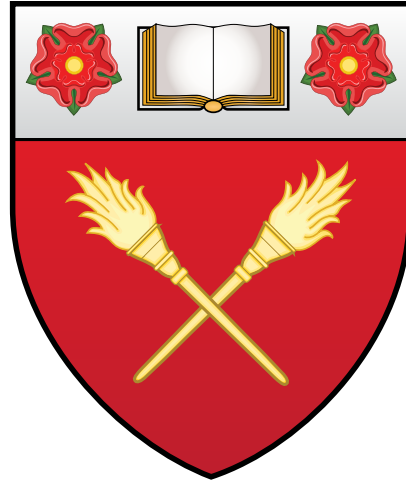


Dislocation Based Modelling of Fusion Relevant Materials.



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Dedication

“Ohana significa familia, y tu familia nunca te abandona, ni te olvida.”

“Ohana means family, family means nobody gets left behind, or forgotten.”

– Stitch, Lilo & Stitch

Malle santa, cuando te hablé desde el trabajo para darte la noticia me dijiste que siempre lo supiste. Recuerdo que lloraste en el teléfono, se te cerró la garganta y a mí también. A diferencia de otras veces, no pudimos platicar mucho porque nos quedamos sin aliento. Así que le hablaste a abuelita Raquel mientras le hablé a papá. Madre, David y yo te debemos tanto que no podemos repagar en mil vidas, pero ten por seguro que eres nuestro ejemplo a seguir. Nuestra madre chingona y chambeadora, solo hay una y como ella no hay ninguna. Sin tí el mundo sería un lugar más cruel y pobre, no merece un ángel tan grande y puro como tú.

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David, “Te mamaste we.” me dijiste cuando me abrazaste al llegar a casa después del trabajo el día que me aceptaron. Me has hecho un chingo de falta ojete, te extraño mucho we. Perdón por no hablar tan seguido, pero me duele colgar we. No estoy tan chisqueado como mamá pero siento feo cuando terminamos de platicar. Me gusta mucho ver tus streams porque me recuerda un poco a sentarme a tu lado a verte jugar... o a cuando veíamos Twitch juntos y veíamos a TB missing legal.

Esto es para mi Ohana por sangre y por elección. Ustedes creyeron en mí cuando yo no lo hacía. Me empujaron a ser mejor. Me extendieron la mano cuando nadie más lo hizo. Me hicieron reír cuando solo sabía llorar. Gracias por hacerme quien soy.

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“You can always judge a man by the quality of his enemies.”

– Oscar Wilde

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I guess I'll content myself with being judged unfavourably by the world.

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Outline

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2. DDD-FEM
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 - (a) Traction errors
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 - (c) Simulations
 - (d) Recommendations and Conclusions
4. Parallelising tractions
 - (a) Algorithms
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5. EasyDD v2.0
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 - (d) Collisions
 - (e) Code redesign
6. Future work
 - (a) Julia redesign
 - (b) Preliminary comparisons
7. Conclusions

Chapter 0

Preface

0.1 Notation

For the sake of clarity the following notation conventions have been used:

- Tensors are denoted by sans serif bold italics, \mathbf{T} .
- Matrices are denoted by serif bold roman, \mathbf{M} .
- Vectors are denoted by serif bold italics, \mathbf{V} .
- Scalars are denoted by serif italics, S .
- Host variables are denoted by $^{\text{h}}a$.
- Global device variables are denoted by $^{\text{d}}a$.
- Thread variables are denoted by $^{\text{t}}a$.

0.2 Conventions

All angles are given in radians unless stated otherwise. All pseudo-code is C-style and assuming C conventions (row-order, 0-start indexing, memory allocation).

Elements, be they Surface Elements (SEs), Finite Elements (FEs), dislocation line segments or else, are denoted by e and their total as E . Similarly nodes are denoted by n and their total N . Parallel indices are denoted by idx to distinguish them from regular indices/counters i , j , k . For disambiguation purposes the exponential function will be denoted as $\exp(x)$.

0.3 Typesetting

This document was typeset using a custom-made \LaTeX document class created by the author, compiled with \XeLaTeX , and the bibliography was produced with \BibTeX . The custom document class can be found in https://github.com/dcelisgarza/latex_template.

0.4 Diagrams

All diagrams are vector graphics drawn using the open-source image editing software [InkScape](#) for its ease of use, elegant simplicity, high quality outputs, and “Draw Freely” philosophy.

Chapter 1

Introduction

1.1 Section 1

1.1.1 Subsection 1

Chapter 2

Coupling Discrete Dislocation Dynamics to Finite Element Methods

2.1 Superposition Scheme

Coupling Discrete Dislocation Dynamics (DDD) to Finite Element Methods (FEMs) [1] is important to properly simulate micromechanical tests because DDD provides us with a more precise set of inputs and greater granularity for solving the FE problem. This can be achieved by using a so-called superposition scheme fig. 2.1 that enables the independent solution of both problems, whilst feeding information from one to the other in a continuous feedback loop.

2.2 Extracting Surface Nodes

The FEM coupler arranges the nodes starting on the xz -plane where $y = 0, \dots, ndy$, $n \in \mathbb{N}$. However in order to couple DDD to FEM we only require the surface nodes where displacements are not calculated. Because we're working with rectangular prisms, we can easily pick out the surface nodes using a search algorithm with a logical mask. MATLAB and Fortran provide vector intrinsics that allow one to do so. Figure 2.2a illustrates only the surface nodes according to our implementation's node arrangement—which is the xz -plane going from $y = y_{\min} \rightarrow y = y_{\max}$. However, due to the nature of the analytical solutions in chapter 3, we need all the surface nodes of the rectangular faces for which there are no displacements. This means that edge nodes are shared between 2 adjacent faces and corner nodes between 3 adjacent faces. In order to properly apply the logical mask to find *only* the surface nodes need to know that each FEs' nodes are numbered according to

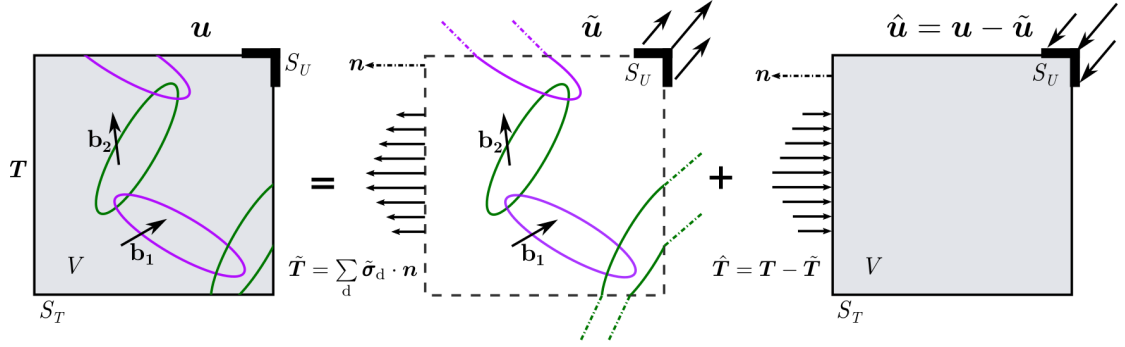
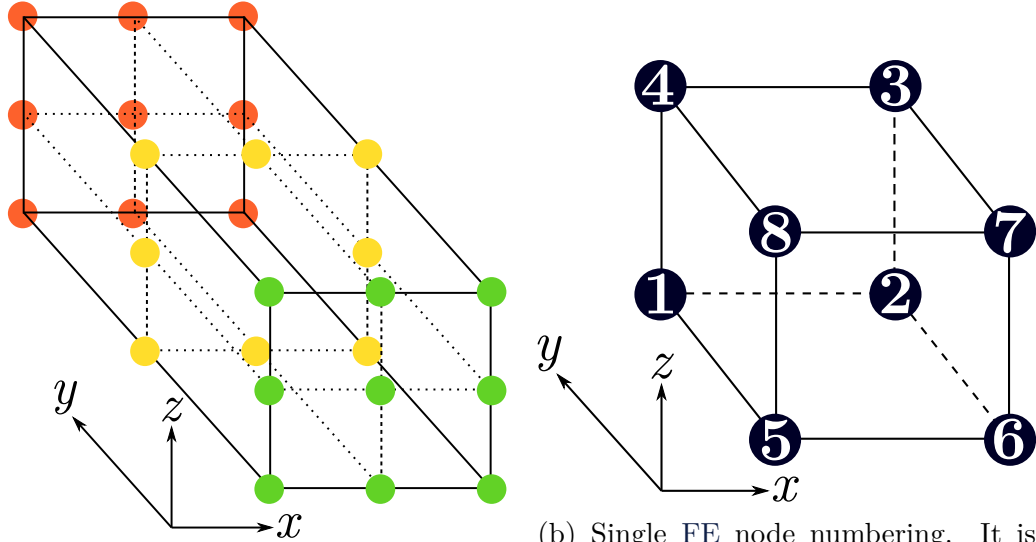


Figure 2.1: The dislocation ensemble in a volume V is bounded by surface S . First, the traction field $\sum_d \tilde{\sigma}_d$ due to the dislocation ensemble is evaluated at the surface. Then, a traditional FEM or Boundary Element Method (BEM) calculates the image traction field $\hat{\sigma} \times \mathbf{n}$. Which is then fed back to the DDD problem to evolve the dislocation positions and repeat the cycle. Image edited from [1].



(a) Arrangement of the surface nodes of our FE model. FE nodes are arranged in chunks of $\Delta x \Delta z$ nodes in our implementation, but we only want the surface nodes.

(b) Single FE node numbering. It is necessary to know which FE plane corresponds to which node labels. This lets us design the auxiliary matrix that selects nodes according to the planes we want to extract.

Figure 2.2: FE node arrangement for coupling to DDD.

fig. 2.2b.

Using fig. 2.2 one can work out which nodes are of interest to whichever surface is being extracted. The ordering of the nodes in the final array will depend on the definition of the problems in chapter 3.

Node selection remains an expensive operation and minimising array indexing is of the utmost importance for the best performance. Selecting nodes in the traditional sense, i.e. with code branching such as `if statements` or `case selection` is unmanageable, verbose and very prone to mistakes. The issue was solved by introducing an auxiliary matrix which defines various parameters that aid node selection and greatly reduces code size, improves readability, and eliminates the need for code branching. The matrix can be constructed utilising fig. 2.2 in order to know which nodes correspond to which FE planes. The p^{th} column of the matrix¹ corresponds to the p^{th} plane (according to an arbitrary plane numbering) and is defined as,

$$\mathbf{V}_p^T = [L_{1p} \quad L_{2p} \quad \cdots \quad L_{Np} \quad A_p \quad C_p], \quad (2.1)$$

where L_{np} is the numeric label for node n as given by fig. 2.2, A_p is the area of the plane, and C_p is the numeric label of the orthogonal coordinate to the plane $C_p = 1, 2, 3$ for the x, y, z coordinates respectively. A_p lets us segment our output and transitional arrays so that the only data being modified is that which corresponds to the correct plane and C_p lets us know which coordinate we must use in our selection criteria. Using our particular node labelling scheme (with dimensions $\Delta x, \Delta y, \Delta z$ respectively in the x, y, z directions), the matrix is defined as,

$$\mathbf{V} = \begin{bmatrix} 5 & 2 & 6 & 1 & 5 & 4 \\ 1 & 6 & 5 & 2 & 6 & 3 \\ 8 & 3 & 7 & 4 & 1 & 8 \\ 4 & 7 & 8 & 3 & 2 & 7 \\ \Delta y \Delta z & \Delta y \Delta z & \Delta x \Delta z & \Delta x \Delta z & \Delta x \Delta y & \Delta x \Delta y \\ 1 & 1 & 2 & 2 & 3 & 3 \end{bmatrix}. \quad (2.2)$$

The information codified in eq. (2.2) lets us index and process only the necessary columns to extract the surface nodes we're interested in. The advantage of this setup over a naïve implementation is that it can be relatively easily expanded, maintained, and is general enough that it lends itself to a variety of selection criteria. The columns from left to right (1 to 6) represent: face 1 $\equiv \min(x)$, yz -

¹MATLAB uses column-major ordering, so this gives us the best performance for vectorised code.

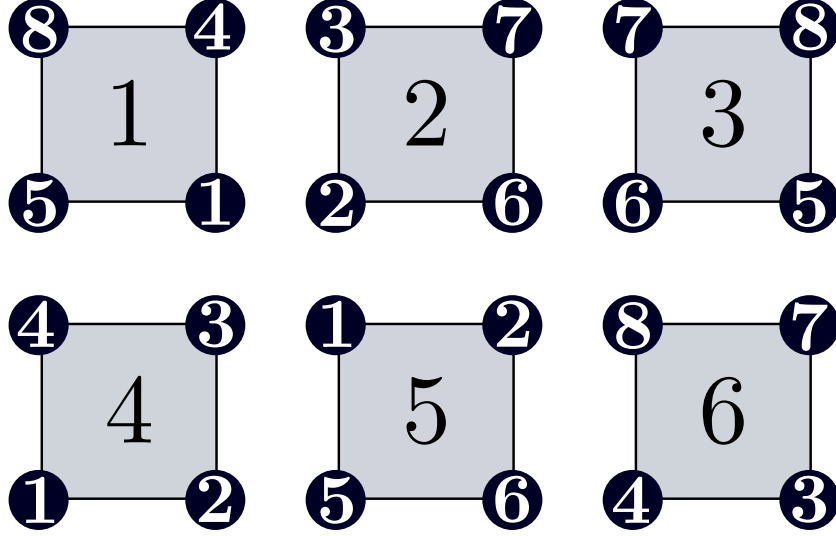


Figure 2.3: Self-consistent, chirality preserving surface planes. Plane and node numbering according to the columns of eq. (2.2) and fig. 2.2.

plane; face 2 $\equiv \max(x)$, yz -plane; face 3 $\equiv \min(y)$, xz -plane; face 4 $\equiv \max(y)$, xz -plane; face 5 $\equiv \min(z)$, xy -plane; face 6 $\equiv \max(z)$, xy -plane.

It is important to make sure the nodes are extracted in a self consistent manner because the problem definitions in chapter 3 assume a specific node ordering. In particular, the labelling chirality must remain the same. This is due to the fact that the problems in chapter 3 utilise internal coordinates derived from a set of basis vectors, and scalar projections onto them. If the relationship between the vectors is kept, but the chirality is different (opposite) the calculated quantities will have the wrong sign. If however the relationship between the vectors is different to the problem formulation, the program will crash due to division by 0. Self-consistency was achieved by placing an imaginary observer inside the FE model facing the $\min(x)$, yz -plane (face 1) and rotating it to view all 6 planes according to fig. 2.3.

2.3 Mapping Forces

$\hat{\mathbf{F}}$ only cares about the global node numbers, of which there are M . The columns of \mathbf{N}_L , \mathbf{F}_N represent the 4 nodes of a SE, the rows represent the SE they are part of. The column vector \mathbf{x} represents the x , y , z -coordinates of the node-element or global node corresponding to their subscript. The column vector $\boldsymbol{\gamma}_t$ is the list

Algorithm 2.1 If $\hat{\mathbf{F}}$'s columns are arranged the same way as γ_t .

```

1:  ▷ Loop through the array containing the node labels of the relevant surface
   nodes.
2:  for  $i = 0; i < \text{length}(\gamma_t); i++$  do
3:      ▷ Save the global node label for the current iteration.
4:       $n \leftarrow \gamma_t[i]$ 
5:      ▷ Use the node label to find a vector with
   the linearised index of all surface nodes whose labels correspond to the node
   whose forces we want to pass on to the FEM coupler.
6:       $\mathbf{L} \leftarrow \text{find}(\mathbf{N}_{\mathbf{L}} == n)$ 
7:      ▷ Loop over coordinates.
8:      for  $k = 0; k < 3; k++$  do
9:          ▷ Use global node label vector to index the
   force array from the analytical force calculation due to dislocations on surface
   nodes. Multiplied by 3 because there are three coordinates per node. We sum
   the forces from the analytical calculation because the same global node can be
   part of multiple surface elements. We add  $k$  because the  $x, y, z$  coordinates
   are consecutively stored in  $\mathbf{F}_{\mathbf{n}}$ .
10:          $\hat{\mathbf{F}}[3 \times \mathbf{L} + k] \leftarrow \hat{\mathbf{F}}[3 \times \mathbf{L} + k] + \sum \mathbf{F}_{\mathbf{n}}[3 \times \mathbf{L} + k]$ 
11:     end for
12: end for

```

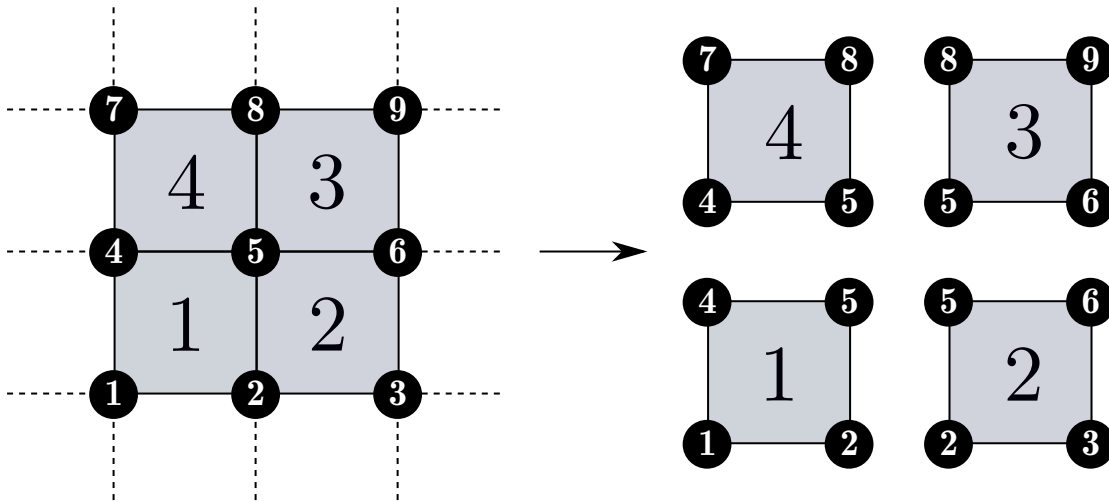


Figure 2.4: FE nodes shared by multiple SEs.

of global nodes for which the forces induced by dislocations must be calculated.

$$\mathbf{x}_{e,n}^\top \equiv \begin{bmatrix} x_{en} & y_{en} & z_{en} \end{bmatrix}, \quad \hat{\mathbf{x}}_n^\top \equiv \begin{bmatrix} x_n & y_n & z_n \end{bmatrix} \quad (2.3)$$

$$\mathbf{N}_L = \begin{bmatrix} l_{1,1} & l_{1,2} & l_{1,4} & l_{1,5} \\ l_{2,2} & l_{2,3} & l_{2,5} & l_{2,6} \\ l_{3,5} & l_{3,6} & l_{3,8} & l_{3,9} \\ l_{4,4} & l_{4,5} & l_{4,7} & l_{4,8} \end{bmatrix}, \quad \boldsymbol{\gamma} = \begin{bmatrix} l_1 \\ l_2 \\ \vdots \\ l_9 \end{bmatrix} \quad (2.4)$$

$$\mathbf{F}_e = \begin{bmatrix} \mathbf{x}_{1,1} & \mathbf{x}_{1,2} & \mathbf{x}_{1,3} & \mathbf{x}_{1,4} \\ \mathbf{x}_{2,1} & \mathbf{x}_{2,2} & \mathbf{x}_{2,3} & \mathbf{x}_{2,4} \\ \mathbf{x}_{3,1} & \mathbf{x}_{3,2} & \mathbf{x}_{3,3} & \mathbf{x}_{3,4} \\ \mathbf{x}_{4,1} & \mathbf{x}_{4,2} & \mathbf{x}_{4,3} & \mathbf{x}_{4,4} \end{bmatrix}, \quad \hat{\mathbf{F}} = \begin{bmatrix} \hat{\mathbf{x}}_1 \\ \hat{\mathbf{x}}_2 \\ \hat{\mathbf{x}}_3 \\ \hat{\mathbf{x}}_4 \end{bmatrix}. \quad (2.5)$$

Chapter 3

Analytical Forces Induced by Dislocations on Linear Rectangular Surface Elements

3.1 Forces Exerted by a Dislocation Line Segment on Linear Rectangular Surface Elements

Coupling DDD to FEM requires the traction field $\boldsymbol{\sigma}^\infty(\mathbf{x}) \cdot \mathbf{n}(\mathbf{x})$ to be distributed among the set of relevant discrete nodes of a FE or Boundary Element (BE) model. This is usually achieved as follows,

$$\mathbf{F}^{(m)} = \int_{S_e} [\boldsymbol{\sigma}^\infty(\mathbf{x}) \cdot \mathbf{n}(\mathbf{x})] N_m(\mathbf{x}) \, dS_e, \quad (3.1)$$

where dS_e is the infinitesimal surface element with surface area S_e . $N_m(\mathbf{x})$ are so-called shape functions (interpolation functions) that distribute the traction field among the surface element's nodes.

The problematic singularity associated with the classical Volterra dislocation is avoided by using the non-singular formulation of Cai et al. [2]. The stress field of a dislocation in a homogenous infinite linear elastic domain can be calculated as a contour integral along the loop [3],

$$\sigma_{ij}^\infty(\mathbf{x}) = C_{ijkl} \oint \epsilon_{lnh} C_{pqmn} \frac{\partial G_{kp}(\mathbf{x} - \mathbf{x}')}{\partial x_q} b_m dx'_h, \quad (3.2)$$

where C_{ijkl} is the elastic stiffness matrix, ϵ_{lnh} the permutation operator, \mathbf{b} the Burgers vector, \mathbf{x}' the coordinate that spans the dislocation, and $G_{kp}(\mathbf{x} - \mathbf{x}')$ is Green's function of elasticity [3]. $G_{kp}(\mathbf{x} - \mathbf{x}')$ is defined as the displacement component in the x_k direction at point \mathbf{x} due to a force applied in the x_p direction

at point \mathbf{x}' . The traditional singularity comes from taking the Burgers vector distribution as a delta function. Cai et al. [2] proposed an alternative definition of $G_{kp}(\mathbf{x} - \mathbf{x}')$ which has a wider isotropic spread mainly localised in a radius a around the dislocation core,

$$G_{ij}(\mathbf{x} - \mathbf{x}') = \frac{1}{8\pi\mu} \left[\delta_{ij} \partial_{pp} - \frac{1}{2(1-\nu)} \partial_{ij} \right] R_a, \quad (3.3)$$

where μ, ν are the isotropic shear modulus and Poisson's ratio respectively, δ_{ij} is the Kronecker Delta, $\partial_{x_1 \dots x_n} \equiv \frac{\partial^n}{\partial x_1 \dots \partial x_n}$. R_a is defined as,

$$\begin{aligned} R_a(\mathbf{x}) &= R(\mathbf{x}) * w(\mathbf{x}) = \int R(\mathbf{x} - \mathbf{x}') w(\mathbf{x}') d^3 \mathbf{x}' \\ &= \sqrt{R^2 + a^2} \end{aligned} \quad (3.4a)$$

$$w(\mathbf{x}) = \frac{15a^4}{8\pi(R^2 + a^2)^{7/2}}, \quad (3.4b)$$

where $w(\mathbf{x})$ is the isotropic Burgers vector distribution derived in the appendix of [2], $\mathbf{x} = (x, y, z)$ and $R(\mathbf{x}) = \sqrt{x^2 + y^2 + z^2}$.

For two dislocation nodes (1, 2) connected by straight line segments eq. (3.2) becomes,

$$\begin{aligned} \boldsymbol{\sigma}^{(12)}(\mathbf{x}) &= -\frac{\mu}{8\pi} \int_{x_1}^{x_2} \left(\frac{2}{R_a^3} + \frac{3a^2}{R_a^5} \right) [(\mathbf{R} \times \mathbf{b}) \otimes d\mathbf{x}' + d\mathbf{x}' \otimes (\mathbf{R} \times \mathbf{b})] \\ &+ \frac{\mu}{4\pi(1-\nu)} \int_{x_1}^{x_2} \left(\frac{1}{R_a^3} + \frac{3a^2}{R_a^5} \right) [(\mathbf{R} \times \mathbf{b}) \cdot d\mathbf{x}'] \mathbf{I}_2 \\ &- \frac{\mu}{4\pi(1-\nu)} \int_{x_1}^{x_2} \frac{1}{R_a^3} [(\mathbf{b} \times d\mathbf{x}') \otimes \mathbf{R} + \mathbf{R} \otimes (\mathbf{b} \times d\mathbf{x}')] \\ &+ \frac{\mu}{4\pi(1-\nu)} \int_{x_1}^{x_2} \frac{3}{R_a^5} [(\mathbf{R} \times \mathbf{b}) \cdot d\mathbf{x}'] \mathbf{R} \otimes \mathbf{R} \end{aligned} \quad (3.5)$$

3.1.1 Resolving Singularities when Dislocation Line Segments are Parallel to Surface Elements

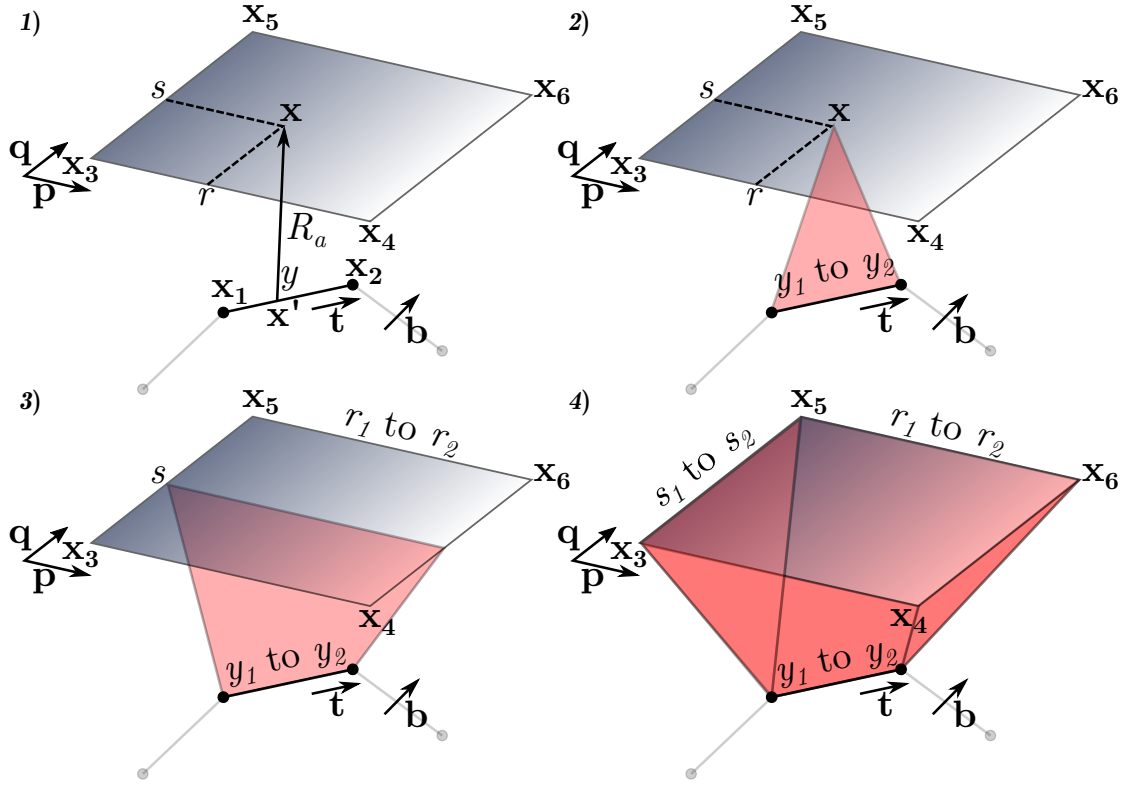


Figure 3.1: Diagram of the line integral method used to find analytical expressions for the forces exerted by dislocations on linear rectangular SEs [1]. 1) For any given point \mathbf{x} on the SE and any given point \mathbf{x}' on the dislocation line segment, define distance R_a . 2) Integrate from $x_1 \rightarrow x_2$ along line direction \mathbf{t} . 3) Integrate from $r_1 \rightarrow r_2$ along vector \mathbf{p} . 4) Integrate from $s_1 \rightarrow s_2$ along vector \mathbf{q} .

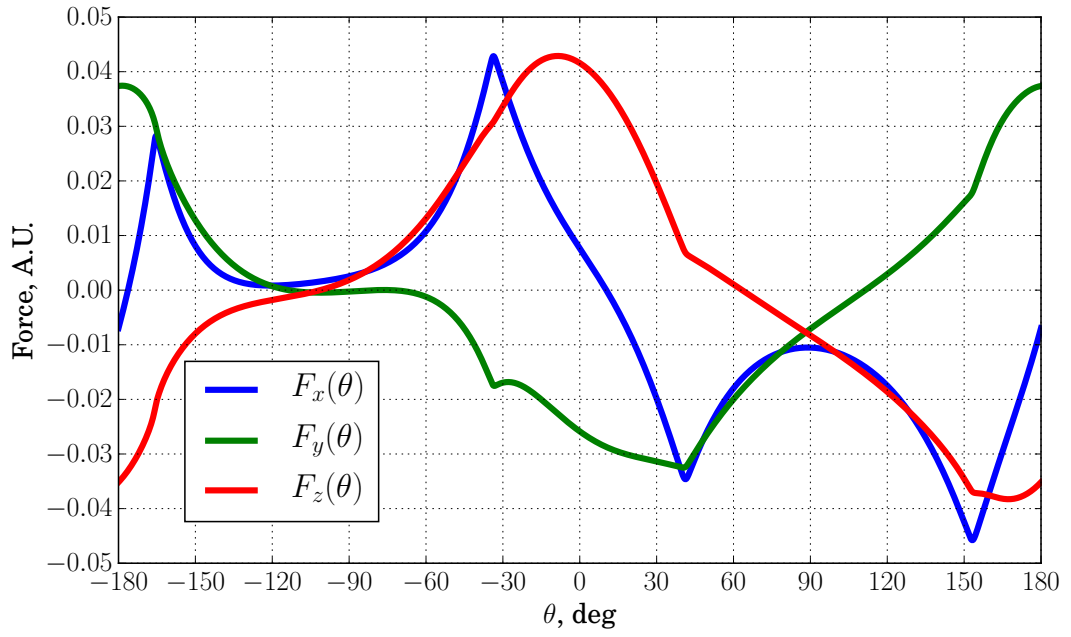


Figure 3.2: Effects of rotating a single dislocation line segment on the forces exerted by it on a linear rectangular SE. The specific values of this function are not known *a priori*, all that is known is that it must be periodic ($T = 2\pi$) and have finite maximum and minimum values. The singularity is avoided by perturbing the angle $\theta = 0 \rightarrow \theta = \pm\epsilon, \epsilon \gtrsim 0$.

Chapter 4

Parallelisation of the Analytical Forces Induced by Dislocations on Surface Elements

4.1 Parallelisation on Graphics Processing Units

Efficient parallelisation requires [coalesced memory access](#), which means we have to be extremely careful when mapping CPU memory to device memory. The fact that threads work “simultaneously”¹ means that in order to obtain good performance, data which is to be “simultaneously” loaded into each thread must be contiguous. This maximises cache memory use and therefore reduces slow memory fetch operations to global or shared memory.

The parallelisation was done only over the [SEs](#) in order to avoid the undesirable and inefficient GPU branching that would occur under other schemes.

The most natural form of parallelisation is to have blocks of $4n$ threads where $n \leq 8 \in \mathbb{N}$. This also fits nicely into the 32 thread per warp paradigm.

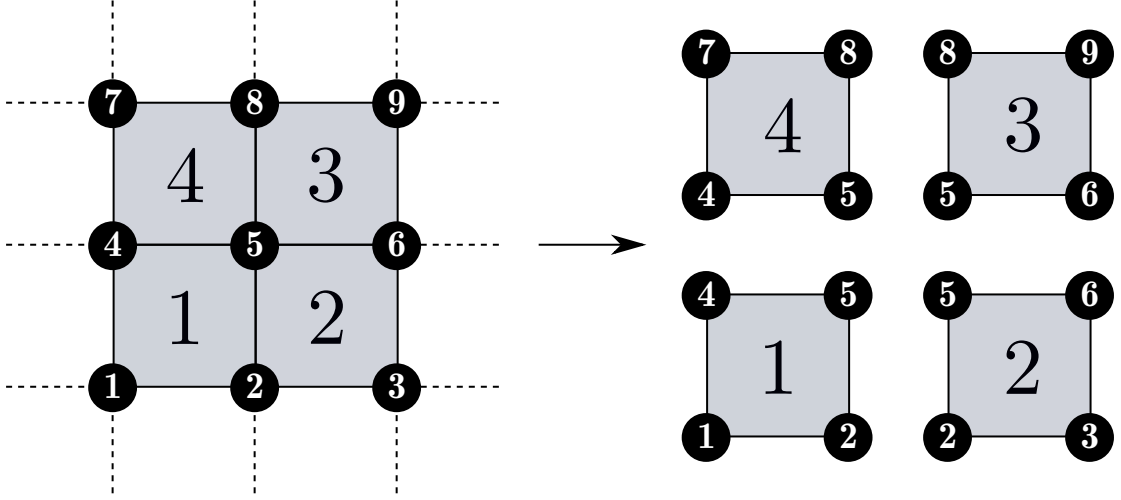


Figure 4.1: Each linear rectangular SE is mapped to one thread.

Algorithm 4.1 Elements in host \mapsto device.

```

1: function ELEMENT_HOST_DEVICE_MAP( $^h\mathbf{X}[N][3 \times E]$ )
2:      $\triangleright$  Input, temporary, and output indices set to zero.
3:      $i, j, k \leftarrow 0$ 
4:      $\triangleright$   $^d\mathbf{X}$  accomodates all 3 coordinates of all  $N$  nodes in all  $E$  elements.
5:      $^d\mathbf{X} \leftarrow \text{malloc}(3 \times E \times N)$ 
6:     for ( $n = 0; n < N; n++$ ) do  $\triangleright$  Loop over nodes in element.
7:          $\triangleright$  Set output index to point at the  $n^{\text{th}}$  node of the first coordinate of
            the first element.
8:          $k \leftarrow j$ 
9:         for ( $c = 0; c < 3; c++$ ) do  $\triangleright$  Loop over coordinates.
10:             $\triangleright$  Set input index to point at the  $c^{\text{th}}$  coordinate of the first element
                of the  $n^{\text{th}}$  node.
11:             $i \leftarrow c$ 
12:            for ( $e = 0; e < E; e++$ ) do  $\triangleright$  Loop over elements.
13:                 $^d\mathbf{X}[k + e] \leftarrow ^h\mathbf{X}[n][i]$ 
14:                 $\triangleright$  Advance output index to point at the  $n^{\text{th}}$  node of the  $c^{\text{th}}$ 
                    coordinate of the first element.
15:                 $i \leftarrow i + 3$ 
16:            end for
17:             $\triangleright$  Advance output index to point at the  $n^{\text{th}}$  node of the  $c^{\text{th}}$ 
                coordinate of the first element.
18:             $k \leftarrow k + E \times N$ 
19:        end for
20:         $\triangleright$  Advance temporary index to point at the  $(n + 1)^{\text{th}}$  node of the first
            coordinate of the first element.
21:         $j \leftarrow j + E$ 
22:    end for
23:    return  $^d\mathbf{X}$ 
24: end function

```

Algorithm 4.2 Elements in device \mapsto thread.

```
1: GPU function ELEMENT_DEVICE_THREAD_MAP( ${}^d\mathbf{X}[3 \times E \times N]$ ,  
    ${}^t\mathbf{X}[N][3]$ )  
2:                                      $\triangleright$  Thread, input indices.  
3:    $\text{idx}, i \leftarrow \text{threadIdx}.x + \text{blockIdx}.x \times \text{blockDim}.x$   
4:   for ( $c = 0; c < 3; c++$ ) do                                      $\triangleright$  Loop over elements.  
5:     for ( $n = 0; n < N; n++$ ) do                                      $\triangleright$  Loop over nodes in element.  
6:        ${}^t\mathbf{X}[n][c] \leftarrow {}^d\mathbf{X}[i + n \times E]$   
7:     end for  
8:      $i \leftarrow i + E \times N$   $\triangleright$  Advance input index to the  $(n + 1)^{\text{th}}$  node of element  
    $i$ .  
9:   end for  
10:  return  ${}^t\mathbf{X}[n][c]$   
11: end GPU function
```

Algorithm 4.3 Force in thread \mapsto device.

```
1: GPU function ADD_FORCE_THREAD_DEVICE( ${}^t\mathbf{F}_n[N][3]$ ,  ${}^t\mathbf{F}_e[3]$ ,  ${}^d\mathbf{F}_n[3 \times$   
    $E \times N]$ ,  
    ${}^d\mathbf{F}_e[3 \times E]$ ,  $\text{idx}$ )  
2:                                      $\triangleright$  Nodal force.  
3:    $\triangleright$  Set output index to whatever parallelisation index is given to the  
   function. In this case the same index we use for the main parallelisation.  
4:    $i \leftarrow \text{idx}$   
5:   for ( $n = 0; n < N; n++$ ) do                                      $\triangleright$  Loop over nodes.  
6:     for ( $c = 0; c < 3; c++$ ) do                                      $\triangleright$  Loop over coordinates.  
7:        $\triangleright$  Ensure nodal forces are correctly added and mapped from local  
       thread memory to global device memory.  
8:        $\text{atomicAdd}({}^d\mathbf{F}_n[i + c \times E \times N], {}^t\mathbf{F}_n[n][c])$   
9:     end for  
10:     $\triangleright$  Displace output index to point at the first coordinate of the  
     $(n + 1)^{\text{th}}$  node of the  $\text{idx}^{\text{th}}$  SE.  
11:     $i \leftarrow i + E$   
12:  end for  
13:                                      $\triangleright$  Total force.  
14:   $i \leftarrow \text{idx}$                                       $\triangleright$  Reset output index.  
15:  for ( $c = 0; c < 3; c++$ ) do                                      $\triangleright$  Loop over coordinates.  
16:     $\text{atomicAdd}({}^d\mathbf{F}_e[i], {}^t\mathbf{F}_e[c])$   
17:     $\triangleright$  Advance the output index to point at the  $(i + 1)^{\text{th}}$  coordinate of the  
     $\text{idx}^{\text{th}}$  SE.  
18:     $i \leftarrow i + E$   
19:  end for  
20:  return  ${}^d\mathbf{F}_n, {}^d\mathbf{F}_e$   
21: end GPU function
```

Algorithm 4.4 Force (parallelise over dislocations) in thread $\xrightarrow{+}$ device.

```

1: GPU function DLN_ADD_FORCE_THREAD_DEVICE( ${}^t\mathbf{F}_n[N][3]$ ,  ${}^t\mathbf{F}_e[3]$ ,
    ${}^d\mathbf{F}_n[3 \times E \times N]$ ,  ${}^d\mathbf{F}_e[3 \times E]$ ,  $k$ )
2:                                      $\triangleright$  Nodal force.
3:                                      $\triangleright$  Set output index to correspond to whatever SE we're on.
   By parallelising over dislocation lines, we must loop through SEs in the main
   code, this is the value of  $k$  we provide. Multiply by 3 and  $N$  because each SE
   has 3 coordinates and  $N$  nodes.
4:    $i \leftarrow 3 \times N \times k$ 
5:    $j \leftarrow 0$                                       $\triangleright$  Auxiliary index.
6:   for ( $n = 0$ ;  $n < N$ ;  $n++$ ) do                        $\triangleright$  Loop over nodes.
7:     for ( $c = 0$ ;  $c < 3$ ;  $c++$ ) do                      $\triangleright$  Loop over coordinates.
8:        $\triangleright$  Ensure nodal forces are correctly added and mapped from local
       thread memory to global device memory.
9:       atomicAdd( ${}^d\mathbf{F}_n[i + j + c]$ ,  ${}^t\mathbf{F}_n[n][c]$ )
10:    end for
11:     $\triangleright$  Displace auxiliary index to point at the first coordinate of the
     $(n + 1)^{\text{th}}$  node of the  $k^{\text{th}}$  SE.
12:     $j \leftarrow j + 3$ 
13:  end for
14:                                      $\triangleright$  Total force.
15:     $\triangleright$  Set auxiliary index to point at the first coordinate of the  $k^{\text{th}}$  SE.
16:     $j \leftarrow 3 \times k$ 
17:    for ( $c = 0$ ;  $c < 3$ ;  $c++$ ) do                        $\triangleright$  Loop over coordinates.
18:      atomicAdd( ${}^d\mathbf{F}_e[j + c]$ ,  ${}^t\mathbf{F}_e[c]$ )
19:    end for
20:    return  ${}^d\mathbf{F}_n$ ,  ${}^d\mathbf{F}_e$ 
21: end GPU function

```

Algorithm 4.5 Nodal force in device \mapsto host.

```

1: function FX_DEVICE_HOST_MAP(d $\mathbf{F}_n[3 \times E \times N]$ , h $\mathbf{F}_n[N][3 \times E]$ )
2:    $i, j \leftarrow 0$  ▷ Set input and output indices to zero.
3:   for  $n = 0; n < N; n++$  do ▷ Loop over nodes.
4:      $j \leftarrow 0$  ▷ Reset output index to point at the first element.
5:     for  $e = 0; e < E; e++$  do ▷ Loop over elements.
6:       for  $c = 0; c < 3; c++$  do ▷ Loop over coordinates.
7:         h $\mathbf{F}_n[n][j + c] =$ d $\mathbf{F}_n[i + e + c \times E \times N]$ 
8:       end for
9:       ▷ Advance output index to point at the first coordinate of the
       $(e + 1)^{\text{th}}$  element.
10:       $j \leftarrow j + 3$ 
11:    end for
12:    ▷ Advance input index to point at the first coordinate of the  $(n + 1)^{\text{th}}$ 
    node of the first element.
13:     $i \leftarrow i + E$ 
14:  end for
15:  return h $\mathbf{F}_n$ 
16: end function

```

4.1.1 Data Mapping

4.1.1.1 Elements: Host \mapsto Device

4.1.1.2 Elements: Device \mapsto Thread

4.1.1.3 Force: Thread $\overset{+}{\mapsto}$ Device

4.1.1.4 Force (Parallelise over Dislocations): Thread $\overset{+}{\mapsto}$ Device

4.1.1.5 Nodal Force: Device \mapsto Host

4.1.1.6 Total Force: Device \mapsto Host

$$\mathbf{X}_{en} := [x_{en}, y_{en}, z_{en}] \quad (4.1a)$$

$$\mathbf{X}_{(1 \rightarrow E)n} \mapsto \mathbf{X}_n$$

$$\mathbf{X}_n := [x_{1n}, y_{1n}, z_{1n}, \dots, x_{En}, y_{En}, z_{En}] \quad (4.1b)$$

$$\mathbf{X}_{1 \rightarrow N} \mapsto \mathbf{X}^{\text{SE}}$$

$$\begin{aligned} \mathbf{X}^{\text{SE}} := & [x_{11}, \dots, x_{E1}, x_{12}, \dots, x_{E2}, \dots, x_{1N}, \dots, x_{EN}, \\ & y_{11}, \dots, y_{E1}, y_{12}, \dots, y_{E2}, \dots, y_{1N}, \dots, y_{EN}, \\ & z_{11}, \dots, z_{E1}, z_{12}, \dots, z_{E2}, \dots, z_{1N}, \dots, z_{EN}] \end{aligned} \quad (4.1c)$$

where e is the SE label, n the node label, E the number of SEs in the scope. N

¹Not quite but essentially simultaneously.

Algorithm 4.6 Total force in device \mapsto host.

```

1: function FTOT_DEVICE_HOST_MAP(d $\mathbf{F}_e[3 \times E]$ , h $\mathbf{F}_e[3 \times E]$ )
2:    $i \leftarrow 0$  ▷ Set output index to zero.
3:   for  $e = 0; e < E; e++$  do ▷ Loop over elements.
4:     for  $c = 0; c < 3; c++$  do ▷ Loop over coordinates.
5:       h $\mathbf{F}_e[i + c] = \text{d\mathbf{F}_e}[e + c \times E]$ 
6:     end for
7:     ▷ Advance the output index to point at the first coordinate of the
        $(e + 1)^{\text{th}}$  surface element.
8:      $i \leftarrow i + 3$ 
9:   end for
10:  return h $\mathbf{F}_e$ 
11: end function

```

the number of nodes in a SE.

Data-mapping according to ?? and relabelling the nodes so they go from $0 \rightarrow N - 1$, the data from fig. 4.1 would be arranged in device memory like eq. (4.2),

$$\begin{aligned}
\mathbf{X}^{\text{SE}} = & \left[\underbrace{h_x, i_x, f_x, g_x}_{\mathbf{x}_0}, \underbrace{a_x, b_x, i_x, h_x}_{\mathbf{x}_1}, \underbrace{i_x, d_x, e_x, f_x}_{\mathbf{x}_2}, \underbrace{b_x, c_x, d_x, i_x}_{\mathbf{x}_3}, \right. \\
& \quad \dots \text{ } y\text{-coord } \dots, \\
& \quad \left. \dots \text{ } z\text{-coord } \dots \right]
\end{aligned} \tag{4.2}$$

In the GPU, each thread will cater to one SE at a time. This means that each thread will have to extract the relevant data from the 1D array with length $3 \times E \times N$ into four 1D arrays of length 3. The purpose of CNE mapping is to provide the warp with coalesced memory access. This is achieved via ??.

Since ?? is performed in a CUDA GPU, threads in a single block execute sequentially from,

$$\text{threadIdx}.x = 0 \rightarrow \text{threadIdx}.x = \text{blockDim}.x - 1, \tag{4.3}$$

while threads in different blocks execute in parallel. Coalesced memory access is ensured by having each block load a cache line whose entries are contiguously accessed by the threads in the block. Using the same notation as ??, full cache line utilisation (optimal cache use) is achieved if cache lines can accomodate l entries

$N = 4, E = 4$

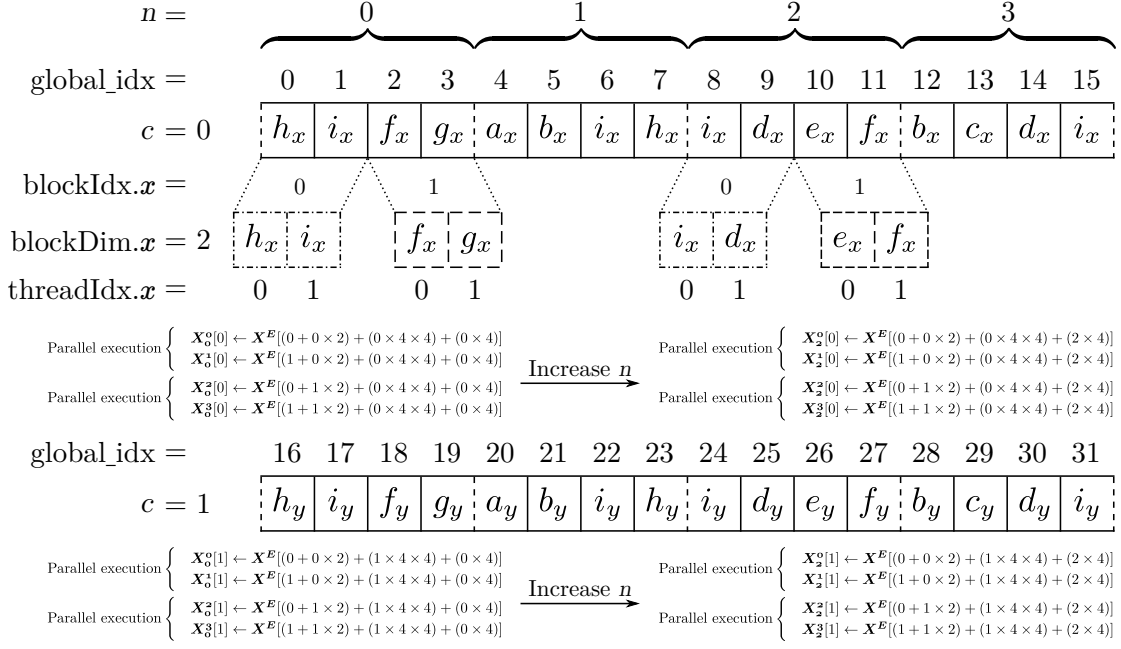


Figure 4.2: Minimum working complete example of the backward coordinate-node-element-map. The explicit calculations of global indices correspond to substituting the values for “idx” and “idxi” as in ?? . \mathbf{X}_a^b denotes a 1D array of length three containing the xyz -coordinates of node a of element b . Each thread concerns itself with only one element at a time. The dash-dot and dashed boxes represent cache lines for a thread block, the dotted line represents a memory fetch, and the dashed lines in \mathbf{X}^E represent steps of E entries (the start of the data for the next node of the element type we’re dealing with). The memory operations are not shown twice to minimise redundancy.

given by eq. (4.4),

$$l = \begin{cases} a \times N \times E & , a > 0 \in \mathbb{N} \\ \text{or} & \\ \frac{1}{2^a} \times N \times E & , a \geq 0 \in \mathbb{N}, N \times E \equiv 0 \pmod{2^a} . \end{cases} \quad (4.4)$$

Figure 4.2 shows an example of ?? up to the y -coordinate.

4.1.1.7 Node Coordinate Element Map

The node-coordinate-element (NCE) data mapping in eq. (4.5) is carried out by algorithm 4.7, each thread looks after a given SE.

Algorithm 4.7 NCE data mapping.

```

for all n nodes  $\in$  surface element do
  for all c coordinates  $\in [x, y, z]$  do
    for all e surface elements  $\in$  surface mesh section do
      list.append(data of the  $n^{\text{th}}$  node with  $c^{\text{th}}$  coordinate of the  $e^{\text{th}}$  SE)
    end for
  end for
end for

```

$$\mathbf{X}_{en} := [x_{en}, y_{en}, z_{en}] \quad (4.5a)$$

$$\mathbf{X}_{(1 \rightarrow E)n} \mapsto \mathbf{X}_n$$

$$\mathbf{X}_n := [x_{1n}, \dots, x_{En}, y_{1n}, \dots, y_{En}, z_{1n}, \dots, z_{En}] \quad (4.5b)$$

$$\mathbf{X}_{1 \rightarrow N} \mapsto \mathbf{X}^{\text{SE}}$$

$$\mathbf{X}^{\text{SE}} := \begin{matrix} [x_{11}, \dots, x_{E1}, y_{11}, \dots, y_{E1}, z_{11}, \dots, z_{E1} \\ \dots, \dots, \\ x_{1M}, \dots, x_{EN}, y_{1N}, \dots, y_{EN}, z_{1N}, \dots, z_{EN}] \end{matrix} \quad (4.5c)$$

where e is the surface element, n the node, E the total number of SEs in scope, N the total number of nodes in each SE.

Data-mapping according to algorithm 4.7 and relabelling the nodes so they go from $0 \rightarrow N - 1$, the data from fig. 4.1 would be arranged in device memory like eq. (4.6),

$$\mathbf{X}^{\text{SE}} = \begin{bmatrix} \underbrace{h_x, i_x, f_x, g_x, h_y, i_y, f_y, g_y, h_z, i_z, f_z, g_z}_{\mathbf{x}_0}, \\ \dots \mathbf{x}_1, xyz\text{-coords} \dots, \\ \dots \mathbf{x}_2, xyz\text{-coords} \dots, \\ \dots \mathbf{x}_3, xyz\text{-coords} \dots \end{bmatrix} \quad (4.6)$$

4.1.1.8 Resolving Data Write Conflicts

Data write conflicts can be a problem in parallel applications where global data is changed by multiple threads within the same clock tick. This can be avoided with atomic operations.

4.1.2 Resolving Parallel Dislocation Line Segments to Surface Elements

Dealing with the special case when the dislocation line segment \mathbf{t} is parallel to the SE is relatively trivial when in a serial program. By the mean value theorem, we

can slightly perturb \mathbf{t} by rotating it by a small angle around the midpoint of \mathbf{t} with respect to the axis of rotation defined by $\mathbf{t} \times \mathbf{n}$. In contrast to serial code, program branches can have a serious effect on parallel performance due to warp divergence. Due to the complexity of the force calculation and relative rarity of the edge case, there is no branching behaviour in the parallel code until *after* the calculation is performed. Where algorithm 4.8 is performed.

Algorithm 4.8 Resolving cases when $\mathbf{t} \parallel \mathbf{n}$ on GPUs.

```

for all surface elements and line segments do
    ...
    if  $\mathbf{t}_{\text{thread}} \parallel \mathbf{n}_{\text{thread}}$  then
         $\mathbf{x}_{\text{buffer}} \leftarrow \mathbf{x}_{\text{thread}}$ 
         $\mathbf{t}_{\text{buffer}} \leftarrow \mathbf{t}_{\text{thread}}$ 
    else
         $\mathbf{F}_{\text{total}} += \mathbf{F}_{\text{thread}}$ 
    end if
end for
return to serial code
if  $\mathbf{x}_{\text{buffer}} \neq \text{empty}$  then
    for all surface elements and line segments do
        perform calculation for  $\mathbf{t} \parallel \mathbf{n}$ 
    end for
end if

```

4.1.2.1 Implementation

The node mapping would be the same as ???. The parallelisation would be on individual packets of a SE with a slightly rotated dislocation line segment, as in a single iteration of the serial code where the dislocation line segment is rotated. The forces would be averaged with atomic operations and `__syncthreads()` before they are added to the total nodal forces.

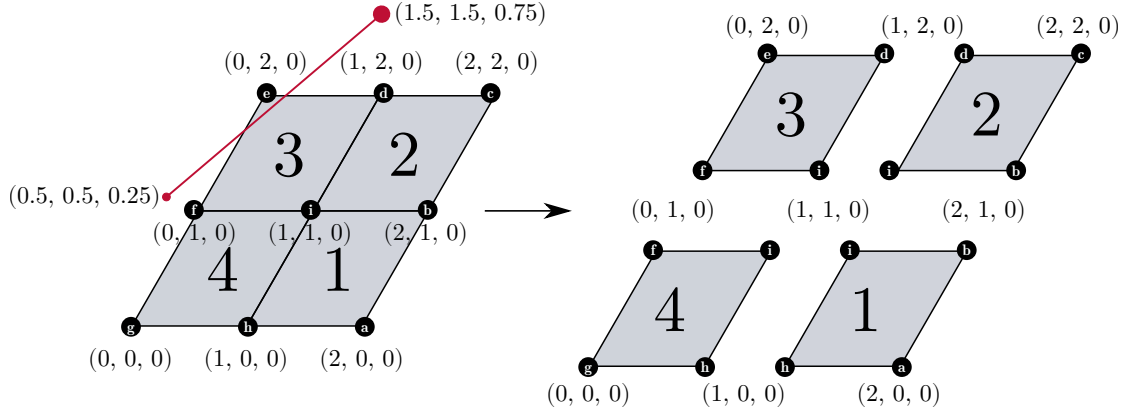


Figure 4.3: Test case for CUDA development. Dislocation line not shown in node separation as the parallelisation happens across surface elements. Not to perspective or scale.

4.1.3 Test Case

According to ?? and the node labelling scheme of fig. 3.1 the data mapping would look like:

$$\mathbf{X}^{\text{SE}} = \begin{bmatrix} \overbrace{1, 1, 0, 0, 2, 2, 1, 1, 1, 1, 0, 0, 2, 2, 1, 1}^{x\text{-coord}}, \\ \overbrace{0, 1, 1, 0, 0, 1, 1, 0, 1, 2, 2, 1, 1, 2, 2, 1}^{y\text{-coord}}, \\ \overbrace{0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0}^{z\text{-coord}} \end{bmatrix} \quad (4.7)$$

In general the number of indices in the array would be $3 \times e \times n$ where e is the number of surface elements, n the number of nodes per element and 3 the number of spatial dimensions.

4.2 Thread Block Size Optimisation

<http://docs.nvidia.com/cuda/cuda-c-best-practices-guide/index.html#memory-optimizations>
<http://docs.nvidia.com/cuda/cuda-c-best-practices-guide/index.html#execution-configuration-optimizations>

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Appendix A

Best Practices

This chapter defines the best practices that will make improve the development and testing pipeline by defining rules and standards that facilitate collaboration.

A.1 Filesystem

All files must be within `/project_name`.

1. All files must be within `/project_name`,
 `~/ = project_name/`,
 `~/~/ = project_name/src/` or `project_name/dev/`.
2. Release code must be within `~/src/`.
3. Development code must be within `~/dev/`.
4. Test data must be within `~/tests/`.
5. Documentation must be within `~/~/doc/`.
6. Examples must be within `~/~/exmp/`.
7. External libraries must be within `~/lib/`.
8. Generated images must be within `~/~/images/`.
9. Old versions recordkeeping must be within `~/prv/va.b.c/`.

A.2 Versioning

1. Use a version control system like [GitHub](#) or [PasteBin](#).

2. There must be a master branch that is only changed when the code is stable and bug free.
3. Development branches should be exploited as seen fit without making things overly convoluted.
4. Commits must be as bug free and regular as possible. When to commit is left to the developer's discretion.
5. Commit messages should be as descriptive as possible.
6. Versions should be specified as `va.b.c` where `a`, `b`, `c` = integers. The three levels are `a` = release version (usable, bug free code), `b` = beta version (code that is undergoing testing), `c` = alpha version (code that is under active development).

A.3 Documentation

A.3.1 Commenting

Every codefile must be appropriately commented by meeting the following guidelines.

1. The start of each codefile must have a heading detailing the creator, date of creation and edit history (date and name of editor).
2. Below the heading there must be a general explanation of the code. It must state any procedures, structures, objects and how they are to be utilised. Any backward or forward dependencies must be stated.
3. Below the description and edit history any relevant literature must be mentioned (dois are preferred). Must be as detailed as possible, include equation numbers/ranges if necessary.
4. The start of every procedure has an explanation of its purpose, inputs, outputs and inputs-outputs.
5. Particularly complicated code blocks must have an in-depth explanation of what it does. Comment each line if necessary.
6. Corrections or additions must be explicitly bounded by comments at the start and end of the change. Both bounding comments must have the author's name and the date. Below the starting comment, there should be an explanation of the change. Any punctual comments can be made as normal.

A.3.2 README

Every codefile must have an associated README `.tex` document that documents the codefile's contents. It must meet the following guidelines as appropriate.

1. The name must be that of the file it documents (minus the extension of course).
2. Description and overall explanation of the codefile's purpose.
3. Overall flow chart or pseudo code describing the file's purpose.
4. Document the codefile's procedures. This means describing and explaining their corresponding inputs, outputs, inputs-outputs, forwards and backwards dependencies, and flow charts or pseudo codes.
5. Unit test designs and results for each procedure. If appropriate also include those of integral tests.

The codefile may also be appended at the end of documentation if desired (the `minted` package is highly recommended).

A.4 Modularisation

1. Code repetition must be kept to a *strict* minimum. Any piece of code that will be reused must be modularised.
2. Procedures must be as self-sufficient as possible *without* repeating code. If repeating code is necessary, replace it with a procedure that is to be repeatedly called instead. Minimising repeated code \ggg procedure self-sufficiency.

A.5 Coding Style

All names must meet the following guidelines.

1. Indent appropriately. Four space tabs are a good compromise between code necking and readability.
2. Minimise the use of nested code blocks, use intrinsics, libraries or create procedures instead.
3. Break up lines that are uncomfortably long, typically anything over 80–100 characters.

4. Names should be appropriately descriptive and human readable.
5. All code and names must be systematic and logical.
6. the use of upper cases should be reserved for parameters (`const` variables in C).
7. Delimit words with “_” *not* case changes.
8. Long and descriptive \ggg short and cryptic.

A.5.1 Filenames

1. If old versions are to be kept, the old *stable* versions of file must have the date of last modification appended *suffixed* after the file extension in the `_yyyymmdd` format. For example, if the stable version of the release code `hello_world_parallel.c` was last modified on April 25, 2017 it should be archived as
`hello_world_parallel.c_20170425`.

A.5.2 Variables, Structures and Objects

1. Only counters and indices can be single letter variables.
2. Structure and object *definitions* are *suffixed* with `_s` and `_o` respectively.
3. Inputs, outputs and input-outputs to procedures must be *prefixed* with `i_`, `o_` and `io_` respectively.

A.5.3 Procedures

1. Functions and subroutines must be *prefixed* with `f_` and `s_` respectively.

Appendix B

Coupling Discrete Dislocation Dynamics to Finite Element Methods

```
1  function node_plane = extract_node_plane(fem_nodes ,
    ↪ fem_node_cnct,...
2
    node_labels, dim, filter
    ↪ ,...
3
    node_plane)
4
    ↪ %%=====%%
5
    ↪ %-----%
6  % Written by famed MATLAB hater and fan of compiled languages,
7  % Daniel Celis Garza, in 13/11/17.
8  %
9  % This code has been optimised as much as I can and has been
    ↪ thoroughly
10 % tested as of 15/11/17.
11
    ↪ %-----%
12 %
13 % Extracts a single node plane according to given filtering
    ↪ criteria
14 % and puts it into a 2D array arranged in xyz-coordinates like
    ↪ so.
15 %
```

```

16      %      Node 1      Node 2      Node 3 ... Node N
17      %
18      %      x_11      x_12      x_13      x_1N      /
19      %      y_11      y_12      y_13      y_1N      / Element 1
20      %      z_11      z_12      z_13      ... z_1N      /_____
21      %      x_21      x_22      x_23      x_2N      /
22      %      y_21      y_22      y_23      y_2N      / Element 2
23      %      z_21      z_22      z_23      z_2N      /_____
24      %
25      %
26      %
27      %
28      %      x_E1      x_E2      x_E3      x_EN      /
29      %      y_E1      y_E2      y_E3      ... y_EN      / Element E
30      %      z_E1      z_E2      z_E3      z_EN      /
31      %
32
33      ↪ %=====
34      % Inputs
35
36      ↪ %=====
37      %
38      % fem_nodes := dimension(:, 3). Assumed shape 2D array with 3
39      ↪ columns.
40      % Describes the coordinates of the nodes of the ft_inite
41      ↪ element model.
42      %
43      % fem_node_cnct := dimension(:, 8). Assumed shape 2D array
44      ↪ with 8
45      % columns. Describes the node connectivity of the finite
46      ↪ elements in
47      % the model.
48      %
49      % node_labels := dimension(:). The labels of the nodes that
50      ↪ are to be
51      % extracted from fem_nodes.
52      %
53      % dim := three times the area of the plane to be extracted (3
54      ↪ dimensions).

```



```

48      %
49      % filter := the selection criteria used to filter which nodes
      ↪ we want
50      % to obtain.
51      %
52      ↪ %=====
53      % Dummy variables
54
55      ↪ %=====
56      %
57      % n_nodes := number of nodes per planar element.
58      %
59      % tmp_nodes := dimension(n_nodes, 3). Contains the
      ↪ xyz-coordinates of
60      % all nodes with the labels specified by node_label.
61      %
62      % mask := same dimension as tmp_nodes. Logical mask whose
      ↪ entries are 1
63      % only when a given condition is met. It is used to index
      ↪ tmp_nodes to
64      % find only the nodes that correspond to a given plane.
65      %
66      ↪ %=====
67      % Input/output variables
68
69      ↪ %=====
70      %
71      % node_plane := dimension(dim/3, n_nodes). It is the slice of
      ↪ the array
72      % that holds the nodes of the plane. Only pass the array slice
73      % that corresponds to the current plane of nodes to be
      ↪ extracted,
74      % a la Fortran intent(in out).
75      %
      ↪ %%=====

```

```

76     %% Error checking
77     n_nodes      = size(node_labels, 2);
78     if(n_nodes ~= size(node_plane , 2))
79         error('extract_node_plane:: n_nodes = %d, size(node_plane,
            ↪ 2) = %d; they must be the same i.e. the number of nodes
            ↪ per planar element', n_nodes, size(node_plane,2))
80     end %if
81
82     %% Finding plane nodes
83     % Find the all the nodes labelled as stated in node_labels
84     % (corresponding to a plane family)
85     tmp_nodes = fem_nodes(fem_node_cnct(:, node_labels), 1:3);
86     % Find an array of entries that meet the filter criteria in
            ↪ tmp_nodes
87     mask = tmp_nodes(:, filter(1)) == filter(2);
88     % Reshape the array to a column array containing only the
            ↪ nodes which
89     % meet the filter criteria.
90     tmp_nodes = reshape(tmp_nodes(mask, :)', n_nodes * dim, 1);
91
92     %% Passing data to the slice of the array that contains the
            ↪ nodes for every plane
93     step = 0;
94     for i = 1: n_nodes
95         node_plane(1: dim, i) = tmp_nodes(1 + step: dim + step);
96         step = step + dim_3;
97     end %for
98
99     %% Cleanup
100    clear dim_3; clear idxf; clear tmp_nodes; clear mask; clear
            ↪ step;
101    end %function

1    function [node_plane_lbl, node_plane] = extract_node_planes(
            ↪ ...
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4                                     n_elem          , n_nodes)
5
6
7                                     %%=====%%
8
9                                     %-----%
10
11 % Written by famed MATLAB hater and fan of compiled languages,
12 % Daniel Celis Garza, in 13/11/17.
13 %
14 % This code has been optimised as much as I can and has been
15 %   thoroughly
16 % tested as of 15/11/17.
17
18                                     %-----%
19
20 %
21 % _lbl := variables suffixed with it are the global label
22 %   equivalent
23 %   of those without.
24 %
25 % Extracts a single node plane according to given filtering
26 %   criteria
27 % and puts it into a 2D array arranged in xyz-coordinates like
28 %   so.
29
30 %
31 %      Node 1      Node 2      Node 3 ... Node N
32 %
33 %      x_11      x_12      x_13      x_1N      /
34 %      y_11      y_12      y_13      y_1N      / Element 1
35 %      z_11      z_12      z_13      ... z_1N      /_____
36 %      x_21      x_22      x_23      x_2N      /
37 %      y_21      y_22      y_23      y_2N      / Element 2
38 %      z_21      z_22      z_23      z_2N      /_____
39 %
40 %      .
41 %      .
42 %      .
43 %
44 %      -----
45 %      x_E1      x_E2      x_E3      x_EN      /
46 %      y_E1      y_E2      y_E3      ... y_EN      / Element E
47 %      z_E1      z_E2      z_E3      z_EN      /
48 %

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```
→ %=====
% Inputs
→ %=====
%
% fem_nodes := dimension(:, 3). Assumed shape 2D array with 3
→ columns.
% Describes the coordinates of the nodes of the ft_inite
→ element model.
%
% fem_node_cnct := dimension(:, 8). Assumed shape 2D array
→ with 8
% columns. Describes the node connectivity of the finite
→ elements in
% the model.
%
% surf_node_util := dimension(n_nodes+2, 6). Each column is
→ laid out as
% follows:
% [node_label_initial; ... ; node_label_final;
% plane_area; coordinate_number].
% Each row corresponds to one plane.
%
% fem_planes := dimension(:). Contains the planes whose nodes
→ are to be
% extracted.
%
% n_elem := total number of elements to be extracted.
%
% n_nodes := number of nodes per planar element.
%
→ %=====
% Dummy variables
→ %=====
%
```

```

64 % plane_idx := the index of the fem plane to be extracted
    ↪ during the
65 % current iteration.
66 %
67 % dim := plane_area, there are this many surface elements per
    ↪ plane.
68 %
69 % dim_3 := 3*plane_area, there are three coordinates per plane
70 % node, plane_area is the total number of nodes in a plane
    ↪ and is found
71 % in surf_node_utils.
72 %
73 % tmp_nodes := dimension(n_nodes, 3). Contains the
    ↪ xyz-coordinates of
74 % all nodes with the labels specified by node_label.
75 %
76 % mask := same dimension as tmp_nodes. Logical mask whose
    ↪ entries are 1
77 % only when a given condition is met. It is used to index
    ↪ tmp_nodes to
78 % find only the nodes that correspond to a given plane.
79 %
80 % coord := orthogonal coordinate to the plane (for node
    ↪ filtering
81 % purposes).
82 %
83 % idxi, idxf := initial and final index that slice the
    ↪ coordinate
84 % output array into sections for each extracted plane.
85 %
86 % idxl, idxm := initial and final index that slice the label
87 % output array into sections for each extracted plane.
88 %
89 % filter := value for filtering data.
90 %
91 % step := step that ensures proper array indexing for each
    ↪ section of
92 % the linear array that contains the nodes of a given plane.
93 %

```

```

94     % n_nodes_p1 := n_nodes + 1
95     %
96     % n_nodes_p2 := n_nodes + 2
97     %
98
99     ↪ %=====
100    % Output variables
101
102    ↪ %=====
103    %
104    % node_plane := dimension(3*n_elem, n_nodes). It is the array
105    ↪ that
106    % holds the planes' nodes' Cartesian coordinates.
107    %
108    ↪ %%=====
109
110    %% Allocating node plane coordinates and labels
111    node_plane      = zeros(3 * n_elem, n_nodes);
112    node_plane_lbl = zeros(    n_elem, n_nodes);
113
114    %% Looping through the planes to be extracted
115    idxl = 1;
116    idxi = 1;
117    n_nodes_p_1 = n_nodes + 1;
118    n_nodes_p_2 = n_nodes + 2;
119    for i = 1: size(fem_planes, 1)
120        %% Assigning control variables for the iteration
121        plane_idx = fem_planes(i);
122        dim      = surf_node_util(n_nodes_p_1, plane_idx);
123        dim_3    = 3 * dim;
124        coord    = surf_node_util(n_nodes_p_2, plane_idx);
125        % Set final index of the output array slice containing the
126        ↪ nodes
127        % for this iteration's plane
128        idxm = idxl + dim - 1;
129        idxf = idxi + dim_3 - 1;
130        %% Finding global node labels and node coordinates of the
131        ↪ plane

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127     % Extracting all nodes whose labels correspond to the
    ↪ plane family
128     % to be extracted this iteration
129     tmp_nodes_lbl = fem_node_cnct(:, surf_node_util(1:n_nodes,
    ↪ plane_idx));
130     tmp_nodes      = fem_nodes(tmp_nodes_lbl, 1:3);
131     % If the plane index is odd it is a face where its
    ↪ orthogonal
132     % coordinate is at a minimum.
133     if(mod(plane_idx, 2) ~= 0)
134         filter = min(tmp_nodes(:, coord));
135     % Otherwise find the face whose orthogonal coordinate is
    ↪ at a maximum
136     else
137         filter = max(tmp_nodes(:, coord));
138     end %if
139     % Find an array of entries that meet the filter criteria
    ↪ in tmp_nodes
140     mask = tmp_nodes(:, coord) == filter;
141     % Reshape the array to a column array containing only the
    ↪ nodes which
142     % meet the filter criteria.
143     tmp_nodes = reshape(tmp_nodes(mask,:)', n_nodes * dim_3,
    ↪ 1);
144     % Reshape the array with global node labels so it can be
    ↪ filtered
145     % using the same mask.
146     tmp_nodes_lbl_size = size(tmp_nodes_lbl, 1);
147     tmp_nodes_lbl      = reshape(tmp_nodes_lbl ,...
148                                 tmp_nodes_lbl_size * n_nodes,
    ↪ 1);
149     tmp_nodes_lbl = tmp_nodes_lbl(mask);
150
151     %% Passing data to the slice of the array that contains
    ↪ the nodes labels and coords for every plane
152     % Pass the node labels of this iteration to the global
    ↪ array.
153     % Set the slice of the tmp_nodes array to be indexed
154     step      = 0;

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```

155     step_lbl = 0;
156     for j = 1: n_nodes
157         % Grab the array slice that corresponds to node j in
           ↪ the
158         % tmp_nodes array and send it off to the corresponding
           ↪ column
159         % of the slice of the output array containing all
           ↪ extracted nodes
160         node_plane(idxi: idxf, j) = tmp_nodes(...
161                                     1 + step: dim_3 +
           ↪ step...
162                                     );
163         node_plane_lbl(idxl: idxm, j) = tmp_nodes_lbl(...
164                                     1 + step_lbl: dim +
           ↪ step_lbl...
165                                     );
166         % Advance the step to move onto the array slice that
           ↪ contains
167         % node j+1 in tmp_nodes
168         step = step + dim_3;
169         step_lbl = step_lbl + dim ;
170     end %for
171     % Set the initial index to shift the output array slice
           ↪ onto the
172     % next plane to be extracted.
173     idxi = idxf + 1;
174     idxl = idxm + 1;
175     clear tmp_nodes; clear tmp_nodes_lbl; clear mask;
176 end %for
177
178 %% Cleanup
179 clear idxl; clear idxi; clear plane_idx; clear dim; clear
           ↪ dim_3;
180 clear coord; clear idxm; clear idxf; clear filter; clear
           ↪ tmp_nodes_lbl_size;
181 clear step; clear step_lbl; clear n_nodes_p1; clear n_nodes_p2;
182 end %function

```



```

1  function [f_dln, f_dln_se] = analytic_traction(
    ↪ ...
2
    se_node_coord, dln_node_coord, burgers      ,
    ↪ n_nodes,...
3
    n_nodes_t      , n_se          , n_dln      ,
    ↪ idxf          ,...
4
    idxi           , f_dln_node     , f_dln_se    ,
    ↪ f_dln         ,...
5
    mu, nu, a      , use_gpu        , n_threads   ,
    ↪ para_scheme)
6
    ↪ %%=====%%
7
    ↪ %-----%
8  % Written by famed MATLAB hater and fan of compiled languages,
9  % Daniel Celis Garza, in 12/11/17--16/11/17.
10 % Refactored into independent subroutines in
    ↪ 11/12/2017--15/12/2017.
11
    ↪ %-----%
12 % Calculates analytical forces to the surfaces of a
    ↪ rectangular
13 % cantilever whose nodes are labelled as follows:
14 %
15 % 4. ----- .3
16 % /\          /\
17 % / \         / \
18 % 1. -\----- .2 \
19 %   \ \       \ \
20 %    \ 8. ----\ - .7
21 %     \ /       \ /
22 %      \/\       \/\
23 %       5. ----- .6
24 % ^z
25 % | y
26 % | /
27 % | /
28 % |----->x
29 %

```

```

30  % Everything has been arranged to conform with the analytical
    ↪ solutions
31  % for forces exerted on linear rectangular surface elements
    ↪ published by:
32  % S. Queyreau, J. Marian, B.D. Wirth, A. Arsenlis, MSMSE,
    ↪ 22(3):035004, (2014).
33  %
34
    ↪ %=====
35  % Inputs
36
    ↪ %=====
37  %
38  % se_node_coord := dimension(n_se*n_nodes, 3). 2D array
39  % with the cartesian coordinates of the finite element nodes
    ↪ subject
40  % to traction boundary conditions. The analytical solution
    ↪ requires
41  % they be arranged element by element in a specific order.
42  %
43  % dln_node_coord := dimension(n_dln*2, 3). 2D array with
44  % the cartesian coordinates of the dislocation line nodes.
    ↪ The
45  % analytical solution requires they be arranged line segment
    ↪ by
46  % line segment.
47  %
48  % burgers := dimension(n_dln, 3). 2D array with the
49  % individual dislocation line segments' Burgers vector. The
50  % analytical solution needs the Burgers vector for each line
    ↪ segment.
51  %
52  % n_nodes := number of nodes per surface element.
53  %
54  % n_nodes_t := total number of nodes with traction boundary
    ↪ conditions.
55  % It is needed to map the analytical solutions' forces to
    ↪ the force

```

```

56      %   array used to account for the forces exerted by the
      ↪   dislocation
57      %   ensemble on the finite element nodes. This does not
      ↪   necessarily
58      %   contain a factor of n_nodes less entries than
      ↪   se_node_coord because
59      %   not all nodes in an element may have traction boundary
      ↪   conditions,
60      %   but still need to be included in the calculation due to
      ↪   the
61      %   problem's formulation.
62      %
63      % n_se := number of surface elements
64      %
65      % n_dln := number of dislocation line segments
66      %
67      % idxf := dimension(:). 1D array containing the indices of the
      ↪   force
68      %   array used to account for the forces exerted by the
      ↪   dislocation
69      %   ensemble on the finite element nodes. It is needed to map
      ↪   the
70      %   analytical solutions' forces to the right index. It is
      ↪   usually
71      %   equal to gamma*3 because only the nodes with traction
      ↪   boundary
72      %   conditions are needed.
73      %
74      % idxi := dimension(n_nodes_t*n_nodes). 1D assumed shape array
      ↪
75      %   containing the global label indices of the nodes whose
      ↪   tractions
76      %   were calculated by the analytical solution.
77      %
78      % mu := shear modulus of material.
79      %
80      % nu := poisson's ratio.
81      %

```

```

82      % a := dislocation core size parameter (non-singular
      ↪      dislocations).
83      %
84
85      ↪      %-----
86      %
87      % Flags:
88      % use_gpu := Flag to use the Graphics Processing Unit (GPU) in
      ↪      the
89      % calculation. If use_gpu == 1, the code is run by a
      ↪      CUDA-enabled
90      % NVidia GPU. The code will crash if there is none, or if
      ↪      there is no
91      % compiled MEX CUDA file. If use_gpu == 0, a serial C code
      ↪      will be
92      % executed. The code will crash if there is no compiled MEX
      ↪      C file.
93      % If use_gpu is any other number, a MATLAB version will be
      ↪      used. This
94      % is much slower than either of the others.
95      %
96
97      ↪      %-----
98      %
99      % Optional parameters:
100     %
101     % n_threads := when use_gpu ==1, it is the number of threads
      ↪      per block
102     % to be used in the parallelisation. This may require some
103     % experimentation to optimise. Defaults to 512. Does nothing
      ↪      if a
104     % GPU isn't used.
105     %
106     % para_scheme := when use_gpu is enabled, it dicates the
107     % parallelisation scheme used. If para_scheme == 1
      ↪      parallelise over
108     % dislocation line segments. If para_scheme == 0 parallelise
      ↪      over

```

```

108 % surface elements. Other parallelisation schemes can be
    ↪ added in the
109 % CUDA C code.
110 %
111
    ↪ %=====
112 % Dummy variables
113
    ↪ %=====
114 %
115 % f_dln_node := dimension(3*n_se, 4). 3D force on each
116 % surface element node.
117 %
118 % f_dln_se := dimension(3*n_se, 1). Total 3D force on each
    ↪ surface
119 % element (sum over all nodes of an element). This isn't
    ↪ used to
120 % couple the forces to the finite element forces but can be
    ↪ used to
121 % provide the evolution of the forces exerted by the
    ↪ dislocations on
122 % the surface elements.
123 %
124 % tmp := dimension(n_nodes). 1D array containing the at-most 4
125 % instances where a surface element node appears (a single
    ↪ node can
126 % be shared by up to 4 surface elements).
127 %
128 % tmp2 := dimension(:). Assumed shape 1D array containing the
    ↪ at-most 4
129 % instances where a surface element node appears (a single
    ↪ node can
130 % be shared by up to 4 surface elements) but without any
    ↪ zero-valued
131 % entries because 0 indices crash MATLAB.
132 %
133 % k := used to traverse idxi to the next block of 4 entries
    ↪ which
134 % contain data relevant to the node in the

```

```

135 %
136
137 ↪ %=====
138 % Inputs/Outputs
139
140 ↪ %=====
141 %
142 % f_dln := dimension(n_nodes_t*3). 1D array containing the
143 ↪ forces
144 % exerted by the dislocations on the surface nodes. It is
145 ↪ used to
146 % correct the finite element forces.
147 %
148 ↪ %%=====
149
150 %% Analytical force calculation.
151 % Parallel CUDA C calculation.
152 if use_gpu == 1
153     % Provide a default number of threads in case none is
154     ↪ given.
155     if ~exist('n_threads', 'var')
156         n_threads = 512;
157     end %if
158     % Provide a default parallelisaion scheme in case none is
159     ↪ given.
160     if ~exist('para_scheme', 'var')
161         % Parallelise over dislocations.
162         para_scheme = 1;
163     end %if
164     [f_dln_node(:, 1), f_dln_node(:, 2),...
165     f_dln_node(:, 3), f_dln_node(:, 4),...
166     f_dln_se] = nodal_surface_force_linear_rectangle_mex_cuda(
167     ↪ ...
168
169     dln_node_coord(:, 1),
170     ↪ dln_node_coord(:, 2) ,...
171
172     se_node_coord(:, 1), se_node_coord
173     ↪ (:, 2) ,...

```

```

163         se_node_coord (:, 3), se_node_coord
           ↪ (:, 4) ,...
164         burgers(:), mu, nu, a, n_se, n_dln,
           ↪ n_threads,...
165         para_scheme);
166 % Serial force calculation in C.
167 elseif use_gpu == 0
168     [f_dln_node(:, 1), f_dln_node(:, 2),...
169     f_dln_node(:, 3), f_dln_node(:, 4),...
170     f_dln_se] = nodal_surface_force_linear_rectangle_mex(
           ↪ ...
171         dln_node_coord(:, 1),
           ↪ dln_node_coord(:, 2),...
172         se_node_coord (:, 1), se_node_coord
           ↪ (:, 2),...
173         se_node_coord (:, 3), se_node_coord
           ↪ (:, 4),...
174         burgers(:), mu, nu, a, n_se,
           ↪ n_dln);
175 % Matlab version.
176 else
177     [f_dln_node(:, 1), f_dln_node(:, 2),...
178     f_dln_node(:, 3), f_dln_node(:, 4),...
179     f_dln_se] = NodalSurfForceLinearRectangle2(
           ↪ ...
180         dln_node_coord(:, 1),
           ↪ dln_node_coord(:, 2),...
181         se_node_coord (:, 1), se_node_coord
           ↪ (:, 2),...
182         se_node_coord (:, 3), se_node_coord
           ↪ (:, 4),...
183         burgers(:), mu, nu, a, n_se,
           ↪ n_dln);
184 end %if
185
186 %% Map analytical nodal forces into a useful form for the
           ↪ force superposition scheme.
187 % Loop through the number of nodes.
188 k = 0;

```

```

189     tmp = zeros(n_nodes, 1);
190     for i = 1: n_nodes_t
191         % Populate tmp array with the indices corresponding to
192         ↪ nodes of
193         % the surface relevant surface element.
194         tmp = idxi(1 + k: k + n_nodes);
195         % Obtain only the indices which are non-zero. Zero indices
196         ↪ mean
197         % those nodes are not under traction boundary conditions.
198         tmp2 = tmp(idxi(1+k: k + n_nodes, 1) ~= 0);
199         for j = 2:-1:0
200             % The index is displaced by -2, -1, 0, corresponding
201             ↪ to the
202             % indices of the x, y, z coordinate respectively.
203             % We add the force contributions from all surface
204             ↪ elements a node
205             % is part of. This gives us the total x,y,z forces on
206             ↪ each node.
207             f_dln(idxf(i) - j) = sum(f_dln_node(tmp2 - j));
208         end %for
209         % Step to the block in idxi where the next surface node's
210         ↪ indices
211         % are found.
212         k = k + 4;
213     end %for
214     clear tmp; clear tmp2; clear k;
215 end % function

```


Appendix C

Implementation of Analytical Forces Induced by Dislocations on Linear Rectangular Surface Elements

C.1 Serial C Code MEX File

C.2 Parallel CUDA C Code MEX File

Appendix D

Talks

D.1 Durham July 12–14 2017

Bridging the gap between the microscopic and macroscopic world is an on-going challenge for science and technology. If we hope to understand complex emergent phenomena we need to study systems that blur the line between micro and macro. In materials science, one such area is the study of extended defects called dislocations; whose nucleation and movement mediate the permanent deformation of materials. These large, high energy defects present very complex and long ranged interactions with each other, crystal boundaries, impurities, free surfaces, and themselves. As such, their dynamics are difficult to study experimentally. So we make justified assumptions and simplifications and create models that let us study them in detail. However if our models are to prove useful in real applications, they must be continually refined and improved by weakening assumptions and removing simplifications. Unfortunately, with increased refinement comes increased computational cost and new challenges. Therefore, finding faster alternatives that do not sacrifice accuracy is of the utmost importance—better yet if the alternatives are more accurate or exact. With the advent of increasingly accessible graphics processor units (GPUs) typically used in video gaming, the power of parallel processing is no longer exclusive to researchers with access to supercomputers. In this talk we will discuss the GPU implementation of exact solutions for the forces dislocations exert on the surfaces of materials.

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