gamma' \vdash e_1 : t' \\ \Theta; \Delta; \Gamma' \vdash e_2 : t' \end{array}}{\Theta; \Delta; \Gamma, \Gamma' \vdash \text{if } e \text{ then } e_1 \text{ else } e_2 : t} \text{TY_BOOL_ELIM}$

751 $\frac{}{\Theta; \Delta; \cdot \vdash k : \text{int}} \text{TY_INT_INTRO}$

752 $\frac{}{\Theta; \Delta; \cdot \vdash el : \text{elt}} \text{TY_ELT_INTRO}$

753 $\frac{\begin{array}{l} \Theta; \Delta; \cdot \vdash v : t \\ v \neq l \end{array}}{\Theta; \Delta; \cdot \vdash \text{Many } v : !t} \text{TY_BANG_INTRO}$

754 $\frac{\begin{array}{l} \Theta; \Delta; \Gamma \vdash e : !t \\ \Theta; \Delta, x : t; \Gamma' \vdash e' : t' \end{array}}{\Theta; \Delta; \Gamma, \Gamma' \vdash \text{let Many } x = e \text{ in } e' : t'} \text{TY_BANG_ELIM}$

755 $\frac{\begin{array}{l} \Theta; \Delta; \Gamma \vdash e : t \\ \Theta; \Delta; \Gamma' \vdash e' : t' \end{array}}{\Theta; \Delta; \Gamma, \Gamma' \vdash (e, e') : t \otimes t'} \text{TY_PAIR_INTRO}$

23:26 NumLin: Linear Types for Linear Algebra

$$\frac{\Theta; \Delta; \Gamma \vdash e_{12} : t_1 \otimes t_2 \quad \Theta; \Delta; \Gamma', a : t_1, b : t_2 \vdash e : t}{\Theta; \Delta; \Gamma, \Gamma' \vdash \mathbf{let} (a, b) = e_{12} \mathbf{in} e : t} \text{TY_PAIR_ELIM}$$

$$\frac{\Theta \vdash t' \text{Type} \quad \Theta; \Delta; \Gamma, x : t' \vdash e : t}{\Theta; \Delta; \Gamma \vdash \mathbf{fun} x : t' \rightarrow e : t' \multimap t} \text{TY_LAMBDA}$$

$$\frac{\Theta; \Delta; \Gamma \vdash e : t' \multimap t \quad \Theta; \Delta; \Gamma' \vdash e' : t'}{\Theta; \Delta; \Gamma, \Gamma' \vdash e e' : t} \text{TY_APP}$$

$$\frac{\Theta, fc; \Delta; \Gamma \vdash e : t}{\Theta; \Delta; \Gamma \vdash \mathbf{fun} 'fc \rightarrow e : 'fc.t} \text{TY_GEN}$$

$$\frac{\Theta \vdash f \text{Perm} \quad \Theta; \Delta; \Gamma \vdash e : 'fc.t}{\Theta; \Delta; \Gamma \vdash e[f] : t[f/fc]} \text{TY_SPC}$$

$$\frac{\Theta; \Delta, g : t \multimap t'; \cdot, x : t \vdash e : t'}{\Theta; \Delta; \cdot \vdash \mathbf{fix} (g, x : t, e : t') : t \multimap t'} \text{TY_FIX}$$

762

763 A.2 Dynamic Semantics

764 $\langle \sigma, e \rangle \rightarrow \text{Config}$ Operational semantics

$$\frac{}{\langle \sigma, \mathbf{let} () = () \mathbf{in} e \rangle \rightarrow \langle \sigma, e \rangle} \text{OP_LET_UNIT}$$

$$\frac{}{\langle \sigma, \mathbf{let} x = v \mathbf{in} e \rangle \rightarrow \langle \sigma, e[v/x] \rangle} \text{OP_LET_VAR}$$

$$\frac{}{\langle \sigma, \mathbf{if} (\mathbf{Many true}) \mathbf{then} e_1 \mathbf{else} e_2 \rangle \rightarrow \langle \sigma, e_1 \rangle} \text{OP_IF_TRUE}$$

$$\frac{}{\langle \sigma, \mathbf{if} (\mathbf{Many false}) \mathbf{then} e_1 \mathbf{else} e_2 \rangle \rightarrow \langle \sigma, e_2 \rangle} \text{OP_IF_FALSE}$$

$$\frac{}{\langle \sigma, \mathbf{let Many} x = \mathbf{Many} v \mathbf{in} e \rangle \rightarrow \langle \sigma, e[v/x] \rangle} \text{OP_LET_MANY}$$

$$\frac{}{\langle \sigma, \mathbf{let} (a, b) = (v_1, v_2) \mathbf{in} e \rangle \rightarrow \langle \sigma, e[v_1/a][v_2/b] \rangle} \text{OP_LET_PAIR}$$

$$\frac{}{\langle \sigma, (\mathbf{fun} 'fc \rightarrow v)[f] \rangle \rightarrow \langle \sigma[f/fc], v[f/fc] \rangle} \text{OP_FRAC_PERM}$$

$$\frac{}{\langle \sigma, \mathbf{fix} (g, x : t, e : t') v \rangle \rightarrow \langle \sigma, e[v/x][\mathbf{fix} (g, x : t, e : t')/g] \rangle} \text{OP_APP_FIX}$$

$$\frac{}{\langle \sigma, (\mathbf{fun} x : t \rightarrow e) v \rangle \rightarrow \langle \sigma, e[v/x] \rangle} \text{OP_APP_LAMBDA}$$

773

- 774 $\frac{\langle \sigma, e \rangle \rightarrow \langle \sigma', e' \rangle}{\langle \sigma, C[e] \rangle \rightarrow \langle \sigma', C[e'] \rangle} \quad \text{OP_CONTEXT}$
- 775 $\frac{\langle \sigma, e \rangle \rightarrow \mathbf{err}}{\langle \sigma, C[e] \rangle \rightarrow \mathbf{err}} \quad \text{OP_CONTEXT_ERR}$
- 776 $\frac{0 \leq k_1, k_2 \quad l \text{ fresh}}{\langle \sigma, \mathbf{matrix} \ k_1 \ k_2 \rangle \rightarrow \langle \sigma + \{l \mapsto_1 M_{k_1, k_2}\}, l \rangle} \quad \text{OP_MATRIX}$
- 777 $\frac{k_1 < 0 \text{ or } k_2 < 0}{\langle \sigma, \mathbf{matrix} \ k_1 \ k_2 \rangle \rightarrow \mathbf{err}} \quad \text{OP_MATRIX_NEG}$
- 778 $\frac{}{\langle \sigma + \{l \mapsto_1 m_{k_1, k_2}\}, \mathbf{free} \ l \rangle \rightarrow \langle \sigma, () \rangle} \quad \text{OP_FREE}$
- 779 $\frac{}{\langle \sigma + \{l \mapsto_f m_{k_1, k_2}\}, \mathbf{share}[f] \ l \rangle \rightarrow \langle \sigma + \{l \mapsto_{\frac{1}{2}f} m_{k_1, k_2}\} + \{l \mapsto_{\frac{1}{2}f} m_{k_1, k_2}\}, (l, l) \rangle} \quad \text{OP_SHARE}$
- 780 $\frac{\sigma' \equiv \sigma + \{l \mapsto_{\frac{1}{2}f} m_{k_1, k_2}\} + \{l \mapsto_{\frac{1}{2}f} m_{k_1, k_2}\}}{\langle \sigma', \mathbf{unshare}[f] \ l \ l \rangle \rightarrow \langle \sigma + \{l \mapsto_f m_{k_1, k_2}\}, l \rangle} \quad \text{OP_UNSHARE_EQ}$
- 781 $\frac{l \neq l'}{\langle \sigma + \{l \mapsto_{\frac{1}{2}f} m_{k_1, k_2}\} + \{l' \mapsto_{\frac{1}{2}f} m'_{k_1, k_2}\}, \mathbf{unshare}[f] \ l \ l' \rangle \rightarrow \mathbf{err}} \quad \text{OP_UNSHARE_NEQ}$
- 782 $\frac{\begin{array}{l} \sigma' \equiv \sigma + \{l_1 \mapsto_{fc_1} m_{1k_1, k_2}\} + \{l_2 \mapsto_{fc_2} m_{2k_2, k_3}\} \\ \sigma_1 \equiv \sigma' + \{l_3 \mapsto_1 m_{3k_1, k_3}\} \\ \sigma_2 \equiv \sigma' + \{l_3 \mapsto_1 (m_1 \ m_2 + m_3)_{k_1, k_3}\} \end{array}}{\langle \sigma_1, \mathbf{gemm}[fc_1] \ l_1[fc_2] \ l_2 \ l_3 \rangle \rightarrow \langle \sigma_2, ((l_1, l_2), l_3) \rangle} \quad \text{OP_GEMM_MATCH}$
- 783 $\frac{k_2 \neq k'_2 \quad \sigma' \equiv \sigma + \{l_1 \mapsto_{fc_1} m_{1k_1, k_2}\} + \{l_2 \mapsto_{fc_2} m_{2k'_2, k_3}\}}{\langle \sigma' + \{l_3 \mapsto_1 m_{1k_1, k_3}\}, \mathbf{gemm}[fc_1] \ l_1[fc_2] \ l_2 \ l_3 \rangle \rightarrow \mathbf{err}} \quad \text{OP_GEMM_MISMATCH}$
- 784

785 **B** Interpretation

786 **B.1** Definitions

787 Operationally, $Heap \sqsubseteq Loc \times Permission \times Matrix$ (a multiset), denoted with a σ .
 788 Define its *interpretation* to be $Loc \rightarrow Permission \times Matrix$ with $\star : Heap \times Heap \rightarrow Heap$ as
 789 follows:

$$790 \quad (\varsigma_1 \star \varsigma_2)(l) \equiv \begin{cases} \varsigma_1(l) & \text{if } l \in \text{dom}(\varsigma_1) \wedge l \notin \text{dom}(\varsigma_2) \\ \varsigma_2(l) & \text{if } l \in \text{dom}(\varsigma_2) \wedge l \notin \text{dom}(\varsigma_1) \\ (f_1 + f_2, m) & \text{if } (f_1, m) = \varsigma_1(l) \wedge (f_2, m) = \varsigma_2(l) \wedge f_1 + f_2 \leq 1 \\ \text{undefined} & \text{otherwise} \end{cases}$$

791 Commutativity and associativity of \star follows from that of $+$.
 792 $\varsigma_1 \star \varsigma_2$ is *defined* if it is for all $l \in \text{dom}(\varsigma_1) \cup \text{dom}(\varsigma_2)$.
 793 Define $\mathcal{H}[\![\sigma]\!] = \star_{(l,f,m) \in \sigma} [l \mapsto_f m]$ and **implicitly denote** $\varsigma \equiv \mathcal{H}[\![\theta(\sigma)]\!]$.
 794
 795 The n -fold iteration for the \rightarrow (functional) relation, is also a (functional) relation:

$$796 \quad \forall n. \mathbf{err} \rightarrow^n \mathbf{err} \quad \langle \sigma, v \rangle \rightarrow^n \langle \sigma, v \rangle \quad \langle \sigma, e \rangle \rightarrow^0 \langle \sigma, e \rangle \quad \langle \sigma, e \rangle \rightarrow^{n+1} ((\langle \sigma, e \rangle \rightarrow) \rightarrow^n)$$

797

798 Hence, all bounded iterations end in either an **err**, a heap-and-expression or a heap-and-value.

799 B.2 Interpretation

$$800 \quad \mathcal{V}_k[\mathbf{unit}] = \{(\emptyset, *)\}$$

801

$$802 \quad \mathcal{V}_k[\mathbf{bool}] = \{(\emptyset, true), (\emptyset, false)\}$$

803

$$804 \quad \mathcal{V}_k[\mathbf{int}] = \{(\emptyset, n) \mid 2^{-63} \leq n \leq 2^{63} - 1\}$$

805

$$806 \quad \mathcal{V}_k[\mathbf{elt}] = \{(\emptyset, f) \mid f \text{ a IEEE Float64 } \}$$

807

$$808 \quad \mathcal{V}_k[f \mathbf{mat}] = \{(\{l \mapsto_{2^{-f}} _ \}, l)\}$$

809

$$810 \quad \mathcal{V}_k[!t] = \{(\emptyset, \mathbf{Many} \ v) \mid (\emptyset, v) \in \mathcal{V}_k[t]\}$$

811

$$812 \quad \mathcal{V}_k[f.c. \ t] = \{(\varsigma, \mathbf{fun} \ f.c. \rightarrow v) \mid \forall f. (\varsigma[f/f.c], v[f/f.c]) \in \mathcal{V}_{k-1}[t[f/f.c]]\}$$

813

$$814 \quad \mathcal{V}_k[t_1 \otimes t_2] = \{(\varsigma_1 \star \varsigma_2, (v_1, v_2)) \mid (\varsigma_1, v_1) \in \mathcal{V}_k[t_1] \wedge (\varsigma_2, v_2) \in \mathcal{V}_k[t_2]\}$$

815

$$816 \quad \mathcal{V}_k[t' \multimap t] = \{(\varsigma_v, v) \mid (v \equiv \mathbf{fun} \ x : t' \rightarrow e \vee v \equiv \mathbf{fix}(g, x : t', e : t)) \wedge$$

$$817 \quad \forall j \leq k, (\varsigma_{v'}, v') \in \mathcal{V}_j[t']. \varsigma_v \star \varsigma_{v'} \text{ defined} \Rightarrow (\varsigma_v \star \varsigma_{v'}, v \ v') \in \mathcal{C}_j[t]\}$$

818

$$819 \quad \mathcal{C}_k[t] = \{(\varsigma_s, e_s) \mid \forall j < k, \sigma_r. \varsigma_s \star \varsigma_r \text{ defined} \Rightarrow \langle \sigma_s + \sigma_r, e_s \rangle \rightarrow^j \mathbf{err} \vee \exists \sigma_f, e_f.$$

$$820 \quad \langle \sigma_s + \sigma_r, e_s \rangle \rightarrow^j \langle \sigma_f + \sigma_r, e_f \rangle \wedge (e_f \text{ is a value} \Rightarrow (\varsigma_f \star \varsigma_r, e_f) \in \mathcal{V}_{k-j}[t])\}$$

821

$$822 \quad \mathcal{I}_k[\cdot]\theta = \{\emptyset\}$$

823

$$824 \quad \mathcal{I}_k[\Delta, x : t]\theta = \{\delta[x \mapsto v_x] \mid \delta \in \mathcal{I}_k[\Delta]\theta \wedge (\emptyset, v_x) \in \mathcal{V}_k[\theta(t)]\}$$

825

$$826 \quad \mathcal{L}_k[\cdot]\theta = \{(\emptyset, [])\}$$

827

$$828 \quad \mathcal{L}_k[\Gamma, x : t]\theta = \{(\varsigma \star \varsigma_x, \gamma[x \mapsto v_x]) \mid (\varsigma, \gamma) \in \mathcal{L}_k[\Gamma]\theta \wedge (\varsigma_x, v_x) \in \mathcal{V}_k[\theta(t)]\}$$

829

$$830 \quad \mathcal{H}[\sigma] = \star_{(l,f,m) \in \sigma} [l \mapsto_f m]$$

$$831 \quad \varsigma \equiv \mathcal{H}[\theta(\sigma)]$$

832

$$833 \quad {}_k[\Theta; \Delta; \Gamma \vdash e : t] = \forall \theta, \delta, \gamma, \sigma. \Theta = \text{dom}(\theta) \wedge (\varsigma, \gamma) \in \mathcal{L}_k[\Gamma]\theta \wedge \delta \in \mathcal{I}_k[\Delta]\theta \Rightarrow$$

$$834 \quad (\varsigma, \theta(\delta(\gamma(e)))) \in \mathcal{C}_k[\theta(t)]$$

836 **C** Lemmas

837 **C.1** $\forall \sigma_s, \sigma_r, e. \varsigma_s \star \varsigma_r \text{ defined} \Rightarrow \forall n. \langle \sigma_s, e \rangle \rightarrow^n = \langle \sigma_s + \sigma_r, e \rangle \rightarrow^n$

838 SUFFICES: By induction on n , consider only the cases $\langle \sigma_s, e \rangle \rightarrow \langle \sigma_f, e_f \rangle$ where $\sigma_s \neq \sigma_f$.

839
840 PROOF SKETCH: Only `OP_FREE`, `MATRIX`, `SHARE`, `UNSHARE_EQ`, `GEMM_MATCH`
841 change the heap: the rest are either parametric in the heap or step to an **err**.

842
843 PROVE: $\langle \sigma_s + \sigma_r, e \rangle \rightarrow \langle \sigma_f + \sigma_r, e_f \rangle$.

844
845 $\langle 1 \rangle 1$. CASE: `OP_FREE`, $\sigma_s \equiv \sigma' + \{l \mapsto_1 m\}$, $\sigma_f = \sigma'$.
846 PROOF: Instantiate `OP_FREE` with $(\sigma' + \sigma_r) + \{l \mapsto_1 m\}$,
847 valid because $l \notin \text{dom}(\varsigma_r)$ by $\varsigma' \star [l \mapsto_1 m] \star \varsigma_r$ defined (assumption).

848 $\langle 1 \rangle 2$. CASE: `OP_MATRIX`
849 PROOF: Rule has no requirements on σ_s so will also work with $\sigma_s + \sigma_r$.

850 $\langle 1 \rangle 3$. CASE: `OP_SHARE`, $\sigma_s \equiv \sigma' + \{l \mapsto_f m\}$, $\sigma_f = \sigma' + \{l \mapsto_{\frac{1}{2}.f} m\} + \{l \mapsto_{\frac{1}{2}.f} m\}$.
851 PROOF: Union-ing σ_r does not remove $l \mapsto_f m$, so that can be split out of $\sigma_s + \sigma_r$ as
852 before.

853 $\langle 1 \rangle 4$. CASE: `OP_UNSHARE_EQ`, $\sigma_s \equiv \sigma' + \{l \mapsto_{\frac{1}{2}.f} m\} + \{l \mapsto_{\frac{1}{2}.f} m\}$, $\sigma_f = \sigma' + \{l \mapsto_f m\}$.
854 $\langle 2 \rangle 1$. Union-ing σ_r does not remove $l \mapsto_{\frac{1}{2}.f} m$, so that can still be split out of $\sigma_s + \sigma_r$.
855 $\langle 2 \rangle 2$. There may also be other valid splits introduced by σ_r .
856 $\langle 2 \rangle 3$. However, by assumption of $\varsigma_s \star \varsigma_r$ defined, any splitting of $\sigma_s + \sigma_r$ will satisfy
857 $f \leq 1$.

858 $\langle 1 \rangle 5$. CASE: `OP_GEMM_MATCH`

859 $\langle 2 \rangle 1$. By assumption of $\varsigma_s \star \varsigma_r$ defined, either l_1 (or l_2 , or both) are not in σ_r , or they
860 are and the matrix values they point to are the same.

861 $\langle 2 \rangle 2$. The permissions (of l_1 and/or l_2) may differ, but `OP_GEMM_MATCH` universally
862 quantifies over them and leaves them unchanged, so they are irrelevant.

863 $\langle 2 \rangle 3$. Only the pointed to matrix value at l_3 changes.

864 $\langle 2 \rangle 4$. SUFFICES: $l_3 \notin \pi_1[\sigma_r]$.

865 $\langle 2 \rangle 5$. By assumption of $\varsigma_s \star \varsigma_r$ defined, $l_3 \notin \text{dom}(\varsigma_r)$.

866 $\langle 2 \rangle 6$. Hence $l_3 \notin \pi_1[\sigma_r]$.

867 **C.2** $\forall k, t. \mathcal{V}_k[t] \subseteq \mathcal{C}_k[t]$

868 Follows from definition of $\mathcal{C}_k[t]$, \rightarrow^j ($\forall n. \langle \sigma, v \rangle \rightarrow^n \langle \sigma, v \rangle$) for arbitrary $j \leq k$ and C.1.

869 **C.3** $\forall \theta, \delta, \gamma, v. \theta(\delta(\gamma(v)))$ is a value.

870 θ is irrelevant because it only maps fractional permission variables to fractional permissions.

871 By construction, δ and γ only map variables to values, and values are closed under substitution.

872 **C.4** $\forall k, \sigma, \sigma', e, e', t. (\varsigma', e') \in \mathcal{C}_k[t] \wedge \langle \sigma, e \rangle \rightarrow \langle \sigma', e' \rangle \Rightarrow (\varsigma, e) \in \mathcal{C}_{k+1}[t]$

873 In the lemma, and for the rest of its proof, $\varsigma = \mathcal{H}[\sigma]$.

874 ASSUME: arbitrary $j < k + 1$, and σ_r such that $\varsigma \star \varsigma_r$ defined.

875

876 $\langle 1 \rangle 1$. CASE: $j = 0$. Clearly $\sigma_f = \sigma_s + \sigma_r$ and $e' = e$.

877 Remains to show that if e is a value then $(\varsigma_s \star \varsigma_r, e) \in \mathcal{V}_k[t]$.

878 This is true vacuously, because by assumption, e is not a value.

879 $\langle 1 \rangle 2$. CASE: $j \geq 1$. We have $\langle \sigma, e \rangle \rightarrow^j = \langle \sigma', e' \rangle \rightarrow^{j-1}$.

880 Instantiate $(\varsigma', e') \in \mathcal{C}_k[t]$, with $j - 1 < k$ and σ_r to conclude the required conditions.

881 **C.5** $j \leq k \Rightarrow _k[\cdot] \subseteq _j[\cdot]$

882 For the rest of this proof, $\varsigma = \mathcal{H}[\sigma]$.

883 Lemma C.4 is the inductive step for this lemma for the $\mathcal{C}[\cdot]$ case.

884 Need to prove for $\mathcal{V}[\cdot]$, by induction on t and then index.

885 SUFFICES: Consider only $t \multimap t'$ case, rest use k directly on structure of type.

886 ASSUME: Arbitrary $j \leq k$ and $(\varsigma_{v'}, v') \in \mathcal{V}_k[t \multimap t']$.

887 PROVE: $(\varsigma_{v'}, v') \in \mathcal{V}_j[t \multimap t']$.

888

889 $\langle 1 \rangle 1$. v' is of the correct syntactic form (lambda or fixpoint) by assumption.

890 $\langle 1 \rangle 2$. ASSUME: arbitrary $j' \leq j$ and $(\varsigma_v, v) \in \mathcal{V}_{j'}[t]$ such that $\varsigma_{v'} \star \varsigma_v$ is defined.

891 $\langle 1 \rangle 3$. SUFFICES: to show $(\varsigma_{v'} \star \varsigma_v, v'v) \in \mathcal{C}_{j'}[t']$.

892 $\langle 1 \rangle 4$. This is true by instantiating $(\varsigma_{v'}, v') \in \mathcal{V}_k[t \multimap t']$ with $j' \leq k$ and $(\varsigma_v, v) \in \mathcal{V}_{j'}[t]$.

893 **C.6** $\forall \Delta, \Gamma, t, k, \theta, \delta, \gamma. \delta \in \mathcal{I}_k[\Delta]\theta \wedge \gamma \in \pi_2[\mathcal{L}_k[\Gamma]\theta] \Rightarrow \text{dom}(\Delta) = \text{dom}(\delta)$
 894 **and** $\text{dom}(\Gamma) = \text{dom}(\gamma)$

895 PROOF: By induction on Δ and Γ .

896 **C.7** $\forall k, \Gamma, \Gamma', \theta, \sigma_+, \gamma_+. (\varsigma_+, \gamma) \in \mathcal{L}_k[\Gamma, \Gamma']\theta \wedge \Gamma, \Gamma' \text{ disjoint} \Rightarrow$
 897 $\exists \sigma, \gamma, \sigma', \gamma'. \sigma_+ = \sigma + \sigma' \wedge \gamma, \gamma' \text{ disjoint} \wedge \gamma_+ = \gamma \cup \gamma' \wedge (\varsigma, \gamma) \in$
 898 $\mathcal{L}_k[\Gamma] \wedge (\varsigma', \gamma') \in \mathcal{L}_k[\Gamma']$

899 PROOF: By induction on Γ' .

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900 **C.8** $\forall e, \sigma, e', \sigma', \theta. \langle \sigma, e \rangle \rightarrow \langle \sigma', e' \rangle \Rightarrow \langle \theta(\sigma), \theta(e) \rangle \rightarrow \langle \theta(\sigma'), \theta(e') \rangle$

901 PROOF: By induction on \rightarrow .

902 $\langle 1 \rangle 1$. ASSUME: Arbitrary $e, \sigma, e', \sigma', \theta$ such that $\langle \sigma, e \rangle \rightarrow \langle \sigma', e' \rangle$.

903 $\langle 1 \rangle 2$. SUFFICES: To consider only the following rules which mention fractional permission
904 variables.

905 OP_FRAC_PERM, OP_SHARE, OP_UNSHARE_(N)EQ and OP_GEMM_(MIS)MATCH.

906 $\langle 1 \rangle 3$. CASE: OP_FRAC_PERM.

907 Because substitution avoids capture,

908 $\langle \theta(\sigma), \theta(\text{fun } 'fc \rightarrow v) [f] \rangle \rightarrow \langle \theta(\sigma' [f/fc]), \theta(v [f/fc]) \rangle$.

909 $\langle 1 \rangle 4$. The rest of the cases are parametric in their use of fractional permission variables and
910 so will take the same step after any substitution.

911 $\langle 1 \rangle 5$. COROLLARY: If $\langle \sigma [f_1/fc], e [f_1/fc] \rangle \rightarrow^n \langle \sigma_2, e'_2 \rangle$ and $\langle \sigma [f_2/fc], e [f_2/fc] \rangle \rightarrow^n \langle \sigma_2, e'_2 \rangle$,
912 then $\exists \sigma, e'. \sigma_1 = \sigma [f_1/fc] \wedge \sigma_2 = \sigma [f_2/fc] \wedge e'_1 = e' [f_1/fc] \wedge e'_2 = e' [f_2/fc]$.

913 **D** Soundness

914 $\forall \Theta, \Delta, \Gamma, e, t. \Theta; \Delta; \Gamma \vdash e : t \Rightarrow \forall k. {}_k \llbracket \Theta; \Delta; \Gamma \vdash e : t \rrbracket$

915 PROOF SKETCH: Induction over the typing judgements.

916
917 ASSUME: 1. Arbitrary $\Theta, \Delta, \Gamma, e, t$ such that $\Theta; \Delta; \Gamma \vdash e : t$.

918 2. Arbitrary $k, \theta, \delta, \gamma, \sigma$ such that:

919 a. $\Theta = \text{dom}(\theta)$

920 b. $\delta \in \mathcal{I}_k \llbracket \Delta \rrbracket \theta$.

921 c. $(\varsigma, \gamma) \in \mathcal{L}_k \llbracket \Gamma \rrbracket \theta$

922 3. W.l.o.g., all variables are distinct, hence Θ , $\text{dom}(\Delta)$ and $\text{dom}(\Gamma)$ are disjoint so
923 order of θ , δ and γ (as substitutions defined recursively over expressions) is
924 irrelevant.

925
926 PROVE: $(\varsigma, \theta(\delta(\gamma(e)))) \in \mathcal{C}_k \llbracket \theta(t) \rrbracket$.

927 ASSUME: Arbitrary $j < k$ and σ_r , such that $\varsigma \star \varsigma_r$ defined.

928 SUFFICES: $\langle \sigma + \sigma_r, e \rangle \rightarrow^j \text{err} \vee \exists \sigma_f, e_f. \langle \sigma + \sigma_r, e \rangle \rightarrow^j \langle \sigma_f + \sigma_r, e_f \rangle$

929 $\wedge (e_f \text{ is a value} \Rightarrow (\varsigma_f \star \varsigma_r, e_f) \in \mathcal{V}_{k-j} \llbracket t \rrbracket)$.

930 SUFFICES: By C.1, to show $\langle \sigma, e \rangle \rightarrow^j \text{err} \vee \exists \sigma_f, e_f. \langle \sigma, e \rangle \rightarrow^j \langle \sigma_f, e_f \rangle$

931 $\wedge (e_f \text{ is a value} \Rightarrow (\varsigma_f, e_f) \in \mathcal{V}_{k-j} \llbracket t \rrbracket)$

932
933 $\langle 1 \rangle 1$. CASE: TY_LET.

934 $\langle 2 \rangle 1$. By induction,

935 1. $\forall k. {}_k \llbracket \Theta; \Delta; \Gamma \vdash e : t \rrbracket$

936 2. $\forall k. {}_k \llbracket \Theta; \Delta; \Gamma', x : t \vdash e' : t' \rrbracket$.

937 $\langle 2 \rangle 2$. By 2c, 3 and C.7, we know there exists the following (for all k):
 938 1. $(\varsigma_e, \gamma_e) \in \mathcal{L}_k[\Gamma]$
 939 2. $\gamma = \gamma_e \cup \gamma_{e'}$
 940 3. $\sigma = \sigma_e + \sigma_{e'}$.

941 $\langle 2 \rangle 3$. So, using $k, \theta, \delta, \gamma_e, \sigma_e$, we have $(\varsigma_e, \theta(\delta(\gamma_e(e)))) \in \mathcal{C}_k[\theta(t)]$.

942 $\langle 2 \rangle 4$. By $\langle 2 \rangle 2$ ($\gamma = \gamma_e \cup \gamma_{e'}$), have $(\varsigma_e, \theta(\delta(\gamma(e)))) \in \mathcal{C}_k[\theta(t)]$.

943 $\langle 2 \rangle 5$. By definition of $\mathcal{C}_k[\cdot]$ and $\langle 2 \rangle 2$, we instantiate with j and $\sigma_r = \sigma_{e'}$ to conclude
 944 that
 945 $\langle \theta(\sigma), \theta(\delta(\gamma(e))) \rangle$ either takes j steps to **err** or another heap-and-expression
 946 $\langle \sigma_f, e_f \rangle$.

947 $\langle 2 \rangle 6$. CASE: j steps to **err**
 948 By `OP_CONTEXT_ERR`, the whole expression reduces to **err** in $j < k$ steps.

949 $\langle 2 \rangle 7$. CASE: j steps to another heap-and-expression.
 950 If it is not a value, then `OP_CONTEXT` runs j times and we are done.

951 $\langle 2 \rangle 8$. If it is, then $\exists i \leq j. (\varsigma_f, v_1) \in \mathcal{V}_{k-i}[\theta(t_1)] \subseteq \mathcal{V}_{k-j}[\theta(t_1)]$ by C.3 and C.5.
 952 So, `OP_CONTEXT` runs i times, and then we have the following.
 953 SUFFICES: $(\varsigma_f \star \varsigma_{e'}, \mathbf{let } x = v \mathbf{ in } \theta(\delta(\gamma(e')))) \in \mathcal{C}_{k-i}[\theta(t')]$ by C.4 i times.
 954 SUFFICES: $(\varsigma_f \star \varsigma_{e'}, \theta(\delta(\gamma(e')))[v/x]) \in \mathcal{C}_{k-i-1}[\theta(t')]$ by C.4.

955 $\langle 2 \rangle 9$. By C.5, $(\varsigma_{e'}, \gamma_{e'}[x \mapsto v]) \in \mathcal{L}_k[\Gamma', x : t]\theta \subseteq \mathcal{L}_{k-i-1}[\Gamma', x : t]\theta$.

956 $\langle 2 \rangle 10$. Instantiate 2 of step $\langle 2 \rangle 1$ with $k - i - 1, \theta, \delta, \gamma_{e'}[x \mapsto v], \sigma_{e'}$ to conclude
 957 $(\varsigma_{e'}, \theta(\delta(\gamma_{e'}[x \mapsto v](e')))) \in \mathcal{C}_{k-i-1}[\theta(t')]$.

958 $\langle 2 \rangle 11$. By 3, we have $\theta(\delta(\gamma(e')))[v/x] = \theta(\delta(\gamma_{e'}[x \mapsto v](e')))$ and
 959 by C.1 we conclude $(\varsigma_f \star \varsigma_{e'}, \theta(\delta(\gamma(e')))[v/x]) \in \mathcal{C}_{k-i-1}[\theta(t')]$.

960 $\langle 1 \rangle 2$. CASE: `TY_PAIR_ELIM`.
 961 PROOF SKETCH: Similar to `TY_LET`, but with the following key differences.

962 $\langle 2 \rangle 1$. When $(\varsigma_f, v) \in \mathcal{V}_{k-i}[\theta(t_1) \otimes \theta(t_2)]$, we have $v = (v_1, v_2)$.

963 $\langle 2 \rangle 2$. SUFFICES: $(\varsigma_{e'}, \theta(\delta(\gamma(e')))) \in \mathcal{C}_{k-i-1}[\theta(t')]$ by C.4 $i + 1$ times.

964 $\langle 2 \rangle 3$. By C.5, $(\varsigma_{e'}, \gamma_{e'}[a \mapsto v_1, b \mapsto v_2]) \in \mathcal{L}_k[\Gamma', a : t_1, b : t_2]\theta \subseteq \mathcal{L}_{k-i-1}[\Gamma', a : t_1, b : t_2]\theta$.
 965

966 $\langle 2 \rangle 4$. Instantiate $_{k-i-1}[\Theta; \Delta; \Gamma', a : t_1, b : t_2 \vdash e' : t']$ with $\theta, \delta, \gamma_{e'}[a \mapsto v_1, b \mapsto v_2], \sigma_{e'}$.

967 $\langle 2 \rangle 5$. By 3 (for $\gamma = \gamma_e \cup \gamma_{e'}$ and a, b), conclude $(\varsigma_{e'}, \theta(\delta(\gamma(e'[v_1/a][v_2/b]))) \in \mathcal{C}_{k-i-1}[\theta(t')]$.

968 $\langle 1 \rangle 3$. CASE: `TY_BANG_ELIM`.
 969 PROOF SKETCH: Similar to `TY_LET`, but with the following key differences.

970 $\langle 2 \rangle 1$. When $(\varsigma_f, v) \in \mathcal{V}_{k-i}[\theta(!t)]$, since $\mathcal{V}_{k-i}[\theta(!t)] = \mathcal{V}_{k-i}[\theta(t)]$,
 971 we have $\varsigma_f = \emptyset$ and $v = \mathbf{Many } v'$ for some $(\emptyset, v') \in \mathcal{V}_{k-i}[\theta(t)]$.

- 972 $\langle 2 \rangle 2$. SUFFICES: $(\varsigma_{e'}, \mathbf{let\ Many\ } x = \mathbf{Many\ } v' \mathbf{ in } \theta(\delta(\gamma(e')))) \in \mathcal{C}_{k-i}[\![\theta(t)]\!]$.
- 973 $\langle 2 \rangle 3$. SUFFICES: $(\varsigma_{e'}, \theta(\delta(\gamma(e')))[v/x]) \in \mathcal{C}_{k-i-1}[\![\theta(t)]\!]$ by C.4 $i + 1$ times.
- 974 $\langle 2 \rangle 4$. Instantiate $_{k-i-1}[\![\Theta; \Delta, x : t, \Gamma' \vdash e' : t']\!]$ with $\theta, \delta_{e'} = \delta[x \mapsto v'], \gamma_{e'}, \sigma_{e'}$.
- 975 $\langle 2 \rangle 5$. By 3, $(\varsigma_{e'}, \theta(\delta(\gamma(e')))[v/x]) \in \mathcal{C}_{k-i-1}[\![\theta(t)]\!]$.
- 976 $\langle 1 \rangle 4$. CASE: TY_UNIT_ELIM.
 977 PROOF SKETCH: Similar to TY_LET, but with the following key differences.
- 978 $\langle 2 \rangle 1$. When $(\varsigma_f, v) \in \mathcal{V}_{k-i}[\![\mathbf{unit}]\!]$, we have $\varsigma_f = \emptyset$ and $v = ()$.
- 979 $\langle 2 \rangle 2$. SUFFICES: $(\varsigma_{e'}, \theta(\delta(\gamma(e')))) \in \mathcal{C}_{k-i-1}[\![\theta(t')]\!]$ by C.4 $i + 1$ times.
- 980 $\langle 2 \rangle 3$. By C.5, $(\varsigma_{e'}, \gamma_{e'}) \in \mathcal{L}_k[\![\Gamma']\!]\theta \subseteq \mathcal{L}_{k-i-1}[\![\Gamma']\!]\theta$.
- 981 $\langle 2 \rangle 4$. Instantiate $_{k-i-1}[\![\Theta; \Delta; \Gamma' \vdash e' : t']\!]$ with $\theta, \delta, \gamma_{e'}, \sigma_{e'}$.
- 982 $\langle 2 \rangle 5$. By 3 $(\varsigma_{e'}, \theta(\delta(\gamma(e')))) \in \mathcal{C}_{k-i-1}[\![\theta(t')]\!]$.
- 983 $\langle 1 \rangle 5$. CASE: TY_BOOL_ELIM.
 984 PROOF SKETCH: Similar to TY_UNIT_ELIM but with $\mathbf{OP_IF_}\{\mathbf{TRUE}, \mathbf{FALSE}\}$, $\varsigma_f = \emptyset$
 985 and $v = \mathbf{Many\ true}$ or $v = \mathbf{Many\ false}$.
- 986 $\langle 1 \rangle 6$. CASE: TY_BANG_INTRO.
- 987 $\langle 2 \rangle 1$. We have, $e = v$ for some value $v \neq l$, $\Gamma = \emptyset$ and so
 988 $\forall k. _k[\![\Theta; \Delta; \cdot \vdash v : t]\!]$ by induction.
- 989 $\langle 2 \rangle 2$. SUFFICES: $(\emptyset, \mathbf{Many\ } \theta(\delta(v))) \in \mathcal{C}_k[\![\theta(t)]\!]$ by 2c ($\varsigma = \emptyset, \gamma = []$).
- 990 $\langle 2 \rangle 3$. Instantiate $_{k-1}[\![\Theta; \Delta; \cdot \vdash v : t]\!]$ with $\theta, \delta, \gamma = [], \sigma = \emptyset$ to obtain $(\emptyset, \theta(\delta(v))) \in$
 991 $\mathcal{C}_k[\![\theta(t)]\!]$.
- 992 $\langle 2 \rangle 4$. Instantiate $(\emptyset, \theta(\delta(v))) \in \mathcal{C}_k[\![\theta(t)]\!]$ with $j = 0$, $\sigma_r = \emptyset$ and C.3 ($\theta(\delta(v))$ is a value),
 993 to conclude $(\emptyset, \theta(\delta(v))) \in \mathcal{V}_k[\![\theta(t)]\!]$.
- 994 $\langle 2 \rangle 5$. By definition of $\mathcal{V}_k[\![\theta(t)]\!]$, C.3 and C.2 we have $(\emptyset, \mathbf{Many\ } \theta(\delta(v))) \in \mathcal{C}_k[\![\theta(t)]\!]$.
- 995 $\langle 1 \rangle 7$. CASE: TY_PAIR_INTRO.
- 996 $\langle 2 \rangle 1$. By 2c, 3 and C.7, we know there exists the following (for all k):
 997 1. $(\varsigma_1, \gamma_1) \in \mathcal{L}_k[\![\Gamma_1]\!]$
 998 2. $(\varsigma_2, \gamma_2) \in \mathcal{L}_k[\![\Gamma_2]\!]$
 999 3. $\gamma = \gamma_1 \cup \gamma_2$
 1000 4. $\sigma = \sigma_1 + \sigma_2$.
- 1001 $\langle 2 \rangle 2$. By induction,
 1002 1. $\forall k. _k[\![\Theta; \Delta; \Gamma_1 \vdash e_1 : t_1]\!]$
 1003 2. $\forall k. _k[\![\Theta; \Delta; \Gamma_2 \vdash e_2 : t_2]\!]$.
- 1004 $\langle 2 \rangle 3$. Instantiate the first with $k, \theta, \delta, \gamma_1, \sigma_1$.

- 1005 $\langle 2 \rangle 4$. By that and $\langle 2 \rangle 1$, $(\varsigma_1, \theta(\delta(\gamma_1(e_1)))) = (\varsigma_1, \theta(\delta(\gamma(e_1)))) \in \mathcal{C}_k[\![\theta(t)]\!]$.
- 1006 $\langle 2 \rangle 5$. So, $\langle \theta(\sigma_1 + \sigma_2), \theta(\delta(\gamma_1(e_1))) \rangle$ either takes j steps to **err** or a heap-and-expression
 1007 $\langle \sigma_{1f}, e_{1f} \rangle$.
- 1008 $\langle 2 \rangle 6$. CASE: j steps to **err**
 1009 By `OP_CONTEXT_ERR`, the whole expression reduces to **err** in $j < k$ steps.
- 1010 $\langle 2 \rangle 7$. CASE: j steps to another heap-and-expression.
 1011 If it is not a value, then `OP_CONTEXT` runs j times and we are done.
- 1012 $\langle 2 \rangle 8$. If it is, then $\exists i_1 \leq j. (\varsigma_{1f}, v_1) \in \mathcal{V}_{k-i_1}[\![\theta(t_1)]\!]$ by C.3 and C.5.
 1013 So, `OP_CONTEXT` runs i_1 times, and then we have the following.
 1014 SUFFICES: By C.4, $(\varsigma_{1f} \star \varsigma_2, (v_1, e_2)) \in \mathcal{C}_{k-i_1}[\![\theta(t_1) \otimes t_2]\!]$.
- 1015 $\langle 2 \rangle 9$. Instantiate the second IH with $k, \theta, \delta, \gamma_2, \sigma_2$.
- 1016 $\langle 2 \rangle 10$. So, $\langle \theta(\sigma_{1f} + \sigma_2), \theta(\delta(\gamma_2(e_2))) \rangle$ either takes j steps to **err** or a heap-and-expression
 1017 $\langle \sigma_{2f}, e_{2f} \rangle$.
- 1018 $\langle 2 \rangle 11$. CASE: j steps to **err**
 1019 By `OP_CONTEXT_ERR`, the whole expression reduces to **err** in $j < k$ steps.
- 1020 $\langle 2 \rangle 12$. CASE: j steps to another heap-and-expression.
 1021 If it is not a value, then `OP_CONTEXT` runs j times and we are done.
- 1022 $\langle 2 \rangle 13$. If it is, then $\exists i_2 \leq j. (\varsigma_{2f}, v_2) \in \mathcal{V}_{k-i_2}[\![\theta(t_2)]\!]$ by C.3 and C.5.
 1023 So, `OP_CONTEXT` runs i_2 times, and then we have the following.
 1024 SUFFICES: By C.4, $(\varsigma_{1f} \star \varsigma_{2f}, (v_1, v_2)) \in \mathcal{V}_{k-i_1-i_2}[\![\theta(t_1) \otimes \theta(t_2)]\!]$.
- 1025 $\langle 2 \rangle 14$. By C.5 and $k - i_1 - i_2 \leq k - i_1, k - i_2$, have
 1026 $(\varsigma_{1f}, v_1) \in \mathcal{V}_{k-i_1}[\![\theta(t_1)]\!]$ and
 1027 $(\varsigma_{2f}, v_2) \in \mathcal{V}_{k-i_2}[\![\theta(t_2)]\!]$ as needed.
- 1028 $\langle 1 \rangle 8$. CASE: `TY_LAMBDA`.
 1029 SUFFICES: By C.2, to show $(\varsigma, \theta(\delta(\gamma(\mathbf{fun} x : t \rightarrow e)))) \in \mathcal{V}_k[\![\theta(t \rightarrow t')]\!]$.
 1030 ASSUME: Arbitrary $j \leq k$, $(\varsigma_v, v) \in \mathcal{V}_j[\![\theta(t)]\!]$ such that $\varsigma \star \varsigma_v$ is defined.
 1031 SUFFICES: $(\varsigma \star \varsigma_v, \theta(\delta(\gamma(\mathbf{fun} x : t \rightarrow e))) v) \in \mathcal{C}_j[\![\theta(t')]\!]$.
 1032 SUFFICES: $(\varsigma \star \varsigma_v, \theta(\delta(\gamma(e)))[v/x]) \in \mathcal{C}_{j-1}[\![\theta(t')]\!]$ by C.4.
- 1033 $\langle 2 \rangle 1$. By induction, $\forall k. {}_k[\![\Theta; \Delta; \Gamma, x : t \vdash e]\!]$.
- 1034 $\langle 2 \rangle 2$. Instantiate it $j - 1, \theta, \delta, \gamma[x \mapsto v], \sigma + \sigma_v$.
- 1035 $\langle 2 \rangle 3$. Hence, $(\varsigma \star \varsigma_v, \theta(\delta(\gamma[x \mapsto v](e)))) \in \mathcal{C}_{j-1}[\![\theta(t')]\!]$.
- 1036 $\langle 2 \rangle 4$. By 3, $\theta(\delta(\gamma[x \mapsto v](e))) = \theta(\delta(\gamma(e)))[v/x]$, we are done.
- 1037 $\langle 1 \rangle 9$. CASE: `TY_APP`.
- 1038 $\langle 2 \rangle 1$. By 2c, 3 and C.7, we know there exists the following (for all k):
 1039 1. $(\varsigma_e, \gamma_e) \in \mathcal{L}_k[\![\Gamma_e]\!]$

- 1040 2. $(\varsigma_{e'}, \gamma_{e'}) \in \mathcal{L}_k[\Gamma_{e'}]$
 1041 3. $\gamma = \gamma_e \cup \gamma_{e'}$
 1042 4. $\sigma = \sigma_e + \sigma_{e'}$.
- 1043 $\langle 2 \rangle 2$. By induction,
 1044 1. $\forall k. {}_k[\Theta; \Delta; \Gamma \vdash e : t' \multimap t]$
 1045 2. $\forall k. {}_k[\Theta; \Delta; \Gamma' \vdash e' : t']$.
- 1046 $\langle 2 \rangle 3$. Instantiate the first with $k, \theta, \delta, \gamma_e, \sigma_e$ to conclude $(\varsigma_e, \theta(\delta(\gamma_e(e)))) \in \mathcal{C}_k[\theta(t') \multimap \theta(t)]$.
 1047
- 1048 $\langle 2 \rangle 4$. Instantiate *this* with j and $\sigma_{e'}$ and use $\langle 2 \rangle 1$ to conclude $\langle \theta(\sigma_e + \sigma_{e'}), \theta(\delta(\gamma(e))) \rangle$
 1049 either takes j steps to **err** or a heap-and-expression $\langle \sigma_f + \sigma_{e'}, e_f \rangle$.
- 1050 $\langle 2 \rangle 5$. CASE: j steps to **err**
 1051 By OP_CONTEXT_ERR, the whole expression reduces to **err** in $j < k$ steps.
- 1052 $\langle 2 \rangle 6$. CASE: j steps to another heap-and-expression.
 1053 If it is not a value, then OP_CONTEXT runs j times and we are done.
- 1054 $\langle 2 \rangle 7$. If it is, then $\exists i_e \leq j. (\varsigma_f, e_f) \in \mathcal{V}_{k-i_e}[\theta(t') \multimap \theta(t)] \subseteq \mathcal{V}_{k-j}[\dots]$ by C.3 and C.5.
 1055 So, OP_CONTEXT runs i_e times, and then we have the following.
 1056 SUFFICES: By C.4 i_e times, $(\varsigma_f \star \varsigma_{e'}, e_f e') \in \mathcal{C}_{k-i_e}[\theta(t')]$.
- 1057 $\langle 2 \rangle 8$. By C.5, $(\varsigma_{e'}, \gamma_{e'}) \in \mathcal{L}_k[\Gamma'] \theta \subseteq \mathcal{L}_{k-i_e}[\Gamma'] \theta$.
- 1058 $\langle 2 \rangle 9$. So, instantiate the second IH with $k - i_e, \theta, \delta, \gamma_{e'}, \sigma_{e'}$ to conclude
 1059 $(\varsigma_{e'}, \theta(\delta(\gamma_{e'}(e')))) \in \mathcal{C}_{k-i_e}[\theta(t')]$.
- 1060 $\langle 2 \rangle 10$. Instantiate *this* with $j - i_e$ and σ_f to conclude $\langle \theta(\sigma_f + \sigma_{e'}), \theta(\delta(\gamma_{e'}(e')))) \rangle$
 1061 either takes $j - i_e$ steps to **err** or $\langle \sigma_f + \sigma'_f, e'_f \rangle$.
- 1062 $\langle 2 \rangle 11$. CASE: $j - i_e$ steps to **err**
 1063 By OP_CONTEXT_ERR, the whole expression reduces to **err** in $j - i_e < k - i_e$
 1064 steps.
- 1065 $\langle 2 \rangle 12$. CASE: $j - i_e$ steps to another heap-and-expression.
 1066 If it is not a value, then OP_CONTEXT runs $j - i_e$ times and we are done.
- 1067 $\langle 2 \rangle 13$. If it is, then $\exists i_{e'} \leq j - i_e. (\varsigma'_f, v_{e'}) \in \mathcal{V}_{k-i_e-i_{e'}}[\theta(t')]$ by C.3.
 1068 So, OP_CONTEXT runs $i_{e'}$ times, and then we have the following.
 1069 SUFFICES: By C.4 $i_{e'}$ times, $(\varsigma_f \star \varsigma'_f, e_f e'_f) \in \mathcal{C}_{k-i_e-i_{e'}}[\theta(t')]$.
- 1070 $\langle 2 \rangle 14$. Instantiate $(\varsigma_f, e_f) \in \mathcal{V}_{k-i_e}[\theta(t') \multimap \theta(t)]$ with $k - i_e - i_{e'} \leq k - i_e$ and
 1071 $(\varsigma_{v'}, v_{e'}) \in \mathcal{V}_{k-i_e-i_{e'}}[\theta(t')]$, to conclude $(\varsigma_f \star \varsigma'_f, e_f e'_f) \in \mathcal{C}_{k-i_e-i_{e'}}[\theta(t)]$ as
 1072 needed.
- 1073 $\langle 1 \rangle 10$. CASE: TY_GEN.
- 1074 $\langle 2 \rangle 1$. By induction, $\forall k. {}_k[\Theta, fc; \Delta; \Gamma \vdash e : t]$.
- 1075 $\langle 2 \rangle 2$. LET: f be arbitrary; $\theta' \equiv \theta[fc \mapsto f]$.

1076 Instantiate induction hypothesis with $k - 1, \theta', \delta, \gamma, \sigma$,
 1077 to conclude $(\varsigma, \theta'(\gamma(\delta(e)))) \in \mathcal{C}_{k-1}[\![\theta'(t)]\!]$ (for all f , by C.8).

1078 $\langle 2 \rangle 3$. Instantiate *this* with j and \emptyset to conclude $\langle \theta(\sigma), \theta'(\gamma(\delta(e))) \rangle$
 1079 either takes j steps to **err** or a heap-and-expression $\langle \sigma', e' \rangle$ (for all f , by C.8).

1080 $\langle 2 \rangle 4$. CASE: j steps to **err**.
 1081 By `OP_CONTEXT_ERR`, whole expression reduces to **err** in $j < k - 1$ steps
 1082 (for $f = fc$).

1083 $\langle 2 \rangle 5$. CASE: j steps to another heap-and-expression.
 1084 If it is not a value, then for $f = fc$, `OP_CONTEXT` runs j times and we are
 1085 done.

1086 $\langle 2 \rangle 6$. If it is, then $\exists i_e \leq j. (\varsigma', e') \in \mathcal{V}_{k-1-i_e}[\![\theta'(t)]\!] \subseteq \mathcal{V}_{k-1-j}[\![\dots]\!]$
 1087 by C.3 and C.5 (for all f , by C.8).

1088 $\langle 2 \rangle 7$. So, `OP_CONTEXT` runs i_e times, and then we have the following.
 1089 SUFFICES: By C.4 i_e times, $(\varsigma', \mathbf{fun}'fc \rightarrow e') \in \mathcal{V}_{k-i_e}[\![\theta('fc. t)]\!]$ (for $f = fc$).

1090 $\langle 2 \rangle 8$. ASSUME: Arbitrary f' .
 1091 SUFFICES: $(\varsigma', e'[f'/fc]) \in \mathcal{V}_{k-1-i_e}[\![\theta(t)[f'/fc]]\!]$ (for $f = fc$).

1092 $\langle 2 \rangle 9$. This is true by instantiating $\langle 2 \rangle 6$ with $f = f'$.

1093 $\langle 1 \rangle 11$. CASE: `TY_SPC`.

1094 $\langle 2 \rangle 1$. By induction, $\forall k. {}_k[\![\Theta; \Delta; \Gamma \vdash e : 'fc. t]\!]$.

1095 $\langle 2 \rangle 2$. Instantiate with $k, \theta, \delta, \gamma, \sigma$ to conclude $(\varsigma, \theta(\delta(\gamma(e)))) \in \mathcal{C}_k[\![\theta('fc. t)]\!]$.

1096 $\langle 2 \rangle 3$. Instantiate *this* with j and \emptyset and to conclude $\langle \theta(\sigma), \theta(\delta(\gamma(e))) \rangle$
 1097 either takes j steps to **err** or a heap-and-expression $\langle \sigma_f, e_f \rangle$.

1098 $\langle 2 \rangle 4$. CASE: j steps to **err**.
 1099 By `OP_CONTEXT_ERR`, the whole expression reduces to **err** in $j < k$ steps.

1100 $\langle 2 \rangle 5$. CASE: j steps to another heap-and-expression.
 1101 If it is not a value, then `OP_CONTEXT` runs j times and we are done.

1102 $\langle 2 \rangle 6$. If it is, then $\exists i_e \leq j. (\varsigma_f, e_f) \in \mathcal{V}_{k-i_e}[\![\theta('fc. t)]\!] \subseteq \mathcal{V}_{k-j}[\![\dots]\!]$ by C.3 and C.5.
 1103 So $e_f \equiv \mathbf{fun}'fc \rightarrow v$ for some v .

1104 $\langle 2 \rangle 7$. So, `OP_CONTEXT` runs i_e times, and then we have the following.
 1105 SUFFICES: By C.4 i_e times, $(\varsigma_f, (\mathbf{fun}'fc \rightarrow v)[f]) \in \mathcal{C}_{k-i_e}[\![\theta(t[f/fc])]\!]$.
 1106 SUFFICES: By C.4 once more, $(\varsigma_f, v[f/fc]) \in \mathcal{C}_{k-i_e-1}[\![\theta(t[f/fc])]\!]$.

1107 $\langle 2 \rangle 8$. This is true by instantiating $\langle 2 \rangle 6$ with f and C.2.

1108 $\langle 1 \rangle 12$. CASE: `TY_FIX`.
 1109 SUFFICES: $(\emptyset, \theta(\delta(\mathbf{fix}(g, x : t, e : t')))) \in \mathcal{V}_k[\![\theta(t \multimap t')]\!]$, by C.2 ($\sigma = \{\}$, $\gamma = []$).
 1110 ASSUME: Arbitrary $j \leq k$, $(\varsigma_v, v) \in \mathcal{V}_j[\![\theta(t)]\!]$ ($\varsigma = \emptyset$, so $\varsigma \star \varsigma_v$ is defined).

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1111 LET: $\tilde{e} \equiv \theta(\delta(e))$.
 1112 SUFFICES: $(\varsigma_v, \mathbf{fix}(g, x : t, \tilde{e} : t') \ v) \in \mathcal{C}_j[\![\theta(t')]\!]$.
 1113 SUFFICES: $(\varsigma_v, \tilde{e} \ [v/x] \ [\mathbf{fix}(g, x : t, \tilde{e} : t')/g]) \in \mathcal{C}_{j-1}[\![\theta(t')]\!]$ by C.4.

1114 $\langle 2 \rangle 1$. By induction, $\forall k. _k[\![\Theta; \Delta, g : t \multimap t'; x : t \vdash e : t']]\!]$.
 1115 $\langle 2 \rangle 2$. Instantiate this with $j - 1, \delta[g \mapsto \mathbf{fix}(g, x : t, \tilde{e} : t')], \gamma = [x \mapsto v], \sigma_v$.
 1116 $\langle 2 \rangle 3$. We have $(\emptyset, \mathbf{fix}(g, x : t, \tilde{e} : t')) \in \mathcal{V}_{j-1}[\![\theta(t \multimap t')]\!]$.

1117 $\langle 3 \rangle 1$. Again by induction (over k), $(\emptyset, \mathbf{fix}(g, x : t, \tilde{e} : t')) \in \mathcal{C}_{j-1}[\![\theta(t \multimap t')]\!]$.
 1118 $\langle 3 \rangle 2$. Instantiate *this* with $j = 0$ and \emptyset and we are done.

1119 $\langle 2 \rangle 4$. We have $(\varsigma_v, v) \in \mathcal{V}_{j-1}[\![\theta(t)]\!]$ by assumption and C.5.
 1120 $\langle 2 \rangle 5$. So we conclude $(\varsigma_v, \theta(\delta'(\gamma(e)))) \in \mathcal{C}_{j-1}[\![\theta(t')]\!]$ as required.

1121 $\langle 1 \rangle 13$. CASE: TY_VAR_LIN.
 1122 PROVE: $(\varsigma, \theta(\delta(\gamma(x)))) \in \mathcal{C}_k[\![\theta(t)]\!]$.

1123 $\langle 2 \rangle 1$. $\Gamma = \{x : t\}$ by assumption of TY_VAR_LIN.
 1124 $\langle 2 \rangle 2$. SUFFICES: $(\varsigma, \gamma(x)) \in \mathcal{C}_k[\![\theta(t)]\!]$ by 3 (θ and δ irrelevant).
 1125 $\langle 2 \rangle 3$. By 2c, there exist $(\varsigma_x, v_x) \in \mathcal{V}_k[\![\theta(t)]\!]$, such that $\varsigma = \varsigma_x$ and $\gamma = [x \mapsto v_x]$.
 1126 $\langle 2 \rangle 4$. Hence, $(\varsigma_x, v_x) \in \mathcal{C}_k[\![\theta(t)]\!]$, by C.2.

1127 $\langle 1 \rangle 14$. CASE: TY_VAR.
 1128 PROVE: $(\varsigma, \theta(\delta(\gamma(x)))) \in \mathcal{C}_k[\![\theta(t)]\!]$.

1129 $\langle 2 \rangle 1$. $x : t \in \Delta$ and $\Gamma = \emptyset$ by assumption of TY_VAR.
 1130 $\langle 2 \rangle 2$. SUFFICES: $(\emptyset, \delta(x)) \in \mathcal{C}_k[\![\theta(t)]\!]$ by 3.
 1131 $\langle 2 \rangle 3$. By 2b, there exists v_x such that $(\emptyset, v_x) \in \mathcal{V}_k[\![\theta(t)]\!]$ (θ irrelevant and γ empty).
 1132 $\langle 2 \rangle 4$. Hence, $(\emptyset, v_x) \in \mathcal{C}_k[\![\theta(t)]\!]$, by C.2.

1133 $\langle 1 \rangle 15$. CASE: TY_UNIT_INTRO.
 1134 True by C.2 and definition of $\mathcal{V}_k[\![\mathbf{unit}]\!]$.

1135 $\langle 1 \rangle 16$. CASE: TY_BOOL_TRUE, TY_BOOL_FALSE, TY_INT_INTRO, TY_ELT_INTRO.
 1136 Similar to TY_UNIT_INTRO.

1137 D.1 Well-formed types

1138 $\boxed{\Theta \vdash f \text{ Perm}}$ Well-formed fractional permissions

1139 $\frac{fc \in \Theta}{\Theta \vdash fc \text{ Perm}}$ WF_PERM_VAR

1140	$\frac{}{\Theta \vdash 1 \text{ Perm}}$	WF_PERM_ZERO
1141	$\frac{\Theta \vdash f \text{ Perm}}{\Theta \vdash \frac{1}{2}f \text{ Perm}}$	WF_PERM_SUCC
1142	$\boxed{\Theta \vdash t \text{ Type}}$	Well-formed types
1143	$\frac{}{\Theta \vdash \text{unit Type}}$	WF_TYPE_UNIT
1144	$\frac{}{\Theta \vdash \text{bool Type}}$	WF_TYPE_BOOL
1145	$\frac{}{\Theta \vdash \text{int Type}}$	WF_TYPE_INT
1146	$\frac{}{\Theta \vdash \text{elt Type}}$	WF_TYPE_ELT
1147	$\frac{\Theta \vdash f \text{ Perm}}{\Theta \vdash f \text{ arr Type}}$	WF_TYPE_ARRAY
1148	$\frac{\Theta \vdash t \text{ Type}}{\Theta \vdash !t \text{ Type}}$	WF_TYPE_BANG
1149	$\frac{\Theta, fc \vdash t \text{ Type}}{\Theta \vdash 'fc.t \text{ Type}}$	WF_TYPE_GEN
1150	$\frac{\Theta \vdash t \text{ Type} \quad \Theta \vdash t' \text{ Type}}{\Theta \vdash t \otimes t' \text{ Type}}$	WF_TYPE_PAIR
1151	$\frac{\Theta \vdash t \text{ Type} \quad \Theta \vdash t' \text{ Type}}{\Theta \vdash t \multimap t' \text{ Type}}$	WF_TYPE_LOLLY
1152		

1153 **E** NumLin Grammar

	m	::=		matrix expressions
			M	matrix variables
			$m + m'$	matrix addition
			$m m'$	matrix multiplication
1154			(m)	S
	f	::=		fractional permission
			fc	variable
			1	whole permission
			$\frac{1}{2}f$	

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1155	t	::=		linear type
		unit		unit
		bool		boolean (true/false)
		int		63-bit integers
		elt		array element
		f arr		arrays
		f mat		matrices
		$!t$		multiple-use type
		$'fc.t$	bind fc in t	frac. perm. generalisation
		$t \otimes t'$		pair
		$t \multimap t'$		linear function
		(t)	S	parentheses
	p	::=		primitive
		not		boolean negation
		$(+)$		integer addition
		$(-)$		integer subtraction
		$(*)$		integer multiplication
		$(/)$		integer division
		$(=)$		integer equality
		$(<)$		integer less-than
		$(+.)$		element addition
		$(-.)$		element subtraction
		$(*.)$		element multiplication
		$(/.)$		element division
		$(=.)$		element equality
		$(<.)$		element less-than
		set		array index assignment
		get		array indexing
		share		share array
		unshare		unshare array
		free		free array
		array		Owl: make array
		copy		Owl: copy array
		sin		Owl: map sine over array
		hypot		Owl: $x_i := \sqrt{x_i^2 + y_i^2}$
		asum		BLAS: $\sum_i x_i $
		axpy		BLAS: $x := \alpha x + y$
		dot		BLAS: $x \cdot y$
		rotmg		BLAS: see its docs
		scal		BLAS: $x := \alpha x$
		amax		BLAS: $\operatorname{argmax} i : x_i$
		setM		matrix index assignment
		getM		matrix indexing

		shareM		share matrix
		unshareM		unshare matrix
		freeM		free matrix
		matrix		Owl: make matrix
		copyM		Owl: copy matrix
		copyM_to		Owl: copy matrix onto another
		sizeM		dimension of matrix
		trnsp		transpose matrix
		gemm		BLAS: $C := \alpha A^{T?} B^{T?} + \beta C$
		symm		BLAS: $C := \alpha AB + \beta C$
		posv		BLAS: Cholesky decomp. and solve
		potrs		BLAS: solve with given Cholesky
		syrk		BLAS: $C := \alpha A^{T?} A^{T?} + \beta C$
v	::=			values
		<i>p</i>		primitives
		<i>x</i>		variable
		()		unit introduction
		true		true
		false		false
		<i>k</i>		integer
		<i>l</i>		heap location
1156		<i>el</i>		array element
		Many <i>v</i>		!-introduction
		fun <i>'fc</i> $\rightarrow v$		frac. perm. abstraction
		(<i>v</i> , <i>v'</i>)		pair introduction
		fun <i>x</i> : <i>t</i> $\rightarrow e$	bind <i>x</i> in <i>e</i>	abstraction
		fix (<i>g</i> , <i>x</i> : <i>t</i> , <i>e</i> : <i>t'</i>)	bind <i>g</i> $\cup x$ in <i>e</i>	fixpoint
		(<i>v</i>)	S	parentheses
e	::=			expression
		<i>p</i>		primitives
		<i>x</i>		variable
		let <i>x</i> = <i>e</i> in <i>e'</i>	bind <i>x</i> in <i>e'</i>	let binding
		()		unit introduction
		let () = <i>e</i> in <i>e'</i>		unit elimination
		true		true
		false		false
		if <i>e</i> then <i>e</i> ₁ else <i>e</i> ₂		if
		<i>k</i>		integer
		<i>l</i>		heap location
		<i>el</i>		array element
		Many <i>e</i>		!-introduction
		let Many <i>x</i> = <i>e</i> in <i>e'</i>		!-elimination
		fun <i>'fc</i> $\rightarrow e$		frac. perm. abstraction

		$e[f]$		frac. perm. specialisation
		(e, e')		pair introduction
		let $(a, b) = e$ in e'	bind $a \cup b$ in e'	pair elimination
		fun $x : t \rightarrow e$	bind x in e	abstraction
		$e e'$		application
		fix $(g, x : t, e : t')$	bind $g \cup x$ in e	fixpoint
		(e)	S	parentheses
C	$::=$			evaluation contexts
		let $x = [-]$ in e	bind x in e	let binding
		let $() = [-]$ in e		unit elimination
		if $[-]$ then e_1 else e_2		if
		Many $[-]$!-introduction
		let Many $x = [-]$ in e		!-elimination
		fun $'fc \rightarrow [-]$		frac. perm. abstraction
		$[-][f]$		frac. perm. specialisation
		$([-], e)$		pair introduction
		$(v, [-])$		pair introduction
		let $(a, b) = [-]$ in e	bind $a \cup b$ in e	pair elimination
		$[-]e$		application
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Θ	$::=$			fractional permission environment
		\cdot		
		Θ, fc		
Γ	$::=$			linear types environment
		\cdot		
		$\Gamma, x : t$		
		Γ, Γ'		
Δ	$::=$			intuitionistic types environment
		\cdot		
		$\Delta, x : t$		
σ	$::=$			heap (multiset of triples)
		$\{\}$		empty heap
		$\sigma + \{l \mapsto_f m_{k_1, k_2}\}$		location l points to matrix m
$Config$	$::=$			result of small step
		$\langle \sigma, e \rangle$		heap and expression
		err		error

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F Primitives

```

module Arr = Owl.Dense.Ndarray.D
type z = Z
type 'a s = Succ
type 'a arr = A of Arr.arr [@@unboxed]
type 'a mat = M of Arr.arr [@@unboxed]
type 'a bang = Many of 'a [@@unboxed]
module Prim :
sig
  val extract : 'a bang -> 'a
  (** Boolean *)
  val not_ : bool bang -> bool bang
  (** Arithmetic, many omitted for brevity *)
  val addI : int bang -> int bang -> int bang
  val eqI : int bang -> int bang -> bool bang
  (** Arrays *)
  val set : z arr -> int bang -> float bang -> z arr
  val get : 'a arr -> int bang -> 'a arr * float bang
  val share : 'a arr -> 'a s arr * 'a s arr
  val unshare : 'a s arr -> 'a s arr -> 'a arr
  val free : z arr -> unit
  (** Owl *)
  val array : int bang -> z arr
  val copy : 'a arr -> 'a arr * z arr
  val sin : z arr -> z arr
  val hypot : z arr -> 'a arr -> 'a arr * z arr
  (** Level 1 BLAS *)
  val asum : 'a arr -> 'a arr * float bang
  val axpy : float bang -> 'a arr -> z arr -> 'a arr * z arr
  val dot : 'a arr -> 'b arr -> ('a arr * 'b arr) * float bang
  val scal : float bang -> z arr -> z arr
  val amax : 'a arr -> 'a arr * int bang
  (* Matrix, some omitted for brevity *)
  val matrix : int bang -> int bang -> z mat
  val eye : int bang -> z mat
  val copy_mat : 'a mat -> 'a mat * z mat
  val copy_mat_to : 'a mat -> z mat -> 'a mat * z mat
  val size_mat : 'a mat -> 'a mat * (int bang * int bang)
  val transpose : 'a mat -> 'a mat * z mat
  (* Level 3 BLAS/LAPACK *)
  val gemm : float bang -> ('a mat * bool bang) -> ('b mat * bool bang) ->
    float bang -> z mat -> ('a mat * 'b mat) * z mat
  val symm : bool bang -> float bang -> 'a mat -> 'b mat -> float bang ->
    z mat -> ('a mat * 'b mat) * z mat
  val gesv : z mat -> z mat -> z mat * z mat
  val posv : z mat -> z mat -> z mat * z mat
  val potrs : 'a mat -> z mat -> 'a mat * z mat
  val syrks : bool bang -> float bang -> 'a mat -> float bang -> z mat ->
    'a mat * z mat
end

```

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G Kalman Filters from NumLin and C

```

let kalman sigma h mu r_1 data_1 =
  let h, _p_k_n_p_ = Prim.size_mat h in
  let k, n = _p_k_n_p_ in
  let sigma_hT = Prim.matrix n k in
  let (sigma, h), sigma_hT =
    Prim.gemm (Many 1.) (sigma, Many false) (h, Many true) (Many 0.) sigma_hT in
  let (h, sigma_hT), r_2 =
    Prim.gemm (Many 1.) (h, Many false) (sigma_hT, Many false) (Many 1.) r_1 in
  let k_by_k, x = Prim.posv_flip r_2 sigma_hT in
  let (h, mu), data_2 =
    Prim.gemm (Many 1.) (h, Many false) (mu, Many false) (Many (-1.)) data_1 in
  let (x, data_2), new_mu =
    Prim.gemm (Many 1.) (x, Many false) (data_2, Many false) (Many 1.) mu in
  let x_h = Prim.matrix n n in
  let (x, h), x_h =
    Prim.gemm (Many 1.) (x, Many false) (h, Many false) (Many 0.) x_h in
  let () = Prim.free_mat x in
  let sigma, sigma2 = Prim.copy_mat sigma in
  let (sigma, x_h), new_sigma =
    Prim.symm (Many true) (Many (-1.)) sigma x_h (Many 1.) sigma2 in
  let () = Prim.free_mat x_h in
  ((sigma, h), (new_sigma, (new_mu, (k_by_k, data_2))))

```

■ **Figure 19** OCaml code for a Kalman filter, generated (at *compile time*) from the code in Figure 12, passed through `ocamlformat` for presentation.

```

static void kalman( const int n,          const int k,
                  const double *sigma, /* n,n */ const double *h, /* k,n */
                  const double *mu, /* n,1 */ double *r, /* k,k */
                  double *data, /* k,1 */ double **ret_sigma /* n,n */ ) {
  double* n_by_k = (double *) malloc(n * k * sizeof(double));
  cblas_dgemm(RowMajor, NoTrans, Trans, n, k, n, 1., sigma, n, h, n, 0., n_by_k, k);
  cblas_dgemm(RowMajor, NoTrans, NoTrans, k, k, n, 1., h, n, n_by_k, k, 1., r, k);
  LAPACKE_dposv(LAPACK_COL_MAJOR, 'U', k, n, r, k, n_by_k, k);
  cblas_dgemm(RowMajor, NoTrans, NoTrans, k, 1, n, 1., h, n, mu, 1, -1., data, 1);
  cblas_dgemm(RowMajor, NoTrans, NoTrans, n, 1, k, 1., n_by_k, k, data, 1, 1., mu, 1);
  double* n_by_n = (double *) malloc(n * n * sizeof(double));
  cblas_dgemm(RowMajor, NoTrans, NoTrans, n, n, k, 1., n_by_k, k, h, n, 0., n_by_n, n);
  free(n_by_k);
  double* n_by_n2 = (double *) malloc(n * n * sizeof(double));
  cblas_dcopy(n*n, sigma, 1, n_by_n2, 1);
  cblas_dsymm(RowMajor, Right, Upper, n, n, -1., sigma, n, n_by_n, n, 1., n_by_n2, n);
  free(n_by_n);
  *ret_sigma = n_by_n2; }

```

■ **Figure 20** CBLAS/LAPACKE implementation of a Kalman filter. I used C instead of Fortran because it is what Owl uses under the hood and OCaml FFI support for C is better and easier to use than that for Fortran. A distinct ‘`measure_kalman`’ function that sandwiches a call to this function with `getrusage` is omitted for brevity.