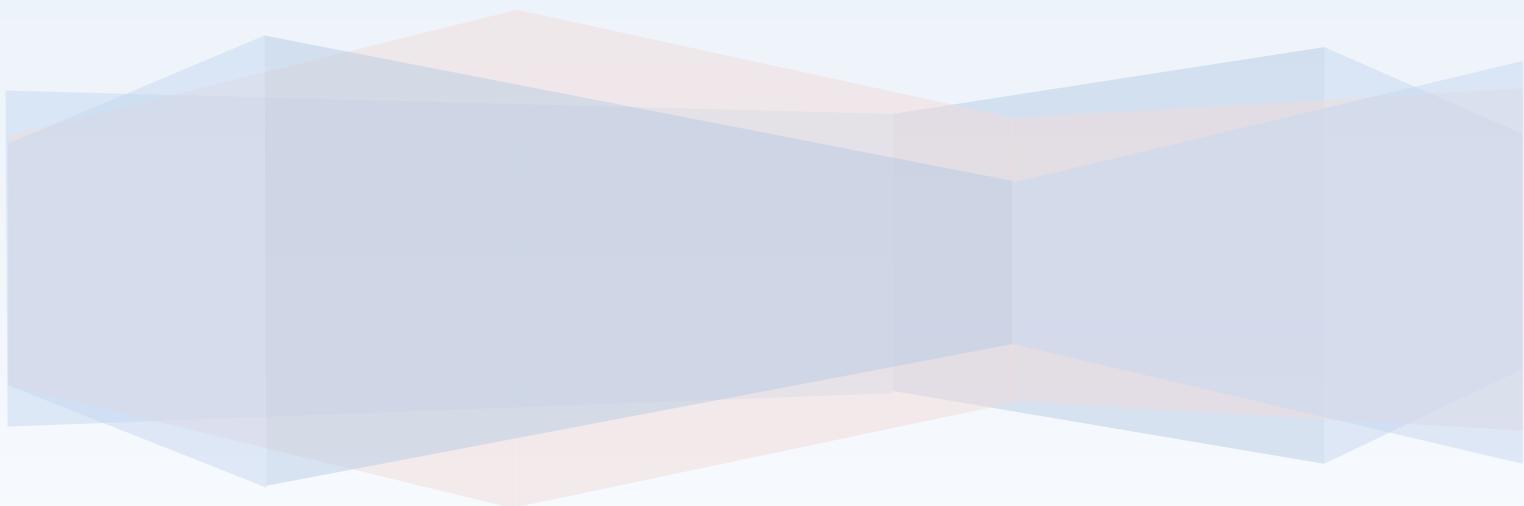


TNE30019 – Unix for Telecommunications

Learning Summary Report

Michael Reeb



Introduction

This report summarises what I have learnt in Unix for Telecommunications. It includes a summary of the completed tasks within the Unit Portfolio and a reflection on my learning.

Assessment Summary

Please complete the tables below to summarise completion of individual tasks in this Portfolio

	Pass	Credit	Distinctio n	High Distinction		
				HD1	HD2	HD3
Lab Report 1	x					
Lab Report 2		x				
Lab Report 3			x			
Jails Report			x			
Research Report					x	
Project						x
Presentation			x			

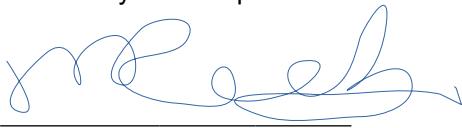
Graded Portfolio Tasks

	Included (please tick)
Pass Tasks (14 tasks)	
Credit Tasks (3 tasks)	
Distinction Tasks (5 tasks)	

Non-Graded Portfolio Tasks

Declaration

I declare that this portfolio is my individual work. I have not copied from any other student's work or from any other source except where due acknowledgment is made explicitly in the text, nor has any part of this submission been written for me by another person.

Signature: 

Reflection

The most important things I learnt:

This unit taught me the fundamentals of Unix from command line. It has removed apprehension and uncertainty that I have faced when using *nix previously.

Entering this unit I had never used Python or Bash before, so learning these two scripting languages was very useful and has increased the amount of programming tools in my tool-chest.

The things that helped me most were:

Google, various textbooks, and Jonathon.

I found the following topics particularly challenging:

CUPS; printers are evil.

I found the following topics particularly interesting:

I really enjoyed learning about the underlying operating system structure. I learnt a lot about how computer operating systems work this semester.

I also really enjoyed the Nmap and TCPDump modules. It was interesting doing some basic port-scanning and network analysis. I might even do some further reading into network security when I find the time.

It was also interesting learning about all the different Unix systems. I am amazed about how easy it is to get applications to run (once you know what to do of course), and how compatible everything is.

That said, I also found that it can be very tedious trying to navigate poorly documented obfuscated manuals and the like, especially when it comes to niche libraries and modules (I am looking at you Apple).

I also found the research papers (Jails and my own topic) to be very interesting. I really enjoy learning, so reading papers is fun.

I feel I learnt these topics, concepts, and/or tools really well:

- I learnt to use Python to beginner level, and learnt about some of its limitations.
- I think that I learnt the course content well, but I am well aware that the course was just a brief introduction to different Unix applications and systems. I am not going to claim expertise in any of these topics.

I still need to work on the following areas:

Everything really. We only had introductions to all of the topics, and any more understanding of these would be useful.

My progress in this unit was ...:



This unit will help me in the future:

This unit has made me comfortable enough and confident to now to start using command line *nix in my own projects. I have a few raspberry Pis that I intend on using to their limits.

I also want to build a home server system, and find an openwrt compatible router so I can play around with it and see if I can optimise my home network.

I really am happy to have had a course on command-line Unix; exposure to this will benefit me greatly.

If I did this unit again I would do the following things differently:

I don't think I would do much differently. I should have put a bit more effort into the Research Report to get the 3rd tier HD, but I wouldn't know what to improve without pinpoint feedback. I am also livid that I got the question wrong on Test 1 about the FreeBSD '&&' vs '| |', as I got them the wrong way around in the test (I even wrote them both down on the paper).

I should have done better on Test 2. Next time I would have to study more.

I really gave this unit as much effort as I was willing to commit. As the saying goes: "I did my best, I have no regrets".

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

MICHAEL REEB

Portfolio Submission

Submitted By:

Michael REEB
4936094

Tutor:

Jonathan KUA

June 5, 2019



Contents

1 Learning Summary Report	1
2 Overall Task Status	2
3 Learning Outcomes	3
3.1 ULO1	3
3.2 ULO2	3
3.3 ULO3	3
3.4 ULO4	3
3.5 ULO5	4
3.6 ULO6	4
3.7 ULO7	4
4 Practical - Lab 2: LaTeX (D)	5
5 Practical - Lab 2: LaTeX (P)	9
6 Practical - Lab 3: Bind (P)	13
7 Communications - Lab Report 1	17
8 Practical - Lab 4: Basic Apache (C)	22
9 Practical - Lab 4: Basic Apache (P)	27
10 Theory - Test 1 (P)	31
11 Communications - Research Report	34
12 Practical - Lab 5: Advanced Apache (D)	43
13 Practical - Lab 5: Advanced Apache (P)	47
14 Communications - Lab Report 2	52
15 Practical - Lab 6: PCAP Programming (D)	62
16 Practical - Lab 6: PCAP Programming (P)	66
17 Practical - Lab 7: NMap (C)	74
18 Practical - Lab 7: NMap (P)	77
19 Practical - Programming Tutorial	81
20 Theory - Jails	86
21 Practical - Lab 8: TCPDump (C)	91
22 Practical - Lab 8: TCPDump (D)	95
23 Practical - Lab 8: TCPDump (P)	101
24 Communications - Speaker Reflection 1	105
25 Communications - Lab Report 3	108

26 Practical - Lab 9: Samba (D)	113
27 Practical - Lab 9: Samba (P)	118
28 Communications - Project Presentation	123
29 Communications - Speaker Reflection 2	129
30 Theory - Test 2 (P)	132
31 Practical - Lab 10: CUPS (D)	135
32 Practical - Lab 10: CUPS (P)	140
33 Practical - Project	144

2 Overall Task Status

Task	Status	Times Assessed
Communications - Research Report	Complete	1
Practical - Lab 1: Introduction to Unix	Complete	1
Theory - Plagiarism Test (P)	Complete	1
Practical - Lab 2: LaTeX (D)	Complete	1
Practical - Lab 2: LaTeX (P)	Complete	1
Theory - Operating Systems (P)	Complete	1
Practical - Lab 3: Bind (P)	Complete	1
Communications - Lab Report 1	Complete	1
Practical - Lab 4: Basic Apache (C)	Complete	1
Practical - Lab 4: Basic Apache (P)	Complete	1
Theory - Test 1 (P)	Complete	1
Practical - Lab 5: Advanced Apache (D)	Complete	1
Practical - Lab 5: Advanced Apache (P)	Complete	1
Theory - Jails	Complete	1
Communications - Lab Report 2	Complete	1
Practical - Lab 6: PCAP Programming (D)	Complete	1
Practical - Lab 6: PCAP Programming (P)	Complete	1
Practical - Project	Complete	1
Practical - Lab 7: NMap (C)	Complete	1
Practical - Lab 7: NMap (P)	Complete	1
Practical - Lab 8: TCPDump (C)	Complete	1
Practical - Lab 8: TCPDump (D)	Complete	1
Practical - Lab 8: TCPDump (P)	Complete	1
Practical - Programming Tutorial	Complete	2
Communications - Lab Report 3	Complete	1
Communications - Project Presentation	Complete	1
Practical - Lab 9: Samba (D)	Complete	1
Practical - Lab 9: Samba (P)	Complete	1
Communications - Speaker Reflection 1	Complete	1
Communications - Speaker Reflection 2	Complete	1
Practical - Lab 10: CUPS (D)	Complete	1
Practical - Lab 10: CUPS (P)	Complete	1
Theory - Test 2 (P)	Complete	1

3 Learning Outcomes

3.1 ULO1

Use navigation processes in an installed Unix System

Task	Rating	Status	Times Assessed
Practical - Lab 1: Introduction to Unix	♦♦♦◊◊	Complete	1
Theory - Plagiarism Test (P)	♦♦◊◊◊	Complete	1
Theory - Test 1 (P)	♦♦◊◊◊	Complete	1
Practical - Project	♦♦♦♦◊	Complete	1
Theory - Test 2 (P)	♦♦◊◊◊	Complete	1

3.2 ULO2

Appreciate the operation of a Unix-based Operating System

Task	Rating	Status	Times Assessed
Theory - Operating Systems (P)	♦♦♦♦◊	Complete	1
Practical - Lab 3: Bind (P)	♦♦♦◊◊	Complete	1
Practical - Lab 4: Basic Apache (C)	♦♦♦◊◊	Complete	1
Practical - Lab 4: Basic Apache (P)	♦♦♦◊◊	Complete	1
Theory - Test 1 (P)	♦♦◊◊◊	Complete	1
Practical - Project	♦♦♦♦◊	Complete	1
Theory - Test 2 (P)	♦♦◊◊◊	Complete	1

3.3 ULO3

Conduct the administration of a Unix server or workstation

Task	Rating	Status	Times Assessed
Theory - Test 1 (P)	♦♦◊◊◊	Complete	1
Practical - Project	♦♦♦◊◊	Complete	1
Practical - Lab 6: PCAP Programming (D)	♦♦♦◊◊	Complete	1
Practical - Lab 6: PCAP Programming (P)	♦♦♦◊◊	Complete	1
Practical - Lab 7: NMap (C)	♦♦♦◊◊	Complete	1
Practical - Lab 7: NMap (P)	♦♦♦◊◊	Complete	1
Practical - Lab 8: TCPDump (C)	♦♦♦◊◊	Complete	1
Practical - Lab 8: TCPDump (D)	♦♦♦◊◊	Complete	1
Practical - Lab 8: TCPDump (P)	♦♦♦◊◊	Complete	1
Theory - Test 2 (P)	♦♦◊◊◊	Complete	1

3.4 ULO4

Configure common network services, devices and security

Task	Rating	Status	Times Assessed
Theory - Test 1 (P)	♦♦◊◊◊	Complete	1
Practical - Lab 5: Advanced Apache (D)	♦♦♦♦♦	Complete	1
Practical - Lab 5: Advanced Apache (P)	♦♦♦♦♦	Complete	1
Practical - Project	♦♦♦♦◊	Complete	1
Practical - Programming Tutorial	♦♦♦♦◊	Complete	2
Communications - Project Presentation	♦♦♦♦◊	Complete	1
Practical - Lab 9: Samba (D)	♦♦♦♦◊	Complete	1
Practical - Lab 9: Samba (P)	♦♦♦♦◊	Complete	1
Theory - Test 2 (P)	♦♦◊◊◊	Complete	1

3.5 ULO5

Demonstrate the use of administration tools on Unix systems

Task	Rating	Status	Times Assessed
Theory - Test 1 (P)	♦♦◊◊◊	Complete	1
Communications - Lab Report 2	♦♦♦♦◊	Complete	1
Theory - Jails	♦♦♦♦♦	Complete	1
Practical - Project	♦♦♦♦◊	Complete	1
Theory - Test 2 (P)	♦♦◊◊◊	Complete	1

3.6 ULO6

Design and construct unfamiliar services

Task	Rating	Status	Times Assessed
Practical - Lab 10: CUPS (P)	♦♦♦♦♦	Complete	1
Practical - Lab 10: CUPS (D)	♦♦♦♦♦	Complete	1
Practical - Project	♦♦♦♦◊	Complete	1
Theory - Test 2 (P)	♦♦◊◊◊	Complete	1
Theory - Test 1 (P)	♦♦◊◊◊	Complete	1

3.7 ULO7

Generate documentation for laboratory work and a research assignment

Task	Rating	Status	Times Assessed
Communications - Research Report	♦♦♦♦♦	Complete	1
Practical - Lab 2: LaTeX (D)	♦♦♦◊◊	Complete	1
Practical - Lab 2: LaTeX (P)	♦♦♦◊◊	Complete	1
Theory - Operating Systems (P)	♦♦♦♦◊	Complete	1
Communications - Lab Report 1	♦♦♦♦◊	Complete	1
Communications - Speaker Reflection 1	♦♦♦♦◊	Complete	1
Theory - Test 1 (P)	♦♦◊◊◊	Complete	1
Communications - Lab Report 2	♦♦♦♦◊	Complete	1
Practical - Project	♦♦♦♦◊	Complete	1
Communications - Lab Report 3	♦♦♦♦◊	Complete	1
Communications - Project Presentation	♦♦♦◊◊	Complete	1
Communications - Speaker Reflection 2	♦♦♦♦◊	Complete	1
Theory - Test 2 (P)	♦♦◊◊◊	Complete	1

4 Practical - Lab 2: LaTeX (D)

Lab task - LaTeX (Distinction)

Outcome	Weight
ULO7	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 2: LaTeX (D)

Submitted By:

Michael REEB

4936094

2019/03/20 11:41

Tutor:

Jonathan KUA

March 20, 2019



Sample Title

Sum Gai*

Ann Otherguy*

sumgai@corporation.com

aotherguy@cooperation.com

FSET Swinburne

Hawthorn, VIC

ABSTRACT

This is the abstract. admire its abstractness. Nam id fermentum dui. Suspendisse sagittis tortor a nulla mollis, in pulvinar ex pretium. Sed interdum orci quis metus euismod, et sagittis enim maximus. Vestibulum gravida massa ut felis suscipit congue. Quisque mattis elit a risus ultrices commodo venenatis eget dui. Etiam sagittis eleifend elementum.

Nam interdum magna at lectus dignissim, ac dignissim lorem rhoncus. Maecenas eu arcu ac neque placerat aliquam. Nunc pulvinar massa et mattis lacinia.

KEYWORDS

component, formatting, style, styling, insert

ACM Reference Format:

Sum Gai and Ann Otherguy. 2018. Sample Title. In *Woodstock '18: ACM Symposium on Neural Gaze Detection, June 03–05, 2018, Woodstock, NY*. ACM, New York, NY, USA, 2 pages. <https://doi.org/10.1145/1122445.1122456>

1 SECTION

This is an example of a section.

1.1 Subsection

This is an example of a subsection

1.1.1 *Subsubsection*. this is an example of a subsubsection. **Now we are in bold font!**

2 LISTS

2.1 Itemised List

- (1) This is the first item
- (2) This is the second item
- (3) This is the third item
- (4) The list goes on
 - (a) this list is nested in another list
 - (b) WOW

*Both authors contributed equally to this research.

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Woodstock '18, June 03–05, 2018, Woodstock, NY

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<https://doi.org/10.1145/1122445.1122456>

- (i) and now we get deeper
- (ii) This text is underlined!
- (5) *This line is in italics*

2.2 Bullet List

- This list uses bullets
- BANG
 - Now we step in again
 - pow
 - * THIS IS IN SMALL CAPS
 - * hello

3 IMAGE/FIGURE INCLUDING REFERENCING

$$x = \sqrt{\pi + e} \int_0^{\infty} \frac{\sinh \theta}{\theta^2} d\theta \quad (1)$$

Figure 1: This is an equation



Figure 2: This is an picture. Seems familiar eh?

This Figure above (Figure 1) can be referenced by using the `\ref{label}` command. As can the picture which can be seen in Figure 2.

4 TABLE

Table 1: This is a table.

Column 1	Column 2	Column 3
Hello but under	I am a pretty this line	nice table is a hard line
now	lets merge	some columns
Merged column takes up 2 spots		purple
		this one takes up all three yeahhhh

We can also reference tables similarly to what we did in Section 3, by labelling and referencing. eg Table 1.

REFERENCING

Acknowledging works similar to referencing, but we use the `\cite{bibID}` command instead [?]. These can be used anywhere in a document. Note the LATEXjargon regarding cite and referencing means different things. referencing references a label, while cite cites the bibliography

ACKNOWLEDGMENTS

To Unix, for being useful. To Jonathan for randomly choosing the same picture as me.

A RESEARCH METHODS

A.1 Part One

Lorem ipsum dolor sit amet, consectetur adipiscing elit. Morbi malesuada, quam in pulvinar varius, metus nunc fermentum urna, id sollicitudin purus odio sit amet enim. Aliquam ullamcorper eu ipsum vel mollis. Curabitur quis dictum nisl. Phasellus vel semper risus, et lacinia dolor. Integer ultricies commodo sem nec semper.

A.2 Part Two

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B ONLINE RESOURCES

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Nam interdum magna at lectus dignissim, ac dignissim lorem rhoncus. Maecenas eu arcu ac neque placerat aliquam. Nunc pulvinar massa et mattis lacinia.

REFERENCES

- [1] D. Doom, M. Evil, and I. Dominate, “Using LaTex for world domination,” Soc. World Enslavement For Evil Masterminds. Bermuda, vol. A247, pp. 529–551, April 2009.

5 Practical - Lab 2: LaTeX (P)

Lab task - LaTeX

Outcome	Weight
ULO7	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 2: LaTeX (P)

Submitted By:

Michael REEB

4936094

2019/03/15 14:38

Tutor:

Jonathan KUA

March 15, 2019



Sample Title

1st Author 1
Author 1 Organisation
name of organisation
City, Country
email address

Author 2
Author 2 Organisation
name of organisation
City, Country
email address

Abstract—This is the abstract. admire its boldness. Lorem ipsum dolor sit amet, consectetur adipiscing elit. Etiam lobortis facilisis sem. Nullam nec mi et neque pharetra sollicitudin. Praesent imperdiet mi nec ante. Donec ullamcorper, felis non sodales commodo, lectus velit ultrices augue, a dignissim nibh lectus placerat pede. Vivamus nunc nunc, molestie ut, ultricies vel, semper in, velit. Ut porttitor. Praesent in sapien. Lorem ipsum dolor sit amet, consectetur adipiscing elit. Duis fringilla tristique neque. Sed interdum libero ut metus. Pellentesque placerat. Nam rutrum augue a leo. Morbi sed elit sit amet ante lobortis sollicitudin. Praesent blandit blandit mauris. Praesent lectus tellus, aliquet aliquam, luctus a, egestas a, turpis. Mauris lacinia lorem sit amet ipsum. Nunc quis urna dictum turpis accumsan semper.

Index Terms—component, formatting, style, styling, insert

I. SECTION

This is an example of a section.

A. Subsection

This is an example of a subsection

1) *Subsubsection*: this is an example of a subsubsection.

Now we are in bold font!

II. LISTS

A. Itemised List

- 1) This is the first item
- 2) This is the second item
- 3) This is the third item
- 4) The list goes on
 - a) this list is nested in another list
 - b) WOW
 - i) and now we get deeper
 - ii) This text is underlined!
- 5) *This line is in italics*

B. Bullet List

- This list uses bullets
- BANG
 - Now we step in again
 - pow
 - * THIS IS IN SMALL CAPS
 - * hello

III. IMAGE/FIGURE INCLUDING REFERENCING

$$x = \sqrt{\pi + e} \int_0^{\infty} \frac{\sinh \theta}{\theta^2} d\theta \quad (1)$$

Fig. 1. This is an equation



Fig. 2. This is an picture. Seems familiar eh?

This Figure above (Figure 1) can be referenced by using the `\ref{label}` command. As can the picture which can be seen in Figure 2.

IV. TABLE

TABLE I
THIS IS A TABLE.

Column 1	Column 2	Column 3
Hello but under	I am a pretty this line	nice table is a hard line
now	lets merge	some columns
Merged column takes up 2 spots		purple
this one takes up all three yeahhhh		

We can also reference tables similarly to what we did in Section III, by labelling and referencing. eg Table I.

REFERENCING

Acknowledging works similar to referencing, but we use the `\cite{bibID}` command instead [1]. These can be used anywhere in a document. Note the L^TE_Xjargon regarding cite and referencing means different things. referecning refernces a label, while cite cites the bibliography

REFERENCES

- [1] D. Doom, M. Evil, and I. Dominate, “Using LaTex for world domination,” Soc. World Enslavement For Evil Masterminds. Bermuda, vol. A247, pp. 529–551, April 2009.

6 Practical - Lab 3: Bind (P)

Lab task - Bind

Outcome	Weight
ULO2	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 3: Bind (P)

Submitted By:

Michael REEB

4936094

2019/03/20 21:10

Tutor:

Jonathan KUA

March 20, 2019



```
Marking P-Lab-03-Bind-P on RULE host rule61
TNE30019/TNE80014 - Portfolio Task - P-Lab03 - BIND - Pass
=====
*****
Portfolio Task: P-Lab-03-Bind, Pass Task

Configure your rule host (rule61) with a forwarding DNS server

Tasks:
Query for: labbind.unix      ; Correct answer: 136.186.230.61
Query for: labnmap.unix      ; Correct answer: 136.186.230.21
Query for: caia.swin.edu.au   ; Correct answer: 136.186.229.16
Query for: 136.186.230.61    ; Correct answer: ns1.unix
*****
```

Date: 20190320_2108

Last 5 Logons:

```
student pts/0 vpn247-80.cc.swin.edu. Wed Mar 20 20:34 still logged in
student pts/0 vpn247-80.cc.swin.edu. Wed Mar 20 20:10 - 20:34 (00:23)
root pts/0 vpn247-80.cc.swin.edu. Wed Mar 20 14:42 - 17:51 (03:09)
student pts/1 vpn247-80.cc.swin.edu. Wed Mar 20 12:03 - 15:13 (03:09)
student pts/0 vpn247-80.cc.swin.edu. Wed Mar 20 11:08 - 14:17 (03:09)
```

RULE host Socket information

```
bind named 24347 20 tcp4 136.186.230.61:53 *:*
bind named 24347 21 tcp4 136.186.230.61:953 *:*
bind named 24347 512 udp4 136.186.230.61:53 *:*
bind named 24347 513 udp4 136.186.230.61:53 *:*
bind named 24347 514 udp4 136.186.230.61:53 *:*
bind named 24347 515 udp4 136.186.230.61:53 *:*
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bind named 24347 534 udp4 136.186.230.61:53 *:*
tempuser sshd 18410 3 tcp4 136.186.230.61:22 136.186.247.80:57007
tempuser sshd 18410 7 tcp4 136.186.230.61:6010 *:*
root sshd 18406 3 tcp4 136.186.230.61:22 136.186.247.80:57007
root sshd 97762 3 tcp4 136.186.230.61:22 *:*
root ntpd 1207 61 udp4 136.186.230.61:123 *:*
```

Performing DNS lookups on 136.186.230.61

```
=====
Querying for (labbind.unix) using DNS server (136.186.230.61) - Correct answer (136.186.230.61)
Using domain server:
```

Name: 136.186.230.61
Address: 136.186.230.61#53
Aliases:

labbind.unix has address 136.186.230.61
Test Result: PASSED

Querying for (labnmap.unix) using DNS server (136.186.230.61) - Correct answer (136.186.230.21)
Using domain server:
Name: 136.186.230.61
Address: 136.186.230.61#53
Aliases:

labnmap.unix has address 136.186.230.21
Test Result: PASSED

Querying for (caia.swin.edu.au) using DNS server (136.186.230.61) - Correct answer (136.186.229.16)
Using domain server:
Name: 136.186.230.61
Address: 136.186.230.61#53
Aliases:

caia.swin.edu.au has address 136.186.229.16
Test Result: PASSED

Querying for (136.186.230.61) using DNS server (136.186.230.61) - Correct answer (ns1.unix.)
Using domain server:
Name: 136.186.230.61
Address: 136.186.230.61#53
Aliases:

61.230.186.136.in-addr.arpa domain name pointer ns1.unix.
Test Result: PASSED

PORTFOLIO TASK RESULT: PASSED

7 Communications - Lab Report 1

Required Lab Report 1

Outcome	Weight
ULO7	◆◆◆◆◇

Date	Author	Comment
2019/04/10 09:18	Jonathan Kua	Great work! Well-written results and discussion sections. The content has achieved Distinction level. One thing you might try for future reports is to develop your own diagrams in EPS format, if applicable.

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Communications - Lab Report 1

Submitted By:

Michael REEB

4936094

2019/04/09 22:12

Tutor:

Jonathan KUA

April 9, 2019



TNE30019 Unix for Telecommunications

Lab 4 - Basic-Apache

Michael Reeb (ID: 4936094)
Swinburne University of Technology
Melbourne, Australia
4936094@student.swin.edu.au

Abstract—Apache is a web server application that can be used to build complex web servers. The main configuration file is `httpd.conf` which is used to change basic settings and be used to call other configuration files, such as the `.htaccess` file called in the later sections to create a password authenticated section of the site.

I. INTRODUCTION TO APACHE

Apache, or Apache HTTP Server is an open-source web server application. It supports many different compiled modules which can be used to extend Apaches functionality. Apache is widely used throughout the world, and accounts for 67% of all web-servers. Apache is used to serve, or handle the HTTP requests from connecting client devices.

II. AIM

In this lab, a basic Apache web-server was implemented, and a private page section was implemented demonstrating more of Apaches functionality. This sever was to serve a basic HTML homepage, and the private segment of the site was to be password protected.

III. EQUIPMENT

RULE host 61 was used as the server. SSH access was though PuTTy running on an external PC connected to Swinburne via VPN.

IV. METHODOLOGY

As per lab handout.

V. RESULTS

Unix for Telecommunications - Lab 4 - Michael Reeb - 4936094



Welcome to my webpage. Peace among worlds!

Fig. 1. Capture of the main index page from host RULE62



Fig. 2. Authentication prompt before being able to access the private section of the site

Oh wow, hello there.



SHhhhhh this is my private part Michael Reeb - 4936094

Fig. 3. Capture of the private section. Authentication is needed before being able to see this page

```
root@rule61:/usr/local/www/apache24/data # sockstat -4
USER COMMAND PID FD PROTO LOCAL ADDRESS FOREIGN ADDRESS
student sshd 62606 3 tcp4 136.186.230.61:22 136.186.247.49:64438
root sshd 62595 3 tcp4 136.186.230.61:22 136.186.247.49:64438
student sshd 61859 3 tcp4 136.186.230.61:22 136.186.247.49:64430
student sshd 61859 7 tcp4 136.186.230.61:6010 *:*
root sshd 61844 3 tcp4 136.186.230.61:22 136.186.247.49:64430
www httpd 89149 3 tcp4 136.186.230.61:80 *:*
www httpd 87505 3 tcp4 136.186.230.61:80 *:*
www httpd 87504 3 tcp4 136.186.230.61:80 *:*
www httpd 87503 3 tcp4 136.186.230.61:80 *:*
www httpd 87502 3 tcp4 136.186.230.61:80 *:*
www httpd 87501 3 tcp4 136.186.230.61:80 *:*
www httpd 87500 3 tcp4 136.186.230.61:80 *:*
root httpd 87499 3 tcp4 136.186.230.61:80 *:*
bind named 51658 20 tcp4 136.186.230.61:53 *:*
bind named 51658 21 tcp4 136.186.230.61:953 *:*
bind named 51658 512 udp4 136.186.230.61:53 *:*
bind named 51658 513 udp4 136.186.230.61:53 *:*
bind named 51658 514 udp4 136.186.230.61:53 *:*
bind named 51658 515 udns4 136.186.230.61:53 *:*
```

Fig. 4. Output of the command SOCKSTAT -4 while Apache is running. Note that Apache service is identified as HTTPD, and it is running a www service from port 80

```
root@rule61:/usr/local/www/apache24/data # cat /var/log/httpd-access.log
136.186.247.203 - - [31/Mar/2019:14:33:08 +1000] "GET /HTTP/1.1" 200 268
136.186.247.203 - - [31/Mar/2019:14:33:08 +1000] "GET /peace.png HTTP/1.1" 200 302893
136.186.247.203 - - [31/Mar/2019:14:33:08 +1000] "GET /favicon.ico HTTP/1.1" 404 209
136.186.247.203 - - [31/Mar/2019:14:35:32 +1000] "GET / HTTP/1.1" 200 326
136.186.247.203 - - [31/Mar/2019:14:35:49 +1000] "GET /test HTTP/1.1" 404 202
root@rule61:/usr/local/www/apache24/data #
```

Fig. 5. Output of the access log of the Apache server.

```
root@rule61:/usr/local/www/apache24/data # cat /var/log/httpd-error.log
[Sun Mar 31 16:33:13.405879 2019] [mpm_prefork:notice] [pid 22182] AH00163: Apache/2.4.33 (FreeBSD) configured -- resuming normal operation
[Sun Mar 31 16:33:13.405879 2019] [core:notice] [pid 22182] AH00094: Command line: "/usr/local/sbin/httpd"
[Sun Mar 31 16:33:13.405879 2019] [mpm_prefork:notice] [pid 22182] AH00163: Apache/2.4.33 (FreeBSD) configured -- resuming normal operation
[Sun Mar 31 15:41:24.329947 2019] [core:notice] [pid 34247] AH00163: Apache/2.4.33 (FreeBSD) configured -- resuming normal operation
[Sun Mar 31 16:23:17.726989 2019] [mpm_prefork:notice] [pid 34247] AH00169: caught SIGTERM, shutting down
[Sun Mar 31 16:23:17.726989 2019] [mpm_prefork:notice] [pid 34247] AH00163: Apache/2.4.33 (FreeBSD) configured -- resuming normal operation
[Sun Mar 31 16:23:17.726989 2019] [core:notice] [pid 42982] AH00094: Command line: "/usr/local/sbin/httpd -D NOHTTUPACCEPT"
[Sun Mar 31 20:27:17.726989 2019] [mpm_prefork:notice] [pid 42982] AH00169: caught SIGTERM, shutting down
[Sun Mar 31 20:27:17.726989 2019] [mpm_prefork:notice] [pid 42982] AH00094: Command line: "/usr/local/sbin/httpd -D NOHTTUPACCEPT"
[Sun Mar 31 20:27:17.726989 2019] [core:notice] [pid 42982] AH00163: Apache/2.4.33 (FreeBSD) configured -- resuming normal operation
[Sun Mar 31 20:36:19.784312 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user pass not found /private/
[Sun Mar 31 20:36:19.784312 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user not found /private/
[Sun Mar 31 20:36:34.506118 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user not found /private/
[Sun Mar 31 20:36:34.506118 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user not found /private/
[Sun Mar 31 20:36:32.508074 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:32.508074 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user not found /private/
[Sun Mar 31 20:36:32.508074 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:37.142872 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:39.033393 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:43.737030 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:43.737030 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:44.370269 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:44.370269 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:44.400418 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:44.400418 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:44.670075 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:44.670075 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:44.768433 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:44.768433 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:44.802105 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:44.802105 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:45.198770 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:45.198770 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:45.559967 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:46.004500 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:46.004500 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:46.004500 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:46.366817 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:46.596061 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:46.596061 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:46.902454 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:46.902454 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
[Sun Mar 31 20:36:47.084377 2019] [auth_basicauth:error] [pid 87500] [client 136.186.247.203:63596] AH01618: user asdf not found /private/
```

Fig. 6. Output of the access log of the Apache server.

VI. DISCUSSION

A. Apache Introduction

The `/etc/services` is the service file that contains a list of all the common services that are run on Unix, the port-number they common use, the protocol the particular service uses, as well as any applicable aliases that the service has. For example, when `sockstat -4` command is ran, this `services` file is queried so the aliases of the services can be seen. Figure 4 is an example of this [3].

B. Apache Configuration File

The Apache configuration `httpd.conf` file determines the basic configuration of the Apache service. This file is main configuration file and can be used to change basic settings, as well as call other modules and other configuration files. Changes to this configuration file require Apache (or the `httpd` service) to restarted to take effect [4].

In this file, we can see that the port that Apache will listen on when running is configured to port 80, which is the well-known port associated with HTTP services. Also it can be seen that the directory that documents will be served from is: `/usr/local/www/apache24/data`. The error log and access log can be found at: `/var/log/httpd-error.log` and `/var/log/httpd-access.log` respectively.

C. Installing a Simple Web Page

A basic page with the name `index.html` was place in the server directory `/usr/local/www/apache24/data`. The page can be seen in Figure 1.

D. Starting Apache

With Apache running, the `sockstat -4` command will show the open IPv4 sockets and which process they relate to. The Apache services have the name `httpd`, and the local address is the RULE IP connected to socket 80 as outlined previously. These processes have no associated "foreign" addresses, which means that they are listening for incoming connections.

The Access Log file as seen in Figure 5 contains the HTTP requests, the IP address of the requesting client, and at what time, and the status code of what the server sends back to the client. In this log for example, we see that client with IP address of 136.186.247.203 requested ((GET/HTTP/1.1)) that the server should send ("GET" part) the domain root file (the / by itself means the domain root which refers to in this case `index.html`) using the protocol HTTP 1.1 (HTTP/1.1). The server responds with the status code (200 in this case) and the size of the returned object (268 bytes). Requesting a page that does not exist would result in a 404 error as the servers response [5].

E. Researching Simple Authentication

`.htaccess` is the configuration file used by Apache to change configurations of different directories so web-server behaviour can be changed. The `.htaccess` file operates on a directory level, and when enabled in the main configuration file (`httpd.conf`), Apache will override the main configuration with the configuration in `.htaccess`. This is done on a per-directory basis, so a server with multiple directories could have many `.htaccess` files change behaviour per directory [1].

F. Create the password file

It is important not to store the password file in a publicly served directory so it is not illegally accessed or modified in anyway. Storing the password file inside the server private directory does not guarantee safety either. By placing the file outside of the Apache serving directory, there is no chance that it may accidentally be accessed by clients, and the Apache software will still be able to access to the password file. The directory chosen for this section was: `/usr/local/etc/apache24/.htpasswd`, which is not inside the Apache serving directory.

G. Create the `.htaccess` file

the `.htaccess` file was created in the Private directory with the following commands:

- `AuthUserFile /usr/local/etc/apache24/.htpasswd`
- `AuthName "No Unauthorised access."`

- `AuthType Basic`
- `Require valid-user`

The `AuthUserFile` command directs Apache to the username/password database to use.

The `AuthName` command is the prompt that is to the client that will appear on the login box.

`AuthType` tell Apache what kind of Authentication to use. Here just the basic authentication is used.

`Require` tell Apache which authenticated users will be granted access [2].

H. Testing

As can be seen in Figure 2, upon trying to access the private section of the page access was not given until a valid user id and password was given. Upon doing so, the private index was served and we were able to see Figure 3. Invalid credentials were also tested, and the page was redirected to the 401 Unauthorised error page. This can also been seen in the `.htaccess-error` log file in Figure 6. It can be seen that many different attempts of login without valid credentials were attempted.

VII. CONCLUSION

Apache is a popular web serving platform that is used for building web servers. It is highly configurable, and can be easily configured for individual-to-individual basis. In this lab, we explored basic configuration, and simple directory security. The `httpd.conf` file is the main configuration file, which can be overwritten on a per-directory basis by `.htaccess` configurations. Password access to private directories is also achievable, and it is important that passwords are stored outside of the Apache serving directories so the password database is not illegally accessed or modified.

REFERENCES

- [1] What is .htaccess? Retrieved from <http://www.htaccess-guide.com/>
- [2] PasswordBasicAuth. (n.d.). Retrieved from <https://wiki.apache.org/httpd/PasswordBasicAuth>
- [3] Haas, J. (2018, December 18). What Is /etc/services in Linux/Unix? Retrieved from <https://www.lifewire.com/what-is-etc-services-2196940>
- [4] Configuration Files. (n.d.). Retrieved from <https://httpd.apache.org/docs/2.4/configuring.html>
- [5] Understanding the Apache Access Log - KeyCDN Support. (n.d.). Retrieved from <https://www.keycdn.com/support/apache-access-log>

8 Practical - Lab 4: Basic Apache (C)

Lab task - Basic Apache (Credit)

Outcome	Weight
ULO2	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 4: Basic Apache (C)

Submitted By:

Michael REEB

4936094

2019/03/31 20:49

Tutor:

Jonathan KUA

March 31, 2019



```
Marking P-Lab-04-Basic-Apache-C on RULE host rule61
TNE30019/TNE80014 - Laboratory - Basic Apache
=====
*****
Portfolio Task: P-Lab-04-Basic-Apache, Credit Task
```

Configure your rule host (rule61) with Apache to serve a single web site

Tasks:

- Apache running on rule61
- Browsing to http://rule61/private will fail with error code 401
- Browsing to http://rule61/private with username(ruleperson)
password(ruleperson) will return a valid web page

Date: 20190331_2048

Last 5 Logons:

student	pts/1	vpn247-203.cc.swin.edu	Sun Mar 31 19:57	still logged in
student	pts/0	vpn247-203.cc.swin.edu	Sun Mar 31 18:59	still logged in
student	pts/3	vpn247-203.cc.swin.edu	Sun Mar 31 14:27 - 18:36	(04:09)
student	pts/1	vpn247-136.cc.swin.edu	Sun Mar 31 11:54 - 14:45	(02:51)
student	pts/8	atc330-19.ict.swin.edu	Mon Mar 25 16:10 - 16:14	(00:04)

RULE host Socket information

www	httpd	89149	3	tcp4	136.186.230.61:80	*:*
www	httpd	87505	3	tcp4	136.186.230.61:80	*:*
www	httpd	87504	3	tcp4	136.186.230.61:80	*:*
www	httpd	87503	3	tcp4	136.186.230.61:80	*:*
www	httpd	87502	3	tcp4	136.186.230.61:80	*:*
www	httpd	87501	3	tcp4	136.186.230.61:80	*:*
www	httpd	87500	3	tcp4	136.186.230.61:80	*:*
root	httpd	87499	3	tcp4	136.186.230.61:80	*:*
tempuser	sshd	83724	3	tcp4	136.186.230.61:22	136.186.247.203:63507
root	sshd	83721	3	tcp4	136.186.230.61:22	136.186.247.203:63507
tempuser	sshd	82089	3	tcp4	136.186.230.61:22	136.186.247.203:63502
tempuser	sshd	82089	7	tcp4	136.186.230.61:6011	*:*
root	sshd	82086	3	tcp4	136.186.230.61:22	136.186.247.203:63502
tempuser	sshd	71710	3	tcp4	136.186.230.61:22	136.186.247.203:63419
tempuser	sshd	71710	7	tcp4	136.186.230.61:6010	*:*
root	sshd	71706	3	tcp4	136.186.230.61:22	136.186.247.203:63419
bind	named	51658	20	tcp4	136.186.230.61:53	*:*
bind	named	51658	21	tcp4	136.186.230.61:953	*:*
bind	named	51658	512	udp4	136.186.230.61:53	*:*
bind	named	51658	513	udp4	136.186.230.61:53	*:*
bind	named	51658	514	udp4	136.186.230.61:53	*:*
bind	named	51658	515	udp4	136.186.230.61:53	*:*
bind	named	51658	516	udp4	136.186.230.61:53	*:*
bind	named	51658	517	udp4	136.186.230.61:53	*:*
bind	named	51658	518	udp4	136.186.230.61:53	*:*
bind	named	51658	519	udp4	136.186.230.61:53	*:*
bind	named	51658	520	udp4	136.186.230.61:53	*:*
bind	named	51658	521	udp4	136.186.230.61:53	*:*
bind	named	51658	522	udp4	136.186.230.61:53	*:*
bind	named	51658	523	udp4	136.186.230.61:53	*:*
bind	named	51658	524	udp4	136.186.230.61:53	*:*
bind	named	51658	525	udp4	136.186.230.61:53	*:*
bind	named	51658	526	udp4	136.186.230.61:53	*:*
bind	named	51658	527	udp4	136.186.230.61:53	*:*
bind	named	51658	528	udp4	136.186.230.61:53	*:*
bind	named	51658	529	udp4	136.186.230.61:53	*:*

```
bind    named      51658 530  udp4   136.186.230.61:53      *:*
bind    named      51658 531  udp4   136.186.230.61:53      *:*
bind    named      51658 532  udp4   136.186.230.61:53      *:*
bind    named      51658 533  udp4   136.186.230.61:53      *:*
bind    named      51658 534  udp4   136.186.230.61:53      *:*
root    sshd       97762  3   tcp4   136.186.230.61:22      *:*
root    ntpd       1207   61  udp4   136.186.230.61:123     *:*
?      ?          ?      ?   tcp4   136.186.230.61:51564   136.186.230.61:80
?      ?          ?      ?   tcp4   136.186.230.61:14865   136.186.230.61:80
?      ?          ?      ?   tcp4   136.186.230.61:80      136.186.247.203:63645
```

```
HTTP QUERY http://rule61/private
=====
--2019-03-31 20:48:23-- http://rule61/private
Resolving rule61 (rule61)... 136.186.230.61
Connecting to rule61 (rule61)|136.186.230.61|:80... connected.
HTTP request sent, awaiting response... 401 Unauthorized
```

Username/Password Authentication Failed.

Test Result: PASSED

```
HTTP QUERY http://rule61/private
- Authentication parameters (--user=ruleperson --password=ruleperson)
=====
--2019-03-31 20:48:23-- http://rule61/private
Resolving rule61 (rule61)... 136.186.230.61
Connecting to rule61 (rule61)|136.186.230.61|:80... connected.
HTTP request sent, awaiting response... 401 Unauthorized
Authentication selected: Basic realm="You wish to enter the private part. Please Enter Password"
Reusing existing connection to rule61:80.
HTTP request sent, awaiting response... 301 Moved Permanently
Location: http://rule61/private/ [following]
--2019-03-31 20:48:23-- http://rule61/private/
Reusing existing connection to rule61:80.
HTTP request sent, awaiting response... 200 OK
Length: 259 [text/html]
Saving to: '/tmp/wget_save.boy7QkPe'
```

OK 100% 29.6M=0s

2019-03-31 20:48:23 (29.6 MB/s) - '/tmp/wget_save.boy7QkPe' saved [259/259]

```
=====
Downloaded Page contents
=====
<html>
<head>
<title> private part </title>
</head>

<body bgcolor="white" text="blue">

<h1> Oh wow, hello there. </h1>
SHhhhhh this is my private part
Michael Reeb - 4936094
```

```
  
</body>
```

```
</html>=====
```

```
Test Result: PASSED
```

```
*****
```

```
PORTFOLIO TASK RESULT: PASSED
```

9 Practical - Lab 4: Basic Apache (P)

Lab task - Basic Apache

Outcome	Weight
ULO2	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 4: Basic Apache (P)

Submitted By:

Michael REEB

4936094

2019/03/31 15:44

Tutor:

Jonathan KUA

March 31, 2019



```
Marking P-Lab-04-Basic-Apache-P on RULE host rule61
TNE30019/TNE80014 - Laboratory - Basic Apache
=====
*****
Portfolio Task: P-Lab-04-Basic-Apache, Pass Task
```

Configure your rule host (rule61) with Apache to serve a single web site

Tasks:

- Apache running on rule61
- Browsing to http://rule61 will return a valid web page

Date: 20190331_1542

Last 5 Logons:

student	pts/3	vpn247-203.cc.swin.edu	Sun Mar 31 14:27	still logged in
student	pts/1	vpn247-136.cc.swin.edu	Sun Mar 31 11:54 - 14:45	(02:51)
student	pts/8	atc330-19.ict.swin.edu	Mon Mar 25 16:10 - 16:14	(00:04)
student	pts/16	atc330-19.ict.swin.edu	Mon Mar 25 14:48 - 15:28	(00:39)
student	pts/0	vpn247-80.cc.swin.edu	Wed Mar 20 23:47 - 23:50	(00:02)

RULE host Socket information

www	httpd	34253	3	tcp4	136.186.230.61:80	*:*
www	httpd	34252	3	tcp4	136.186.230.61:80	*:*
www	httpd	34251	3	tcp4	136.186.230.61:80	*:*
www	httpd	34250	3	tcp4	136.186.230.61:80	*:*
www	httpd	34249	3	tcp4	136.186.230.61:80	*:*
www	httpd	34248	3	tcp4	136.186.230.61:80	*:*
root	httpd	34247	3	tcp4	136.186.230.61:80	*:*
tempuser	sshd	20964	3	tcp4	136.186.230.61:22	136.186.247.203:63191
root	sshd	20961	3	tcp4	136.186.230.61:22	136.186.247.203:63191
tempuser	sshd	20954	3	tcp4	136.186.230.61:22	136.186.247.203:63189
tempuser	sshd	20954	7	tcp4	136.186.230.61:6011	*:*
root	sshd	20951	3	tcp4	136.186.230.61:22	136.186.247.203:63189
bind	named	51658	20	tcp4	136.186.230.61:53	*:*
bind	named	51658	21	tcp4	136.186.230.61:953	*:*
bind	named	51658	512	udp4	136.186.230.61:53	*:*
bind	named	51658	513	udp4	136.186.230.61:53	*:*
bind	named	51658	514	udp4	136.186.230.61:53	*:*
bind	named	51658	515	udp4	136.186.230.61:53	*:*
bind	named	51658	516	udp4	136.186.230.61:53	*:*
bind	named	51658	517	udp4	136.186.230.61:53	*:*
bind	named	51658	518	udp4	136.186.230.61:53	*:*
bind	named	51658	519	udp4	136.186.230.61:53	*:*
bind	named	51658	520	udp4	136.186.230.61:53	*:*
bind	named	51658	521	udp4	136.186.230.61:53	*:*
bind	named	51658	522	udp4	136.186.230.61:53	*:*
bind	named	51658	523	udp4	136.186.230.61:53	*:*
bind	named	51658	524	udp4	136.186.230.61:53	*:*
bind	named	51658	525	udp4	136.186.230.61:53	*:*
bind	named	51658	526	udp4	136.186.230.61:53	*:*
bind	named	51658	527	udp4	136.186.230.61:53	*:*
bind	named	51658	528	udp4	136.186.230.61:53	*:*
bind	named	51658	529	udp4	136.186.230.61:53	*:*
bind	named	51658	530	udp4	136.186.230.61:53	*:*
bind	named	51658	531	udp4	136.186.230.61:53	*:*
bind	named	51658	532	udp4	136.186.230.61:53	*:*
bind	named	51658	533	udp4	136.186.230.61:53	*:*
bind	named	51658	534	udp4	136.186.230.61:53	*:*
root	sshd	97762	3	tcp4	136.186.230.61:22	*:*

```
root      ntpd        1207  61  udp4    136.186.230.61:123    *:*
?          ?          ?      ?  tcp4    136.186.230.61:80     136.186.247.203:63226
?          ?          ?      ?  tcp4    136.186.230.61:80     136.186.247.203:63230
?          ?          ?      ?  tcp4    136.186.230.61:57343   136.186.230.61:80
```

```
HTTP QUERY http://rule61
=====
--2019-03-31 15:42:14--  http://rule61/
Resolving rule61 (rule61)... 136.186.230.61
Connecting to rule61 (rule61)|136.186.230.61|:80... connected.
HTTP request sent, awaiting response... 200 OK
Length: 326 [text/html]
Saving to: '/tmp/wget_save.NxJcb4NA'

OK                                         100% 34.7M=0s
```

```
2019-03-31 15:42:14 (34.7 MB/s) - '/tmp/wget_save.NxJcb4NA' saved [326/326]
```

```
=====
Downloaded Page contents
=====
<html>

<head>
<title> favorites / bookmark title goes here </title>
</head>

<body bgcolor="white" text="blue">

<h1>Unix for Telecommunications - Lab 4 - Michael Reeb - 4936094</h1>
Welcome to my webpage. Peace amoung worlds!


</body>
```

```
</html>
=====
```

```
Test Result: PASSED
```

```
*****
PORTFOLIO TASK RESULT: PASSED
```

10 Theory - Test 1 (P)

In class test run during tutorial in week 5 of semester

Outcome	Weight
ULO3	♦♦◊◊◊
ULO6	♦♦◊◊◊
ULO7	♦♦◊◊◊
ULO2	♦♦◊◊◊
ULO4	♦♦◊◊◊
ULO5	♦♦◊◊◊
ULO1	♦♦◊◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Theory - Test 1 (P)

Submitted By:

Michael REEB

4936094

2019/04/12 14:43

Tutor:

Jonathan KUA

April 12, 2019



Unix for Telecommunications



Portfolio Task - T-Test-1

Pass Level Task

Factors per 100

Student ID

4936094

Student Name

Michael Reeb

Test Date

Week 5

Time Allowed

30 minutes

Assessment – Staff Use Only

In order to pass the test you MUST score a minimum of **80%**

Question:	1	2	3	4	5	Total
Points:	4	6	5	6	7	28
Score:	4	6	4	6	5.5	25.5

91%
PASS

11 Communications - Research Report

Research report required on Assignment topic

Outcome	Weight	
ULO7	◆◆◆◆◆	
Date	Author	Comment
2019/04/30 15:11	Jonathan Kua	Excellent work! Great job!

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Communications - Research Report

Submitted By:

Michael REEB

4936094

2019/04/17 10:54

Tutor:

Jonathan KUA

April 17, 2019



A Brief Overview of LTE Video Broadcast Technologies

Michael Reeb - 4936094
TNE30019 Unix for Telecommunications
Communications Report
4936094@student.swin.edu.au

Abstract—Currently the standard technology used is in cellular mobile networks is LTE, and its successor LTE-A, a Fourth Generation technology that offers bandwidth up-to 1Gbps in the down link and multicast/broadcast capabilities via eMBMS. Modern mobile devices are being released with chipsets that are compatible with this broadcast technology, and coupled with the benefits of the H.264 encoding which allows for the implementation of DASH (and the related HLS) media transport systems, the future for efficient broadcasting high quality, live video over cellular networks with great user experience, is a certainty.

I. INTRODUCTION

With the widespread adoption of portable, powerful smart devices, the demand for high-quality, flexible streaming services has also increased. Currently video content counts for almost 58% of global internet traffic [1], and this consumer video demand has placed a large strain upon cellular telecommunication services who are currently implementing the latest cellular technology to improve network efficiency, improve customer Quality of Service (QoS), and maximise bandwidth capabilities.

The streaming of events, particularly live sport and entertainment, has become a profitable and marketable service for many mobile telecommunication providers who are now moving away from traditional services (such as hardwired phone and voice), and focusing on media streaming services and platforms. This paper overviews the different technology that is needed for broadcasting live-video via a cellular network, and explores cellular technology and future enhancements, video codecs, and media distribution systems.

II. ON-DEMAND STREAMING OPERATION

There are two main models for serving streaming media: the push method where the transmission of data is requested by the server, and the pull method where the request for transmission is performed by the client devices. Each method has its advantages and disadvantages with respect to performance, and application.

A. Push Streaming

A push service is one where the request of the transmission of payload data is performed by the server. Once the client and server have established a connection, the server pushes the new data to the client until the client stops or interrupts the session [8]. This is useful for streaming services, as

it allows the server to perform multi-cast (ie, sending the same payload data to multiple services) which can lower the bandwidth requirements of the server thus reducing redundant transmission. An example a push based service is Internet Radio. A client subscribes to the stream, and data is pushed as it is created by the server. There are no “rewind” type capabilities that are performed by the server(however they may be some local caching on the client device that allows short rewind capabilities) as the client only receives what the server sends.

In these systems, the client device application can be considered passive; there is very little complexity in the client application as it only needs to be able to connect to the server, receive, decode, and buffer the data as before using it. The server holds the complexity and sends commands to the user device regarding QoS and other session changes [8]. Example protocols are SIP (Session Initiation Protocol, used in telephony applications including video conferencing) and RTP (Real Time Protocol, which is used often in live-streaming video and audio, as it also supports multi-cast) [8].

B. Pull Streaming

A pull service is a service where the request for transmission of data is performed on the client side. The content is pre-encoded and retrieved piecemeal by requests from the client. Multiple users each have a unique connection to the server and request the data that they would like to be served. This is useful for streaming services , including live streaming services, as it allows each client to watch the same media as when it is convenient for client, have rewind/playback capabilities (for example, YouTube live has limited rewind capabilities for live streams), and dynamic quality control which can be adjusted on the fly . With pull-type services, the complexity lies in the client application; the server simply serves the requests of the client devices. While less efficient than Push based services this is not a large issue if large bandwidth is available to user devices.

Modern pull based streaming services use HTTP for as the transport protocol, which itself uses TCP to guarantee packet delivery. Unlike the Push streaming examples, there is no built in broadcast or multi-cast support for HTTP/TCP as each user device acts streams the media independently from each other.

This can particularly inefficient if devices are streaming the same data as the server will send redundant copies of the media to each device independently.

Example services of Pull type services are YouTube, (including YouTube Live), Netflix, and Spotify [8].

C. Streaming Implementations

Streaming protocols are used to transport the encoded video data (explored further in Section II-F) from the sever to the client mobile devices. These implementations wrap the encoded video, perform QoS, and live-video quality switching. Historically, RTP and other UDP-based services were used as the transport protocol overhead was small, and thus efficient as bandwidth was often limited and inconsistent. However with modern network technology offering larger bandwidth and faster throughput, this low-overhead has primary become a less of concern with reliability of delivery becoming more of a focus. This focus shift has lead to HTTP over TCP becoming the most common media streaming transport protocol used in current devices [9].

While there are many different HTTP streaming implementations, including Adobe HTTP Dynamic Streaming, and Microsoft's Smooth Streaming, this paper will primary focus on the two most popular: Apples HTTP Live Streaming (HLS), and the MPEG standardised Dynamic Adaptive Streaming over HTTP (MPEG-DASH, or DASH).

D. HLS

HLS is an adaptive bitrate streaming protocol that was developed by Apple and natively supported in QuickTime, Safari, iOS and other Apple related devices and technologies. The popularity of HLS has increased with the popularity of Apple Products and Devices,

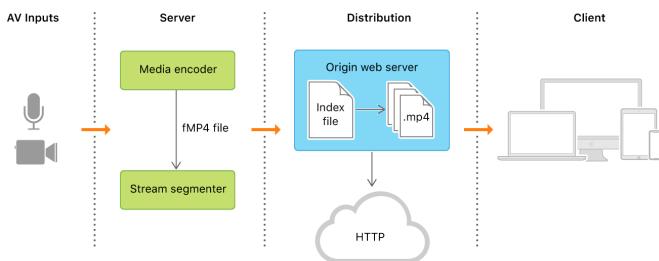


Fig. 1. Flow chart of Apple's HLS [11]

The HLS Media Preparation Server is responsible for taking input streams of audio and video input and digitally encoding them using an encoding scheme (as explored in Section II-F). The encoded content is then placed into an MPEG Transport Stream container, where it is delivered to the Stream Segmenter, where the stream is segmented into chunks approximately 5-10 seconds long of MPEG-Transport Stream video files (.ts files). It is here that the video also can be encoded at different qualities for adaptive stream capabilities. The Stream Segmenter also produces the index files that describe the available chunked .ts files at the

differing quality levels. This chunking and indexing process can also be performed in real-time, for live-streaming content. [9]

The .ts files and index files are then pushed to a HTTP web sever where they are published. The web server is only responsible for delivering the requested chunks [9].

Finally, the client device establishes a connection with the server, where it requests the index and .ts files to download. The client is responsible for selection which quality tier to choose from and adapting the quality during the stream as bandwidth and throughput fluctuates [9].

E. MPEG-DASH

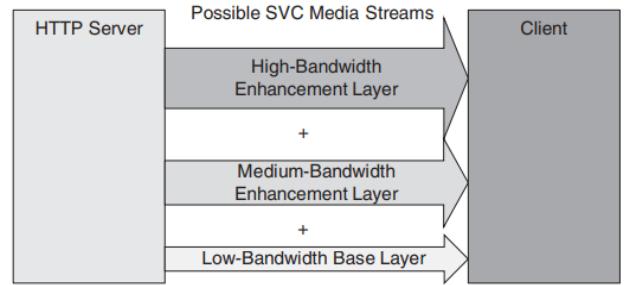


Fig. 2. Example of a 3-Layer SVC. To achieve highest possible quality, the base layer and the two additional enhancement layer streams are needed [9].

DASH is term that describes any HTTP streaming service that allows the client devices to select various quality tiers offered by the HTTP server. DASH is not a protocol, or system itself, but a set of formats that can be used with HTTP to deliver streamable media. Thus, a total analysis of DASH is outside the scope of this paper will not be performed, but a brief overview of MPEG-DASH will be analysed; a specific instance of DASH defined in the ISO/IEC FCD 23009-1 standard [9].

Similarly to HLS, MPEG-DASH utilises HTTP over TCP for transport, and supports various encoding schemes. MPEG-DASH also segments the input audio/video stream into individual chunks of differing quality, and clients connected to the MPEG-DASH web server decide which video quality tier to stream. However, differences lie in the operation of this quality selection.

With MPEG-DASH the adjustment of the quality of the stream is performed by the addition of Scalable Video Codec (SVC) layers. This means that a client will always receive the base SVC layer, and increasing the quality is achieved by the addition of higher enhancement layers. An example of this can be seen in Figure 2.

F. Encoder and Decoder Algorithms, and Protocols

Another important aspect of efficient transmission is the encoder/decoder (codec) algorithm and structure. As wireless broadcast technology is inherently lossy, its is highly likely that during streaming there is a likelihood that packets

will be lost. Now, particularly if this occurs during an LTE-Broadcast (which is inherently a push based service) there is no backwards channel for the mobile devices to request re-transmission of the lost packet. Thus, the codec used has to find a balance between data compression, and ability to withstand packet losses in transmission. [7]. Historically, many different codecs have existed, such as MPEG-2, and MPEG-4. These have been supplanted by modern codecs that are doubly as effective (ie for a given bandwidth are able to provide twice the video quality [9]). The commonly used codec for live-streaming is H.264. H.264 is an industry standard video compression system that commonly is used for video streaming services, being both used in HLS and DASH delivery systems outlined in Section II-D and Section II-E.

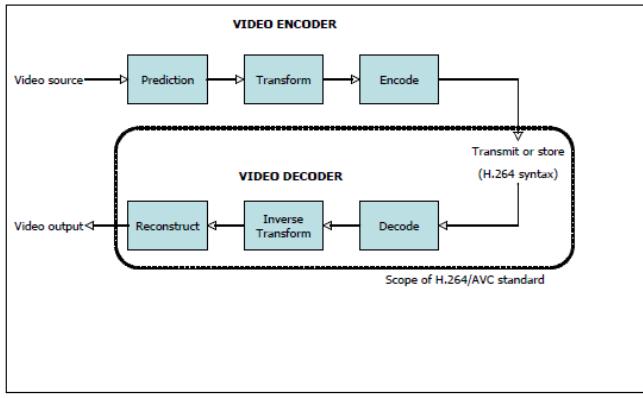


Fig. 3. H264 encoding and decoding process [13].

The encoding and decoding process may vary between implementations of H.264, but all will follow a similar process to that in Figure 3. First, the encoder processes breaks down the frame of the video in what is known as a Macroblock (a 16×16 pixel box) which is compared to previous coded data of the current frame (known as Intra prediction), or previously coded frames (known as Inter prediction). The encoder subtracts the prediction from the current macro-blocks to obtain the residual. A block of these residuals is then transformed into a set of weighted co-coefficients which is then quantised to some value, called the Quantisation Parameter (QP) [13]. This QP can be changed depending on the desired Quality/Compression ratio (ie higher compression means lower quality, and high quality means low compression) [9].

The output of the encoding processes is then used to create a bit stream. This bit stream is the information that is sent to servers and received by client devices. The decoding process is the reverse of this, and a video image is reconstructed.

A H.264 stream consists of different frames, categorised by the amount of prediction used by the encoder when encoding said frame. An I-frame contains no prediction at all, and are sent at certain intervals to ensure that the codec reconstruction

process has something to synchronise to. If I-frames are lost, any frames that reference this I-frame cannot be reconstructed. A P-frame is a frame that has prediction from a previous I-frame, or from another P-frame. And a B-frame contains predictions from both previous and subsequent frames [9]. Reference frames will be significantly smaller than non-referenced frames, and as such I-frames will cause peak-traffic while P-frames and B-frames will only cause minimal amounts of traffic [9].

The compression performance of H.264 is significantly greater than that of its predecessors. H.264 can deliver significantly better performance than that of MPEG-2 and MPEG-4, however at the price of greater computational power. As computational power has increased in user devices, H.264 has seen adaptation throughout video technology from media streaming to media storage.

III. LONG TERM EVOLUTION

A. LTE Operation Overview

Long Term Evolution (LTE) is a communications standard for wireless designed for cellular networks to simplify the network architecture than of the 3G (and previous) protocols by being based on an all-IP packet switched network [3, pg.470]. This allowed LTE to have significantly higher bit-rates, whilst still maintaining simultaneous access for multiple users.

In LTE the core network and air interface was separated into two distinct entities: eNodeB and LTE EUTRAN respectively. This modularisation allows for cost-effective improvement and advancements to be implemented in either of these systems without having the whole network having to be reconfigured. This simple network structure results in simpler operation, and thus greater performance [3, pg.473].

B. LTE Air Interface

LTE uses Orthogonal Frequency Division Multiple Access (OFDMA) in the down-link channels [3, pg.473], and Frequency Division Multiple Access (FDMA) in the up-link. These modulation technologies were chosen for high bit-rates, lower latency, and lower power consumption (particularly in the uplink as this transmission would be powered by the mobile device which often use battery power). With 20Mhz wide channels, the performance for LTE was 100Mbps in the down-link, and 50Mbps in the uplink.

The LTE network architecture was greatly simplified from the 3G technologies, so all QoS traffic was now packet-switched (including voice and multimedia), unlike the older 3G technologies whose infrastructure was comprised of both packet and circuit-switched services.

The implementation of all LTE transmissions being apart of a QoS system allows for prioritisation of different transmission packets, which greatly improves throughput for high resource real time applications, such as live-video and live-audio.

LTE also supports multiple antenna transmission (MIMO)

which improves spectral efficiency, bandwidth, and lowers interference.

C. LTE Advanced: True 4G

LTE Advanced was an upgrade to LTE that met the standards to be a true 4G technology. These advancements allows 1 Gbps peak rates in the down link, and 500Mbps peak rates in the uplink [3, pg.502]. This increase in data rates was achieved through many different advancements in various technological aspects of the system, namely: increasing the bandwidths of the channels to a maximum of 100Mhz via carrier aggregation, MIMO enhancements, introduction of smaller mobile-cell subdivisions and relays, and enhanced interference modelling [3].

IV. LTE OPERATION AND REAL TIME STREAMING

Video traffic comprises a significantly large portion of current mobile traffic which requires larger capacities on the network back-haul, but also requires reliable connectivity over the radio interface. LTE and LTE-Advanced both have QoS built into them allowing the priority of certain services over others. This is important for Real Time Streaming over LTE, as packet delay and losses over the radio link can cause user experience to decline when streaming media.

LTE also has issues with bandwidth consumption with multiple concurrent streaming clients. Each mobile device has its own assigned channel (a piece of bandwidth allocated to this device for data transmission) and due to the intensive data needed and QoS requirements for video streaming, the LTE transmitter and backhaul network can quickly reach maximum capacity. This will cause any new customers trying to connect to the network to either fail, or force current customers to share some of their allocated bandwidth which can effect the video streaming quality, and reliability.

If multiple clients are streaming the exact same content, (say a live sport game) then a lot of these transmissions are going to be replicas of each other and thus redundant. Instead of sending multiple transmissions over separate channels of the same data, and given broadcast nature of cellular networks, a more efficient way of transmitting would be to have all the mobile devices subscribe to the same channel to receive the media data. This technique is known as cellular-broadcasting, was introduced by the LTE group in its Evolved Multimedia Broadcast/Multicast Services (eMBMS) specification, also commonly known as LTE-Broadcast.

V. EVOLVED MULTIMEDIA BROADCAST MULTICAST SERVICE

eMBMS is the LTE implementation of the of Multimedia Broadcast/Multicast Service (MBMS) which brings the advantages of LTE to MBMS allowing HD broadcast service by taking advantage of the higher bandwidth channels that LTE offers.

eMBMS is a “point-to-multi-point” service where data is transmitted from a single-source transmitter to multiple receivers [4] allowing for multiple devices to receive the same broadcast

transmission within the same cell. eMBMS reduces cellular network bandwidth by eliminating the redundancy caused by having multiple devices streaming the same data over many uni-cast links. Thus for live-streaming services (such as live-event broadcasting or television over LTE) this can radically free the cellular network bandwidth, and reduces the load on the cellular back-haul

Up to 60 percent of an LTE frequency carrier can be used for MBMS transmissions, and using 20MHz LTE eMBMS allows for approximately 5 HD streams to be broadcast, parallel with normal unicast services [4].

Currently eMBMS needs unique chip set on the device so backwards compatibility is not possible with older devices. This is currently a limiting factor of adoption for using eMBMS services as software decoding is CPU intensive and given most mobile devices are battery powered, a software fix is not viable [4]. However, the latest mobile devices do support eMBMS, so as consumers upgrade their devices the rollout of eMBMS services will be able to increase.

A. Efficiency with eMBMS

Efficiency and resource savings for eMBMS is not as simple as a linear-relationship (ie two devices on a broadcast channel does not save half the amount of resources). For a unicast transmission the LTE network is able to adapt the signal modulation and coding rate to suit the current reception conditions, which is not possible to do with broadcast transmission. The broadcast transmitter must use a robust modulation and encoding schemes that are able to reach devices located on the edge of the cellular coverage area. Devices on the edge of the cell will suffer from larger signal losses due to distance, but will also receive interference from neighbouring cell transmissions. Thus, for adequate broadcast transmission a larger amount of resources are required compared to unicast transmission; which equates to many devices uni-casting the same stream before it is more efficient for the network to switch over to multi-cast [4].

B. Single Frequency Network

However, if multiple neighbouring cells transmit the same data fully synchronised so that the neighbouring radio signals are interfering constructively, the received signal (particularly at the cell edges) will increase.

This transmission technique is called “signal frequency network” and can greatly increase the total cellular network efficiency. This is particularly effective for popular live events, such as live sport finals (ie AFL grand final) where many user devices located in different neighbouring cells are casting the same media.

The neighbouring cell radio transmitters purposefully select the same frequency blocks to transmit the media on, such that the signal is indistinguishable from each other. As these frequency blocks are relatively static, the rest of the blocks are able to be used for regular, non-SFN transmissions (including other LTE broadcasts, and regular unicast traffic). For LTE base-stations separated by 500m, the efficiency of MBMS with

SFN transmission was evaluated to be 3bytes/Hz as compared with 0.25 bytes/Hz for MBMS over UMTS (a 3G technology) [4].

C. Stream Authentication with eMBMS

A large issue with eMBMS is ensuring that only paying customers are able to access the eMBMS service. Due to the broadcast nature cellular networks, the eMBMS signal must be adequately encrypted so unauthorised devices cannot listen in the live stream. Secondly, as the encryption keys need to be distributed among a large number of UE devices and need to be changed at regular intervals to prevent unauthorised distribution of them to free loading devices who wish to access the media illegally.

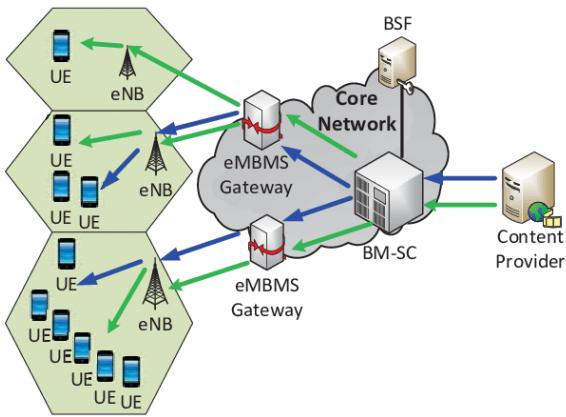


Fig. 4. A simplified example of eMBMS architecture [6]

The Key Management Mechanism, which is used by the multicast network provide forward and backward security, consists of the following parts: a Bootstrapping Server Function (BSF), Broadcast Multi-cast Service Centre (BM-SC), a content provider, an eMBMS gateway, and mobile user devices (UEs). A simplified diagram of this system can be seen in Figure 4.

The BSF is a part of the the Generic Bootstrapping Architecture that is responsible for establishing shared secrets between the mobile devices and the BM-SC. The BM-SC is the entry-point for the content providers and forwards the broadcasts packets to the eMBMS Gateway, where the packets are then further distributed to the eNodeB. The eNodeB then broadcasts the packets to the mobile user equipment (UE in the diagram). To protect the eMBMS data, four security keys are used which are: the MBMS Reqeust Key (MRK), the MBMS Service Key (MSK), the MBMS Traffic Key (MTK), and the MBMS User Key (MUK) [5].

THe MRK is used for authentication a mobile device to a BM-SC when performing the key requests. The MUK is used for securing the distribution the MSK, and the MSK is used for securing eMBMS sessions and the distribution of the MTK. The MTK is used for encrypting and decrypting of the eMBMS packets [5]. Thus, the MTK, MUK, and MSK

are used for protecting data, while the MRK is used for authenticating mobile devices. During the eMBMS service, the MSK and MTK are updated (or rekeyed) to protect the stream [5].

A simplified User Service Join Procedure is as follows, where $UE_{\{1..n\}}$ denote the mobile devices that have already joined [5]:

- 1) The new UE user device, UE_{n+1} , sends a service join request to the BM-SC. BM-SC then will ask that the joining device UE_{n+1} to begin the bootstrap authentication procedure with the BSF [5].
- 2) Device UE_{n+1} performs this bootstrap procedure with the BSF to obtain a MRK_{n+1} . This MRK key is then used by the mobile device to derive the MUK_{n+1} key.
- 3) After deriving MRK_{n+1} and MUK_{n+1} , UE_{n+1} performs authentication with the BM-SC using MRK_{n+1}
- 4) The BM-SC generates a new MSK, MSK_{new} which unicasts this to all existing mobile devices UE_k already on the service, encrypted with the UE_k 's MUK_k key.
- 5) Finally, the BM-SC generates a new MTK, which is encrypted with MSK_{new} and broadcast out to all user equipment devices that have joined the MBMS service.

This solution outlined above for rekeying every-time a device joins or leaves the service group works well for a small number of devices, but degrades quickly if there is a large number of devices in the usergroup (live sporting events may attract thousands of user devices for example). Performance quickly degrades as the devices continuously joining and dropping at seemly random intervals will cause frequent rekeying. This frequent rekeying will cause a large amount of re-authentication traffic, given such a large user-group of user devices can cause swap the Core Network causing it to under-perform.

Also, if the mobile devices were not able to receive the new keys from the broadcast for some reason (say, from momentary small signal fade, or interference) than this device will have to reauthenticate to the user group from the beginning. Frequent rekeying means that the probability that a device will miss the re-keying is increased, and thus can cause an endless loop of rekeying transmissions [5].

Alternatively having a longer fixed time-frame before rekeying will cause the keys to be renewed less often. While this will prevent the Core Network from being overwhelmed with reauthentication traffic and prevent rekeying loops, the likelihood of the multicast being compromised and having freeloading devices access the stream for free is higher, which results in lost revenue for the content provider. This balance between the stream security and performance of the networks is known as the security-performance trade-off, and is an important metric used by the service providers to adjust the network security protocols to find an acceptable threshold between the two criteria.

A solution the security-performance is the Dynamic Rekeying Algorithm (DRA). This algorithm is performed by the network operators to determine when rekeying in necessary. DRA is

used designed to update security keys based on a dynamic time slot, rather than when a mobile user device joins or leaves, or after a set period of time. This algorithm allows the network operators to define a performance metric between acceptable revenue loss of the content provider, and network performance of the cellular network , and find the optimal interval time that performs acceptably on both of these metrics.

An example algorithm for determining the interval time is given as:

$$\phi = \sum_{k=N(t_i-1)+1}^{N(t_i)} \quad (1)$$

where: ϕ is the accumulated free enjoy time interval, $N(t_i)$ is the number of user devices leaving the multi-cast group between $t = 0$ and $t = t_i$,

and s_k is the time when one user leaves the user-group. Thus, the time interval between $[s_k, t_i]$ the user is still able to freely accesses the multi-cast content because it still hold the valid encryption keys.

When the accumulated free-enjoy time exceeds a defined limit, then the keys are updated and rekeyed. Operators are thus able to easily adjust this limit as circumstances change to ensure the optimal network performance and revenue lost metric is achieved.

VI. MOBILE-TO-MOBILE COOPERATION

Another interesting technology that could improve QoS of live streaming is Mobile-to-Mobile cooperation. As video-streaming requires particularly high-data rates to maintain a high QoS, a mobile device connected via LTE has large energy drain on the device battery as the LTE radio system is not as optimised for energy efficient than that of a low power transmission protocols such as Bluetooth [12].

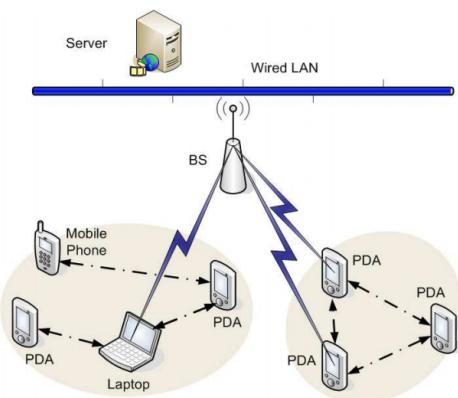


Fig. 5. An example of a clustered network. Only once device connects via the long range link, and the payload is distributed to the other devices via a short-range protocol [14]

As can be seen in Figure 5, the principle of Mobile-to-mobile cooperation is as follows: Mobile devices wanting to

stream the data from the Long Range LTE link base station that are also geographically near each other create an ad-hoc wireless network cluster via a short range energy efficient protocol, such as Bluetooth. Then in each cluster a single device becomes the cluster head, and receives the data packets from Long Range link, and then further broadcasts them to the rest of the UT via the Short Range. Only a single device at a time needs to connect to the LR at anytime, which can greatly reduce LR link bandwidth requirements by reducing the amount of Unicast connections. The average power consumption across all devices in cluster is significantly lower as only one device is using the high power network protocols where the other devices use the low-power devices to receive the relayed data. [12] [14].

VII. CONCLUSION

The future for broadcasting live video over cellular networks is currently being implemented worldwide. The rollout of eMBMS enabled user devices, and user demands for streaming live-video will see subsequent services provided by cellular network providers who aim to efficiently optimise network bandwidth and connectivity. Coupled with the improvement in codec as mobile device computational power has increased; it is near certainty that customers will be able to live-stream high quality, high definition video media on their mobile devices.

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12 Practical - Lab 5: Advanced Apache (D)

Lab task - Advanced Apache (Distinction)

Outcome	Weight
ULO4	♦♦♦♦♦

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 5: Advanced Apache (D)

Submitted By:

Michael REEB

4936094

2019/04/08 15:28

Tutor:

Jonathan KUA

April 8, 2019



```

Marking P-Lab-05-Advanced-Apache-D on RULE host rule61
TNE30019/TNE80014 - Laboratory - Advanced Apache
=====
*****
Portfolio Task: P-Lab-04-Advanced-Apache, Pass Task

Configure your rule host (rule61) with Apache to serve a virtual hosted web site

Tasks:
Set rule61 as DNS Server
Query for: aristocrats.unix ; Correct answer: 136.186.230.61
Apache running on rule61
Browsing to http://aristocrats.unix will return a correct web page
*****

```

Date: 20190408_1527

Last 5 Logons:

student	pts/18	atc330-21.ict.swin.edu	Mon Apr 8 14:56	still logged in
student	pts/17	atc330-21.ict.swin.edu	Mon Apr 8 14:56	still logged in
student	pts/1	vpn247-242.cc.swin.edu	Sun Apr 7 09:49 - 12:59	(03:09)
student	pts/0	vpn247-242.cc.swin.edu	Sun Apr 7 08:52 - 13:01	(04:09)
student	pts/1	vpn247-228.cc.swin.edu	Sat Apr 6 23:17 - 02:27	(03:09)

RULE host Socket information

www	httpd	52418	3	tcp4	136.186.230.61:80	*:*
www	httpd	52417	3	tcp4	136.186.230.61:80	*:*
www	httpd	52416	3	tcp4	136.186.230.61:80	*:*
www	httpd	52415	3	tcp4	136.186.230.61:80	*:*
www	httpd	52414	3	tcp4	136.186.230.61:80	*:*
root	httpd	52413	3	tcp4	136.186.230.61:80	*:*
tempuser	sshd	43937	3	tcp4	136.186.230.61:22	136.186.15.217:55890
root	sshd	43930	3	tcp4	136.186.230.61:22	136.186.15.217:55890
tempuser	sshd	43898	3	tcp4	136.186.230.61:22	136.186.15.217:55876
root	sshd	43891	3	tcp4	136.186.230.61:22	136.186.15.217:55876
bind	named	41933	20	tcp4	136.186.230.61:53	*:*
bind	named	41933	21	tcp4	136.186.230.61:953	*:*
bind	named	41933	512	udp4	136.186.230.61:53	*:*
bind	named	41933	513	udp4	136.186.230.61:53	*:*
bind	named	41933	514	udp4	136.186.230.61:53	*:*
bind	named	41933	515	udp4	136.186.230.61:53	*:*
bind	named	41933	516	udp4	136.186.230.61:53	*:*
bind	named	41933	517	udp4	136.186.230.61:53	*:*
bind	named	41933	518	udp4	136.186.230.61:53	*:*
bind	named	41933	519	udp4	136.186.230.61:53	*:*
bind	named	41933	520	udp4	136.186.230.61:53	*:*
bind	named	41933	521	udp4	136.186.230.61:53	*:*
bind	named	41933	522	udp4	136.186.230.61:53	*:*
bind	named	41933	523	udp4	136.186.230.61:53	*:*
bind	named	41933	524	udp4	136.186.230.61:53	*:*
bind	named	41933	525	udp4	136.186.230.61:53	*:*
bind	named	41933	526	udp4	136.186.230.61:53	*:*
bind	named	41933	527	udp4	136.186.230.61:53	*:*
bind	named	41933	528	udp4	136.186.230.61:53	*:*
bind	named	41933	529	udp4	136.186.230.61:53	*:*
bind	named	41933	530	udp4	136.186.230.61:53	*:*
bind	named	41933	531	udp4	136.186.230.61:53	*:*
bind	named	41933	532	udp4	136.186.230.61:53	*:*
bind	named	41933	533	udp4	136.186.230.61:53	*:*
bind	named	41933	534	udp4	136.186.230.61:53	*:*
root	sshd	97762	3	tcp4	136.186.230.61:22	*:*

```
root      ntpd      1207  61  udp4   136.186.230.61:123    *:*
?        ?       ?    ?  tcp4   136.186.230.61:46601  136.186.230.61:80

CHECKING DNS CONFIGURATION
=====
Querying for (aristocrats.unix) using DNS server (136.186.230.61) - Correct answer (136.186.230.61)
Using domain server:
Name: 136.186.230.61
Address: 136.186.230.61#53
Aliases:

aristocrats.unix has address 136.186.230.61
Test Result: PASSED
```

```
ADDING DNS SERVER 136.186.230.61 TO SYSTEM RESOLV.CONF
=====
HTTP QUERY http://aristocrats.unix
=====
--2019-04-08 15:27:43--  http://aristocrats.unix/
Resolving aristocrats.unix (aristocrats.unix)... 136.186.230.61
Connecting to aristocrats.unix (aristocrats.unix)|136.186.230.61|:80... connected.
HTTP request sent, awaiting response... 200 OK
Length: unspecified [text/html]
Saving to: '/tmp/wget_save.9kNVIFCG'

OK                                30.8K=0.02s
```

```
2019-04-08 15:27:43 (30.8 KB/s) - '/tmp/wget_save.9kNVIFCG' saved [494]

=====
Downloaded Page contents
=====
<html><head><title>Aristocrats Family Theater Restaurant</title></head><body>
<center><h1>Hello!</h1></center><p>
<center><h2>The Aristocrats Family Theater Restaurant site is under construction...</h2></center></p>
<center><h3>Check back soon for updates.</h3></center></p>
<center><h3>Remember: It is only <i>-394079263</i> seconds until our first location opens!</h3><center>
<center><h3>Right now it is Mon Apr 8 15:27:43 EST 2019 and we open on Thu Oct 12 13:00:00 EST 2006
</body></html>
```

```
=====
Checking Downloaded page contents

Test Result: PASS(download) + PASS(contents)
```

```
REMOVING DNS SERVER 136.186.230.61 FROM SYSTEM RESOLV.CONF
=====
```

```
*****
PORTFOLIO TASK RESULT: PASSED
```

13 Practical - Lab 5: Advanced Apache (P)

Lab task - Advanced Apache

Outcome	Weight
ULO4	◆◆◆◆◆

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 5: Advanced Apache (P)

Submitted By:

Michael REEB

4936094

2019/04/07 09:51

Tutor:

Jonathan KUA

April 7, 2019



```

Marking P-Lab-05-Advanced-Apache-P on RULE host rule61
TNE30019/TNE80014 - Laboratory - Advanced Apache
=====
*****
Portfolio Task: P-Lab-04-Advanced-Apache, Pass Task

```

Configure your rule host (rule61) with Apache to serve a virtual hosted web site

Tasks:

```

Set rule61 as DNS Server
Query for: armitagechemicals.unix      ; Correct answer: 136.186.230.61
Apache running on rule61
Browsing to http://armitagechemicals.unix will return a correct web page
*****
```

Date: 20190407_0951

Last 5 Logons:

student	pts/1	vpn247-242.cc.swin.edu	Sun Apr 7 09:49	still logged in
student	pts/0	vpn247-242.cc.swin.edu	Sun Apr 7 08:52	still logged in
student	pts/1	vpn247-228.cc.swin.edu	Sat Apr 6 23:17 - 02:27	(03:09)
student	pts/2	vpn247-228.cc.swin.edu	Sat Apr 6 22:02 - 23:15	(01:13)
student	pts/7	vpn247-49.cc.swin.edu.	Mon Apr 1 14:25 - 17:35	(03:09)

RULE host Socket information

www	httpd	47022	3	tcp4	136.186.230.61:80	*:*
www	httpd	47021	3	tcp4	136.186.230.61:80	*:*
www	httpd	47020	3	tcp4	136.186.230.61:80	*:*
www	httpd	47019	3	tcp4	136.186.230.61:80	*:*
www	httpd	47018	3	tcp4	136.186.230.61:80	*:*
root	httpd	47017	3	tcp4	136.186.230.61:80	*:*
tempuser	sshd	46525	3	tcp4	136.186.230.61:22	136.186.247.242:59761
tempuser	sshd	46525	7	tcp4	136.186.230.61:6011	*:*
root	sshd	46522	3	tcp4	136.186.230.61:22	136.186.247.242:59761
bind	named	41933	20	tcp4	136.186.230.61:53	*:*
bind	named	41933	21	tcp4	136.186.230.61:953	*:*
bind	named	41933	512	udp4	136.186.230.61:53	*:*
bind	named	41933	513	udp4	136.186.230.61:53	*:*
bind	named	41933	514	udp4	136.186.230.61:53	*:*
bind	named	41933	515	udp4	136.186.230.61:53	*:*
bind	named	41933	516	udp4	136.186.230.61:53	*:*
bind	named	41933	517	udp4	136.186.230.61:53	*:*
bind	named	41933	518	udp4	136.186.230.61:53	*:*
bind	named	41933	519	udp4	136.186.230.61:53	*:*
bind	named	41933	520	udp4	136.186.230.61:53	*:*
bind	named	41933	521	udp4	136.186.230.61:53	*:*
bind	named	41933	522	udp4	136.186.230.61:53	*:*
bind	named	41933	523	udp4	136.186.230.61:53	*:*
bind	named	41933	524	udp4	136.186.230.61:53	*:*
bind	named	41933	525	udp4	136.186.230.61:53	*:*
bind	named	41933	526	udp4	136.186.230.61:53	*:*
bind	named	41933	527	udp4	136.186.230.61:53	*:*
bind	named	41933	528	udp4	136.186.230.61:53	*:*
bind	named	41933	529	udp4	136.186.230.61:53	*:*
bind	named	41933	530	udp4	136.186.230.61:53	*:*
bind	named	41933	531	udp4	136.186.230.61:53	*:*
bind	named	41933	532	udp4	136.186.230.61:53	*:*
bind	named	41933	533	udp4	136.186.230.61:53	*:*
bind	named	41933	534	udp4	136.186.230.61:53	*:*
tempuser	sshd	36568	3	tcp4	136.186.230.61:22	136.186.247.242:56913
tempuser	sshd	36568	7	tcp4	136.186.230.61:6010	*:*

```
root      sshd      36565  3  tcp4    136.186.230.61:22      136.186.247.242:56913
tempuser sshd      34472  3  tcp4    136.186.230.61:22      136.186.247.242:56009
root      sshd      34469  3  tcp4    136.186.230.61:22      136.186.247.242:56009
root      sshd      97762  3  tcp4    136.186.230.61:22      *:*
root      ntpd      1207   61  udp4    136.186.230.61:123     *:*
?        ?        ?      ?  tcp4    136.186.230.61:48605  136.186.230.61:80

CHECKING DNS CONFIGURATION
=====
Querying for (armitagechemicals.unix) using DNS server (136.186.230.61) - Correct answer (136.186.230.61)
Using domain server:
Name: 136.186.230.61
Address: 136.186.230.61#53
Aliases:

armitagechemicals.unix has address 136.186.230.61
Test Result: PASSED
```

```
ADDING DNS SERVER 136.186.230.61 TO SYSTEM RESOLV.CONF
=====
HTTP QUERY http://armitagechemicals.unix
=====
--2019-04-07 09:51:07--  http://armitagechemicals.unix/
Resolving armitagechemicals.unix (armitagechemicals.unix)... 136.186.230.61
Connecting to armitagechemicals.unix (armitagechemicals.unix)|136.186.230.61|:80... connected.
HTTP request sent, awaiting response... 200 OK
Length: 303 [text/html]
Saving to: '/tmp/wget_save.19SndrMV'
```

```
OK                                         100% 34.1M=0s

2019-04-07 09:51:07 (34.1 MB/s) - '/tmp/wget_save.19SndrMV' saved [303/303]
```

```
=====
Downloaded Page contents
=====
<html>
<head><title>Welcome to Armitage Chemicals</title></head>
<body>
<center>
<h1>Welcome to Armitage Chemicals</h1>
<img src=armitagechemicals.jpg>
<p>Armitage Chemicals Headquarters<p>
To contact us, please deliver all letters/emails to the guard at the front gate.<br>
</center>
```

```
</body>
</html>
=====
```

```
Checking Downloaded page contents
```

```
Test Result: PASS(download) + PASS(contents)
```

REMOVING DNS SERVER 136.186.230.61 FROM SYSTEM RESOLV.CONF

=====

PORTFOLIO TASK RESULT: PASSED

14 Communications - Lab Report 2

New Description

Outcome	Weight
ULO7	♦♦♦♦◊

Outcome	Weight
ULO5	♦♦♦♦◊

Date	Author	Comment
2019/04/30 15:23	Jonathan Kua	Excellent work! Good to see you have included your program codes in the report, not the easiest thing with LaTeX!

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Communications - Lab Report 2

Submitted By:

Michael REEB

4936094

2019/04/27 18:21

Tutor:

Jonathan KUA

April 27, 2019



TNE30019 Unix for Telecommunications

Lab 6 - PCAP Programming

Michael Reeb (ID: 4936094)
Swinburne University of Technology
Melbourne, Australia
4936094@student.swin.edu.au

Abstract—PCAP is a traffic analysis API that can be used to create programs that can capture network packets and perform analysis on them. dkpt is a Python wrapper of the Unix libpcap implementation of PCAP, and can easily be used to sort and manage the captured traffic packets, and output useful human-readable information.

I. INTRODUCTION TO PCAP

PCAP (Packet Capture) is an application programming interface (API) that is used for capturing network traffic. The PCAP API is implemented in Unix-like systems under the name libpcap, and on Windows under the name Npcap. These PCAP libraries are able to save captured packets to file, read the saved packets from file; and this data can also be analysed and manipulated. These API implementations are written in C [1], and the information in the pcap file is stored in their hexadecimal values.

In this lab, we used the Python module dkpt, which is a wrapper that sits on top of the C library [?].

II. AIM

Dkpt was used for analysis of given PCAP files which contained previously captured TCP and UDP traffic. Relevant IP packet information was pulled from the traffic dump, and printed to the console.

In Part 2 the packets were sorted into flows and information was regarding the flow behaviour was analysed. A flow is simply a unique combination of the following parameters: the Source IP and Source Socket, the Destination IP and Destination Socket, and the protocol type. All this information was then printed to screen in a user-readable manner.

III. EQUIPMENT

RULE host 61 was used as the server. SSH access was through PuTTY running on an external PC connected to Swinburne via VPN. Python2.7 was used as a scripting language to perform the required tasks with the PCAP

libraries. The dpkt python library was used for parsing the PCAP packets.

IV. METHODOLOGY

As per lab handout.

V. DISCUSSION

A. Part 1 : mtcldump

The purpose of mtcldump was to analyse a PCAP file, and for each file in the dump, print the packet information (Timestamp, Source IP and socket, Destination IP and socket, Protocol, Packet-length and TTL) to the console.

The !# was the first line in the file and pointed to the Python2.7 binaries so when executed this script would be able to locate the Python interpreter.

In the main method of the program, the PCAP file had to be opened using the argument string given when the script was executed. This was achieved using the `open()` Python command. This file was then passed to the `dkpt.pcap.Reader()`, which is parses the file extracting the buffered hexadecimal data. This `Reader()` also implements an iterator, which will be used to cycle through to get each the next packet of hexadecimal data.

As the script iterated through the dump, each element in the iterator was extracted as a Ethernet Frame, using `eth=dpkt.ethernet.Ethernet(buf)`. Then, the IP packet was extracted from the frame using `ip=eth.data`.

With the IP packet isolated, it was now possible to extract the information needed for printing to screen. The source IP was extracted with `ip.source` member, and was converted into the usual IPv4 format (A.B.C.D) using the socket module command: `src_ip=socket.inet_ntoa(ip.src)`. The same method was used to extract the destination port, however using `ip.dst` member to access the destination IPv4 address.

The source and destination ports were extracted using: `src_port=ip.data.sport` and `dst_port=ip.data.dport` respectively. The protocol type was extracted using the `ip.p` member which is an integer. This was compared to the `dpkt.ip.IPPROTO_UDP` which is just a enumerator of all the IP protocols. Since all our traffic was either UDP or TCP, a simple if/else statement was used to determine the correct IP protocol, and then assign a correctly formatted string to a local variable to use for printing to screen.

The ttl and packet length fields were accessed using `ip.ttl` and `ip.len`. Finally, at the end of the loop a running packet count would increment per packet, and the total number of bytes would accumulate by adding the `ip.len` to a running total. All this information would be printed to screen.

B. Part 2 : `tcpdump`

The aim of part 2 was to extract the flow information from the PCAP dump, and determine how many packets per flow there were, and the total number of bytes each flow had .

It was decided to contain all the relevant flow information within a class, called `tcpFlow`. This class would manage the flow identification, the timestamp of the beginning of each flow, the total number packets, and the total number of bytes. A constructor was would initialised these class members, importantly noting that the packet counter was initialised as 1, not 0, ensuring no off-by-one errors

In the main loop of the program, similarly to Part 1, the PCAP file was opened, read, and iterated through using the exact same method. However, instead of printing the packets individually, each packets flow was identified and if determined to be unique (ie a new flow), it would be added to a Python dictionary collection. The dictionary was chosen because it maps a key with a value.

In this case, for each packet, a key was generated with a method (`getFlowKeyNiceAndConcise(ip)`) which took an ip packet as an argument and would output a TCP flow identifier (source ip address, source port, destination ip address, destination port, and ip protocol concatenated together) as a string. This key was then checked to see if it already existed in the dictionary.

If this key was unique, then an instance of the `tcpFlow()` was constructed, and the key/instance pair was added in the dictionary. If the flow was seen before, then in that `tcpFlow` the packet count was incremented and the amount of bytes of the packet was added to the running

total (using the class functions `incrementPacket()`, and `incrementBytes(bytes)` which takes the `ip.len` as an argument).

This system would find each unique flow and add them to the dict structure, however as Dictionary structure would not have the data in order needed to be marked by the RULE marking system. To solve this, a quick workaround was achieved by taking each unique flow in the dictionary, mapping the `tcpFlow` instances into another dict with the timestamp (of the first flow packet) as a key, and then sorting them by this by the timestamp. This issue was to do with the choice of the dictionary collection, as their functionality was confused with the C++ std::map structure functionality.

A better system would be to have used either a python-tuple or another class as the key, which would of allowed for sorting the first dictionary; or instead of using a dictionary, an array or other list could have been used, mitigating the limitations of the dictionary structure all together.

C. Making the scripts executable from the directory

To make the scripts executable from the directory, the `chmod+xscriptName` command was used in root privileged. This would allow the script to be executable.

VI. CONCLUSION

PCAP is a useful API for analysing network traffic data. In this lab a Python implementation was used, and a simple PCAP traffic analyser was created. In the second part, this was expanded to be able to analyse traffic flows.

REFERENCES

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VII. RESULTS

Listing 1. `tcpsummary.py`

```
1 #!/usr/local/bin/python2.7
2 #version 1 by Michael Reeb
3 ######
4 # Sample Python program with printing functionaility for P-Lab-06-PCAP-Programming-P
5 ######
6
7 #####
8 # You will probably want all these libraries for your actual program
9 #####
10 import datetime
11 import time
12 import sys
13 import dpkt
14 import socket
15
16 #####
17 # print_flow_summary_header()
18 #
19 # Prints the first lines of output as required by the marking script
20 #####
21 def print_flow_summary_header():
22     print '\n'----- PER-FLOW INFO -----'\n',
23
24 #####
25 # print_flow_summary_footer()
26 #
27 # Prints the last line of output as required by the marking script
28 #####
29 def print_flow_summary_footer():
30     print'-----',
31
32 #####
33 # print_flow_summary_single(u_int isrc_ip, u_short src_port, u_int idst_ip, u_short dst_port,
34 #                           char *protocol, struct timeval *start_time, u_int flow_packets, u_int flow_
35 #
36 # isrc_ip      - 32-bit IP address
37 # src_port     - Source port number
38 # idst_ip      - 32-bit IP address
39 # dst_port     - Destination port number
40 # protocol     - String containing protocol (should be 'tcp' or 'udp')
41 # start_time   - Timestamp of first packet in the nominated flow
42 # flow_packets - Integer count of total packets in this flow
43 # flow_bytes   - Integer sum of all packet sizes in this flow
44 #
45 # Print the lines of data as required by the marking script for ALL packets in ONE flow from the PCAP file
46 # The function will ensure formatting and convert IP addresses to strings for display
47 #####
48 def print_flow_summary_single(src_ip, src_port, dst_ip, dst_port, protocol, \
49                             start_time, flow_packets, flow_bytes):
50
51     print 'Flow: %s:%s-%s:%s' % (src_ip, src_port, dst_ip, dst_port)
52     print 'Protocol: %s' % (protocol)
53     print 'Start time: %s' % (start_time)
54     print 'Total packets: %u\nTotal bytes: %u\n' % (flow_packets, flow_bytes)
55
56 #####
57 # run_example
58 #
```

```

59 # Main function of sample program to demonstrate use of provided functions
60 #
61 # 1) Parse command line parameter. Fail if PCAP filename not provided
62 # 2) Call print_flow_summary_header() to display initial output
63 # 3) Generate fake data for flow 1 to print as an example
64 # 4) Call print_flow_summary_single() to display fake data. You will need to this once for each flow in
65 # 5) Generate more fake data for flow 2 to print as an example
66 # 6) Call print_flow_summary_single() to display flow 2 fake data
67 # 7) Call print_flow_summary_footer() to display final output
68 ######
69
70 def getFlowKeyNiceAndConcise(ip):
71     src_ip = socket.inet_ntoa(ip.src)
72     src_port = ip.data.sport
73     dst_ip = socket.inet_ntoa(ip.dst)
74     dst_port = ip.data.dport
75     if ip.p == dpkt.ip.IP_PROTO_UDP:
76         protocol = "udp"
77     else:
78         protocol = "tcp"
79
80
81     temp = '%s:%s -> %s:%s (%s)' % \
82         (src_ip, src_port, dst_ip, dst_port, protocol)
83     return temp
84
85
86 class tcpFlow:
87 #
88     def __init__(self, ip_packet, time_stamp):
89         self.source_ip = socket.inet_ntoa(ip_packet.src)
90         self.src_port = ip_packet.data.sport
91         self.dst_ip = socket.inet_ntoa(ip_packet.dst)
92         self.dst_port = ip_packet.data.dport
93         self.time_stamp = time_stamp
94         self.total_packets = 1
95         self.total_bytes = ip_packet.len
96         if ip_packet.p == dpkt.ip.IP_PROTO_UDP:
97             self.protocol = "udp"
98         else:
99             self.protocol = "tcp"
100
101
102     # if ip_packet.p == dpkt.ip.IP_PROTO_UDP:
103     #     self.protocol = "udp"
104     # else:
105     #     self.protocol = "tcp"
106
107 #
108     def getTimeStamp(self):
109         return self.time_stamp
110
111
112     def incrementPacket(self):
113         self.total_packets +=1
114 #
115     def incrementBytes(self, bytes):
116         self.total_bytes += bytes
117         #print the summary of the flow
118     def printInstanceFlowSummary(self):
119         print_flow_summary_single(self.source_ip, self.src_port, self.dst_ip, self.dst_port, self.protocol,
120
121

```

```

122 ######
123
124 def run_example():
125
126     global total_packet_count, total_bytes_count
127
128     try:
129         sys.argv[1]
130         dmp_file = sys.argv[1]
131         fp_dmp_file = open(dmp_file)
132     except Exception as e:
133         print 'Error: please supply pcap filename!\n'
134         return
135
136     print_flow_summary_header()
137
138     # need to move this into a class file, then create a list of these new classes for each flow
139     pcap = dpkt.pcap.Reader(fp_dmp_file)
140
141     #declare tcp flows dict to manage flows. dicts use key:flow arrangement
142     tcp_flows = dict()
143
144     for ts, buf in pcap:
145         eth = dpkt.ethernet.Ethernet(buf)
146         ip = eth.data
147         #create key variable here to make below look nicer
148         key = getFlowKeyNiceAndConcise(ip)
149         # print '----- current flow key %s' %(key)
150         #check to see if flow already exists. if it does, increment the counters, if not construct new
151         if key in tcp_flows.keys():
152             tcp_flows[key].incrementPacket()
153             tcp_flows[key].incrementBytes(ip.len)
154         else: #add new instance of a TCP flows to list, have to pass constructor time and ip packet
155             timestamp = str(datetime.datetime.utcnow().timestamp(ts))
156             tcp_flows[key] = tcpFlow(ip,timestamp)
157             #tcp_flows.update(key:tcpFlow(ip,str(datetime.datetime.utcnow().timestamp(ts))))
158
159     #create a second list because python dict sucks and is nothing like cpp map
160     #so we have to find a stupid work around to sort the damn thing.
161     #so we create a new dictionary that associates the timestamp to the class.
162     sorted_flows = dict()
163     for keys in tcp_flows:
164         tempKey = tcp_flows[keys].getTimeStamp()
165         sorted_flows[tempKey]=tcp_flows[keys]
166
167     #now we sort that list by timestamp and print the summary to the screen so we have things in the
168     for k in sorted(sorted_flows):
169         sorted_flows[k].printInstanceFlowSummary()
170
171
172
173
174     # start_time = time.time()
175     # src_ip = '1.2.3.4'
176     # src_port = '99'
177     # dst_ip = '5.6.7.8'
178     # dst_port = '199'
179     # protocol = 'tcp'
180
181     # print_flow_summary_single(src_ip, src_port, dst_ip, dst_port, protocol, start_time, 100, 150000)
182
183     # start_time = time.time()
184     # src_ip = '10.0.0.1'

```

```

185     # src_port = '123'
186     # dst_ip = '10.1.1.200'
187     # dst_port = '456'
188     # protocol = 'udp'
189
190     # print_flow_summary_single(src_ip, src_port, dst_ip, dst_port, protocol, start_time, 8, 12000)
191
192     print_flow_summary_footer()
193
194     fp_dmp_file.close()
195
196 if __name__ == '__main__':
197     run_example()

```

Listing 2. mtcpcap.py

```

1#!/usr/local/bin/python2.7
2#version 1 by Michael Reeb
3#####
4# Sample Python program with printing functionality for P-Lab-06-PCAP-Programming-P
5#####
6#####
7#####
8# You will probably want all these libraries for your actual program
9#####
10import datetime
11import time
12import sys
13import dpkt
14import socket
15
16#####
17# print_packet_info(ts, src_ip, src_port, dst_ip, dst_port, protocol, pkt_len, ttl)
18#
19# ts      - Packet timestamp
20# isrc_ip - 32-bit IP address
21# src_port - Source port number
22# idst_ip - 32-bit IP address
23# dst_port - Destination port number
24# protocol - String containing protocol (should be 'tcp' or 'udp')
25# pkt_len - Packet length as integer
26# ip_ttl  - IP Packet TTL as byte value
27#
28# Print the line of data as required by the marking script for ONE packet in the PCAP file
29# The function will ensure formatting and convert IP addresses to strings for display
30#####
31def print_packet_info(ts, src_ip, src_port, dst_ip, dst_port, protocol, pkt_len, ttl):
32
33    # get UTC timestamp, source/destination IP/port
34    utc_timestamp = str(datetime.datetime.utcfromtimestamp(ts))
35
36    # create log line
37    log_line = '[%s] - %s:%s -> %s:%s (%s, len=%d, ttl=%d)' % \
38        (utc_timestamp, src_ip, src_port, dst_ip, dst_port, protocol, pkt_len, ttl)
39
40    print log_line
41
42#####
43# print_summary(total_packet_count, total_bytes_count, average_pkt_size)
44#
45# total_pkts - Integer count of total packets
46# total_bytes - Integer count of sum of packet sizes
47# avg_pkt    - Integer representation of average packet size

```

```

48 #
49 # Print the three parameters in the format as required by the marking script
50 #####
51 def print_summary(total_packet_count, total_bytes_count, average_pkt_size):
52
53     print '\n----- SUMMARY -----\\n'
54     print 'Total packets: %i' % (total_packet_count)
55     print 'Total bytes (bytes): %i' % (total_bytes_count)
56     print 'Average packet size (bytes): %i' % (average_pkt_size),
57     print '\\n-----'
58
59 #####
60 # run_example
61 #
62 # Main function of sample program to demonstrate use of provided functions
63 #
64 # 1) Parse command line parameter. Fail if PCAP filename not provided and file cannot be opened
65 # 2) Generate fake data to print as an example
66 # 3) Call print_packet_info() to display fake data. You will need to do this once for each packet in the
67 # 4) Call print_summary() to display a fake summary. You will need to do this with your accumulated data
68 #####
69 def run_example():
70
71     global total_packet_count, total_bytes_count
72     pkt_count = 0
73     byte_count = 0
74
75     try:
76         sys.argv[1]
77         dmp_file = sys.argv[1]
78         fp_dmp_file = open(dmp_file)
79     except Exception as e:
80         print 'Error: please supply pcap filename!\\n'
81     return
82
83 #For each packet in the dmp_file
84
85     pcap = dpkt.pcap.Reader(fp_dmp_file)
86
87     for ts, buf in pcap:
88         eth = dpkt.ethernet.Ethernet(buf)
89         ip = eth.data
90         src_ip = socket.inet_ntoa(ip.src)
91         src_port = ip.data.sport
92         dst_ip = socket.inet_ntoa(ip.dst)
93         dst_port = ip.data.dport
94         if ip.p == dpkt.ip.IP_PROTO_UDP:
95             protocol = "udp"
96         else:
97             protocol = "tcp"
98         ttl = ip.ttl
99         len = ip.len
100        print_packet_info(ts, src_ip, src_port, dst_ip, dst_port, protocol, len, ttl)
101        pkt_count += 1
102        byte_count += ip.len
103
104
105
106
107
108
109 #For each packet in the dmp_file
110

```

```
111
112
113
114     print_summary(pkt_count, byte_count, byte_count/ pkt_count)
115
116     fp_dmp_file.close()
117
118 if __name__ == '__main__':
119     run_example()
```

15 Practical - Lab 6: PCAP Programming (D)

Lab task - PCAP Programming (Distinction)

Outcome	Weight
ULO3	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 6: PCAP Programming (D)

Submitted By:

Michael REEB

4936094

2019/04/16 14:58

Tutor:

Jonathan KUA

April 16, 2019



```
Marking P-Lab-06-PCAP-Programming-D on RULE host rule61
TNE30019/TNE80014 - Laboratory - PCAP Programming
=====
*****
Portfolio Task: P-Lab-06-PCAP-Programming, Distinction Task
```

```
Develop a tcpsummary program to output a summary of all flows within a tcpdump
PCAP file
```

Tasks:

```
Create an executable as /home/student/lab06/tcpsummary
```

```
Executable takes one parameter - name of PCAP file
```

```
Output standard summary block of text for each unique flow in PCAP file
```

```
Date: 20190416_1455
```

Last 5 Logons:

```
student pts/8 vpn247-130.cc.swin.edu Tue Apr 16 12:05 still logged in
student pts/4 vpn247-202.cc.swin.edu Mon Apr 15 21:59 - 01:09 (03:09)
student pts/4 vpn247-202.cc.swin.edu Mon Apr 15 21:25 - 21:59 (00:34)
student pts/5 ses236-63.cc.swin.edu. Mon Apr 15 17:18 - 17:30 (00:12)
student pts/5 ses236-63.cc.swin.edu. Mon Apr 15 17:15 - 17:16 (00:00)
```

```
CHECKING EXECUTABLE
```

```
=====
```

```
Checking executable (/usr/home/student/pcap_lab/tcpsummary), permissions(rwx??????)...
```

```
-rwxr-xr-x 1 0 1001 6942 Apr 16 14:52 /usr/home/student/pcap_lab/tcpsummary
```

```
RUNNING TESTS
```

```
=====
```

```
=====
```

```
Running test: /usr/home/student/pcap_lab/tcpsummary test1.pcap
```

```
--- Program output -----
```

```
----- PER-FLOW INFO -----
```

```
Flow: 172.16.11.2:61016-172.16.10.2:80
```

```
Protocol: tcp
```

```
Start time: 2018-08-16 02:48:10.238506
```

```
Total packets: 6
```

```
Total bytes: 466
```

```
Flow: 172.16.10.2:80-172.16.11.2:61016
```

```
Protocol: tcp
```

```
Start time: 2018-08-16 02:48:10.238523
```

```
Total packets: 5
```

```
Total bytes: 986
```

```
Flow: 172.16.11.3:37269-172.16.10.2:12345
```

```
Protocol: udp
```

```
Start time: 2018-08-16 02:48:13.537505
```

```
Total packets: 5
```

```
Total bytes: 7490
```

```
Flow: 172.16.10.2:12345-172.16.11.3:37269
```

```
Protocol: udp
```

```
Start time: 2018-08-16 02:48:13.591438
```

```
Total packets: 1
```

```
Total bytes: 1498
```

```
-----  
-----
```

```
Test Result: PASSED
```

```
=====  
Running test: /usr/home/student/pcap_lab/tcpsummary test2.pcap
```

```
--- Program output -----
```

```
----- PER-FLOW INFO -----
```

```
Flow: 172.16.11.72:33711-172.16.10.62:5001  
Protocol: tcp  
Start time: 2018-08-01 03:08:34.575832  
Total packets: 156  
Total bytes: 234000
```

```
Flow: 172.16.10.62:5001-172.16.11.72:33711  
Protocol: tcp  
Start time: 2018-08-01 03:08:34.576840  
Total packets: 78  
Total bytes: 4056
```

```
-----  
-----
```

```
Test Result: PASSED
```

```
=====  
Running test: /usr/home/student/pcap_lab/tcpsummary test3.pcap
```

```
Program output: Hidden
```

```
Test Result: PASSED
```

```
*****  
PORTFOLIO TASK RESULT: PASSED
```

16 Practical - Lab 6: PCAP Programming (P)

Lab task - PCAP Programming

Outcome	Weight
ULO3	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 6: PCAP Programming (P)

Submitted By:

Michael REEB

4936094

2019/04/15 22:33

Tutor:

Jonathan KUA

April 15, 2019



```
Marking P-Lab-06-PCAP-Programming-P on RULE host rule61
TNE30019/TNE80014 - Laboratory - PCAP Programming
=====
*****
Portfolio Task: P-Lab-06-PCAP-Programming, Pass Task
```

```
Develop a mini-tcpdump program to output a summary line for each packet in a
tcpdump PCAP file
```

Tasks:

```
Create an executable as /home/student/lab06/mtcpdump
Executable takes one parameter - name of PCAP file
Output standard summary line for each packet in PCAP file
Output summary of total bytes and packets captured, and average packet size
*****
```

Date: 20190415_2227

Last 5 Logons:

```
student pts/4 vpn247-202.cc.swin.edu Mon Apr 15 21:59 still logged in
student pts/4 vpn247-202.cc.swin.edu Mon Apr 15 21:25 - 21:59 (00:34)
student pts/5 ses236-63.cc.swin.edu. Mon Apr 15 17:18 - 17:30 (00:12)
student pts/5 ses236-63.cc.swin.edu. Mon Apr 15 17:15 - 17:16 (00:00)
student pts/15 ses236-63.cc.swin.edu. Mon Apr 15 17:02 - 17:14 (00:11)
```

CHECKING EXECUTABLE

```
=====
Checking executable (/usr/home/student/pcap_lab/mtcpdump), permissions(rwx??????)...
-rwxr-xr-x 1 0 1001 3722 Apr 15 22:03 /usr/home/student/pcap_lab/mtcpdump
```

RUNNING TESTS

```
=====
Running test: /usr/home/student/pcap_lab/mtcpdump test1.pcap
```

```
--- Program output -----
[2018-08-16 02:48:10.238506] - 172.16.11.2:61016 -> 172.16.10.2:80 (tcp, len=60, ttl=63)
[2018-08-16 02:48:10.238523] - 172.16.10.2:80 -> 172.16.11.2:61016 (tcp, len=60, ttl=64)
[2018-08-16 02:48:10.288710] - 172.16.11.2:61016 -> 172.16.10.2:80 (tcp, len=52, ttl=63)
[2018-08-16 02:48:10.288738] - 172.16.11.2:61016 -> 172.16.10.2:80 (tcp, len=198, ttl=63)
[2018-08-16 02:48:10.288746] - 172.16.10.2:80 -> 172.16.11.2:61016 (tcp, len=52, ttl=64)
[2018-08-16 02:48:10.288860] - 172.16.10.2:80 -> 172.16.11.2:61016 (tcp, len=770, ttl=64)
[2018-08-16 02:48:10.288865] - 172.16.10.2:80 -> 172.16.11.2:61016 (tcp, len=52, ttl=64)
[2018-08-16 02:48:10.339317] - 172.16.11.2:61016 -> 172.16.10.2:80 (tcp, len=52, ttl=63)
[2018-08-16 02:48:10.339331] - 172.16.11.2:61016 -> 172.16.10.2:80 (tcp, len=52, ttl=63)
[2018-08-16 02:48:10.339563] - 172.16.11.2:61016 -> 172.16.10.2:80 (tcp, len=52, ttl=63)
[2018-08-16 02:48:10.339586] - 172.16.10.2:80 -> 172.16.11.2:61016 (tcp, len=52, ttl=64)
[2018-08-16 02:48:13.537505] - 172.16.11.3:37269 -> 172.16.10.2:12345 (udp, len=1498, ttl=63)
[2018-08-16 02:48:13.549001] - 172.16.11.3:37269 -> 172.16.10.2:12345 (udp, len=1498, ttl=63)
[2018-08-16 02:48:13.560213] - 172.16.11.3:37269 -> 172.16.10.2:12345 (udp, len=1498, ttl=63)
[2018-08-16 02:48:13.571413] - 172.16.11.3:37269 -> 172.16.10.2:12345 (udp, len=1498, ttl=63)
[2018-08-16 02:48:13.582614] - 172.16.11.3:37269 -> 172.16.10.2:12345 (udp, len=1498, ttl=63)
[2018-08-16 02:48:13.591438] - 172.16.10.2:12345 -> 172.16.11.3:37269 (udp, len=1498, ttl=64)
```

----- SUMMARY -----

```
Total packets: 17
Total bytes (bytes): 10440
Average packet size (bytes): 614
```

Test Result: PASSED

```
=====  
Running test: /usr/home/student/pcap_lab/mtcpdump test2.pcap
```



```
[2018-08-01 03:08:34.729868] - 172.16.11.72:33711 -> 172.16.10.62:5001 (tcp, len=1500, ttl=63)
[2018-08-01 03:08:34.730868] - 172.16.11.72:33711 -> 172.16.10.62:5001 (tcp, len=1500, ttl=63)
[2018-08-01 03:08:34.730884] - 172.16.10.62:5001 -> 172.16.11.72:33711 (tcp, len=52, ttl=64)
```

```
----- SUMMARY -----
```

```
Total packets: 234
Total bytes (bytes): 238056
Average packet size (bytes): 1017
```

```
-----  
-----  
Test Result: PASSED
```

```
=====  
Running test: /usr/home/student/pcap_lab/mtcpdump test3.pcap
```

```
Program output: Hidden
```

```
Test Result: PASSED
```

```
*****  
PORTFOLIO TASK RESULT: PASSED
```

17 Practical - Lab 7: NMap (C)

Lab task - Nmap (Credit)

Outcome	Weight
ULO3	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 7: NMap (C)

Submitted By:

Michael REEB

4936094

2019/04/12 16:02

Tutor:

Jonathan KUA

April 12, 2019





This room is fullfilled mit special
electronische equippment. Fingergrabbing
and pressing the cnoepkes from the
computers is allowed for die experts only!

So all the Lefthanders stay away and do
not disturben the brainstorming von here
working intelligencies.

Otherwise you will be out thrown and
kicked anderswhere! Also: please keep
still and only watchen astonished the
blinkenlights.

Secret Message: No spitzensparken for the new RULE system

18 Practical - Lab 7: NMap (P)

Lab task - Nmap

Outcome	Weight
ULO3	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 7: NMap (P)

Submitted By:

Michael REEB

4936094

2019/04/12 15:59

Tutor:

Jonathan KUA

April 12, 2019



```
student@rule61:~ % nmap -v -r -p- -sV rule21.caia.swin.edu.au
Starting Nmap 7.70 ( https://nmap.org ) at 2019-04-12 15:46 EST
NSE: Loaded 43 scripts for scanning.
Initiating Ping Scan at 15:46
Scanning rule21.caia.swin.edu.au (136.186.230.21) [2 ports]
Completed Ping Scan at 15:46, 0.00s elapsed (1 total hosts)
Initiating Parallel DNS resolution of 1 host. at 15:46
Completed Parallel DNS resolution of 1 host. at 15:46, 0.00s elapsed
Initiating Connect Scan at 15:46
Scanning rule21.caia.swin.edu.au (136.186.230.21) [65535 ports]
Discovered open port 7/tcp on 136.186.230.21
Discovered open port 13/tcp on 136.186.230.21
Discovered open port 21/tcp on 136.186.230.21
Discovered open port 22/tcp on 136.186.230.21
Discovered open port 23/tcp on 136.186.230.21
Discovered open port 79/tcp on 136.186.230.21
Discovered open port 80/tcp on 136.186.230.21
Discovered open port 110/tcp on 136.186.230.21
Discovered open port 143/tcp on 136.186.230.21
Increasing send delay for 136.186.230.21 from 0 to 5 due to
max_successful_tryno increase to 4
Connect Scan Timing: About 9.14% done; ETC: 15:52 (0:05:08 remaining)
Connect Scan Timing: About 18.29% done; ETC: 15:52 (0:04:33 remaining)
Connect Scan Timing: About 27.43% done; ETC: 15:52 (0:04:01 remaining)
Connect Scan Timing: About 36.58% done; ETC: 15:52 (0:03:30 remaining)
Connect Scan Timing: About 45.72% done; ETC: 15:52 (0:02:59 remaining)
Connect Scan Timing: About 54.87% done; ETC: 15:52 (0:02:29 remaining)
Connect Scan Timing: About 64.01% done; ETC: 15:52 (0:01:59 remaining)
Connect Scan Timing: About 73.16% done; ETC: 15:52 (0:01:28 remaining)
Connect Scan Timing: About 82.31% done; ETC: 15:52 (0:00:58 remaining)
Discovered open port 53959/tcp on 136.186.230.21
Completed Connect Scan at 15:52, 328.98s elapsed (65535 total ports)
Initiating Service scan at 15:52
Scanning 10 services on rule21.caia.swin.edu.au (136.186.230.21)
Completed Service scan at 15:52, 6.01s elapsed (10 services on 1 host)
NSE: Script scanning 136.186.230.21.
Initiating NSE at 15:52
Completed NSE at 15:52, 0.07s elapsed
Initiating NSE at 15:52
Completed NSE at 15:52, 0.00s elapsed
Nmap scan report for rule21.caia.swin.edu.au (136.186.230.21)
Host is up (0.000073s latency).
rDNS record for 136.186.230.21: passTest.unix
Not shown: 65525 closed ports
PORT      STATE SERVICE      VERSION
7/tcp      open  echo
13/tcp     open  daytime
21/tcp     open  ftp          WU-FTPD or MIT Kerberos ftptd 6.00LS
22/tcp     open  ssh          OpenSSH 7.2 (FreeBSD 20160310; protocol 2.0)
23/tcp     open  telnet       BSD-derived telnetd
79/tcp     open  finger       FreeBSD fingerd
```

```
80/tcp      open  http        thttpd 2.27 19Oct2015
110/tcp     open  tcpwrapped
143/tcp     open  tcpwrapped
53959/tcp   open  http        thttpd 2.27
Service Info: Host: rule21; OSs: Unix, FreeBSD; CPE: cpe:/o:freebsd:freebsd
```

```
Read data files from: /usr/local/share/nmap
Service detection performed. Please report any incorrect results at
https://nmap.org/submit/ .
Nmap done: 1 IP address (1 host up) scanned in 336.26 seconds
```

19 Practical - Programming Tutorial

Successful completion of programming task in tutorial

Outcome	Weight
ULO4	♦♦♦♦◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Programming Tutorial

Submitted By:

Michael REEB

4936094

2019/05/09 13:12

Tutor:

Jonathan KUA

May 9, 2019



```
1 // visit
2 → https://www.binarytides.com/receive-full-data-with-recv-socket-function-in-c/
3
4 #include <stdio.h>
5 #include <stdlib.h>
6
7 #define __USE_GNU
8 #include <string.h>
9
10 #include <sys/socket.h>
11 #include <netinet/in.h>
12 #include <arpa/inet.h>
13 #include <netdb.h>
14
15 #include <errno.h>
16 #include <unistd.h>
17 #define WWW_PORT      80
18
19 //-----
20 // Print an error message to stderr and terminate the program immediately
21 void printerror(char *pcErrStr) {
22     fprintf(stderr, "ERROR: %s\n\n", pcErrStr);
23     exit(1);
24 }
25
26 //-----
27 // Input parameters are the name of the server to query and a pointer to a
28 // sockaddr_in structure to fill. The function performs a lookup of the
29 // specified Internet name and stores the address along with the web port
30 // number (WWW_PORT) in the provided sockaddr_in structure
31 void FillAddress(char *pcURL, struct sockaddr_in *psAddress) {
32     struct hostent      *psHostDetails;
33
34     if ((psHostDetails = gethostbyname(pcURL)) == NULL) printerror("Can't determine
35         ← address of website");
36
37     memcpy(&(psAddress->sin_addr), psHostDetails->h_addr_list[0], 4);
38     psAddress->sin_family = AF_INET;
39     psAddress->sin_port = htons(WWW_PORT);
40
41     printf("Server details:\n");
42     printf("  Server:\t%s\n", pcURL);
43     printf("  IP Address:\t%s\n", inet_ntoa(psAddress->sin_addr));
44 }
45
46 //-----
47 int main(int iArgC, char **ppcArgV) {
48     struct sockaddr_in sAddress;
49     char              *pcWWWServer, *pcPage, *pcPos;
50
51     if (iArgC != 2) printerror("Command Line requires one parameter = website");
```

```

52 // If the provided URL begins with "http://" then set pcWWWServer to point
53 // to the next character otherwise set it to point to the entire provided URL
54 if (pcWWWServer = strstr(ppcArgV[1], "http://")) pcWWWServer+= 7; else
→ pcWWWServer = ppcArgV[1];
55
56 // If there is a "/" in the URL (after the http:// has been stripped) then
57 // replace it with a 0 so the string terminates as just the name. Create
58 // a new string pointer (pcPage) to be a copy of the remaining string OR
59 // just "/" if no page details have been provided
60 if (pcPos = strstr(pcWWWServer, "/"))
{
61     pcPage = strdup(pcPos);
62     pcPos[0] = 0;
63 } else
64     pcPage = strdup("/");
65
66 // Call FillAddress() to fill sAddress with the IP address details (and port
67 // number 80) of the requested URL
68 FillAddress(pcWWWServer, &sAddress);
69
70 // Put rest of code here ----- ]
→ -----
71 //create a socket. note socket returns an integer based upon its success
72 //
73 // int socket(int domain, int type, int protocol);
74 int mySock=socket(AF_INET,SOCK_STREAM,0);
75 if( mySock<0) {
76     printf("Error creating socket\n");
77     exit(1);
78 }
79 //try and connect socket to server
80 //
81 // int connect(int sockfd, const struct sockaddr *addr, socklen_t addrlen);
82 if(connect(mySock,(struct sockaddr*)&sAddress, sizeof(sAddress) ) != 0){
83     printf("Error in connecting socket\n");
84     exit(1);
85 }
86
87 // create a plain-text GET request for HTTP. should included other HTTP versions
88 //
89 // GET <pgName> HTTP/1.1\r\nHost: <hostName>\r\nConnection: close\r\n\r\n"
90 char httpRequest[1024];
91 sprintf(httpRequest, "GET %s HTTP/1.1\r\nHost: %s\r\nConnection: close\r\n\r\n",
→ pcPage, pcWWWServer);
92
93 //http://man7.org/linux/man-pages/man2/send.2.html
94 //send the httpMessage through the socket to the server
95 //ssize_t send(int socket, const void *buffer, size_t length, int flags);
96 //returns the amount of bytes sent.
97 //returns -1 if there is an error
98 ssize_t hasSent = send(mySock,httpRequest,sizeof(httpRequest),0);
99 //check if it has sent without issues. exit on error
100 if(hasSent < 0){
    printf("Error in socket sending data\n");

```

```
102     exit(1);
103 }
104
105
106 //http://man7.org/linux/man-pages/man2/recv.2.html
107 //listen to socket for response.
108 //recv function is used to do this.
109 //
110 // ssize_t recv(int sockfd, void *buf, size_t len, int flags);
111 // until a message is ready to receive, recv returns -1
112 //also need a buffer
113 char recvBuffer[16000];
114 //open a file with fopen() for writing the scrape to
115 FILE* outputTxtFile;
116 if ((outputTxtFile = fopen("scrape.txt","w")) == NULL){
117     printf("Error cannot open file\n");
118     exit(1);
119 }
120
121 while(recv(mySock,recvBuffer, sizeof(recvBuffer),0) > 0 ){
122     //printf("%s\n",recvBuffer);
123     //print buffer to file. add null chars on end and flush
124     // fprintf(outputTxtFile,recvBuffer);
125     fputs(recvBuffer,outputTxtFile);
126     fflush(outputTxtFile);
127 }
128 //fprintf(outputTxtFile,recvBuffer);
129 //printf("%s\n",recvBuffer);
130
131
132 //http://man7.org/linux/man-pages/man3/fopen.3.html
133 // /*fopen(const char *pathname, const char *mode);
134 //mode "w" Truncate file to zero length or create text file for writing.
135 //           The stream is positioned at the beginning of the file.
136 //outputTxtFile
137
138
139 fclose(outputTxtFile);
140 printf("%s\n",recvBuffer);
141 printf("-----Scrape Completed-----\n");
142 close(mySock);
143     return 0;
144 }
```

20 Theory - Jails

Brief report summarising jail operations

Outcome	Weight	
ULO5	◆◆◆◆◆	
Date	Author	Comment
2019/04/30 15:04	Jonathan Kua	Excellent work!

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Theory - Jails

Submitted By:

Michael REEB

4936094

2019/04/27 21:59

Tutor:

Jonathan KUA

April 27, 2019



A Brief Overview of Operating System Virtualisation

Michael Reeb - 4936094
TNE30019 Unix for Telecommunications
Communications Report
4936094@student.swin.edu.au

Abstract—Hardware virtualisation allows a computer to logically divide its resources to operate like separate machines. These separate machines can run differing operating systems each with its own kernel, in the case of a virtual machine, or can be a compartmentalised space that shares its kernel with the other instances but processes are limited to that instances workspace, such is the case with FreeBSD Jails.

I. INTRODUCTION

Hardware virtualisation is the process of logically dividing its resources to allow multiple Virtual Machines (VM) to run simultaneously. These VMs perform like a regular physical machine, and ideally to the Operating System (OS) and/or applications running on them, VMs are indistinguishable from a physical machine.

II. VIRTUALISATION

The simplest form of Virtualisation is emulation: the host computer simulates all the underlying physical hardware (memory, CPU, CPU registers, etc) in the emulator. The Guest Operating system; that is the system that is to be emulated, then runs inside this emulator and behaves as it would if it was running on physical hardware, without any support from the host Operating System.

While hardware emulation works well for modern computers emulating old (and less powerful) hardware; it is not feasible for a computer to emulate an equivalent system.

A Virtual Machine Monitor (VMM), or Hypervisor allows the Guest OS to run directly on the CPU, and when the Guest OS attempts to execute privileged instructions (which it does not have permission to perform as it is operating in user mode) the Hypervisor traps them, and then steps in and emulates the instruction on the virtual machine. Thus, a hypervisor simply monitors the guest OS and emulates the Guest OS privileged instruction calls, such as memory allocation, I/O operation, context switching etc.

A Virtual Machine Monitor, or Hypervisor creates the illusion that there are multiple machines running on the same physical hardware. It is essential that a Hypervisor be able to score well in following three dimensions to ensure that the Virtual Machine behaves identically to a real physical machine:

- 1) **Safety**: the Hypervisor should be in full control of the visualised resources.

- 2) **Fidelity**: the behaviour of a program executing in a VM should be identical to that of one executing on a physical machine.
- 3) **Efficiency** As much as possible of the program code should run without the Hypervisor intervening.

Thus, for a machine to achieve virtualisation, the sensitive instructions must be a subset of the privileged instructions.

There is no clear distinction between the categories of Hypervisors, and certain Hypervisor implementations can fall into multiple categories. However, in this paper we will draw two-distinct categories for hypervisors: type 1, where the virtual machine runs directly on a Hypervisor incorporated with the kernel, and Type-2 hypervisors where the virtual machine runs as an independent application on the host OS like any other application.

A. Type 1 Hypervisor

A Type-1 Hypervisor is similar to an operating system as it is the only process that is running in the most privileged mode and has direct access to the systems hardware. This is colloquially known as a "bare-metal" hypervisor.

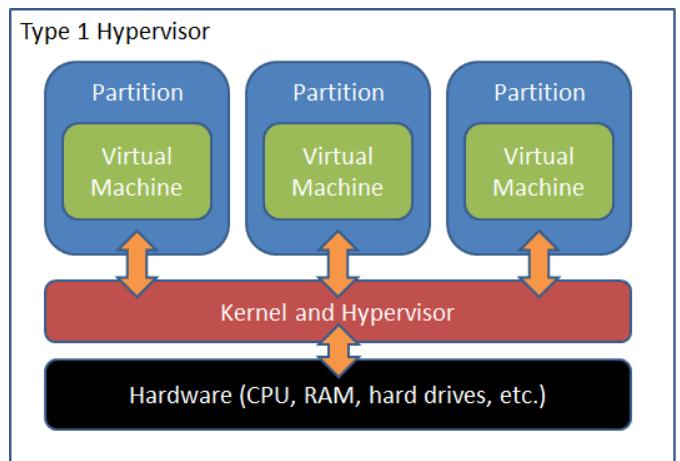


Fig. 1. A Type-1 hypervisor access flow. The Virtual Machine accesses the Hypervisor which then translates these access calls to system hardware [1]

Type 1 Hypervisors directly manage the hardware and memory to create "virtual machines" that guest operating systems

can execute within. An example of the Type-1 Hypervisor flow can be seen in Figure 1.

B. Type 2 Hypervisor

Unlike a Type-1 Hypervisor which runs in the most privileged mode and has access to system hardware, Type-2 Hypervisor is simply an application that runs on a "host" operating system which creates a virtual machine that the guest operating systems can execute in.

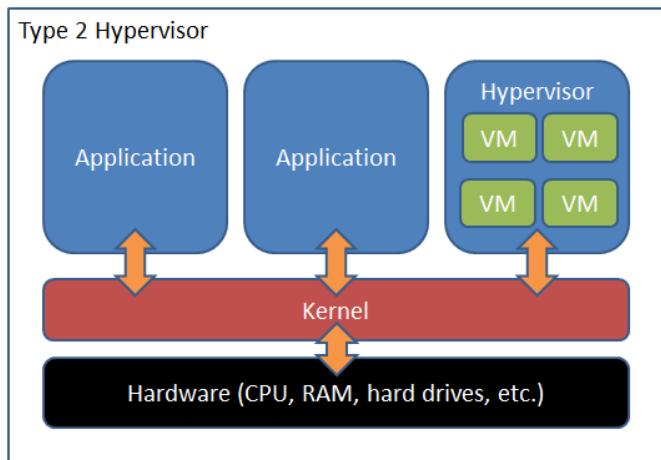


Fig. 2. A Type-1 hypervisor access flow. The Virtual Machine accesses the Hypervisor which then translates these access calls to system hardware [1]

Management of memory and hardware is performed by the Host OS like any other application that is run.

C. Virtual Machine Uses

Virtual machines are useful as they are both able to ensure complete isolation from each other, as well the ability to run different operating systems that can be configured to optimise the running application. An example use of a Type-1 Hypervisor running multiple virtual machines is having a single piece hardware running multiple VMs that each host a different server. That way, if one server is compromised, the hackers bound to only that virtual machine. Secondly, if a server crashes the operating system, only that virtual machine will be affected and the other VMs operate normally.

An example use of a Type-2 system is a software developer who is developing a cross-platform application. Instead of having to have multiple different physical computers, the developer could run the other operating system inside the virtual machine and perform all the code compilation and testing inside the VM.

III. JAILS

A. chroot Jail

chroot is a Unix operation that changes the apparent working directory for a given process and its children. The processes cannot see or access files outside the designated directory tree, and the artificial root directory is called a "chroot jail".

A chroot jail is not a good idea for security purposes, such as isolating server systems from each-other, as is possible for privileged root users to escape the jail via performing a second "chroot", or mounting a file system.

Instead, chroot is used in applications such as testing and development, where isolation of the processes simplifies the operation for the user. An example of this is having to compile software where multiple conflicting libraries can cause linkage errors. Instead, each project can be ran in its own chroot jail, with the necessary libraries present and thus isolating the conflicting libraries from each other.

B. FreeBSD Jail

FreeBSD jail operates in a similar way to a Type 2 Hypervisor, where an underlying Host OS instance of FreeBSD is used for communication to the hardware.

Each jail is partitioned into its own mini-system where applications can operate. Unlike a Virtual machine, a FreeBSD jail shares the FreeBSD OS and kernel with all other jails. Therefore, there is no Virtual Machine that is being ran, so the overhead of the FreeBSD is vastly minimised.

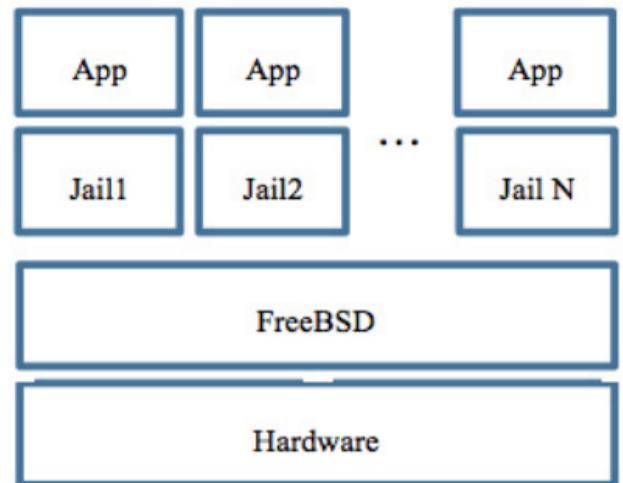


Fig. 3. An example of the structure of a FreeBSD jail. Note that each jail shares the same OS and kernel [2]

Each jail is a sand-boxed virtual environment and has its own files, processes, IP-addresses, and user/superuser accounts. Processes that are ran on one jail are not allowed to interact with those of another jail so that each jail is effectively sealed from one another. Each jail can run different software (even the same software of differing versions and configurations) as decided upon by the jail owner.

Limitations are introduced in Free-BSD jail environment as each jail is bound to the kernel of the base system. It is simply not possible for a jail to run a different kernel or different kernel configuration, which may be needed for optimal application performance. Secondly, this host kernel is a single-point of failure: if some bug or exploit is used to crash

the base kernel, then all services in all jails will also cease to function.

Thus Free-BSD jails can be used to compartmentalised sensitive and vulnerable services from each other. Given that a central-server is used to support multiple services; in this example a FTP service and a Web server are running from the same machine. By placing the separate services in differing jails, if the Web server is compromised and the attacker is able to browse through the directory tree, they will never be able to see the FTP service information.

C. Jail-Like tools offered by other OS

There are many alternative applications on other operating systems that can offer process isolation. For windows, software such as Sandboxie can perform application isolation. Linux Containers allow multiple running isolated Linux systems on a single Linux kernel, and MacOS has applications such as Docker.

REFERENCES

- [1] <https://www.altaro.com/hyper-v/hyper-v-terminology-host-operating-system-or-parent-partition/>
- [2] Performance of Jails versus Virtualization for Cloud Computing Solutions
- [3] <http://www.cs.yale.edu/homes/aspnes/pinewiki/Virtualization.html>
- [4] <https://corensic.wordpress.com/2011/12/12/virtual-machines-the-traps/>

21 Practical - Lab 8: TCPDump (C)

Lab task - TCPDump (Credit)

Outcome	Weight
ULO3	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 8: TCPDump (C)

Submitted By:

Michael REEB

4936094

2019/04/17 22:51

Tutor:

Jonathan KUA

April 17, 2019



TNE30019-Unix For Telecommunications

Lab 8 - TCPDump

Michael Reeb - 4936094

April 17, 2019

1 Results

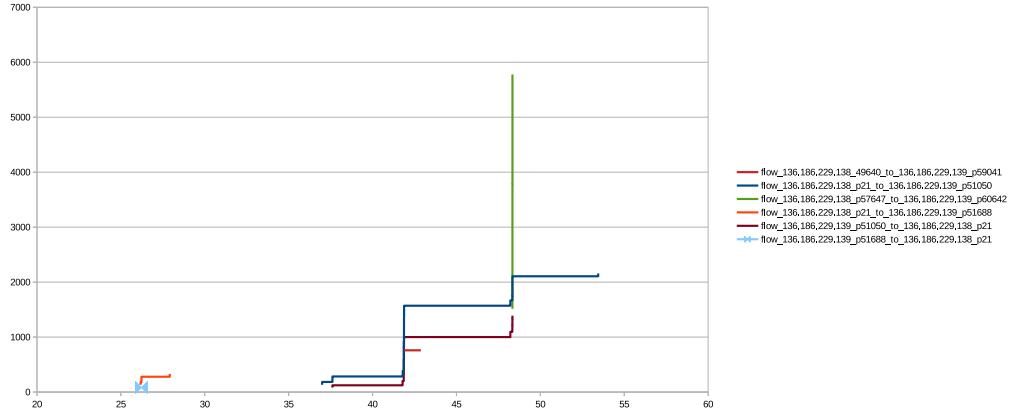


Figure 1: File Transfer Protocol (FTP) flows.

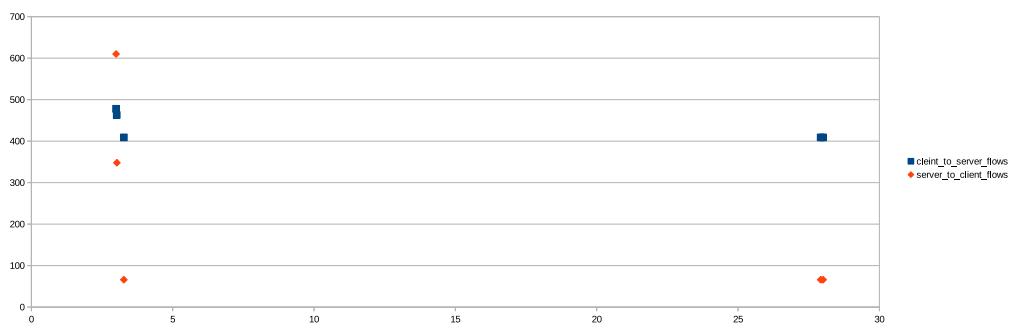


Figure 2: HTTP Flows. Note that each point is a separate flow. Flows from the Client-to-Server are blue, and flows from Server-to-Client are in orange.

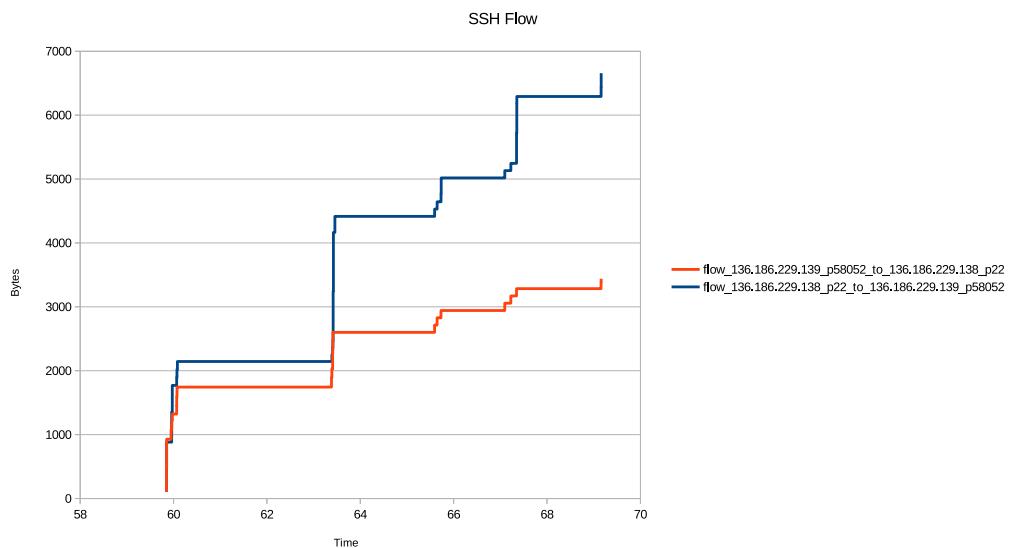


Figure 3: Secure Shell (SSH) flows.

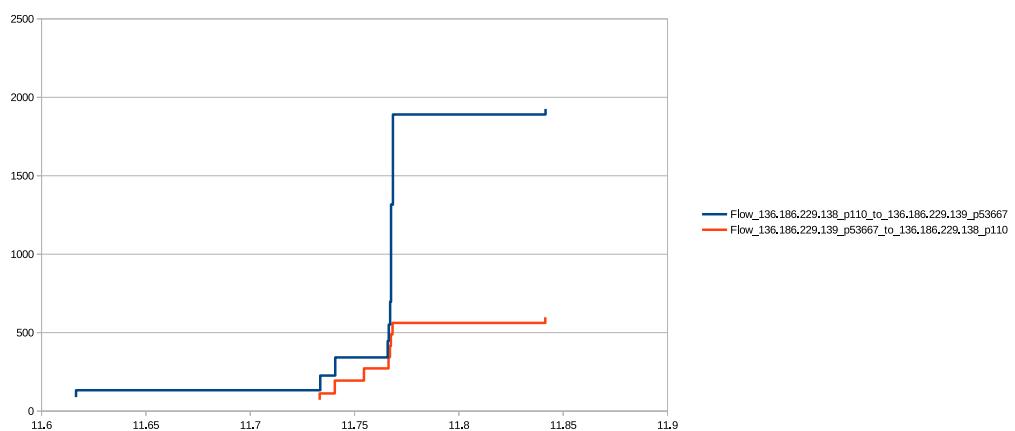


Figure 4: Post Office Protocol (POP) flows.

22 Practical - Lab 8: TCPDump (D)

Lab task - TCPDump (Distinction)

Outcome	Weight
ULO3	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 8: TCPDump (D)

Submitted By:

Michael REEB

4936094

2019/04/17 23:58

Tutor:

Jonathan KUA

April 18, 2019



TNE30019-Unix For Telecommunications

Lab 8 - TCPDump

Michael Reeb - 4936094

April 17, 2019

Contents

1	Results	1
2	Discussion	3
2.1	FTP	3
2.2	HTTP	3
2.3	SSH	3
2.4	POP	4
3	Conclusion	4

1 Results

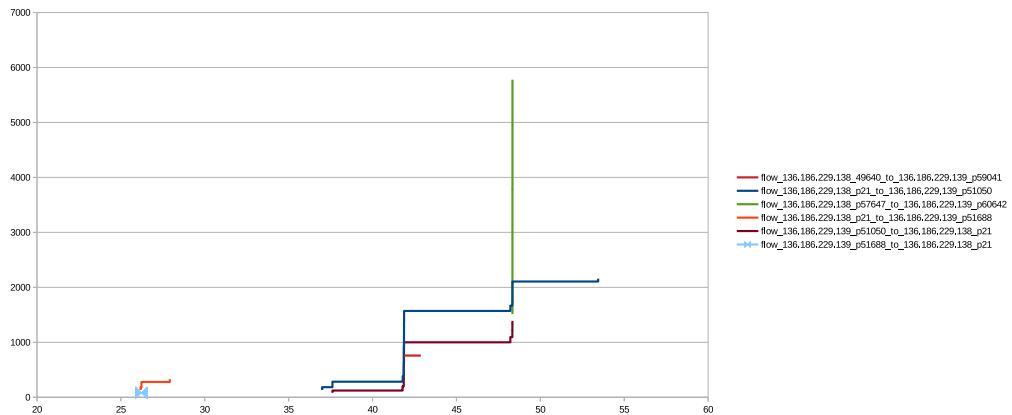


Figure 1: File Transfer Protocol (FTP) flows.

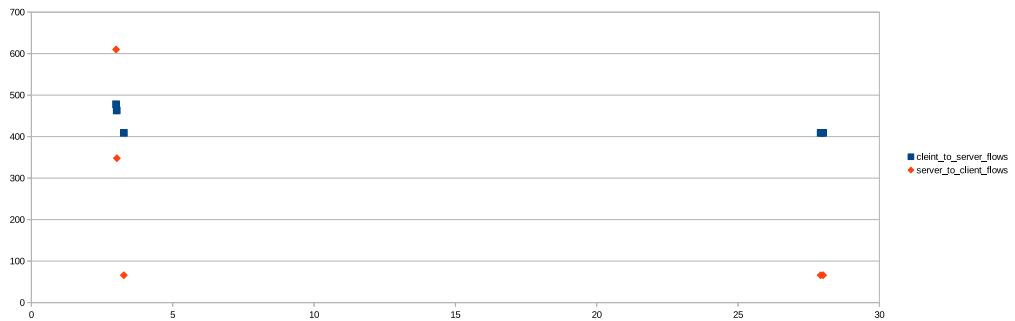


Figure 2: HTTP Flows. Note that each point is a separate flow. Flows from the Client-to-Server are blue, and flows from Server-to-Client are in orange.

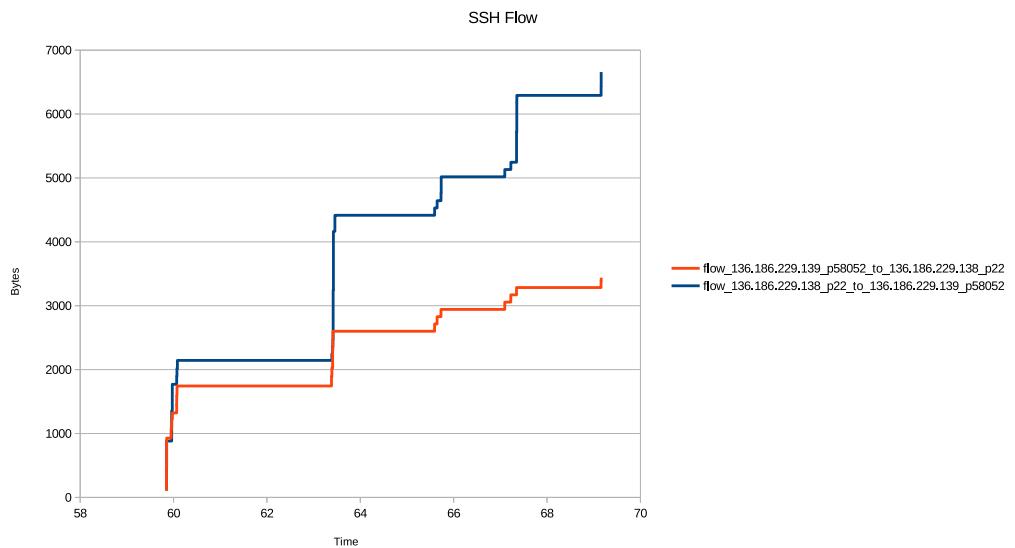


Figure 3: Secure Shell (SSH) flows.

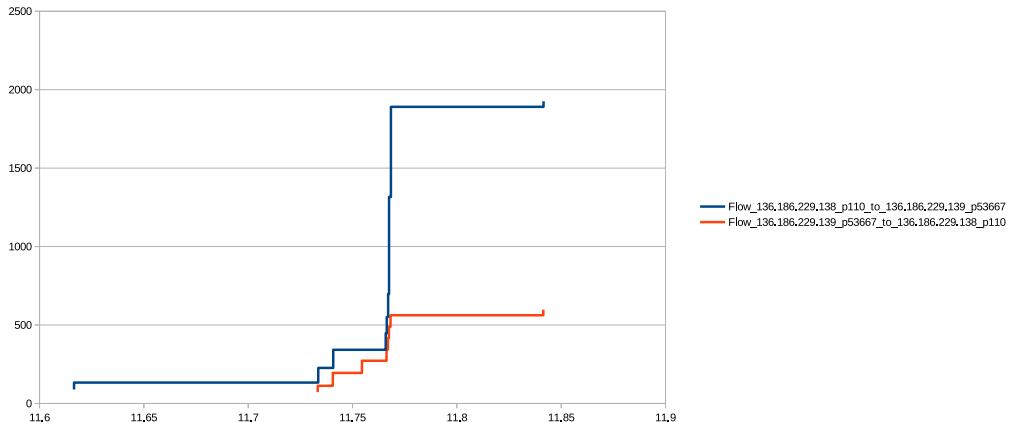


Figure 4: Post Office Protocol (POP) flows.

2 Discussion

In this lab, the Server was located on IP 136.186.229.138, while the client was located on IP 136.186.229.139

2.1 FTP

In Figure 1, it can be seen that there are six distinct flows, 2 of which are from the client device to the server, and four are from the server to the client. The first "burst" of communication was the failed login attempt using `anonymous` as username, and the servers response to the failed login, and when the client abruptly disconnected. In the second burst, the client successfully logged into the FTP server. Here we see that the server has two flows to the client, which is the way FTP operates. FTP sends commands over one connection, and the data over the second connection.

This transmission is bursty, as we can see large peaks of data transfer, followed by periods of inactivity.

2.2 HTTP

As seen in Figure 2 there are many numerous flows of traffic. This is because modern web browsers operate using parallel HTTP connects so web pages are able to be downloaded faster. This can be seen around the first burst of activity, where multiple connections are used to download the web-page. The flows that occur later on have to do with the browser making requests attempting to download the web-page's `favicon.ico`; the little icon that appears next a web-pages name in the tab-bar on modern browsers.

Again, this traffic is bursty as there are periods of high traffic, and then periods of no traffic.

2.3 SSH

In Figure 3 we can see that SSH has two distinct flows: one from the client to server, the other from server to client. we can see that the flows are steady, until there is a large amount

of data transferred from the server to the client, before remaining steady again. This may indicate that a file had been downloaded.

While not as severe as the previous protocols, these SSH flows are still bursty, particularly in the flow from the server to the client.

2.4 POP

In Figure 4 we can see that POP protocol has two distinct flows: one from the client to the server, and the other from the server to the client. The initial traffic is due to the login and verification process (we can see the password is *mypassword*, and this is followed by a large burst in the Server→Client caused by the Client asking to retrieve an email which was then subsequently served.

Again, this traffic is bursty, as can be seen by the large data transmission caused by the email retrieval.

3 Conclusion

Bursty traffic can cause negative user experience in networks with limited bandwidth, such as the Last Mile between ISPs and home networks. There is no simple solution to bursty traffic without increasing bandwidth, which is costly endeavour for ISP.

Due to the nature of operation TCP congestion control which attempts to probe the network for maximum throughput , traffic bursts will be an inevitable part of internet usage. Burst Traffic can be minimised by the use of buffers, which attempt to absorb the burst traffic. However, this can introduce frustrating RTT times into a network as buffered need to be cleared, which can particularly effect Real Time services.

23 Practical - Lab 8: TCPDump (P)

Lab task - TCPDump

Outcome	Weight
ULO3	♦♦♦◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 8: TCPDump (P)

Submitted By:

Michael REEB

4936094

2019/04/12 23:36

Tutor:

Jonathan KUA

April 12, 2019



- 1) Over what timeframe was the trace captured? From 17:45:17.195008 until 17:46:26 (approx 1 min)
 - 2) First was http request, then pop3 (email) request, then FTP transfer, then SSH
 - 3) wharrop2.caia.swin.edu.au virtual host thttpd located on the Hhoang2.caia.swin.edu.au server (Hhoang2.caia.swin.edu.au). This site is the Armitage Chemicals website from Lab 5.

```
7:45:20.191605 IP nakter2.caia.swin.edu.au.58960 > hhoang2.caia.swin.edu.au.http: Flags [P.], seq 1:413, ack 1, win 33304, options [nop,nop,TS val 600647893 ecr 8260862], length 412  
#: .. .8 .. _P..P..E.....  
(. .. GET / HTTP/1.1  
host: wharzrp2.caia.swin.edu.au  
user-agent: Mozilla/5.0 (X11; FreeBSD amd64; en-US; rv:1.8.0.1) Gecko/20061002 Firefox/1.5.0.1  
accept: text/xml,application/xml,application/xhtml+xml,text/html;q=0.9,image/png,*/*;q=0.5  
accept-language: en-us,en;q=0.5  
accept-encoding: gzip,deflate  
accept-charset: ISO-8859-1,utf-8;q=0.7,*;q=0.7  
keep-alive: 300  
connection: keep-alive
```

```
17:45:20.191881 IP hhongq2.csia.swin.edu.au.http > maktek2.csia.swin.edu.au.58960: Flags [P.], seq 1:545, ack 513, win 33304, options [nop,nop,TS val 8260862 ecr 600647893], length 544
E.T. 0.8.F.....P.B...".e....{.....
...#.({.HTTP/1.1 200 OK
Server: thtppd/2.25D 29dec2003
Content-Type: text/html; charset=iso-8859-1
Date: Sat, 12 Sep 2006 07:45:20 GMT
Last-Modified: Sat, 12 Sep 2006 14:23:59 GMT
Accept-Ranges: bytes
Connection: close
Content-Length: 303

<html>
<head><title>Welcome to Armitage Chemicals</title></head>
<body>
<center>
<h1>Welcome to Armitage Chemicals</h1>

<>Armitage Chemicals Headquarters<br>
To contact us, please deliver all letters/emails to the guard at the front gate.<br>
</center>
</body>
</html>
```

4) He typed the wrong username, user anonymous does not exist

```
17:45:43.422714 IP nakte2.caia.swin.edu.au.51688 > hhoang2.caia.swin.edu.au.ftp: Flags [P.], seq 1:17, ack 67, win 33304, options [nop,nop,TS val 600671122 ecr 8263179], length 16  
E0.D#0..0.4.....CE.T!....2....  
#.## User anonymous  
  
17:45:43.423660 IP hhoang2.caia.swin.edu.au.ftp > nakte2.caia.swin.edu.au.51688: Flags [P.], seq 67:96, ack 17, win 33304, options [nop,nop,TS val 600671122 ecr 8263185], length 29  
E..Q.C#8.H.....CE.T!....  
.##.## User anonymous unknown.
```

5) Laurence retrieve a copy (RETR) of the blkmarble-desktopImage.jpg image

```
7:46:05.532002 IP nakter2.caia.swin.edu.au.51050 -> hohoang2.caia.swin.edu.au.ftp: Flags [P.], seq 283:335, ack 670, win 33304, options [nop,nop,TS val 600693230 ecr 8265395], length 52  
0..h@e@e@:.....)....OH.....  
....~-RETR /usr/home/lslowman/Blkmarble-desktopImage.jpg
```

6) It is not possible to see the contents of the SSH session as they are encrypted. However we were able to see SSH perform the public-key exchange part of the Diffie-Hellman key exchange.

7)the contents of the email was:

```
17:45:28.963395 IP hhoang2.caia.swin.edu.au.pop3 > nakter2.caia.swin.edu.au.53667: Flags
E..'.[@.D.....n.....;t.....
.~.K#.K.+OK sending message ending with a '.' on a line by itself
From wharrop@wharrop2.caia.swin.edu.au Tue Oct 3 17:08:30 2006
Return-Path: <wharrop@wharrop2.caia.swin.edu.au>
Received: from wharrop2.caia.swin.edu.au (localhost.caia.swin.edu.au [127.0.0.1])
    by wharrop2.caia.swin.edu.au (8.13.3/8.13.3) with ESMTP id k9378UJ7005779
    for <lslowman@wharrop2.caia.swin.edu.au>; Tue, 3 Oct 2006 17:08:30 +1000 (EST)
    (envelope-from wharrop@wharrop2.caia.swin.edu.au)
Received: (from wharrop@localhost)
    by wharrop2.caia.swin.edu.au (8.13.3/8.13.3/Submit) id k9378Ums005778
    for lslowman; Tue, 3 Oct 2006 17:08:30 +1000 (EST)
    (envelope-from wharrop)
Date: Tue, 3 Oct 2006 17:08:30 +1000 (EST)
From: Warren Harrop <wharrop@wharrop2.caia.swin.edu.au>
Message-Id: <200610030708.k9378Ums005778@wharrop2.caia.swin.edu.au>
To: lslowman@wharrop2.caia.swin.edu.au
Subject: Labs?

Lawrence Slowman, you have yet to hand in any labs. Is there any reason for this? Are
you still enrolled?

Warren Harrop
```

It was sent by wharrop@wharrop2.caia.swin.edu.au (Warren Harrop) a 17:08:30
8) Laurences' password is: mypassword

```
17:45:28.949529 IP nakter2.caia.swin.edu.au.53667 > hhoang2.caia.swin.edu.au.pop3: Flags [P.], seq 22:39, ack 79, win 33304, options [nop,nop,TS val 600656650 ecr 8261736], length 17
.0.E6.0.7.....n...;I.....
.~.K
.~.hPASS mypassword
Transcript end
```

24 Communications - Speaker Reflection 1

Reflection of presentation of first Industry Invited Speaker

Outcome	Weight
ULO7	♦♦♦♦◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Communications - Speaker Reflection 1

Submitted By:

Michael REEB

4936094

2019/05/20 15:39

Tutor:

Jonathan KUA

May 20, 2019





Unix for Telecommunications

Portfolio Task – C-Speaker-Reflection-Form

Name: Michael Heeb Date: 20/5/2019

- 1) Reflect on one thing you learnt today that you were completely surprised by

Job turn around . industry has no loyalty? or do employees have no loyalty?

- 2) How do you think todays presentation changed the way you look at your career following graduation?

Going to focus on big data, AI/machine learning
as a ~~career~~ career path wrt software + comp sci

(P)

25 Communications - Lab Report 3

New Description

Outcome	Weight
ULO7	◆◆◆◆◇

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Communications - Lab Report 3

Submitted By:

Michael REEB

4936094

2019/05/02 22:58

Tutor:

Jonathan KUA

May 2, 2019



TNE30019 Unix for Telecommunications

Lab 9 - Samba

Michael Reeb (ID: 4936094)
Swinburne University of Technology
Melbourne, Australia
4936094@student.swin.edu.au

Abstract—The software package Samba allows for the creation of network directories, or shares. This allows for cross-platform file sharing (ie from Unix to Windows), and the creation of public share drives.

I. INTRODUCTION TO PCAP

Samba is a software package that implements the SMB protocol, and can be used to enable file and print sharing across multiple Operating Systems (OS) who implement different file structures. A Samba "share" is simply a directory on Unix that is offered by Samba via the SMB protocol to a connecting client to use a regular directory.

II. AIM

In this lab, Samba was used to create networked drives that could be remotely accessed by a Windows OS from Windows File Explorer using the Network Drive wizard. These shares were private, so only the account owner would be able to access them. Any files transferred or modified on this drive would transfer to the Unix local host drive of the account.

In the second part of the lab, a public shared drive was implemented so all accounts were able create, and modify files located within this share.

III. EQUIPMENT

RULE host 61 was used as the server. SSH access was though PuTTY running on an external PC connected to Swinburne via VPN. Samba version 4.8.2 was used to implement the service.

IV. METHODOLOGY

As per lab handout.

V. DISCUSSION

A. Basic Samba Configuration

Initially, Samba was configured in the `rc.conf` script so that it would run on startup, and so that the service commands would operate. This was performed by appending `samba_server_enable="YES"` to the `rc.conf` file located in `/etc` directory. The [global] section of the configuration file determines the server-wide settings, as well as the options that will apply to all shares. It was under this section that the work-group was assigned (using the `workgroup=MSHOME` option), server string was set (using `serverstring=rule61.caia.swin.edu.au`), and to ensure that authentication is used with encrypted passwords, both the options `security=user` and `encryptpasswords=yes` were activated. To ensure that only clients from the Swinburne were able to connect, the option `hostsallow=136.186.*.*` was implemented, and `*` is the wild card character meaning that any IP in the `136.186.0.0` domain is allowed to connect. Finally, the `ntlmauth=ntlmv1-permitted` was enable under the global settings to ensure that NTMLv1 and above is used for all clients.

B. Create Users

Create new accounts on the Rule host is accomplished using the FreeBSD `adduser` command ran from the shell. Two accounts were created, one called **samba** with the password *samba*, the second username was **autocollector** with the password being *autocollector*. Next, a Samba password had to be created for each of the above accounts. This was done using the `smbpasswd-aUSERNAME` shell command, where `USERNAME` the account name (`samba` and `autocollector`) that was created above. To ensure these accounts were correctly created, the

shell command pdbedit -L was used. This command generates a list of the users that are located in the Samba database.

C. Create Shares

In the `smb4.conf` file, the section denoted `[homes]` is used to configure services for connecting clients to their home directories. These are created at run-time by the Samba server, by using the Samba variable substitutions `%u` which replaces the variable with the current Samba session username.

Thus, the home shares were created with the following configuration option: `path=/home/%u` which sets the Share path depending on the user connecting (ie for the account autocollector, the Samba server will change this option to: `path=/home/autocollector`).

The `readonly=no` allows the creation and modification of files on the share. Finally, the `browsable=no` option hides the

`path[homes]` share from being able to be accessed, but will allow the user to autoconnect to their directory as explained above. If this option is not set, then any user will be able to see all the shares of the all accounts in the `homes` share, albeit a user will still only be able to access their own share.

Functionality was then tested by logging in to each account on Windows, creating a file, and then checking the accounts home directory on Rule 61 to ensure that the file is there too.

D. Creation of the public share

A public share can be created in the `smb4.conf` so all users can access. This was created in the `[public]` section. As given in the lab specification the location shared folder was set with the `path=/home/samba/pubStuff`. Again the `readonly=no` option was used to files can be created and modified. The `browsable=yes` option is used to ensure that the share is seen in the list of available shares by the client.

The `guestok=yes` ensures that no password is needed for connecting to this service.

The `createmode=777` determines the permission files created in this share. This value is bit-wise "And" with created files permissions. The `forcecreatemode=777` also is used to determine the permissions and is bit-wise ORed onto the file. 777 means that any user of any group can read, write, and execute these files created in this share.

Finally, to have this share work correctly, either the

share directory permission in Unix must be changed (by using the command `chmod777 -R /home/samba/pubStuff` which will also apply this to any directories), or by using the `forceuser=samba` option is the `[public]` share which will change any account accessing this share to use the share as if they were the user `samba`

VI. CONCLUSION

Samba is a software package that can be used to create shared network drives across multiple OS platforms. In this lab it was demonstrated that Samba can be used to generate private shares at run time on a per account basis, and that a public share directory can be setup that allows any user to create, edit and modify files.

REFERENCES

- [1] Configuration Manual for smb.conf. (n.d.). Retrieved from <https://www.samba.org/samba/docs/4.7/man-html/smb.conf.5.html>

VII. RESULTS

Listing 1. smb4.conf

```
1 [global]
2 workgroup = MSHOME
3 server string = rule61.caia.swin.edu.au
4 security = user
5 encrypt passwords = yes
6 hosts allow 136.186.*.*
7 ntlm auth = ntlmv1-permitted
8
9 [home]
10 comment = %u Home Directory
11 path = /home/%u
12 read only = no
13 #valid users = %u
14 #—————add this line to prevent the "homes" from appearing
15 browseable = no
16
17 [public]
18 comment = Public Shared Directory
19 path = /home/samba/pubStuff
20 read only = no
21 guest ok = yes
22 browseable = yes
23 create mask = 777
24 #
25 #either need to force user as below, or change permissions of public folder to
26 #read/write/executable access (chmod 777)
27 ;force user = samba
28 force create mode = 777
```

26 Practical - Lab 9: Samba (D)

Lab task - Samba (Distinction)

Outcome	Weight
ULO4	♦♦♦♦◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 9: Samba (D)

Submitted By:

Michael REEB

4936094

2019/04/29 22:37

Tutor:

Jonathan KUA

April 29, 2019



```

Marking P-Lab-09-Samba-D on RULE host rule61
TNE30019/TNE80014 - Laboratory - Samba
=====
*****
Portfolio Task: P-Lab-09-Samba, Distinction Task

Configure your rule host (rule61) with Samba running a public share

Tasks:
Two Samba users: samba(password=samba) and autocollector(password=autocollector)
Samba share available to both users: //rule61/public
Share maps to /home/samba/pubStuff
Read/write access to both users
Both users can edit/delete each others files
*****

```

```

Date: 20190429_2236
Last 5 Logons:
student pts/2 vpn247-162.ccs.swin.edu Mon Apr 29 21:52 still logged in
student pts/7 atc330-20.ict.swin.edu Mon Apr 29 15:32 - 16:24 (00:52)
student pts/5 atc330-20.ict.swin.edu Mon Apr 29 14:31 - 16:23 (01:51)
student pts/7 vpn247-53.ccs.swin.edu. Sat Apr 27 15:29 - 18:39 (03:09)
student pts/7 vpn247-53.ccs.swin.edu. Sat Apr 27 15:23 - 15:25 (00:01)

```

```

RULE host Socket information
root smbd 45415 35 tcp4 136.186.230.61:445 *:*
root smbd 45415 36 tcp4 136.186.230.61:139 *:*
root nmbd 45411 16 udp4 136.186.230.61:137 *:*
root nmbd 45411 17 udp4 136.186.230.61:138 *:*
root nmbd 45411 18 udp4 136.186.230.61:137 *:*
root nmbd 45411 19 udp4 136.186.230.61:137 *:*
root nmbd 45411 20 udp4 136.186.230.61:138 *:*
root nmbd 45411 21 udp4 136.186.230.61:138 *:*
www httpd 37620 3 tcp4 136.186.230.61:80 *:*
www httpd 37617 3 tcp4 136.186.230.61:80 *:*
tempuser sshd 35750 3 tcp4 136.186.230.61:22 136.186.247.162:49949
tempuser sshd 35750 7 tcp4 136.186.230.61:6010 *:*
root sshd 35741 3 tcp4 136.186.230.61:22 136.186.247.162:49949
root cupsd 73173 5 tcp4 136.186.230.61:631 *:*
www httpd 59579 3 tcp4 136.186.230.61:80 *:*
www httpd 59578 3 tcp4 136.186.230.61:80 *:*
www httpd 59577 3 tcp4 136.186.230.61:80 *:*
www httpd 59575 3 tcp4 136.186.230.61:80 *:*
www httpd 59574 3 tcp4 136.186.230.61:80 *:*
www httpd 59573 3 tcp4 136.186.230.61:80 *:*
www httpd 59572 3 tcp4 136.186.230.61:80 *:*
www httpd 59571 3 tcp4 136.186.230.61:80 *:*
root httpd 52413 3 tcp4 136.186.230.61:80 *:*
bind named 41933 20 tcp4 136.186.230.61:53 *:*
bind named 41933 21 tcp4 136.186.230.61:953 *:*
bind named 41933 512 udp4 136.186.230.61:53 *:*
bind named 41933 513 udp4 136.186.230.61:53 *:*
bind named 41933 514 udp4 136.186.230.61:53 *:*
bind named 41933 515 udp4 136.186.230.61:53 *:*
bind named 41933 516 udp4 136.186.230.61:53 *:*
bind named 41933 517 udp4 136.186.230.61:53 *:*
bind named 41933 518 udp4 136.186.230.61:53 *:*
bind named 41933 519 udp4 136.186.230.61:53 *:*
bind named 41933 520 udp4 136.186.230.61:53 *:*
bind named 41933 521 udp4 136.186.230.61:53 *:*

```

```
bind    named      41933 522  udp4   136.186.230.61:53      *:*
bind    named      41933 523  udp4   136.186.230.61:53      *:*
bind    named      41933 524  udp4   136.186.230.61:53      *:*
bind    named      41933 525  udp4   136.186.230.61:53      *:*
bind    named      41933 526  udp4   136.186.230.61:53      *:*
bind    named      41933 527  udp4   136.186.230.61:53      *:*
bind    named      41933 528  udp4   136.186.230.61:53      *:*
bind    named      41933 529  udp4   136.186.230.61:53      *:*
bind    named      41933 530  udp4   136.186.230.61:53      *:*
bind    named      41933 531  udp4   136.186.230.61:53      *:*
bind    named      41933 532  udp4   136.186.230.61:53      *:*
bind    named      41933 533  udp4   136.186.230.61:53      *:*
bind    named      41933 534  udp4   136.186.230.61:53      *:*
root    sshd       97762  3   tcp4   136.186.230.61:22      *:*
root    ntpd       1207   61  udp4   136.186.230.61:123     *:*
```

TEMPORARILY EDITING nsmb.conf

=====

Setting access to RULE61

 Username: SAMBA

 Password: samba

RUNNING SAMBA ACCESS TEST

=====

Checking access to //samba@RULE61/public...

Mounting (//samba@RULE61/public) to (/var/tmp/rule/collect/mnt)...

Checking Mount Status...

//SAMBA@RULE61/PUBLIC on /var/tmp/rule/collect/mnt (smbfs)

Creating file (auto_collect_test)...

Checking file access (read)...

=< contents >=====

[RULE61]

addr=rule61

[RULE61:SAMBA]

password=samba

=====

Checking file properties on actual RULE Host...

-rwxrwxrwx 1 1002 1002 52 Apr 29 22:36 /usr/jails/unix/lab/rule61/usr/home/samba/pubStuff/auto_collect_test

Checking contents of written file...

Unmounting (/var/tmp/rule/collect/mnt)

Testing summary:

Mounting share:	Pass
File creation:	Pass
Accessing mounted file:	Pass
File contents correct:	Pass
File location on host:	Pass
File Owner permissions:	Pass
File Group permissions:	Pass

```
File Public permissions: Pass

REINSTATING OLD nsmb.conf
=====

TEMPORARILY EDITING nsmb.conf
=====
Setting access to RULE61
    Username: AUTOCOLLECTOR
    Password: autocollector

RUNNING SAMBA ACCESS TEST
=====

Checking access to //autocollector@RULE61/public...
Mounting (//autocollector@RULE61/public) to (/var/tmp/rule/collect/mnt)...

Checking Mount Status...
//AUTOCOLLECTOR@RULE61/PUBLIC on /var/tmp/rule/collect/mnt (smbfs)

Creating file (auto_collect_test)...

Checking file access (read)...
=< contents >=====
[RULE61]
addr=rule61

[RULE61:AUTOCOLLECTOR]
password=autocollector
=====

Checking file properties on actual RULE Host...
-rwxrwxrwx 1 1003 1002 68 Apr 29 22:36 /usr/jails/unix/lab/rule61/usr/home/samba/pubStuff/auto_collect_test

Checking contents of written file...

Unmounting (/var/tmp/rule/collect/mnt)

Testing summary:
Mounting share: Pass
File creation: Pass
Accessing mounted file: Pass
File contents correct: Pass
File location on host: Pass
File Owner permissions: Pass
File Group permissions: Pass
File Public permissions: Pass

REINSTATING OLD nsmb.conf
=====

*****
PORTFOLIO TASK RESULT: PASSED
```

27 Practical - Lab 9: Samba (P)

Lab task - Samba

Outcome	Weight
ULO4	♦♦♦♦◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 9: Samba (P)

Submitted By:

Michael REEB

4936094

2019/04/29 15:29

Tutor:

Jonathan KUA

April 29, 2019



```
Marking P-Lab-09-Samba-P on RULE host rule61
TNE30019/TNE80014 - Laboratory - Samba
=====
*****
Portfolio Task: P-Lab-09-Samba, Pass Task
```

Configure your rule host (rule61) with Samba running private home shares

Tasks:

```
Two Samba users: samba(password=samba) and autocollector(password=autocollector)
Samba share available to both users: //rule61/home
Each share maps to /home/<username>
Users access their own (respective) home directories only
*****
```

Date: 20190429_1529

Last 5 Logons:

```
student pts/5 atc330-20.ict.swin.edu Mon Apr 29 14:31 still logged in
student pts/7 vpn247-53.cc.swin.edu. Sat Apr 27 15:29 - 18:39 (03:09)
student pts/7 vpn247-53.cc.swin.edu. Sat Apr 27 15:23 - 15:25 (00:01)
student pts/8 vpn247-190.cc.swin.edu Wed Apr 24 23:01 - 01:12 (02:10)
student pts/1 vpn247-58.cc.swin.edu. Tue Apr 23 22:13 - 22:55 (00:41)
```

RULE host Socket information

User	Process	Port	Local Address	Remote Address
root	smbd	61871	39 tcp4	136.186.230.61:445
root	smbd	61866	35 tcp4	136.186.230.61:445
root	smbd	61866	36 tcp4	136.186.230.61:139
root	nmbd	61862	16 udp4	136.186.230.61:137
root	nmbd	61862	17 udp4	136.186.230.61:138
root	nmbd	61862	18 udp4	136.186.230.61:137
root	nmbd	61862	19 udp4	136.186.230.61:137
root	nmbd	61862	20 udp4	136.186.230.61:138
root	nmbd	61862	21 udp4	136.186.230.61:138
www	httpd	59579	3 tcp4	136.186.230.61:80
www	httpd	59578	3 tcp4	136.186.230.61:80
www	httpd	59577	3 tcp4	136.186.230.61:80
www	httpd	59576	3 tcp4	136.186.230.61:80
www	httpd	59575	3 tcp4	136.186.230.61:80
www	httpd	59574	3 tcp4	136.186.230.61:80
www	httpd	59573	3 tcp4	136.186.230.61:80
www	httpd	59572	3 tcp4	136.186.230.61:80
www	httpd	59571	3 tcp4	136.186.230.61:80
www	httpd	59570	3 tcp4	136.186.230.61:80
tempuser	sshd	49718	3 tcp4	136.186.230.61:22
root	sshd	49711	3 tcp4	136.186.230.61:22
root	httpd	52413	3 tcp4	136.186.230.61:80
bind	named	41933	20 tcp4	136.186.230.61:53
bind	named	41933	21 tcp4	136.186.230.61:953
bind	named	41933	512 udp4	136.186.230.61:53
bind	named	41933	513 udp4	136.186.230.61:53
bind	named	41933	514 udp4	136.186.230.61:53
bind	named	41933	515 udp4	136.186.230.61:53
bind	named	41933	516 udp4	136.186.230.61:53
bind	named	41933	517 udp4	136.186.230.61:53
bind	named	41933	518 udp4	136.186.230.61:53
bind	named	41933	519 udp4	136.186.230.61:53
bind	named	41933	520 udp4	136.186.230.61:53
bind	named	41933	521 udp4	136.186.230.61:53
bind	named	41933	522 udp4	136.186.230.61:53
bind	named	41933	523 udp4	136.186.230.61:53

```
bind    named      41933 524  udp4   136.186.230.61:53      *:*
bind    named      41933 525  udp4   136.186.230.61:53      *:*
bind    named      41933 526  udp4   136.186.230.61:53      *:*
bind    named      41933 527  udp4   136.186.230.61:53      *:*
bind    named      41933 528  udp4   136.186.230.61:53      *:*
bind    named      41933 529  udp4   136.186.230.61:53      *:*
bind    named      41933 530  udp4   136.186.230.61:53      *:*
bind    named      41933 531  udp4   136.186.230.61:53      *:*
bind    named      41933 532  udp4   136.186.230.61:53      *:*
bind    named      41933 533  udp4   136.186.230.61:53      *:*
bind    named      41933 534  udp4   136.186.230.61:53      *:*
root    sshd       97762  3   tcp4   136.186.230.61:22      *:*
root    ntpd       1207   61  udp4   136.186.230.61:123     *:*
?       ?          ?       ?   tcp4   136.186.230.61:57615  136.186.230.61:139
?       ?          ?       ?   tcp4   136.186.230.61:57664  136.186.230.61:139
```

TEMPORARILY EDITING nsmb.conf

=====

Setting access to RULE61

 Username: SAMBA

 Password: samba

RUNNING SAMBA ACCESS TEST

=====

Checking access to //samba@RULE61/home...

Mounting (//samba@RULE61/home) to (/var/tmp/rule/collect/mnt)...

Checking Mount Status...

//SAMBA@RULE61/HOME on /var/tmp/rule/collect/mnt (smbfs)

Creating file (auto_collect_test)...

Checking file access (read)...

=< contents >=====

[RULE61]

addr=rule61

[RULE61:SAMBA]

password=samba

=====

Checking file properties on actual RULE Host...

-rwxr--r-- 1 1002 1002 52 Apr 29 15:29 /usr/jails/unix/lab/rule61/usr/home/samba/auto_collect_tes t

Checking contents of written file...

Unmounting (/var/tmp/rule/collect/mnt)

Testing summary:

Mounting share:	Pass
File creation:	Pass
Accessing mounted file:	Pass
File contents correct:	Pass
File location on host:	Pass
File Owner permissions:	Pass
File Group permissions:	Pass

```
File Public permissions: Pass

REINSTATING OLD nsmb.conf
=====

TEMPORARILY EDITING nsmb.conf
=====
Setting access to RULE61
    Username: AUTOCOLLECTOR
    Password: autocollector

RUNNING SAMBA ACCESS TEST
=====

Checking access to //autocollector@RULE61/home...
Mounting (//autocollector@RULE61/home) to (/var/tmp/rule/collect/mnt)... 

Checking Mount Status...
//AUTOCOLLECTOR@RULE61/HOME on /var/tmp/rule/collect/mnt (smbfs)

Creating file (auto_collect_test)...

Checking file access (read)...
=< contents >=====
[RULE61]
addr=rule61

[RULE61:AUTOCOLLECTOR]
password=autocollector
=====

Checking file properties on actual RULE Host...
-rwxr--r-- 1 1003 1003 68 Apr 29 15:29 /usr/jails/unix/lab/rule61/usr/home/autocollector/auto_collect_test

Checking contents of written file...

Unmounting (/var/tmp/rule/collect/mnt)

Testing summary:
Mounting share: Pass
File creation: Pass
Accessing mounted file: Pass
File contents correct: Pass
File location on host: Pass
File Owner permissions: Pass
File Group permissions: Pass
File Public permissions: Pass

REINSTATING OLD nsmb.conf
=====

*****
PORTFOLIO TASK RESULT: PASSED
```

28 Communications - Project Presentation

New Description

Outcome	Weight
ULO4	◆◆◆◆◇

Outcome	Weight
ULO7	◆◆◆◇◇

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Communications - Project Presentation

Submitted By:

Michael REEB

4936094

2019/05/29 12:43

Tutor:

Jonathan KUA

May 29, 2019



Issues with current system

- Human error in entries breaking the DNS
 - Long downtime to fix this
 - Can cause our other customers to also experience down time
- Customers must contact our representatives to perform DNS update
 - Annoying for customer
 - Annoying for our service representatives
- Have to manage reverse DNS separately
 - Human errors
 -

New Implementation

- Removing the potential for human error with automated scripts
 - Script handles all request formatting and execution
- Allow customers to self-managed their resources
 - Implemented with a login system
 - Customers can only change their allocated domain
 - Customers managed with Admin tools.
 - Customers only need to contact Admin for special requests, ie creating of whole domains
- Automatically manage the reverse DNS
 - Remove human error
 - Ensure that domains are not “forgotten” to be added if they are within the managed domain.

Where the project is at

- Proof of concept is complete. However, this software is not ready for web facing customer service
 - Working demonstration
 - Functional (albeit feature limited)
 - Web-interface & console interface
- To go forward to create commercial product needs complete rebuild to ensure robustness
 - Including all record types (AAAA, PTR).
 - More security – Web service must be secure
 - Encryption of databases
 - Authentication of users
 - Better menuing system (current one is not scalable)
 - Improved UX design for both Web Interface and Console interface
 - As many features implemented into the web interface as possible as this the easiest service for both parties to manage.
 - Customer security issue with using SSH ports and commands

Risk Analysis – Likelihood x Impact

- Human error in entries breaking the DNS
 - Current system - $5 \times 4 = 20$
 - Proposed system – $1 \times 4 = 4$
- Customer dissatisfaction with our service
 - Current system – $3 \times 4 = 12$
 - Proposed system – $1 \times 3 = 3$
- Issue with managing reverse DNS
 - Current system $5 \times 4 = 20$
 - Proposed system – $1 \times 3 = 3$

29 Communications - Speaker Reflection 2

Reflection of presentation of second Industry Invited Speaker

Outcome	Weight
ULO7	♦♦♦♦◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Communications - Speaker Reflection 2

Submitted By:

Michael REEB

4936094

2019/05/27 16:27

Tutor:

Jonathan KUA

May 27, 2019



①



Unix for Telecommunications

Portfolio Task – C-Speaker-Reflection-Form

Name:

Michael Reah

Date:

27/5/19

1) Reflect on one thing you learnt today that you were completely surprised by
change in view regarding security; ie attacks ~~were~~ are now
assumed to happen; and that your system is ~~secure~~
assumed to be going to be compromised.

2) How do you think todays presentation changed the way you look at your career following graduation?

Cyber kill chain is an intersecting ~~tool~~ tool for
security design.

"Non-Technical" Management → avoid whenever possible.

NIST cyber security framework.

30 Theory - Test 2 (P)

In class test run during tutorial in week 11 of semester

Outcome	Weight
ULO3	♦♦◊◊◊
ULO6	♦♦◊◊◊
ULO7	♦♦◊◊◊
ULO2	♦♦◊◊◊
ULO4	♦♦◊◊◊
ULO5	♦♦◊◊◊
ULO1	♦♦◊◊◊

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Theory - Test 2 (P)

Submitted By:

Michael REEB

4936094

2019/05/27 15:21

Tutor:

Jonathan KUA

May 27, 2019





Unix for Telecommunications

Portfolio Task – T-Test-2

Pass Level Task

Student ID

4936094

Student Name

Michael Leeb

Test Date

Week 11

Time Allowed

30 minutes

Assessment – Staff Use Only

In order to pass the test you **MUST** score a minimum of **60%**

Question:	1	2	3	4	Total
Points:	4	10	7	8	29
Score:	4	7½	7	6	24½

(P)

31 Practical - Lab 10: CUPS (D)

Lab task - CUPS (Distinction)

Outcome	Weight
ULO6	◆◆◆◆◆

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 10: CUPS (D)

Submitted By:

Michael REEB

4936094

2019/05/16 22:38

Tutor:

Jonathan KUA

May 16, 2019



```
Marking P-Lab-10-CUPS-D on RULE host rule61
TNE30019/TNE80014 - Laboratory - CUPS
=====
*****
Portfolio Task: P-Lab-10-CUPS, Distinction Task
```

Configure your rule host (rule61) with CUPS with a printer called PDFPrinter

Tasks:

Server accepts print jobs to printer PDFPrinter from user student and root
PDF files are created in correct locations AND with correct permissions
PDF files are valid and contain correct information

Date: 20190516_2237

Last 5 Logons:

root	pts/20		Thu May 16 16:43 - 16:43 (00:00)
root	pts/5		Thu May 16 13:59 - 20:08 (06:09)
root	pts/2		Thu May 16 13:52 still logged in
student	pts/17	atc330-17.ict.swin.edu	Thu May 16 12:54 - 13:43 (00:49)
root	pts/16		Tue May 14 13:42 - 19:45 (06:02)

RULE host Socket information

root	cupsd	48320 5	tcp4	136.186.230.61:631	*:*
www	httpd	66695 3	tcp4	136.186.230.61:80	*:*
root	smbd	45415 35	tcp4	136.186.230.61:445	*:*
root	smbd	45415 36	tcp4	136.186.230.61:139	*:*
root	nmbd	45411 16	udp4	136.186.230.61:137	*:*
root	nmbd	45411 17	udp4	136.186.230.61:138	*:*
root	nmbd	45411 18	udp4	136.186.230.61:137	*:*
root	nmbd	45411 19	udp4	136.186.230.61:137	*:*
root	nmbd	45411 20	udp4	136.186.230.61:138	*:*
root	nmbd	45411 21	udp4	136.186.230.61:138	*:*
www	httpd	37620 3	tcp4	136.186.230.61:80	*:*
www	httpd	59579 3	tcp4	136.186.230.61:80	*:*
www	httpd	59578 3	tcp4	136.186.230.61:80	*:*
www	httpd	59577 3	tcp4	136.186.230.61:80	*:*
www	httpd	59575 3	tcp4	136.186.230.61:80	*:*
www	httpd	59574 3	tcp4	136.186.230.61:80	*:*
www	httpd	59573 3	tcp4	136.186.230.61:80	*:*
www	httpd	59572 3	tcp4	136.186.230.61:80	*:*
www	httpd	59571 3	tcp4	136.186.230.61:80	*:*
root	httpd	52413 3	tcp4	136.186.230.61:80	*:*
bind	named	41933 20	tcp4	136.186.230.61:53	*:*
bind	named	41933 21	tcp4	136.186.230.61:953	*:*
bind	named	41933 512	udp4	136.186.230.61:53	*:*
bind	named	41933 513	udp4	136.186.230.61:53	*:*
bind	named	41933 514	udp4	136.186.230.61:53	*:*
bind	named	41933 515	udp4	136.186.230.61:53	*:*
bind	named	41933 516	udp4	136.186.230.61:53	*:*
bind	named	41933 517	udp4	136.186.230.61:53	*:*
bind	named	41933 518	udp4	136.186.230.61:53	*:*
bind	named	41933 519	udp4	136.186.230.61:53	*:*
bind	named	41933 520	udp4	136.186.230.61:53	*:*
bind	named	41933 521	udp4	136.186.230.61:53	*:*
bind	named	41933 522	udp4	136.186.230.61:53	*:*
bind	named	41933 523	udp4	136.186.230.61:53	*:*
bind	named	41933 524	udp4	136.186.230.61:53	*:*
bind	named	41933 525	udp4	136.186.230.61:53	*:*
bind	named	41933 526	udp4	136.186.230.61:53	*:*

```
bind    named      41933 527  udp4   136.186.230.61:53      *:*
bind    named      41933 528  udp4   136.186.230.61:53      *:*
bind    named      41933 529  udp4   136.186.230.61:53      *:*
bind    named      41933 530  udp4   136.186.230.61:53      *:*
bind    named      41933 531  udp4   136.186.230.61:53      *:*
bind    named      41933 532  udp4   136.186.230.61:53      *:*
bind    named      41933 533  udp4   136.186.230.61:53      *:*
bind    named      41933 534  udp4   136.186.230.61:53      *:*
root    sshd       97762  3   tcp4   136.186.230.61:22      *:*
root    ntpd       1207   61  udp4   136.186.230.61:123     *:*
```

SENDING PRINT JOBS

=====

Removing existing output print jobs files on rule host for (student)...

Sending print job as (student)

Pausing to allow output file to be created...

Checking output directory (/usr/jails/unix/lab/rule61/usr/home/PDF/student), owner (1001:student), permissions(rw???????)...

drwxr-x--- 2 1001 0 3 May 16 22:38 student

Checking output file - expected location(/usr/jails/unix/lab/rule61/usr/home/PDF/student), owner(1001:student), permissions(rw???????)...

total 1

-rw-r----- 1 1001 0 15157 May 16 22:38 cups_test.txt-16052019-22:38:00.pdf

Checking PDF...

IF THE FOLLOWING IS EQUAL TO 'PDF%' (or similar) PLEASE LET ME KNOW

=====

%-1

=====

----- No current checking for PDF validity

Checking PDF contents...

==</usr/jails/unix/lab/rule61/usr/home/PDF/student/cups_test.txt-16052019-22:38:00.pdf>=====

#####
TEST OUTPUT FILE FOR CUPS PRINT JOB
#####

#####

Hello world...

=====

Testing summary:

Print job accepted:	Pass
Correct output directory created:	Pass
Output file created:	Pass
Output directory owner:	Pass
Output directory permissions:	Pass
Output file owner:	Pass
Output file permissions:	Pass
Output file is a valid PDF:	Pass
Output file contents:	Pass

Removing existing output print jobs files on rule host for (root)...

```
Sending print job as (root)

Pausing to allow output file to be created...

Checking output directory (/usr/jails/unix/lab/rule61/usr/home/PDF/remroot), owner (0:root), permissions(rw??????)...
drwxr-x--- 2 0      0 3 May 16 22:38 remroot

Checking output file - expected location(/usr/jails/unix/lab/rule61/usr/home/PDF/remroot), owner(0:root), permissions(rw??????)...
total 33
-rw-r----- 1 0  0 15157 May 16 22:38 cups_test.txt-16052019-22:38:13.pdf

Checking PDF...
IF THE FOLLOWING IS EQUAL TO 'PDF%' (or similar) PLEASE LET ME KNOW
=====
%-1
=====
----- No current checking for PDF validity

Checking PDF contents...
==</usr/jails/unix/lab/rule61/usr/home/PDF/remroot/cups_test.txt-16052019-22:38:13.pdf>=====
#####
## TEST OUTPUT FILE FOR CUPS PRINT JOB ##
#####

Hello world...

=====
Testing summary:
Print job accepted:          Pass
Correct output directory created: Pass
Output file created:          Pass
Output directory owner:        Pass
Output directory permissions:  Pass
Output file owner:            Pass
Output file permissions:      Pass
Output file is a valid PDF:   Pass
Output file contents:         Pass

*****
PORTFOLIO TASK RESULT: PASSED
```

32 Practical - Lab 10: CUPS (P)

Lab task - CUPS

Outcome	Weight
ULO6	♦♦♦♦♦

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Lab 10: CUPS (P)

Submitted By:

Michael REEB

4936094

2019/05/02 12:44

Tutor:

Jonathan KUA

May 2, 2019



```
Marking P-Lab-10-CUPS-P on RULE host rule61
TNE30019/TNE80014 - Laboratory - CUPS
=====
*****
Portfolio Task: P-Lab-10-CUPS, Pass Task
```

Configure your rule host (rule61) with CUPS with a printer called PDFPrinter

Tasks:

- Configure a CUPS server on your RULE host
- Create a printer called PDFPrinter
- Configure this device to accept print jobs

```
*****
```

Date: 20190502_1243

Last 5 Logons:

```
student pts/8 atc330-19.ict.swin.edu Thu May 2 11:45 still logged in
```

```
/usr/jails/unix/lab/rule61/var/log/utx.log begins Thu May 2 11:45:00 AEST 2019
```

RULE host Socket information

User	Process	Port	Type	Local Address	Foreign Address	
root	cupsd	29213	5	tcp4	136.186.230.61:631	*:*
tempuser	sshd	16660	3	tcp4	136.186.230.61:22	136.186.15.215:56311
root	sshd	16217	3	tcp4	136.186.230.61:22	136.186.15.215:56311
root	smbd	45415	35	tcp4	136.186.230.61:445	*:*
root	smbd	45415	36	tcp4	136.186.230.61:139	*:*
root	nmbd	45411	16	udp4	136.186.230.61:137	*:*
root	nmbd	45411	17	udp4	136.186.230.61:138	*:*
root	nmbd	45411	18	udp4	136.186.230.61:137	*:*
root	nmbd	45411	19	udp4	136.186.230.61:137	*:*
root	nmbd	45411	20	udp4	136.186.230.61:138	*:*
root	nmbd	45411	21	udp4	136.186.230.61:138	*:*
www	httpd	37620	3	tcp4	136.186.230.61:80	*:*
www	httpd	37617	3	tcp4	136.186.230.61:80	*:*
www	httpd	59579	3	tcp4	136.186.230.61:80	*:*
www	httpd	59578	3	tcp4	136.186.230.61:80	*:*
www	httpd	59577	3	tcp4	136.186.230.61:80	*:*
www	httpd	59575	3	tcp4	136.186.230.61:80	*:*
www	httpd	59574	3	tcp4	136.186.230.61:80	*:*
www	httpd	59573	3	tcp4	136.186.230.61:80	*:*
www	httpd	59572	3	tcp4	136.186.230.61:80	*:*
www	httpd	59571	3	tcp4	136.186.230.61:80	*:*
root	httpd	52413	3	tcp4	136.186.230.61:80	*:*
bind	named	41933	20	tcp4	136.186.230.61:53	*:*
bind	named	41933	21	tcp4	136.186.230.61:953	*:*
bind	named	41933	512	udp4	136.186.230.61:53	*:*
bind	named	41933	513	udp4	136.186.230.61:53	*:*
bind	named	41933	514	udp4	136.186.230.61:53	*:*
bind	named	41933	515	udp4	136.186.230.61:53	*:*
bind	named	41933	516	udp4	136.186.230.61:53	*:*
bind	named	41933	517	udp4	136.186.230.61:53	*:*
bind	named	41933	518	udp4	136.186.230.61:53	*:*
bind	named	41933	519	udp4	136.186.230.61:53	*:*
bind	named	41933	520	udp4	136.186.230.61:53	*:*
bind	named	41933	521	udp4	136.186.230.61:53	*:*
bind	named	41933	522	udp4	136.186.230.61:53	*:*
bind	named	41933	523	udp4	136.186.230.61:53	*:*
bind	named	41933	524	udp4	136.186.230.61:53	*:*
bind	named	41933	525	udp4	136.186.230.61:53	*:*
bind	named	41933	526	udp4	136.186.230.61:53	*:*

```
bind    named      41933 527  udp4   136.186.230.61:53      *:*
bind    named      41933 528  udp4   136.186.230.61:53      *:*
bind    named      41933 529  udp4   136.186.230.61:53      *:*
bind    named      41933 530  udp4   136.186.230.61:53      *:*
bind    named      41933 531  udp4   136.186.230.61:53      *:*
bind    named      41933 532  udp4   136.186.230.61:53      *:*
bind    named      41933 533  udp4   136.186.230.61:53      *:*
bind    named      41933 534  udp4   136.186.230.61:53      *:*
root    sshd       97762  3   tcp4   136.186.230.61:22      *:*
root    ntpd       1207   61  udp4   136.186.230.61:123     *:*
```

CHECKING PRINTER INSTALLATION/AVAILABILITY

```
=====
```

Listing installed printers on rule61...

PDFPrinter accepting requests since Thu May 2 12:26:13 2019

Testing...

Printer (PDFPrinter) is installed: Pass

Printer is accepting jobs: Pass

```
*****
```

PORTFOLIO TASK RESULT: PASSED

33 Practical - Project

project submission

Outcome	Weight
ULO3	♦♦♦◊◊
ULO6	♦♦♦♦◊
ULO7	♦♦♦♦◊
ULO2	♦♦♦♦◊
ULO4	♦♦♦♦◊
ULO5	♦♦♦♦◊
ULO1	♦♦♦♦◊

Date	Author	Comment
2019/06/05 12:30	Jonathan Kua	Hi Michael, maybe I missed something in your document, but I can't seem to find the username and password to login via the Web interface.
2019/06/05 15:26	Michael Reeb	Hi Jon, To launch the webserver, simply login to the webserver account. username is <webserver>; password (should) also be <webserver>. To connect to the webpage (I use Chrome). The webserver is located on: http://136.186.230.101:8000, ie port 8000 (due to privilege issues when using well known ports). I have just connected to the webserver via VPN and it seems to work fine. Thanks, mike
2019/06/05 15:34	Michael Reeb	as for login into the web interface, the accounts you have created in the Admin tools should work to login here
2019/06/05 15:43	Jonathan Kua	Great, thanks Michael.

SWINBURNE UNIVERSITY OF TECHNOLOGY

UNIX FOR TELECOMMUNICATIONS

DOUBTFIRE SUBMISSION

Practical - Project

Submitted By:

Michael REEB

4936094

2019/05/29 21:20

Tutor:

Jonathan KUA

May 29, 2019



SWINBURNE UNIVERSITY OF TECHNOLOGY

TNE30019 - C-Project-Bind - Project Report

Michael REEB
4936094

May 29, 2019

Contents

1	Aim	2
2	Setting Up and Running Project	2
2.1	Dependencies	2
2.2	List of freeBSD accounts and Passwords	2
2.3	List of created files	3
2.3.1	Main python files	3
2.3.2	Database files	3
2.3.3	Other Files	4
2.4	Running and troubleshooting the Web-server	4
2.5	Running and troubleshooting the User CLI interface	4
2.6	Running and troubleshooting the Admin CLI interface	4
3	Self Reflection	5

1 Aim

The aim of this project was to implement a fully operable DDNS server with easy to use User front-end, in both command line interface and a web interface. This would allow for Users to be self-sufficient with their DNS needs and thus would not have to rely upon Administrators to perform basic tasks.

The second part of this project was to implement a suite of Admin tools that would allow for: creating and deleting of User accounts, creating and removing DNS zones, and creating and removing reverse DNS zones. (Note that the reverse DNS control was only partially implemented and thus removed from the demonstration.

2 Setting Up and Running Project

This whole project was written in Python3.6, and used a number of libraries to implement the web interface, database, and IP ranges. Python3.6 was chosen throughout the project such that the scripts would be able to be directly called by the differing interfaces reducing the need for wrapper classes; and so that behaviour would be consistent thought out.

2.1 Dependencies

The dependencies for this project are listed below in "pip install" format for ease of replication for the future user. Note that this list may not be comprehensive and is upon the reader to recognise missing packages and install them. It is also assumed that Python3.6 is installed with the PIP package installer.

```
pip install flask-login  
pip install flask-sqlalchemy  
pip install flask-migrate  
pip install flask-wtf  
pip install python-dotenv  
pip install flask  
pip install ipaddress
```

2.2 List of freeBSD accounts and Passwords

The following is a list of freeBSD users, passwords, and a description of the account usage. The username/password is given in the form of: **username - password**.

student - z2rx3gyp: The student account is the account to log in to to access FreeBSD root privileges. Also has a working copy of all the files

root - bqxplus1 Root can be accessed from If anything breaks, use this account to fix it

ddns - ddns - Login into this account will automatically run the User command line interface.

admin - admin - Login into this account will automatically run the Admin tools command line interface. This account can also be used to gain root privileges.

webserver - webserver - Login into this account will automatically run the Webserver in a background process that will persist after the terminal has been closed.

2.3 List of created files

A large amount of file was created for this project, so I will briefly explain the files I wrote and what they contain. Apart from the Named configuration files and zones files, and the login scripts, all files are found within the `/usr/local/share/webInterface/venv/microblog` directory.

2.3.1 Main python files

- `pythonDNS.py` - This contains all the nsupdate functions, reverse update, and basic querying . This is where checking domains, ip ranges, creating the nsupdate commands and sending them to the server. This also contains all the User console functions, and is responsible for User login, and command parsing.
- `pythonDNS_Admintools.py` - This contains all the Administration scripts - Creating zone files, updating Named.conf, adding Users to the database, adding Domains to the Database, etc. This script also contains all the Admin console functions, and is also able to call the User tools.
- `microblog.py` - The entry point for the web-server. The web-server was built using the Flask framework , and this file is where we can specify the port and ip address that the server will run on.
- `__init__.py` - Initialises the Flask web-server data.
- `config.py` - Contains the configuration info for the database Flask-Alchemy (which uses SQL-Lite)
- `forms.py` - Specifies the forms used in the dynamically generated HTML pages that is used in the User web interface. Allows the reading and writing of data
- `models.py` - Specifies the SQL tables and databases that will be used throughout the programs to query from.
- `templates.py` - Contains the information related to the operation of the html pages, including forms, redirections, database checking, user login etc. This file directly uses the HTML template pages located in the `./templates/` directory

2.3.2 Database files

- `app.db` - contains all the User login data. Passwords are hashed with sha256.
- `domainDb.db` - contains all the Domain data.
- `reverseDb.db` - contains the reverse DNS info (ip range)

2.3.3 Other Files

- `./templates/` - contains all the HTML webpage templates that are then filled in by Python in `templates.py` before being served. They are basic and straightforward, and need no further explanation.
- Each user account has a script in their respective home directory that begins their service.
A template file in `../named/dynamic` directory where the Dynamic Zones are located. This template is used to generate new Zones.

2.4 Running and troubleshooting the Web-server

To begin the web-server simply login to the `webserver` account, and the login script should automatically run an instance of this web-server.

the server can then be accessed on `136.186.230.101:8000`.

If this fails for whatever reason, exit to the shell (ctrl+c) then try manually running the shell script (`./startWebserver`)

If this script fails, then navigate and run `/usr/local/share/webInterface/venv/microblog/microblog.py`

Also check that there is no current running background web server (easiest I found was with `sockstat -4`) and kill if necessary.

Operating the website is self explanatory. Please be mindful as this is not very well sanity checked, so it is not overly difficult to create requests and/or queries that can break or lockup the script.

2.5 Running and troubleshooting the User CLI interface

To begin the User command line script, simply login to the DDNS account, and then after a pause the Login prompt should begin. Here you will be able to enter valid User credentials. Once logged in, you can add and remove DNS entries from the respective domains by following the console prompts.

Likewise, if this fails try executing the login script `loginScript`. If the login script is broken, then run the program directly from `/usr/local/share/webInterface/venv/microblog/pythonDNS.py`. Please be mindful as this is not very well sanity checked, so it is not overly difficult to create requests and/or queries that can break or lockup the script.

2.6 Running and troubleshooting the Admin CLI interface

Like wise, to run the Admin tools, simply sshing to the Admin user should open the Admin-tools CLI. If that does not work, execute the script `loginScript` in the gome directory. Finally, attempt to run the `/usr/local/share/webInterface/venv/microblog/pythonDNS_Admtools.py` script if nothing else works.

Also worth noting, be careful with these Admin tools as it is possible to break things peculiarly. Currently there is no database referencing, so deleting an entry (say a Domain) will not update to other tables (say the User who contains a field of Domains that they are allowed to update to), and thus the User will be able to attempt to NSUpdate non-existent

Zones.

3 Self Reflection

To continue on the is project would require a complete rewrite, as this project was built by tacking on more and more features which has caused a lot of clumsy practises.

Better and structured value checking should be implemented through this program, as often it is done in the lower underlying classes and functions, where values should be checked near to where they are read in from user input.

The database should be constructed in such way that the tables are properly join, as currently information is saved as a string which is parsed and queried in the other database table.

The reverse DNS lookup should be properly implemented, currently it was a huge hassle to attempt to implement, as through out the program the reverse IP range was statically defined when needed, in hind sight this was NOT an easy way of doing this.

The console and command parsing is clumsy, and in the next version a proper command parser and menuing system should be implemented.

The Web interface is rather plain, and could use some enhancements.

Big Brother is Watching You

All in all, I am happy with what I achieved. I entered this subject never having used a Unix os before, and never have done any BASH, Python, or HTML before. This subject has made me familiar with some vital knowledge that has filled a lot of holes I had regarding Operating Systems, networking, computer systems in general. Whilst a lot of work, I thoroughly enjoyed this Unit. Also printers suck