

Laboratory practice No. 1: RECURSION

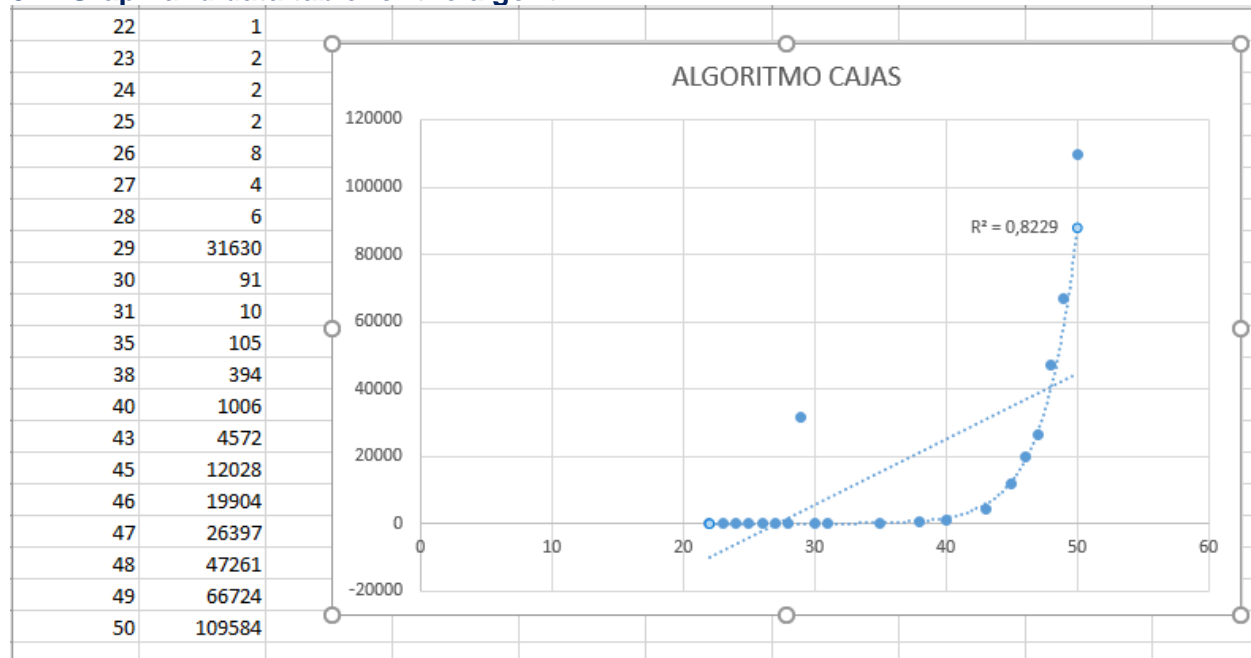
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3) Practice for final project defense presentation

3.1. The asymptotic complexity given for exercise 1.2 can be described with this formula:
 $T(n) = T(n-1) + T(n-2) + C$. Processed by Wolfram app to obtain the solution of this formula we got:
 $t(n) = c_1 2^{(n-1)} + C(2^n - 1)$. Therefore we can conclude that working and analyzing this equation the asymptotic complexity of this algorithm is $O(2^n)$.

3.2 Graph and data table for the algorithm:



Estimated time for the algorithm to calculate the shapes for a size 50x2:

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ESTRUCTURA DE DATOS 1

Código ST0245

For a size of 50x2 cm squared, the time it takes to calculate the ways in which rectangles can be placed is 109584 milliseconds, equivalent to 1.8264 minutes.

3.3 It would not work; the complexity makes the algorithm take time to process the information while the area of the rectangle goes up. It would be necessary to find a way to develop another algorithm with a lower complexity that could allow a faster data processing.

3.4 Algorithm:

```
public boolean groupSum5(int start, int[] nums, int target) {
    if (nums.length == start)
    {
        if(target == 0)
        {
            return true;
        }
        return false;
    }
    if (nums[start] % 5 != 0)
    {
        return groupSum5(start+1, nums, target-nums[start]) || groupSum5(start+1,
        nums, target);
    }
    else
    {
        if(start< nums.length-1 && nums[start+1] == 1)
        {
            return groupSum5(start+2, nums, target-nums[start]);
        }
        else
        {
            return groupSum5(start+1, nums, target-nums[start]);
        }
    }
}
```

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ESTRUCTURA DE DATOS 1

Código ST0245

"Given an array of ints, is it possible to choose a group of some of the ints, such that the group sums to the given target with these additional constraints: all multiples of 5 in the array must be included in the group. If the value immediately following a multiple of 5 is 1, it must not be chosen. (No loops needed.)" -CodingBat

the parameters are an initial position, an array of integers and a target that must be equal to the sum of certain integers of the int array.

first, the stop condition is whether the initial position or position to evaluate is equal to the length of the ints array, in that case, if we reached our target, then the algorithm will return true, otherwise, it will return false.

now, if the integer in the position is not divisible by 5, then the algorithm will evaluate two cases: the first, subtracting the int in the initial position from the target and increasing the initial position in 1, and the second, only increasing the initial position in 1.

if the integer in the position is divisible by 5, there are 2 cases, if the next integer in the array is the number 1 (taking into account that we are not in the last position of the array), the 1 will not be used and the algorithm will run again, subtracting the integer in the position from the target and increasing by 2 the initial position. the other case is just if the integer is just divisible by 5, hence, it will subtract the integer from the target and will run again with the initial position increased in 1.

3.5

- **Recursion 1:**
 - Algorithm Count7:

```
public int count7(int n) {
    int aux = 0;
    if(n<7){
        return 0;
    }else{
        if(n%10 == 7){
            aux = 1;
        }else{
            aux = 0;
        }
        return aux + count7(n/10);
    }
}
```

The asymptotic complexity of this algorithm is given by:

$$T(n) = T(n-1) + C$$

$$T(n) = c1 + C$$

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ESTRUCTURA DE DATOS 1

Código ST0245

complexity = $O(n)$

"n" is the number from which the 7's will be counted

- Algorithm CountX:

```
countX

public int countX(String str) {
    if (str.equals("")){
        return 0;
    }else if(str.charAt(0)=='x'){
        return 1+countX(str.substring(1));
    }
    return countX(str.substring(1));
}
```

$T(n) = T(n-1) + C$

$T(n) = c1 + C$

complexity = $O(n)$

"n" is the string length

- Algorithm PowerN:

```
public int powerN(int base, int n) {
    return n == 0 ? 1: base * powerN(base,n-1);
}
```

The asymptotic complexity of this algorithm is given by:

$T(n) = T(n-1) + C$

$T(n) = c1 + C$

complexity = $O(n)$

"n" is the exponent

- Algorithm noX:

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ESTRUCTURA DE DATOS 1

Código ST0245

```
public String noX(String str) {
    if (str.length() == 0)
    {
        return "";
    }
    else if (str.charAt(0) == 'x')
    {
        return noX(str.substring(1, str.length()));
    }
    else
    {
        return str.substring(0,1) + noX(str.substring(1, str.length()));
    }
}
```

The asymptotic complexity of this algorithm is given by:

$$T(n) = T(n-1) + C$$

$$T(n) = c_1 + Cn$$

$$\text{complexity} = O(n)$$

"n" is the string length

- Algorithm bunnyEars2:

```
public int bunnyEars2(int n) {
    if(n == 0)
    {
        return 0;
    }
    if(n % 2 == 0)
    {
        return 3 + bunnyEars2(n-1);
    }
    else
    {
        return 2 + bunnyEars2(n-1);
    }
}
```

The asymptotic complexity of this algorithm is given by:

$$T(n) = T(n-1) + C$$

$$T(n) = c_1 + Cn$$

$$\text{complexity} = O(n)$$

"n" is the quantity of bunnies

- **Recursion 2:**

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ESTRUCTURA DE DATOS 1

Código ST0245

- Algorithm splitArray:

```
public boolean splitArray(int[] nums) {
    return help(nums, 0, 0, 0);
}

public boolean help(int[] nums, int g1, int g2, int ind)
{
    if(nums.length == ind)
    {
        return (g1 == g2);
    }
    return help(nums, g1 + nums[ind], g2, ind+1) || help(nums, g1, g2+nums[ind],
ind+1);
}
```

The asymptotic complexity of this algorithm is given by:

$$T(n) = T(n-1) + T(n-1) + C$$

$$T(n) = c1 \cdot 2^{(n-1)} + C(2^n - 1)$$

$$\text{complexity} = O(2^n)$$

"n" is the quantity of elements in the array that have not been used yet

- Algorithm split53:

```
public boolean split53(int[] nums) {
    return help(nums, 0, 0, 0);
}

public boolean help(int[] nums, int g1, int g2, int ind)
{
    if(nums.length == ind)
    {
        return (g1 == g2);
    }
    if(nums[ind] % 5 == 0)
    {
        return help(nums, g1+nums[ind], g2, ind+1);
    }
    else if(nums[ind] % 3 == 0)
    {
        return help(nums, g1, g2 + nums[ind], ind+1);
    }
    return help(nums, g1 + nums[ind], g2, ind+1) || help(nums, g1, g2+nums[ind],
ind+1);
}
```

The asymptotic complexity of this algorithm is given by:

$$T(n) = T(n-1) + T(n-1) + C$$

$$T(n) = c1 \cdot 2^{(n-1)} + C(2^n - 1)$$

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ESTRUCTURA DE DATOS 1

Código ST0245

complexity = $O(2^n)$

“n” is the quantity of elements in the array that have not been used yet

- Algorithm splitOdd10:

```
public boolean splitOdd10(int[] nums) {
    return help(nums, 0, 0, 0);
}

public boolean help(int[] nums, int m10, int odd, int ind)
{
    if(nums.length == ind)
    {
        return (m10 % 10 == 0 && odd % 2 != 0);
    }
    return help(nums, m10+nums[ind], odd, ind+1) || help(nums, m10,
        odd+nums[ind], ind+1);
}
```

The asymptotic complexity of this algorithm is given by:

$$T(n) = T(n-1) + T(n-1) + C$$

$$T(n) = c1 \cdot 2^{(n-1)} + C(2^n - 1)$$

complexity = $O(2^n)$

“n” is the quantity of elements in the array that have not been used yet

- Algorithm groupNoAdj:

```
public boolean groupNoAdj(int start, int[] nums, int target) {
    if (nums.length <= start)
    {
        if(target == 0)
        {
            return true;
        }
        return false;
    }

    return groupNoAdj(start+2, nums, target-nums[start]) || groupNoAdj(start+1,
        nums, target);
}
```

The asymptotic complexity of this algorithm is given by:

$$T(n) = T(n-2) + T(n-1) + C$$

$$T(n) = c1Fn + c2Ln + C$$

complexity = $O(2^n)$

“n” is the quantity of elements in the array that have not been used yet

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ESTRUCTURA DE DATOS 1

Código ST0245

- Algorithm groupSum6:


```
public boolean groupSum6(int start, int[] nums, int target) {
    if (nums.length == start)
    {
        if(target == 0)
        {
            return true;
        }
        return false;
    }
    if (nums[start] != 6)
    {
        return groupSum6(start+1, nums, target-nums[start]) || groupSum6(start+1,
        nums, target);
    }
    else
    {
        return groupSum6(start+1, nums, target-nums[start]);
    }
}
```

The asymptotic complexity of this algorithm is given by:

$$T(n) = T(n-1) + T(n-1) + C$$

$$T(n) = c1 \cdot 2^{(n-1)} + C(2^n - 1)$$

$$\text{complexity} = O(2^n)$$

“n” is the quantity of elements in the array that have not been used yet

4) Practice for midterms

- 4.1 (start + 1, nums, target)
- 4.2 a
- 4.3.1 (n-a,a,b,c)
- 4.3.2 (res,solucionar(n,b,c,n+1))
- 4.3.3 (res,solucionar(n,c,n+1,n+1))
- 4.4 e
- 4.5.1 L2: return 1; // L3: n-2 // L4:n-4
- 4.5.2 b
- 4.6.1 sumaAux(n, i+2)
- 4.6.2 sumaAux(n, i+1)
- 4.7.1 S, i+1, t-S[i]
- 4.7.2 S, i+1, t

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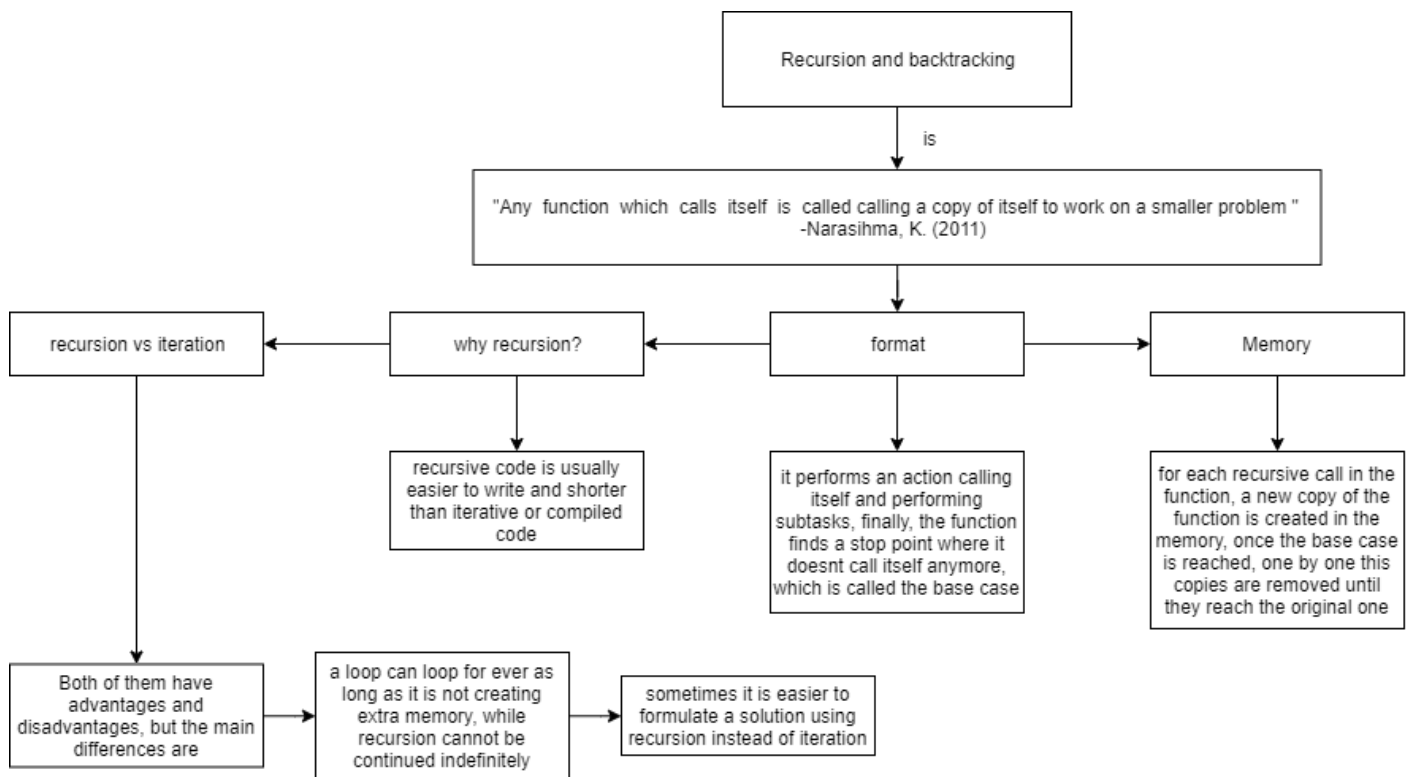
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ESTRUCTURA DE DATOS 1

Código ST0245

4.8.1 return 0
 4.8.2 $n_i + n_j$
 4.9 c.22
 4.10 b. 6
 4.11.1 $(n-1) // \text{lucas}(n-2)$
 4.11.2 c. $O(2^n)$
 4.12.1 sat
 4.12.2 Math.max(fi,fj)
 4.12.3 sat

5) Recommended reading (optional)



Narasimha, K. (2011) Data Structures And Algorithms. Recuperado de: <https://www.docdroid.net/ZPfHmS5/data-structures-and-algorithms-narasimha->

6) Teamwork and gradual progress (optional)

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ESTRUCTURA DE DATOS 1

Código ST0245

Reuniones:

Reuniones

Reunión proyecto

🕒 9 de ago.

Reunión proyecto - 2

🕒 13 de ago.

Reunión Laboratorio

🕒 14 de ago.

Reunión Laboratorio - 2

🕒 15 de ago.

Reunión Laboratorio - 3

🕒 16 de ago.

agosto 2019					
LUN.	MAR.	MIÉ.	JUE.	VIE.	SÁB.
5	6	7	8	9 1 tarjeta Reunión proyecto	10
12	13 1 tarjeta Reunión proyecto - 2	14 1 tarjeta Reunión Laboratorio	15 2 tarjetas Reunión Laboratorio - 2 Simulacro preguntas de sustentación de proyecto p1	16 2 tarjetas Reunión Laboratorio - 3 Lectura: Socialización y sintetización de información	17

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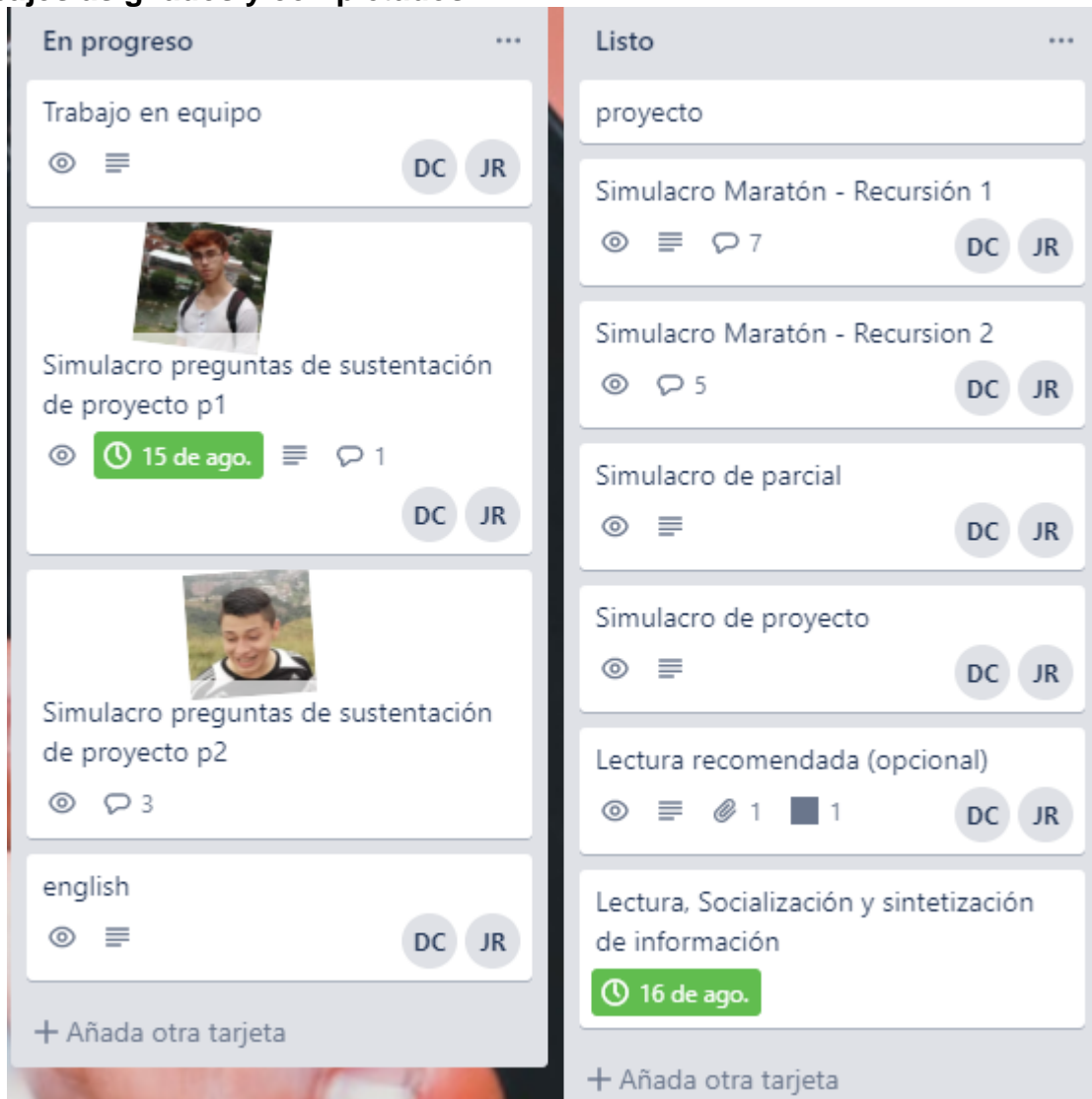
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ESTRUCTURA DE DATOS 1
Código ST0245

Trabajos asignados y completados:



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