Project: "Recognition of digital circle"

Introduction

The objective of this project is to implement a linear programming based algorithm for the recognition of digital circles.

We expect from you:

- A short report with answers to the "formal" questions and a description of your implementation choices and results.
- A C++ project (CMakeLists.txt plus several commented cpp program files).

1 Digital disk

Let $I \subset \mathbb{Z}^2$ be a rectangular digital image. Let $Z \subset I$ be a digital set and $\bar{Z} = I \setminus Z$ be its complement set.

A digital set Z is a digital disk if and only if there exists a circle of center $\omega \in \mathbb{R}^2$ and of radius $\rho \in \mathbb{R}$ such that:

$$\begin{cases}
\forall z \in Z, \ ||z - \omega||^2 \le \rho^2 \\
\forall z \in \bar{Z}, \ ||z - \omega||^2 \ge \rho^2
\end{cases}$$
(1)

Question 1 Let us consider the Gauss digitization of a convex shape $X \subset \mathbb{R}^2$ at gridstep h:

$$Dig(X, h) = X \cap (h \cdot \mathbb{Z}^2).$$

Show that if X is an Euclidean disk, then Dig(X,h) is a digital disk. Show however that the converse is not true.

Question 2 There exists a unique circle passing by three digital points. Show that we can test whether another digital point lies INSIDE, ON or OUTSIDE such a circle with integer-only computations and without explicitly computing its center and radius. You may have a look at "incircle test", broadly used in computational geometry, e.g. to compute the Delaunay triangulation of a point set.

Question 3 Implement a function that checks whether two given digital sets are separated by a given circle passing by three digital points or not.

2 Linear programming

Let us now consider algorithm 1 (which uses routine 2). It is a randomized and recursive algorithm that checks whether two point sets are separable by a circle in expected linear-time (see [?] or [?][chapter4]).

The union of the inner point set, denoted by S^- , and the outer point set, denoted by S^+ , is merely denoted by S. All the points of S are numbered from 1 to |S|, the size of S. The idea consists in maintaining a separating circle while iterating over the points $s_i \in S$ for i from 1 to |S|. For each point s_i , three cases may occur:

• if it belongs to S^- (resp. S^+) and it is located INSIDE (resp. OUTSIDE) or ON the current separating circle, there is nothing to do.

- Otherwise (lines 6-9 of algorithm 2):
 - 1. Either the two input sets are not circularly separable at all,
 - 2. or there exists a separating circle passing by s_i .

In the aim of deciding between these last two alternatives, the set of possible separating circles is restricted to circles passing by s_i and the same algorithm is recursively called from 1 to i (line 9 of algorithm 2). At each recursive call, the set of possible separating circles is restricted so that the base case involves a unique circle passing by three given points and consists in checking whether it separates S^- from S^+ or not (lines 11-17 of algorithm 2).

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Algorithm 1: areCircularlySeparable(S^-, S^+, p_1, p_2, p_3)

Input: S^-, S^+ \subset \mathbb{Z}^2, the inner and outer point sets

p_1, p_2, p_3 \in \mathbb{Z}^2, three points characterizing a circle

Result: "true" if S^- and S^+ are circularly separable, "false" otherwise

Output: p_1, p_2, p_3, three points caracterizing a separating circle if "true"

// initialisation step

1 Construct the set of S = S^- \cup S^+;

2 Randomly permute the points of S;

// points of S are numbered from 1 to |S|, |S| is the size of the set

3 Initialize p_1, p_2, p_3 with three points of S;

// recursive step

4 return areCircularlySeparable(S^-, S^+, S, |S|, p_1, p_2, p_3, 3);
```

Question 4 Implement algorithms 1 and 2. Provide test files.

3 Experiments

Question 5 In order to check whether two connected digital sets Z and \bar{Z} are circularly separable, it is enough to consider only the digital boundaries of Z and \bar{Z} . Implement a function that takes as input the common contour of Z and \bar{Z} and that checks whether Z and \bar{Z} are circularly separable or not. You may use DGTAL and more precisely the class GridCurve that can return the set of boundary digital points (see e.g. IncidentPointsRange).

Question 6 Perform a running time analysis of your recognition function.

- Implement a function that constructs the contour of a disk of a given radius.
- Output the running time of your recognition function for disks of increasing radius.
- Plot the graph of the running times against the size of the contour.
- Do you observe the expected linear-time complexity?

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\overline{\textbf{Algorithm 2}}: are Circularly Separable (S^-, S^+, S, n, p_1, p_2, p_3, k)
   Input: S^-, S^+ \subset \mathbb{Z}^2, the inner and outer point sets, S = S^- \cup S^+
   n, number of points of S to process (1 \le n \le |S|)
   p_1, p_2, p_3 \in \mathbb{Z}^2, three points characterizing a circle
   k, number of variable points among \{p_1, p_2, p_3\} (0 \le k \le 3)
   Result: "true" if S^- and S^+ are circularly separable, "false" otherwise
   Output: p_1, p_2, p_3, three points caracterizing a separating circle if "true"
 1 are Separable \leftarrow TRUE;
 2 if k > 0 then
       for l from 1 to k, initialize p_l with a point of S;
 3
 4
        while are Separable and i < n do
 5
            if (s_i \in S^- \text{ and } s_i \text{ is strictly OUTSIDE the circle passing by } p_1, p_2, p_3)
 6
            or (s_i \in S^+ \text{ and } s_i \text{ is strictly INSIDE the circle passing by } p_1, p_2, p_3) then
 7
 8
               areSeparable \leftarrow areCircularlySeparable(S^-, S^+, i, p_1, p_2, p_3, k-1);
9
11 else
       i \leftarrow 1:
12
13
        while are Separable and i < n do
            if (s_i \in S^- \text{ and } s_i \text{ is strictly OUTSIDE the circle passing by } p_1, p_2, p_3)
14
            or (s_i \in S^+ \text{ and } s_i \text{ is strictly INSIDE the circle passing by } p_1, p_2, p_3) then
             | areSeparable \leftarrow FALSE ;
16
17
18 return areSeparable;
```

4 Extra works

Question 7 Modify your recognition procedure in order to have an on-line algorithm, which takes input points two by two (one belonging to the boundary of Z and one belonging to the boundary of \bar{Z}) and updates the current separating circle on the fly. What is the time complexity of your procedure?

Question 8 Use your on-line procedure to partition a contour into digital circular arcs and to compute the whole set of maximal digital circular arcs.