

RFC-012: Lamport Logical Clocks

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Abstract

This RFC specifies Lamport logical clocks for Cyberspace distributed ordering, providing a happened-before relation across federated vaults without synchronized physical clocks.

Motivation

Physical clocks lie:

- **Clock skew:** Machines disagree on time
- **Clock drift:** Skew grows over time
- **NTP failures:** Synchronization breaks
- **Relativity:** No global “now” anyway

Lamport clocks provide:

1. **Logical ordering:** $a \rightarrow b$ means a happened before b
2. **No synchronization:** No NTP, no GPS, no trusted time
3. **Causality tracking:** If a caused b, then $C(a) < C(b)$
4. **Consistency:** All nodes agree on partial order

From Lamport (1978):

Time, Clocks, and the Ordering of Events in a Distributed System

Specification

The Happened-Before Relation

Definition: $a \rightarrow b$ (a happened before b) if:

1. a and b are in same process and a comes before b, or
2. a is sending a message and b is receiving it, or
3. $\exists c$ such that $a \rightarrow c$ and $c \rightarrow b$ (transitivity)

Lamport Clock Rules

Each process P maintains counter $C(P)$:

Rule 1: Before each event, increment $C(P)$
 $C(P) := C(P) + 1$

Rule 2: When sending message m , attach $C(P)$
 $\text{send}(m, C(P))$

Rule 3: When receiving message (m, T) , update clock
 $C(P) := \max(C(P), T) + 1$

Clock Condition

If $a \rightarrow b$ then $C(a) < C(b)$.

Note: Converse is not guaranteed. $C(a) < C(b)$ does not imply $a \rightarrow b$.

Data Structures

Logical Timestamp

```
(define-record-type <lamport-clock>
  (make-lamport-clock counter node-id)
  lamport-clock?
  (counter clock-counter)
  (node-id clock-node-id))
```

Timestamped Event

```
(define-record-type <timestamped-event>
  (make-timestamped-event clock event)
  timestamped-event?
  (clock event-clock)
  (event event-data))
```

Operations

Increment (Local Event)

```
(define (clock-tick! clock)
  "Increment clock for local event"
```

```

(set! (clock-counter clock)
      (+ 1 (clock-counter clock)))
clock)

```

Send (Message Departure)

```

(define (clock-send clock message)
  "Attach timestamp to outgoing message"
  (clock-tick! clock)
  (make-timestamped-message
   (clock-counter clock)
   (clock-node-id clock)
   message))

```

Receive (Message Arrival)

```

(define (clock-receive! clock timestamped-message)
  "Update clock on message receipt"
  (let ((remote-time (message-timestamp timestamped-message)))
    (set! (clock-counter clock)
          (+ 1 (max (clock-counter clock) remote-time)))
    (message-payload timestamped-message)))

```

Total Ordering

Lamport clocks give partial order. For total order, break ties with node ID:

```

(define (lamport-compare a b)
  "Total order: compare by timestamp, then by node-id"
  (let ((ta (clock-counter (event-clock a)))
        (tb (clock-counter (event-clock b)))
        (na (clock-node-id (event-clock a)))
        (nb (clock-node-id (event-clock b))))
    (cond
     ((< ta tb) -1)
     ((> ta tb) 1)
     ((string<? na nb) -1)
     ((string>? na nb) 1)
     (else 0))))

```

Application to Cyberspace

Audit Trail Ordering

```
(audit-entry
  (lamport-clock 1042 "alice-vault")
  (timestamp "2026-01-06T15:30:00Z") ; Physical (advisory)
  (action (seal-commit "abc123"))
  ...)
```

Logical clock provides causal ordering even if physical clocks differ.

Federation Message Ordering

```
(federation-message
  (logical-time 573)
  (from "bob-vault")
  (type announcement)
  (payload ...))
```

Receivers update their clocks, ensuring consistent view of causality.

Conflict Detection

If events are concurrent (neither $a \rightarrow b$ nor $b \rightarrow a$):

```
(define (concurrent? a b)
  "True if neither event happened before the other"
  (let ((ta (clock-counter (event-clock a)))
        (tb (clock-counter (event-clock b))))
    ;; If timestamps equal and different nodes, concurrent
    ;; More sophisticated: use vector clocks for accurate detection
    (and (= ta tb)
          (not (equal? (clock-node-id (event-clock a))
                        (clock-node-id (event-clock b)))))))
```

Vector Clocks (Extension)

For precise concurrency detection, use vector clocks:

```
(define-record-type <vector-clock>
  (make-vector-clock entries)
  vector-clock?
  (entries vc-entries)) ; Hash: node-id → counter
```

;; a → b iff VC(a) < VC(b) (component-wise)
;; a || b (concurrent) iff neither VC(a) < VC(b) nor VC(b) < VC(a)

Trade-off: $O(N)$ space vs $O(1)$ for Lamport clocks.

Integration Points

With Audit Trail (RFC-003)

```
(audit-entry
  (sequence 42)           ; Local sequence
  (lamport-clock 1042)    ; Logical timestamp
  (physical-time "...")  ; Advisory only
  ...)
```

With Consensus (RFC-011)

```
(consensus-message
  (sequence 573)          ; Consensus sequence
  (lamport-clock 2891)    ; Causal context
  ...)
```

With Replication (RFC-001)

```
(seal-publish "2.0.0"
  (logical-time 892)
  ...)
```

Security Considerations

Clock Manipulation

A Byzantine node could: – Report artificially high clock values – Attempt to order events incorrectly

Mitigations: – Clock values bounded by received values + margin – Signatures prevent retroactive tampering – Audit trails detect anomalies

Denial of Service

Flooding with high-timestamped messages forces clock advancement.

Mitigation: Rate limiting, reputation systems.

Implementation Notes

Thread Safety

Clock updates must be atomic:

```
(define (atomic-tick! clock)
  (mutex-lock! clock-mutex)
  (clock-tick! clock)
  (mutex-unlock! clock-mutex))
```

Persistence

Clock state survives restarts:

```
(define (save-clock clock path)
  (with-output-to-file path
    (lambda ()
      (write `(lamport-clock
                (counter ,(clock-counter clock))
                (node-id ,(clock-node-id clock)))))))
```

References

1. Lamport, L. (1978). Time, Clocks, and the Ordering of Events in a Distributed System.
 2. Fidge, C. J. (1988). Timestamps in Message-Passing Systems.
 3. Mattern, F. (1989). Virtual Time and Global States of Distributed Systems.
 4. RFC-003: Cryptographic Audit Trail
 5. RFC-010: Federation Protocol
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Changelog

- 2026-01-06 – Initial specification

Implementation Status: Proposed **Clock Type:** Lamport scalar, vector clock extension **Ordering:** Partial (causal), total with tie-breaking