

# Lecture1 Discrete Math

## Theorem 1 Rational Root Theorem:

If  $p(x)$  is a monic polynomial with integer coefficients, then the roots of  $p(x)$  are either integers or irrational numbers.

**Theorem2:** If  $n$  is an integer, then  $n$  is either an integer or an irrational number

# Lecture2 Discrete Math-- Propositional Logic

## Propositions:

A Proposition is a statement that is either True (T) or False (F), but not both

T and F are the **truth values** of a Proposition.

Propositions must be **declarative** (must declare a fact) and must have a truth value.

## Propositional Logic aka Propositional Calculus:

### Negation (NOT):

Definition: Let  $p$  be a proposition. The negation of  $p$  denoted  $\neg p$  or  $\bar{p}$  is the statement "not  $p$ ".

### Conjunction (AND):

Definition: Let  $p$  and  $q$  be propositions. The conjunction of  $p$  and  $q$  denoted by  $p \wedge q$  is the proposition " $p$  and  $q$ ".

The conjunction of two propositions is true if and only if both propositions are true and is false otherwise.

### Disjunction (OR):

Definition: Let  $p$  and  $q$  be propositions. The disjunction of  $p$  and  $q$  denoted by  $p \vee q$ , is the proposition " $p$  or  $q$ ".

The disjunction of two propositions is false if and only if both propositions are false and is true otherwise.

### Exclusive Or (XOR):

**Definition:** Let  $p$  and  $q$  be propositions. The 'exclusive or' of  $p$  and  $q$  denoted by  $p \oplus q$  is the proposition " $p$  xor  $q$ ", that is true when exactly one of  $p$  and  $q$  is true and is false otherwise.

### Implication ( $\rightarrow$ ):

**Definition:** Let  $p$  and  $q$  be propositions. The conditional statement  $p \rightarrow q$  is the proposition "if  $p$ , then  $q$ ". The implication  $p \rightarrow q$  is false when  $p$  is true and  $q$  is false, and true otherwise.

If  $(p \rightarrow q)$ , then:

$p$  is sufficient for  $q$

$q$  is necessary for  $p$

### Converse 逆命题, Contrapositive 逆否命题 and Inverse 否命题:

**Definition:** The proposition  $q \rightarrow p$  is called the converse of  $p \rightarrow q$ .

The proposition  $\neg q \rightarrow \neg p$  is called the contrapositive of  $p \rightarrow q$ .

The proposition  $\neg p \rightarrow \neg q$  is called the inverse of  $p \rightarrow q$ .

命题类型	形式	逻辑关系
原命题	$P \rightarrow Q$	原命题的真值决定逆否命题的真值。
逆命题	$Q \rightarrow P$	与原命题的真值无关。
否命题	$\neg P \rightarrow \neg Q$	与逆命题等价。
逆否命题	$\neg Q \rightarrow \neg P$	与原命题等价。

### Biconditionals ( $\leftrightarrow$ )

**Definition:** Let  $p$  and  $q$  be propositions. The biconditional statement  $p \leftrightarrow q$  is the proposition " $p$  if and only if  $q$ ".  
 $p \leftrightarrow q$  is true when  $p$  and  $q$  have the same truth values and is false otherwise.

### Precedence of Logical Operators:

Operator	Precedence
$\neg$	1
$\wedge$	2
$\vee$	3
$\rightarrow$	4
$\leftrightarrow$	5

Methods of Propositional Logic:

Truth Tables for Compound Propositions

Logical Equivalences:

### Tautology 重言式:

Definition: A compound proposition that is always true, irrespective of the truth values of the propositional variables in it, is called a Tautology.

### Contradiction 矛盾:

Definition: A compound proposition that is always false, irrespective of the truth values of the propositional variables in it, is called a Contradiction.

### Logically equivalent:

Definition: The compound propositions  $p$  and  $q$  are called logically equivalent (denoted by  $p \equiv q$ ) if  $p \leftrightarrow q$  is a tautology.

And sometimes we also use  $\Leftrightarrow$  in place of  $\equiv$ .

$p \equiv q$  is a statement, not a compound proposition

Identity laws, Domination laws, Idempotent laws, Negation laws:

$$\begin{array}{cccc} p \wedge T \equiv p & p \vee T \equiv T & p \vee p \equiv p & p \vee \neg p \equiv T \\ p \vee F \equiv p & p \wedge F \equiv F & p \wedge p \equiv p & p \wedge \neg p \equiv F \end{array}$$

Commutativity Rules:

$$\begin{array}{l} p \wedge q \equiv q \wedge p \\ p \vee q \equiv q \vee p \end{array}$$

Associativity Rules:

$$\begin{array}{l} (p \wedge q) \wedge r \equiv p \wedge (q \wedge r) \\ (p \vee q) \vee r \equiv p \vee (q \vee r) \end{array}$$

Double Negation, Bi-Implication and Contrapositive rules:

$$\begin{array}{l} \neg(\neg p) \equiv p \\ (p \leftrightarrow q) \equiv (p \rightarrow q) \wedge (q \rightarrow p) \quad \text{Also: } (p \leftrightarrow q) \equiv \neg p \leftrightarrow \neg q \equiv (p \wedge q) \vee (\neg p \wedge \neg q) \\ (p \rightarrow q) \equiv (\neg q \rightarrow \neg p) \end{array}$$

Conditional-Disjunction Equivalence (aka Implication rule):

$$(p \rightarrow q) \equiv (\neg p \vee q)$$

Observe that  $\neg p \vee q$  is false only when  $p$  is true and  $q$  is false.

$$\neg(p \rightarrow q) \equiv (p \wedge \neg q)$$

Distributive rule of Disjunction over Conjunction. Distributive rule of Conjunction over Disjunction:

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

Also:  $p \rightarrow (q \wedge r) \equiv (p \rightarrow q) \wedge (p \rightarrow r)$

$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

Also:  $(p \rightarrow r) \wedge (q \rightarrow r) \equiv (p \vee q) \rightarrow r$

De Morgan's Laws:

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$\text{In general: } \neg\left(\bigwedge_{j=1}^n p_j\right) \equiv \bigvee_{j=1}^n \neg p_j$$

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

$$\text{In general: } \neg\left(\bigvee_{j=1}^n p_j\right) \equiv \bigwedge_{j=1}^n \neg p_j$$

Satisfiability:

Definition: A compound proposition is satisfiable if there is an assignment of truth values to its variables that makes it true.

Definition: An assignment of truth values that makes a compound proposition true is called a solution of the satisfiability problem.

We also can use Boolean Arithmetic to computing Truth Values

We set 0-False, 1-True and  $\oplus$ - addition modulo 2:

$$\neg P : 1 \oplus P$$

$$P \wedge Q : P \oplus Q \oplus PQ$$

$$P \vee Q : PQ$$

$$P \rightarrow Q : (1 \oplus P)Q$$

$$P \leftrightarrow Q : P \oplus Q$$

(因为布尔代数中 XOR 是完备的)