

# Using PLY For Great Good

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# Whirlwind Tour of Grammars

## Definition: Context Free Grammar

- $V$  a finite set of non-terminals or variables
- $\sigma$  a set of terminals
- $R$  the relation of productions
- $S$  the grammars start style

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## Example: $L_{ab} = \{a^n b^n : n \geq 0\}$

- $V = S$
- $\sigma = a, b$
- $R = \{S \rightarrow aSb\}$

# Using grammars in the real world

## A few things they do

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## BNF: the grammar you know

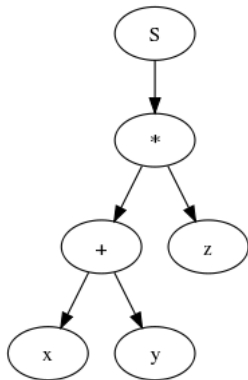
- Developed by John Backus and Peter Naur in the 1960's
- Looks like  $\langle symbol \rangle ::= \_expression\_$
- Also looks like

$list\_display ::= "[expression\_list|list\_comprehension]"$   
 $list\_comprehension ::= expressionlist\_for$

<http://docs.python.org/reference/expressions.html>

# Some more grammar business

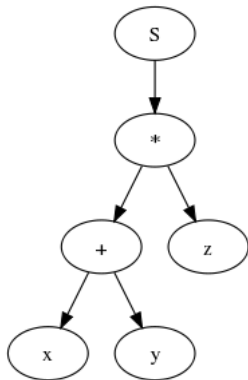
## The Parse Tree



[http://en.wikipedia.org/wiki/Context-free\\_grammar](http://en.wikipedia.org/wiki/Context-free_grammar)

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## The Parse Tree



## The Matching Grammar

$$T \rightarrow x$$

$$T \rightarrow y$$

$$T \rightarrow z$$

$$S \rightarrow S + T$$

$$S \rightarrow S - T$$

$$S \rightarrow S * T$$

$$S \rightarrow S / T$$

$$T \rightarrow (S)$$

$$S \rightarrow T$$

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# Introducing the tools

## PLY

- For building Yacc Style Grammars
- Defining lexer tokens are the terminal set  $\sigma$
- Defining production rules represent  $R$
- Production rules are defined in docstrings
- Always resolves Shift/Reduce with Shift



# Introducing the tools

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## Example: PLY Grammar

```
def p_expression_plus(p):  
    '''expression : expression PLUS term'''  
    p[0] = p[1] + p[3]  
  
def p_expression_term(p):  
    '''expression : term'''  
    p[0] = p[1]
```

# What to avoid when building grammars

## Don't

- use complex lexer tokens
- write the grammar before the semantics
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## Do

- write out the productions for your grammar on paper
- allow the parser to do the syntax not the lexer
- only do one thing in each production
- try optimizing only when your grammar is complete

# What I ended up doing

The output... grrrr

1 subgoal

---

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$$\text{forall } A \ B : \text{Prop}, (A \rightarrow (\sim \sim A) \ \backslash / \ B)$$

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A few problems

- regex won't easily work
- this is still only a simple bit of output

# A bit of grammar the implementation

```
from ply import yacc
from lexer import tokens
__all__ = ['tokens', 'parser', 'precedence']

def p_proofst(p):
    '''proofst : subgoal hyp goal
               | subgoal goal'''
    if len(p) == 4:
        p[0] = dict(subgoal=p[1], hyp=p[2], goal=p[3])
    else:
        p[0] = dict(subgoal=p[1], goal=p[2])

def p_subgoal(p):
    '''subgoal : NUMBER SUBGOAL'''
    p[0] = '_' . join((str(p[1]), str(p[2])))

def p_hyp(p):
    '''hyp : ID COLON PROP hyp
           | ID COLON expr hyp
           | ID COLON PROP
           | ID COLON expr
           ,'''
    hyp = [dict(name=p[1], type=p[3])]
    if len(p) == 5:
        p[0] = hyp + [p[4]]
    else:
        p[0] = dict(name=p[1], type=p[3])

parser = yacc.yacc(debug=True)
```

## Seventh inning stretch: Demo Time



<http://imgur.com/gallery/f5kzo>

# Additional Resources

- [http://en.wikipedia.org/wiki/Context-free\\_grammar](http://en.wikipedia.org/wiki/Context-free_grammar)
- <http://dinosaur.compilertools.net/>
- <http://www.dabeaz.com/ply/>
- <http://www.mostly-decidable.org>
- <http://github.com/dcolish/Cockerel/tree/master/coqd/parser/>