

## TP5: Convex Hulls and Digital Convex Hulls

In this TP, the idea is to implement a convex hull algorithm and to experiment its complexity (number of edges) on digital objects.

### Preliminaries Writing svg files

In this lab, you will need to write images of points and lines. A practical way is to write svg files, which are text files using an xml markup language. They can be directly opened, for instance, in your browser. A simple svg file writing some text and plotting one dot and one line looks like:

```
<svg width="512" height="512" xmlns="http://www.w3.org/2000/svg">
  <line x1="15" y1="5" x2="500" y2="95" stroke="red" />
  <circle cx="256" cy="256" r="5" fill="blue" />
  <text x="210" y="60">Text</text>
</svg>
```

### Exercise 1 Orientation predicate

**Question:** Implement the *Orientation*( $p, q, r$ ) predicate as discussed in the lecture.

- To avoid numerical issues, we encourage you to consider points in  $\mathbb{Z}^2$  on a digital domain (instead of  $\mathbb{R}^2$ ).
- To display a line in DGtal between Points  $p$  and  $q$ , you can use the following method of the Board class: `board.drawLine( p[0], p[1], q[0], q[1]);`

**Question:** Test the *Orientation* predicate with an implementation of the segment-segment intersection detection (cf lecture). Experiment your intersection test on all cases (regular intersection, alignment, no intersection, intersection point is a vertex, ...).

The rest of the TP focuses on the implementation and the experimentation of a *convex hull algorithm*, namely, Graham's scan algorithm. At the end we would like you:

- To test the convex hull construction on point sets defined by the digitization (at a given resolution  $h$ ) of a disc defined as a digital set. You will write an svg file to display your result.
- To plot (using `gnuplot` or `svg`) the number of edges when  $h \rightarrow 0$  in log-scale. The aim is to observe the  $N^{2/3}$  behavior we discussed in the lecture for convex hull in  $N \times N$  domains.

Graham's scan implementation on the point set ( $O(n \log n)$ ) should only use the *Orientation*( $p, q, r$ ) predicate and point coordinate comparisons.

## Exercise 2 - Convex hull

### Questions:

- Implement the Graham's scan algorithm as described in the lecture:
  - First, sort the points by polar angle (cf the `qsort` C function man page or the C++ `std::sort` to do the sort). As discussed in the lecture, you would just have to replace the comparison function/functor by the Orientation predicate with fixed  $p_0$ .
  - The second step consists in a simple stack based removal (using for example `std::queue`).
- Display the number of edges as a function of  $h$  in a graph to observe the  $N^{2/3}$  behavior.

Note that Graham's scan can be sped up just considering border point and not interior points (for experiments on digital shapes).