Discrete Mathematics Question Set 1

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- 1. Simplify the statements so that no negation is outside a quantifier or an expression involving logical connectives. (20 pt.)
- **a.** $\neg \forall x \forall y (P(x,y) \lor Q(x,y))$

$$\neg \forall x \forall y (P(x,y) \lor Q(x,y))$$

We can change the universal quantifier to the existential quantifier because $\neg \forall = \exists$ and factor the \neg into the \land operation. The equation now becomes:

$$\exists x \forall y \neg (P(x,y) \lor Q(x,y))$$

Apply DeMorgans Law:

$$\exists x \forall y (\neg P(x,y) \land \neg Q(x,y))$$

We cannot simplify any more.

b. $\neg(\exists x\exists y\neg P(x,y) \land \forall x\forall yQ(x,y))$ Distribute the outer negation and using DeMorgans Law:

$$\neg \exists x \neg \exists y \neg \neg P(x,y) \lor \neg \forall x \neg \forall y \neg Q(x,y)$$

Simplify statements:

$$\forall x \forall y P(x,y) \lor \exists x \exists y \neg Q(x,y)$$

We cannot simplify any more.

2. Prove the validity of this argument:

Premises:

$$(p \land t) \to (r \lor s) \tag{1}$$

$$q \to (u \land t) \tag{2}$$

$$u \to p$$
 (3)

$$\neg s$$
 (4)

Conclusion:

$$q \rightarrow r$$

- 1) $(p \wedge t) \rightarrow (r \vee s)$ Premise
- 2) $(p \wedge t)$ Premise
- 3) p 2, Specification
- 4) t 2, Specification
- q 2, specime q 5)
- 6) $q \to (u \land t)$ Premise
- 7) $u \wedge t$ 5,6 Modus Ponens
- 8) $u \to p$ Premise
- 9) $u \to (p \land t)$ 3, 4 Conjunction
- 10) $(p \wedge t) \rightarrow (r \vee s)$ Premise
- 11) $u \to (r \lor s)$ 9, 10 Hypothetical Syllogism
- 12) $\neg s$ Premise
- 13) $u \to r$ 11, 12 Disjunctive Syllogism
- 14) $(u \wedge t) \rightarrow r$ 7, 13 Conjunction
- 15) $q \to (u \land t)$ Premise
- 16) $q \rightarrow r$ 14, 15 Hypothetical Syllogism

Other Method:

- 1) q Premise
- 2) $q \to (u \land t)$ Premise
- 3) $u \wedge t$ 1,2 Modues Ponens
- 4) u 3, Specification
- 5) t 3, Specification
- 6) $u \to p$ Premise
- 7) p 4, 6 Modus Ponens
- 8) $p \wedge t$ 5, 7 Specification
- 9) $(p \wedge t) \rightarrow (r \vee s)$ Premise
- 10) $(r \lor s)$ 8,9 Modus Ponens
- 11) $\neg s$ Premise
- 12) r 10, 11 Disjunctive Syllogism
- 13) $q \rightarrow r$ 1,12 Conditional Proof

Because q is implied as a premise, and we have derived r using rules of inference, we can use conditional proof as a method of saying $q \to r$.

3. Is this statement a tautology? (Explain using a truth table) (10 pts.)

$$(\neg q \land (p \to q)) \to \neg p$$

P		Q	not p	not q	p -> q	not q and (p implies q)	(not q and (p implies q)) implies not p
	1	1	0	0	1	0	1
	1	0	0	1	0	0	1
	0	1	1	0	1	0	1
	0	0	1	1	1	1	1

This statement is a tautology.

- **4.** Prove that if n is an integer and $n^3 + 5$ is odd, then n is even using (20 pts.)
 - a) proof by contraposition
 - b) proof by contradiction

Proof 1 Assume n is an odd integer.

So n = 2k + 1 for some integer k.

Plugging in n = 2k + 1 into $n^3 + 5$ we get

$$(2k+1)^3 + 5$$
$$8k^3 + 12k^2 + 6k + 1 + 5$$
$$8k^3 + 12k^2 + 6k + 6$$

We can factor out a 2 from this equation to get

$$2(4k^3 + 6k^2 + 3k + 3)$$

Since k was an integer, $8k^3 + 12k^2 + 6k + 6$ is also an integer. Because we assumed n was odd and $n^3 + 5$ can be represented as 2(j) where

 $j = 4k^3 + 6k^2 + 3k + 3,$

n has to be even.

Proof 2 Suppose that if n is an integer and $n^3 + 5$ is even, then n has to be odd.

If n is odd, it can be represented as 2k + 1 for some integer k.

We can plug it into $n^3 + 5$

$$(2k+1)^3 + 5$$

$$8k^3 + 12k^2 + 6k + 1 + 5$$

$$8k^3 + 12k^2 + 6k + 6$$

$$2(4k^3 + 6k^2 + 3k + 3)$$

Since we assumed that n was an integer, $2(4k^3 + 6k^2 + 3k + 3)$ must also be an integer.

Because $n^3 + 5$ can be represented as 2(j) where $j = (4k^3 + 6k^2 + 3k + 3)$,

this directly contradicts with our claim that n has to be odd.

Therefore, n is even.

- ${f 5.}$ Determine whether each of these statements is true or false and explain the reason briefly. (30 pts.)
 - a) $x \in \{x\}$

This is true. The set $\{x\}$ contains one element which is x.

b) $\{x\} \subseteq \{x\}$

This is true. Every element in $\{x\}$ is part of the set $\{x\}$.

c) $\{x\} \in \{x\}$

This is false. The set $\{x\}$ does not have within itself another set, only an element x.

- 6 Let A, B, and C be sets. Using membership or Venn diagram show that
- a) $(A \cup B) \subseteq (A \cup B \cup C)$

 $A \cup B \cup C$ is a union of all three sets such that it contains all elements of all three sets.

Therefore, the elements of $A \cup B$ are part of the union of all three sets. Therefore, $A \cup B \subseteq (A \cup B \cup C)$ is true.

b) $(A-B)-C\subseteq A-C$

(A-B)-C is the set of all elements that are unique to only A.

The set A-C is the set of all elements that are unique to A, but also includes elements from B as well.

Therefore, the elements in (A - B) - C which only includes elements unique to A is a subset of A - C which includes all elements unique to set A, plus elements of set B.