

$fc$  fractional capability variable  
 $x, g, a, b$  expression variable  
 $k$  integer variable  
 $el$  array-element variable

$symb ::=$   
 $\lambda$   
 $\otimes$   
 $\multimap$   
 $\vdash$   
 $\in$   
 $\forall$   
 $\text{Cap}$   
 $\text{Type}$   
 $!$   
 $\rightarrow$   
 $\text{value}$

$f ::=$   
 $fc$  fractional capability variable  
 $\mathbf{Z}$  zero  
 $\mathbf{S}f$  successor

$t ::=$  linear type  
 $\text{unit}$  unit  
 $\text{bool}$  boolean (true/false)  
 $\text{int}$  63-bit integers  
 $\text{elt}$  array element  
 $f \text{ arr}$  arrays  
 $f \text{ mat}$  matrices  
 $!t$  multiple-use type  
 $\forall fc. t$  bind  $fc$  in  $t$  frac. cap. generalisation  
 $t \otimes t'$  pair  
 $t \multimap t'$  linear function  
 $t\{f/fc\}$  M substitution  
 $(t)$  S parentheses

$e ::=$  expression  
 $p$  primitives (arithmetic, L1 BLAS, Owl)  
 $x$  variable  
 $()$  unit introduction  
 $\text{true}$  true (boolean introduction)  
 $\text{false}$  false (boolean introduction)  
 $\text{if } e \text{ then } e_1 \text{ else } e_2$  if (boolean elimination)  
 $k$  integer  
 $el$  array element  
 $\text{many } e$  packing-up a non-linear value  
 $\text{let many } x = e \text{ in } e'$  using a non-linear value  
 $\text{fun } fc \rightarrow e$  frac. cap. abstraction  
 $e[f]$  frac. cap. specialisation  
 $(e, e')$  pair introduction  
 $\text{let } (a, b) = e \text{ in } e'$  bind  $a \cup b$  in  $e'$  pair elimination

		<b>fun</b> $x : t \rightarrow e$	bind $x$ in $e$	abstraction
		$e\ e'$		application
		<b>fix</b> $(g, x : t, e : t')$	bind $g \cup x$ in $e$	fixpoint
$p$	::=			primitive
		<b>set</b>		array index assignment
		<b>get</b>		array indexing
		$(+)$		integer addition
		$(-)$		integer subtraction
		$(*)$		integer multiplication
		$(/)$		integer division
		$(=)$		integer equality
		$(<)$		integer less-than
		$(+.)$		element addition
		$(-.)$		element subtraction
		$(*.)$		element multiplication
		$(/.)$		element division
		$(=.)$		element equality
		$(<.)$		element comparison (less-than)
		$(\&\&)$		boolean conjunction
		$(  )$		boolean disjunction
		<b>not</b>		boolean negation
		<b>share</b>		share array
		<b>unshare</b>		unshare array
		<b>free</b>		free array
		<b>array</b>		Owl: make array
		<b>copy</b>		Owl: copy array
		<b>sin</b>		Owl: sine of all elements in array
		<b>hypot</b>		Owl: $x_i := \sqrt{x_i^2 + y_i^2}$
		<b>asum</b>		BLAS: $\sum_i  x_i $
		<b>axpy</b>		BLAS: $x := \alpha x + y$
		<b>dot</b>		BLAS: $x \cdot y$
		<b>rotmg</b>		BLAS: gen. mod. Givens rotation
		<b>scal</b>		BLAS: $x := \alpha x$
		<b>amax</b>		BLAS: index of maximum absolute value
		<b>setM</b>		matrix index assignment
		<b>getM</b>		matrix indexing
		<b>shareM</b>		share matrix
		<b>unshareM</b>		unshare matrix
		<b>freeM</b>		free matrix
		<b>matrix</b>		Owl: make matrix
		<b>copyM</b>		Owl: copy matrix
		<b>gemv</b>		BLAS: $y := \alpha A^{T?} x + \beta y$
		<b>symv</b>		BLAS: $y := \alpha A_{\text{sym}} x + \beta y$
		<b>trmv</b>		BLAS: $x := A^{T?} * x$
		<b>trsv</b>		BLAS: $x := A^{-1 \cdot T?} * x$

	$ \begin{array}{ l} \mathbf{ger} \\ \mathbf{gemm} \\ \mathbf{trmm} \\ \mathbf{trsm} \end{array} $	BLAS: $A := \alpha * x * y^T + A$ BLAS: $C := \alpha * A^{T?} * B^{T?} + \beta C$ BLAS: $B := \alpha * A^{T?} * B$ and swapped BLAS: $B := \alpha * A^{-1 \cdot T?} * B$ and swapped
$\Theta$	$ \begin{array}{ l} ::= \\   \cdot \\   \Theta, fc \end{array} $	fractional capability environment
$\Gamma$	$ \begin{array}{ l} ::= \\   \cdot \\   \Gamma, x : t \\   \Gamma, \Gamma' \end{array} $	linear types environment
$\Delta$	$ \begin{array}{ l} ::= \\   \cdot \\   \Delta, x : t \end{array} $	linear types environment
<i>formula</i>	$ \begin{array}{ l} ::= \\   \textit{judgement} \\   x : t \in \Delta \\   x : t \in \Gamma \\   fc \in \Theta \\   \mathbf{value}(e) \end{array} $	
<i>Well_Formed</i>	$ \begin{array}{ l} ::= \\   \Theta \vdash f \mathbf{Cap} \\   \Theta \vdash t \mathbf{Type} \end{array} $	Valid fractional capabilities Valid types
<i>Values</i>	$ \begin{array}{ l} ::= \\   \mathbf{value}(e) \end{array} $	Value restriction for !-introduction
<i>Types</i>	$ \begin{array}{ l} ::= \\   \Theta; \Delta; \Gamma \vdash e : t \end{array} $	Typing rules for expressions (no primitives yet)
<i>judgement</i>	$ \begin{array}{ l} ::= \\   \textit{Well_Formed} \\   \textit{Values} \\   \textit{Types} \end{array} $	
<i>user_syntax</i>	$ \begin{array}{ l} ::= \\   fc \\   x \\   k \\   el \\   symb \\   f \end{array} $	

	$t$
	$e$
	$p$
	$\Theta$
	$\Gamma$
	$\Delta$
	$formula$

$\boxed{\Theta \vdash f \text{ Cap}}$  Valid fractional capabilities

$\frac{fc \in \Theta}{\Theta \vdash fc \text{ Cap}}$	WF_CAP_VAR
$\overline{\Theta \vdash \mathbf{Z} \text{ Cap}}$	WF_CAP_ZERO
$\frac{\Theta \vdash f \text{ Cap}}{\Theta \vdash \mathbf{S} f \text{ Cap}}$	WF_CAP_SUCC

$\boxed{\Theta \vdash t \text{ Type}}$  Valid types

$\overline{\Theta \vdash \mathbf{unit} \text{ Type}}$	WF_TYPE_UNIT
$\overline{\Theta \vdash \mathbf{bool} \text{ Type}}$	WF_TYPE_BOOL
$\overline{\Theta \vdash \mathbf{int} \text{ Type}}$	WF_TYPE_INT
$\overline{\Theta \vdash \mathbf{elt} \text{ Type}}$	WF_TYPE_ELT
$\frac{\Theta \vdash f \text{ Cap}}{\Theta \vdash f \mathbf{arr} \text{ Type}}$	WF_TYPE_ARRAY
$\frac{\Theta \vdash t \text{ Type}}{\Theta \vdash !t \text{ Type}}$	WF_TYPE_BANG
$\frac{\Theta, fc \vdash t \text{ Type}}{\Theta \vdash \forall fc. t \text{ Type}}$	WF_TYPE_GEN
$\frac{\Theta \vdash t \text{ Type} \quad \Theta \vdash t' \text{ Type}}{\Theta \vdash t \otimes t' \text{ Type}}$	WF_TYPE_PAIR
$\frac{\Theta \vdash t \text{ Type} \quad \Theta \vdash t' \text{ Type}}{\Theta \vdash t \multimap t' \text{ Type}}$	WF_TYPE_LOLLY

$\boxed{\mathbf{value}(e)}$  Value restriction for !-introduction

$\overline{\mathbf{value}(p)}$	VAL_PRIM
$\overline{\mathbf{value}(() )}$	VAL_UNIT_INTRO
$\overline{\mathbf{value}(\mathbf{true})}$	VAL_BOOL_TRUE
$\overline{\mathbf{value}(\mathbf{false})}$	VAL_BOOL_FALSE

$$\begin{array}{c}
\frac{}{\mathbf{value}(k)} \quad \text{VAL\_INT\_INTRO} \\
\frac{}{\mathbf{value}(el)} \quad \text{VAL\_ELT\_INTRO} \\
\frac{}{\mathbf{value}(x)} \quad \text{VAL\_VAR} \\
\frac{}{\mathbf{value}(\mathbf{fix}(g, x : t, e : t'))} \quad \text{VAL\_FIX} \\
\frac{}{\mathbf{value}(\mathbf{fun } x : t \rightarrow e)} \quad \text{VAL\_LAMBDA} \\
\frac{\mathbf{value}(e)}{\mathbf{value}(\mathbf{fun } fc \rightarrow e)} \quad \text{VAL\_GEN} \\
\frac{\mathbf{value}(e)}{\mathbf{value}(e[fc])} \quad \text{VAL\_SPC} \\
\frac{\mathbf{value}(e)}{\mathbf{value}(\mathbf{many } e)} \quad \text{VAL\_BANG\_INTRO} \\
\frac{\mathbf{value}(e_1) \quad \mathbf{value}(e_2)}{\mathbf{value}((e_1, e_2))} \quad \text{VAL\_PAIR\_INTRO}
\end{array}$$

$\Theta; \Delta; \Gamma \vdash e : t$

Typing rules for expressions (no primitives yet)

$$\begin{array}{c}
\frac{}{\Theta; \Delta; \cdot, x : t \vdash x : t} \quad \text{TY\_VAR\_LIN} \\
\frac{x : t \in \Delta}{\Theta; \Delta; \cdot \vdash x : t} \quad \text{TY\_VAR} \\
\frac{}{\Theta; \Delta; \cdot \vdash () : !\mathbf{unit}} \quad \text{TY\_UNIT\_INTRO} \\
\frac{}{\Theta; \Delta; \cdot \vdash \mathbf{true} : !\mathbf{bool}} \quad \text{TY\_BOOL\_TRUE} \\
\frac{}{\Theta; \Delta; \cdot \vdash \mathbf{false} : !\mathbf{bool}} \quad \text{TY\_BOOL\_FALSE} \\
\frac{\Theta; \Delta; \Gamma \vdash e : \mathbf{bool} \quad \Theta; \Delta; \Gamma' \vdash e_1 : t' \quad \Theta; \Delta; \Gamma' \vdash e_2 : t'}{\Theta; \Delta; \Gamma, \Gamma' \vdash \mathbf{if } e \mathbf{ then } e_1 \mathbf{ else } e_2 : t} \quad \text{TY\_BOOL\_ELIM} \\
\frac{}{\Theta; \Delta; \cdot \vdash k : !\mathbf{int}} \quad \text{TY\_INT\_INTRO} \\
\frac{}{\Theta; \Delta; \cdot \vdash el : !\mathbf{elt}} \quad \text{TY\_ELT\_INTRO} \\
\frac{\mathbf{value}(e) \quad \Theta; \Delta; \cdot \vdash e : t}{\Theta; \Delta; \cdot \vdash \mathbf{many } e : !t} \quad \text{TY\_BANG\_INTRO} \\
\frac{\Theta; \Delta; \Gamma \vdash e : !t \quad \Theta; \Delta, x : t; \Gamma' \vdash e' : t'}{\Theta; \Delta; \Gamma, \Gamma' \vdash \mathbf{let many } x = e \mathbf{ in } e' : t'} \quad \text{TY\_BANG\_ELIM} \\
\frac{\Theta; \Delta; \Gamma \vdash e : t \quad \Theta; \Delta; \Gamma' \vdash e' : t'}{\Theta; \Delta; \Gamma, \Gamma' \vdash (e, e') : t \otimes t'} \quad \text{TY\_PAIR\_INTRO}
\end{array}$$

$$\begin{array}{c}
\frac{\Theta; \Delta; \Gamma \vdash e_{12} : t_1 \otimes t_2 \quad \Theta; \Delta; \Gamma', a : t_1, b : t_2 \vdash e : t}{\Theta; \Delta; \Gamma, \Gamma' \vdash \mathbf{let} (a, b) = e_{12} \mathbf{in} e : t} \text{TY\_PAIR\_ELIM} \\
\\
\frac{\Theta \vdash t' \text{Type} \quad \Theta; \Delta; \Gamma, x : t' \vdash e : t}{\Theta; \Delta; \Gamma \vdash \mathbf{fun} x : t' \rightarrow e : t' \multimap t} \text{TY\_LAMBDA} \\
\\
\frac{\Theta; \Delta; \Gamma \vdash e : t' \multimap t \quad \Theta; \Delta; \Gamma' \vdash e' : t'}{\Theta; \Delta; \Gamma, \Gamma' \vdash e e' : t} \text{TY\_APP} \\
\\
\frac{\Theta, fc; \Delta; \Gamma \vdash e : t}{\Theta; \Delta; \Gamma \vdash \mathbf{fun} fc \rightarrow e : \forall fc. t} \text{TY\_GEN} \\
\\
\frac{\Theta \vdash f \text{Cap} \quad \Theta; \Delta; \Gamma \vdash e : \forall fc. t}{\Theta; \Delta; \Gamma \vdash e[f] : t\{f/fc\}} \text{TY\_SPC} \\
\\
\frac{\Theta; \Delta, g : t \multimap t'; \cdot, x : t \vdash e : t'}{\Theta; \Delta; \cdot \vdash \mathbf{fix} (g, x : t, e : t') : !(t \multimap t')} \text{TY\_FIX}
\end{array}$$

Definition rules: 19 good 0 bad

Definition rule clauses: 43 good 0 bad