AVL Tree

Outline

Background

Define height balancing

Maintaining balance within a tree

- AVL trees
- Difference of heights
- Maintaining balance after insertions and erases

Background

From previous lectures:

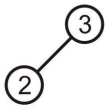
- Binary search trees store linearly ordered data
- Best case height: $\mathbf{O}(\ln(n))$
- Worst case height: $\mathbf{O}(n)$

Requirement:

- Define and maintain a *balance* to ensure $\mathbf{O}(\ln(n))$ height

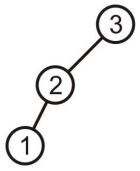
Prototypical Examples

These two examples demonstrate how we can correct for imbalances: starting with this tree, add 1:



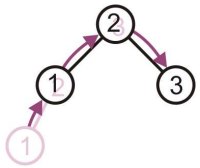
Prototypical Examples

This is more like a linked list; however, we can fix this…



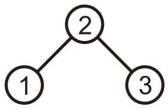
Prototypical Examples

Promote 2 to the root, demote 3 to be 2's right child, and 1 remains the left child of 2



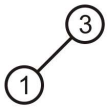
Prototypical Examples

The result is a perfect, though trivial tree



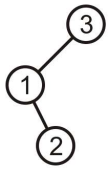
Prototypical Examples

Alternatively, given this tree, insert 2



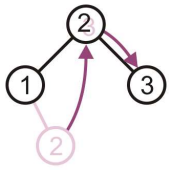
Prototypical Examples

Again, the product is a linked list; however, we can fix this, too



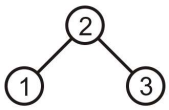
Prototypical Examples

Promote 2 to the root, and assign 1 and 3 to be its children



Prototypical Examples

The result is, again, a perfect tree



These examples may seem trivial, but they are the basis for the corrections in the next data structure we will see: AVL trees

AVL Trees

We will focus on the first strategy: AVL trees

- Named after Adelson-Velskii and Landis

Balance is defined by comparing the height of the two sub-trees

Recall:

- An empty tree has height -1
- A tree with a single node has height 0

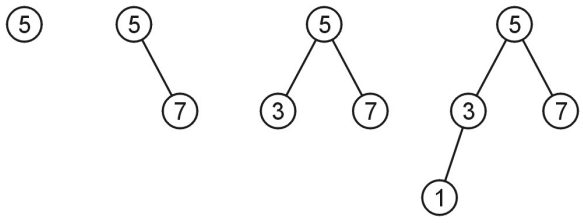
AVL Trees

A binary search tree is said to be AVL balanced if:

- The difference in the heights between the left and right sub-trees is at most 1, and
- Both sub-trees are themselves AVL trees

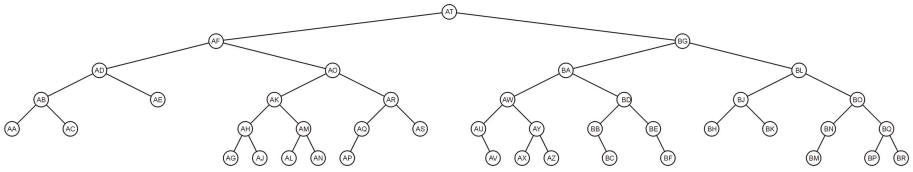
AVL Trees

AVL trees with 1, 2, 3, and 4 nodes:



AVL Trees

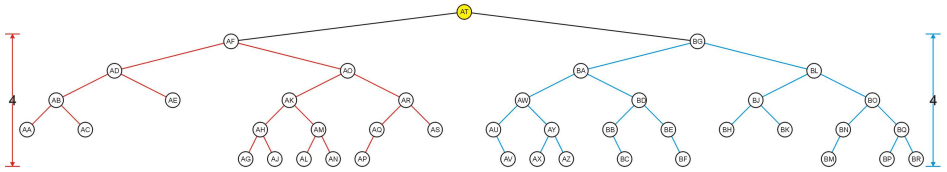
Here is a larger AVL tree (42 nodes):



AVL Trees

The root node is AVL-balanced:

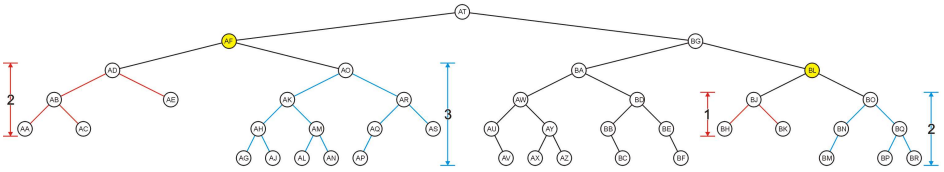
- Both sub-trees are of height 4:



AVL Trees

All other nodes (*e.g.*, AF and BL) are AVL balanced

– The sub-trees differ in height by at most one



Height of an AVL Tree

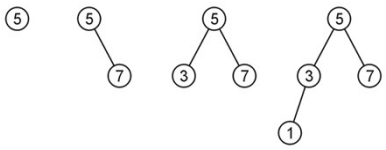
Let $F(h)$ be the fewest number of nodes in a tree of height h

From a previous slide:

$F(0) = 1$

$F(1) = 2$

$F(2) = 4$



Can we find $F(h)$?

Height of an AVL Tree

The worst-case AVL tree of height h would have:

- A worst-case AVL tree of height $h - 1$ on one side,
- A worst-case AVL tree of height $h - 2$ on the other, and
- The root node

We get: $\mathrm{F}(h) = \mathrm{F}(h-1) + 1 + \mathrm{F}(h-2)$

Height of an AVL Tree

This is a recurrence relation:

$$\mathrm{F}(h) = \begin{cases} 1 & h = 0 \\ 2 & h = 1 \\ \mathrm{F}(h-1) + \mathrm{F}(h-2) + 1 & h > 1 \end{cases}$$

The solution?

- Note that $F(h) + 1 = (F(h-1) + 1) + (F(h-2) + 1)$
- Therefore, $F(h) + 1$ is a Fibonacci number:

$$\begin{aligned} F(0) + 1 &= 2 \quad \rightarrow \quad F(0) = 1 \\ F(1) + 1 &= 3 \quad \rightarrow \quad F(1) = 2 \\ F(2) + 1 &= 5 \quad \rightarrow \quad F(2) = 4 \\ F(3) + 1 &= 8 \quad \rightarrow \quad F(3) = 7 \\ F(4) + 1 &= 13 \quad \rightarrow \quad F(4) = 12 \\ F(5) + 1 &= 21 \quad \rightarrow \quad F(5) = 20 \\ F(6) + 1 &= 34 \quad \rightarrow \quad F(6) = 33 \end{aligned}$$

Height of an AVL Tree

This is approximately

$$F(h) \approx 1.8944\ \phi^h - 1 \approx O(\phi^h)$$

where $\phi \approx 1.6180$ is the golden ratio

Given a number of nodes N, the maximum value of h of AVL Tree will be:

$1.8944\ \phi^h - 1 = N$ => h = 1.4404*log(N+1)= O(logN)

Operations in AVL Tree

Searching, Complexity? O(log N)

FindMin, Complexity? O(log N)

Deletion? Insertion?

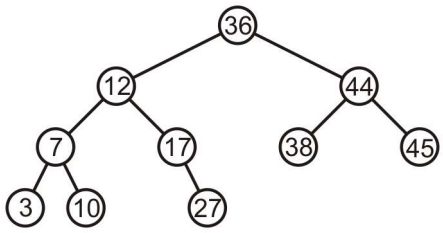
Maintaining Balance

To maintain AVL balance, observe that:

- Inserting a node can increase the height of a tree by at most 1
- Removing a node can decrease the height of a tree by at most 1

Maintaining Balance

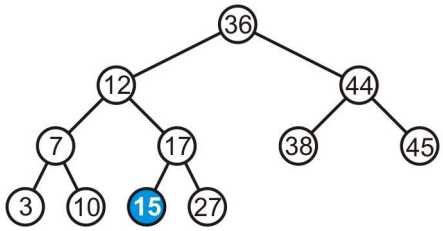
Consider this AVL tree



Maintaining Balance

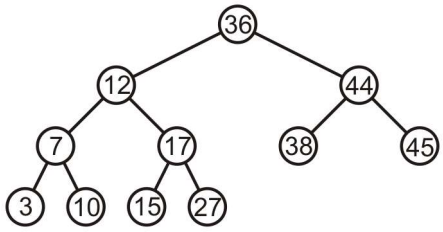
Consider inserting 15 into this tree

– In this case, the heights of none of the trees change



Maintaining Balance

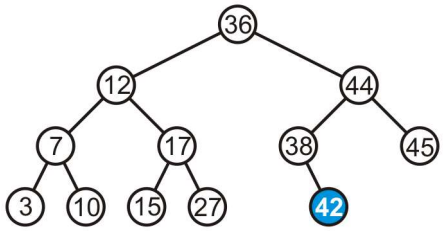
The tree remains balanced



Maintaining Balance

Consider inserting 42 into this tree

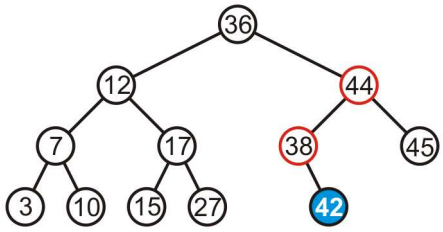
- In this case, the heights of none of the trees change



Maintaining Balance

Consider inserting 42 into this tree

- Now we see the heights of two sub-trees have increased by one
- The tree is still balanced



Maintaining Balance

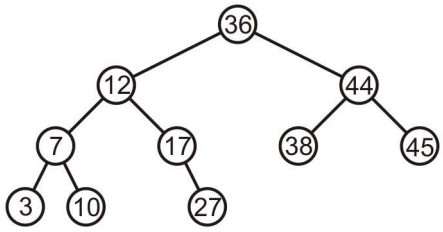
Only insert and erase may change the height

- This is the only place we need to update the height
- These algorithms are already recursive

Maintaining Balance

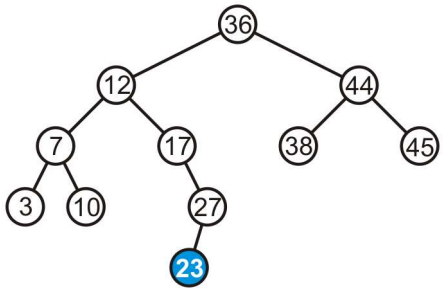
If a tree is AVL balanced, for an insertion to cause an imbalance:

- The heights of the sub-trees must differ by 1
- The insertion must increase the height of the deeper sub-tree by 1



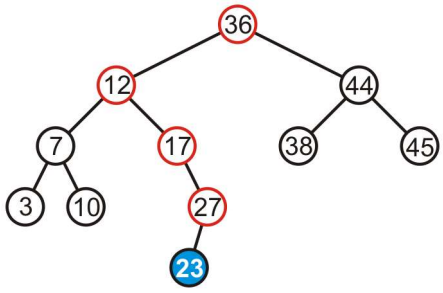
Maintaining Balance

Suppose we insert 23 into our initial tree



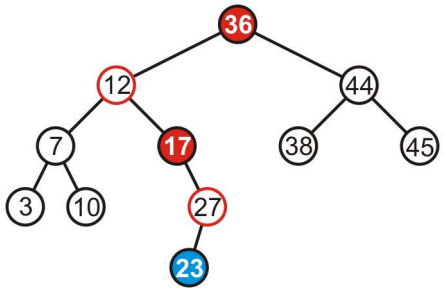
Maintaining Balance

The heights of each of the sub-trees from here to the root are increased by one



Maintaining Balance

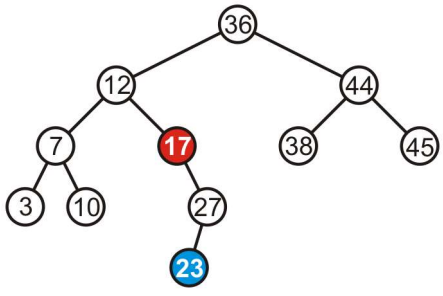
However, only two of the nodes are unbalanced: 17 and 36



Maintaining Balance

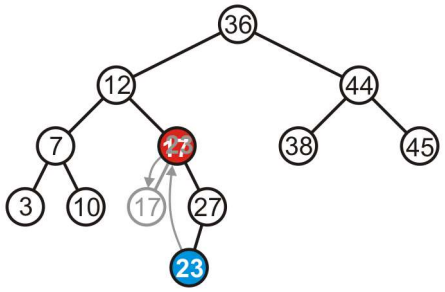
However, only two of the nodes are unbalanced: 17 and 36

- We only have to fix the imbalance at the lowest node



Maintaining Balance

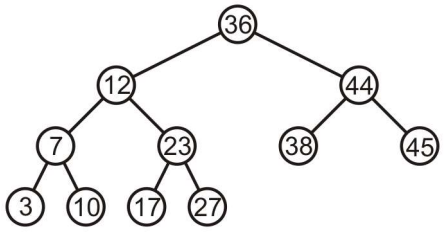
We can promote 23 to where 17 is, and make 17 the left child of 23



Maintaining Balance

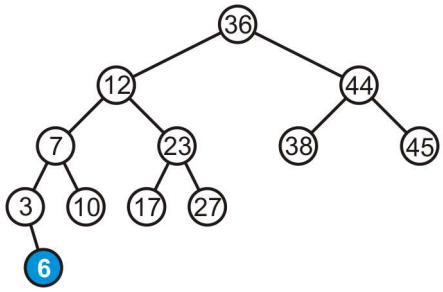
Thus, that node is no longer unbalanced

- Incidentally, neither is the root now balanced again, too



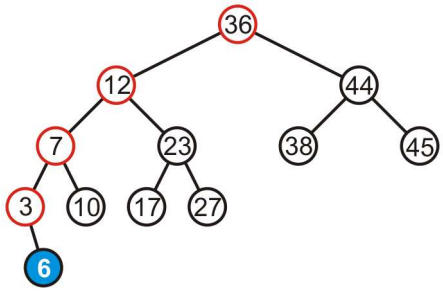
Maintaining Balance

Consider adding 6:



Maintaining Balance

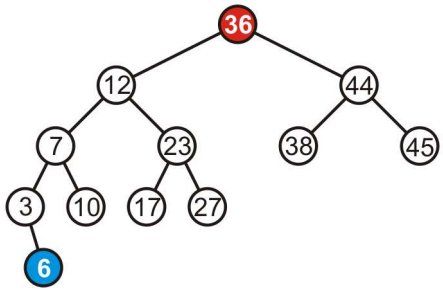
The height of each of the trees in the path back to the root are increased by one



Maintaining Balance

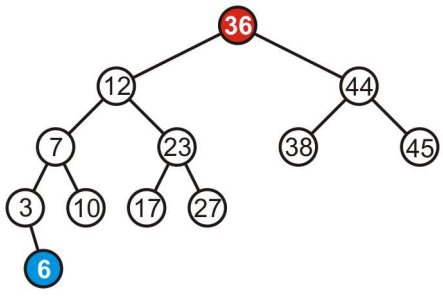
The height of each of the trees in the path back to the root are increased by one

- However, only the root node is now unbalanced



Maintaining Balance

To fix this, we will look at the general case…



Maintaining Balance

If an insertion cause an imbalance, which nodes can be affected?

Nodes on the path of the inserted node.

Let f be the node that must rebalanced.

Case 1: insertion in the left subtree of the left child of f

Case 2: insertion in the right subtree of the left child of f

Case 3: insertion in the left subtree of the right child of f

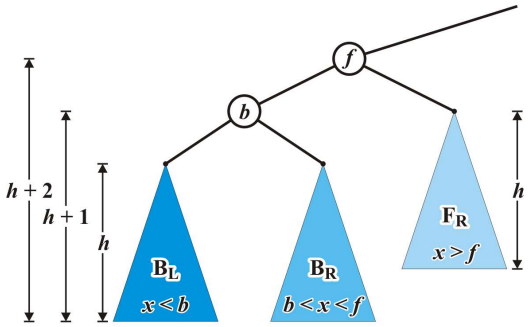
Case 4: insertion in the right subtree of the right child of f

Case 1 and 4 (or Case 2 and 3) are the same. We only solve case 1 and 2

Maintaining Balance: Case 1 Single Rotation

Consider the following setup

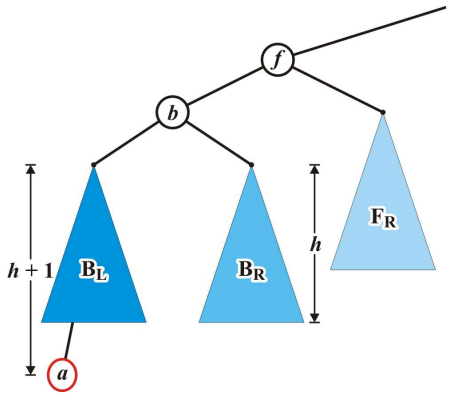
– Each blue triangle represents a tree of height h



Maintaining Balance: Case 1
Single Rotation

Insert a into this tree: it falls into the left subtree $\mathbf{B_L}$ of b $(a<b<f)$

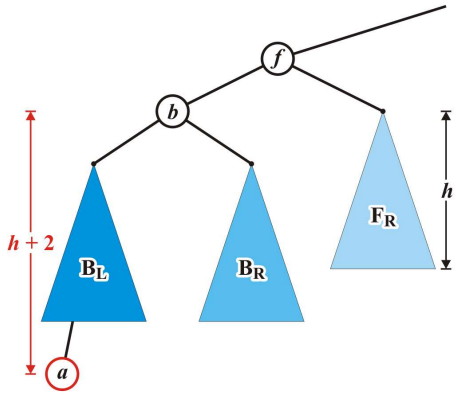
- Assume $\mathbf{B_L}$ remains balanced
- Thus, the tree rooted at b is also balanced



Maintaining Balance: Case 1
Single Rotation

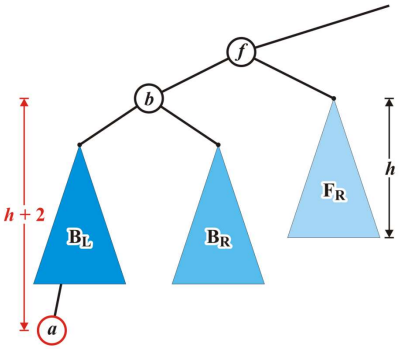
The tree rooted at node f is now unbalanced

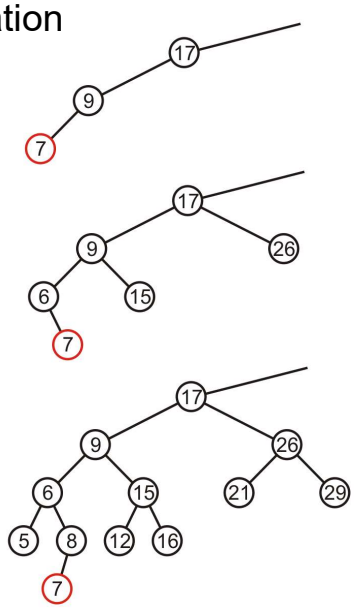
- We will correct the imbalance at this node



Maintaining Balance: Case 1
Single Rotation

Here are examples of when the insertion of 7 may cause this situation when $h = -1, 0,$ and 1

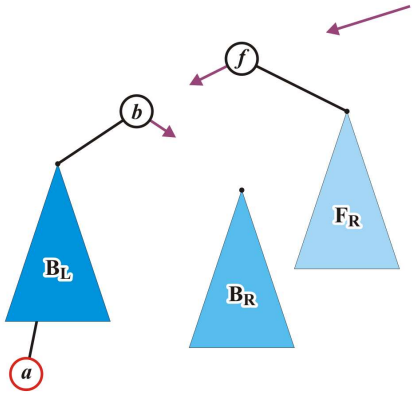




Maintaining Balance: Case 1
Single Rotation

We will modify these three pointers

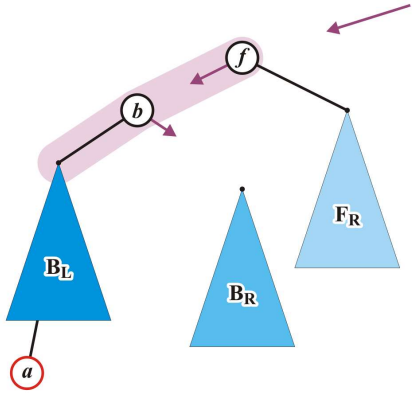
– At this point, `this` references the unbalanced root node *f*



Maintaining Balance: Case 1 Single Rotation

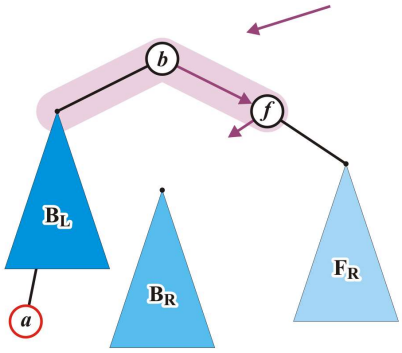
Specifically, we will rotate these two nodes around the root:

- Recall the first prototypical example
- Promote node ***b*** to the root and demote node ***f*** to be the right child of ***b***



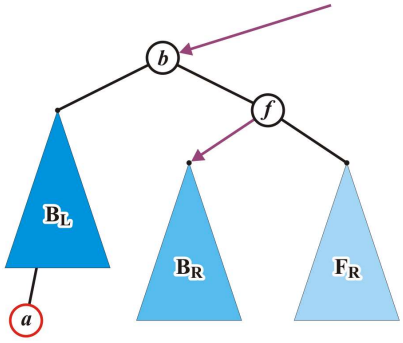
Maintaining Balance: Case 1

This requires the address of node ***f*** to be assigned to the `right tree` member variable of node ***b***



Maintaining Balance: Case 1

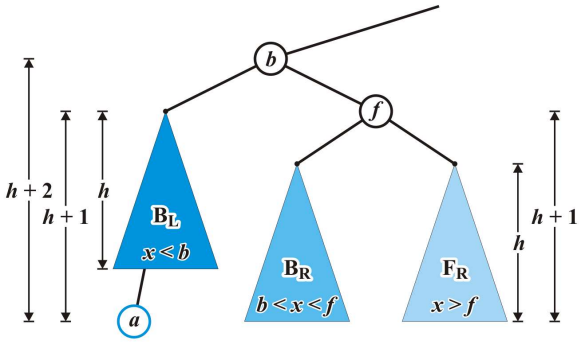
Assign any former parent of node f to the address of node b
Assign the address of the tree $\mathbf{B_R}$ to `left tree` of node f



Maintaining Balance: Case 1
Single Rotation

The nodes b and f are now balanced and all remaining nodes of the subtrees are in their correct positions

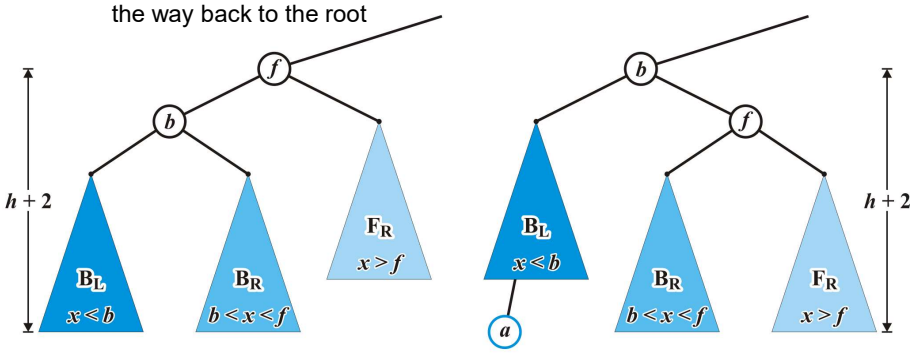
– The height of f is now $h + 1$ while b remains at height $h + 2$



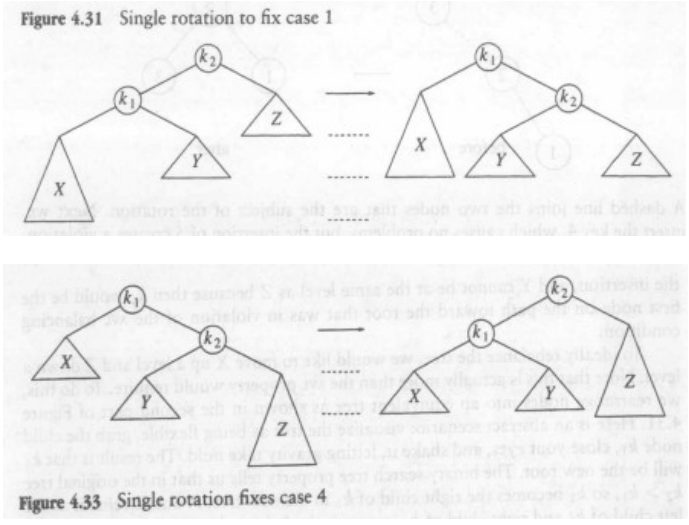
Maintaining Balance: Case 1
Single Rotation

Additionally, height of the corrected tree rooted at b equals the original height of the tree rooted at f

– Thus, this insertion will no longer affect the balance of any ancestors all the way back to the root

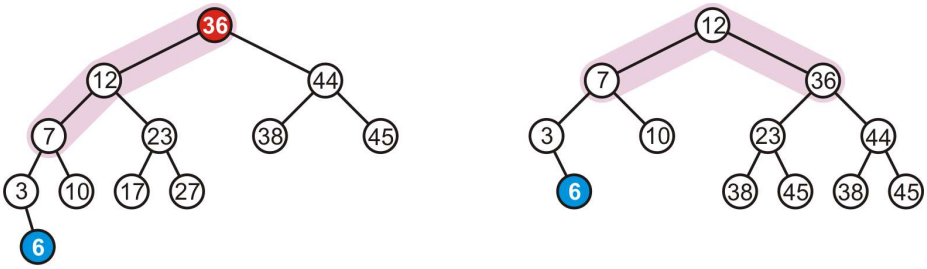


Summary: Case 1&4
Single Rotation



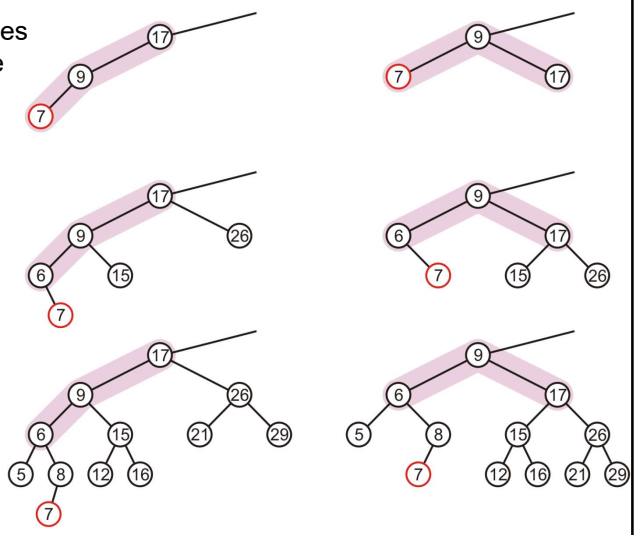
Maintaining Balance: Case 1
Single Rotation

In our example case, the correction is



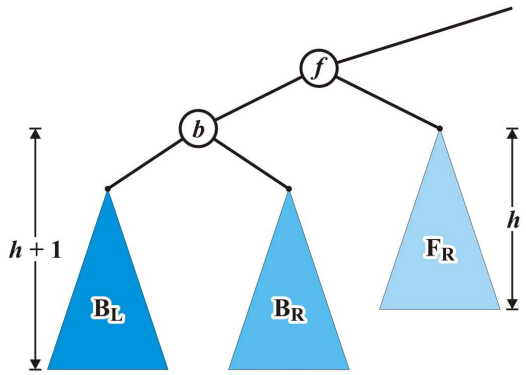
Maintaining Balance: Case 1

In our three sample cases with $h = -1, 0$, and 1, the node is now balanced and the same height as the tree before the insertion



Maintaining Balance: Case 2 Double Rotation

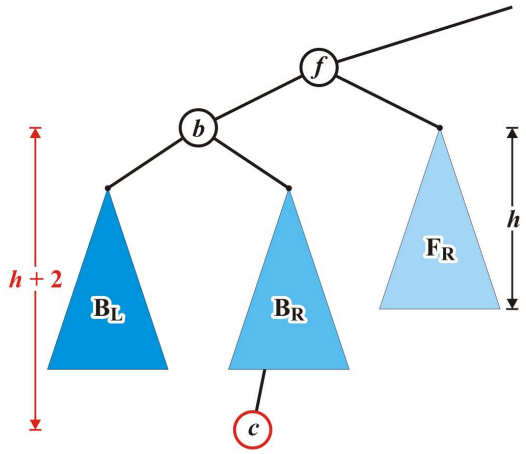
Alternatively, consider the insertion of c where $b < c < f$ into our original tree



Maintaining Balance: Case 2

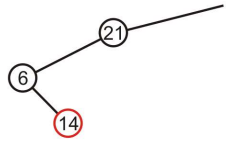
Assume that the insertion of c increases the height of $\mathbf{B_R}$

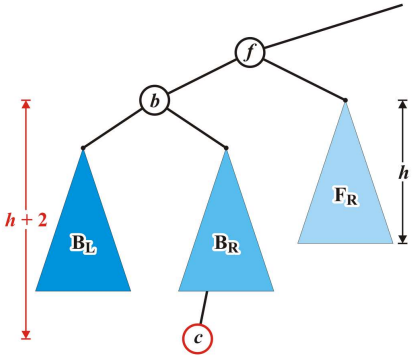
- Once again, f becomes unbalanced

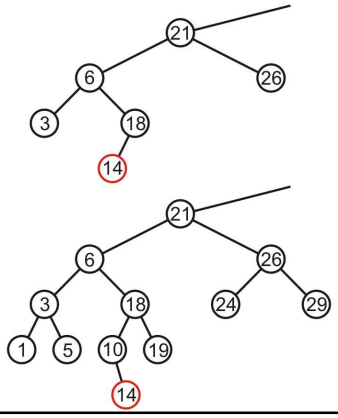


Maintaining Balance: Case 2

Here are examples of when the insertion of 14 may cause this situation when $h = -1, 0$, and 1

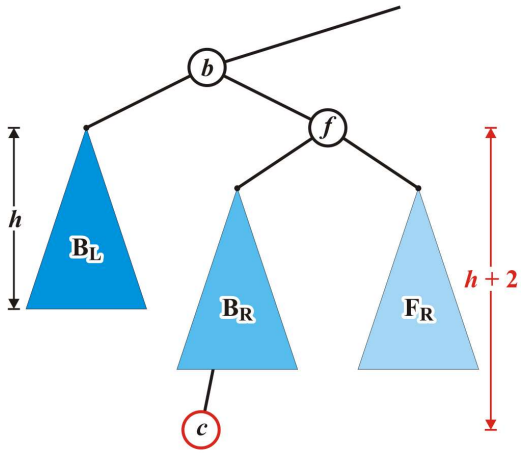






Maintaining Balance: Case 2

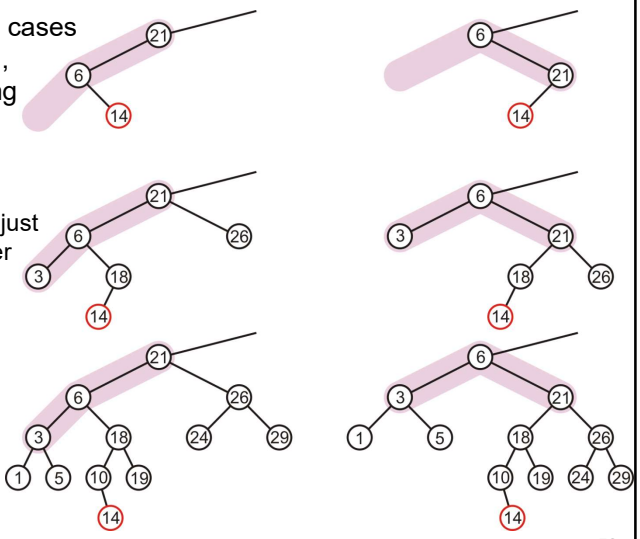
Unfortunately, the previous correction does not fix the imbalance at the root of this sub-tree: the new root, ***b***, remains unbalanced



Maintaining Balance: Case 2

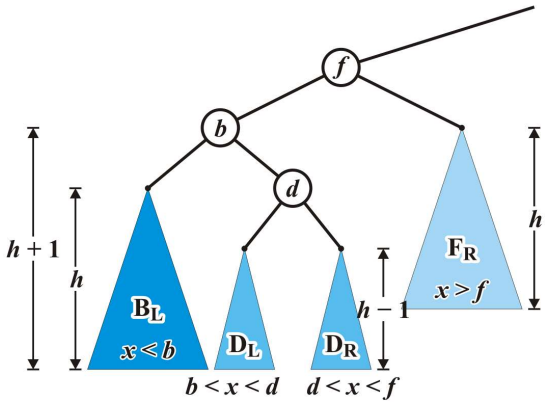
In our three sample cases with $h = -1, 0$, and 1, doing the same thing as before results in a tree that is still unbalanced…

- The imbalance is just shifted to the other side



Maintaining Balance: Case 2 Double Rotation

Re-label the tree by dividing the left subtree of f into a tree rooted at d with two subtrees of height $h - 1$

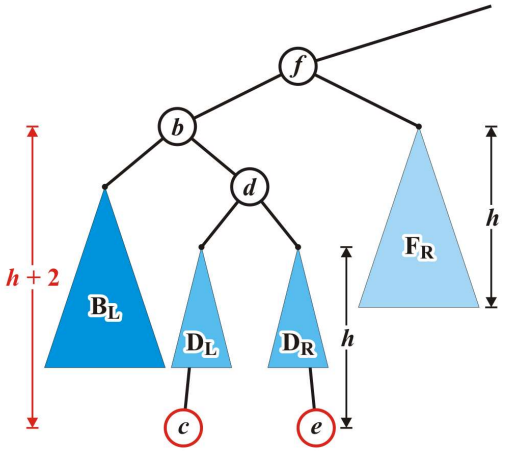


AVL trees

Maintaining Balance: Case 2
Double Rotation

Now an insertion causes an imbalance at *f*

– **The addition of either *c* or *e* will cause this**

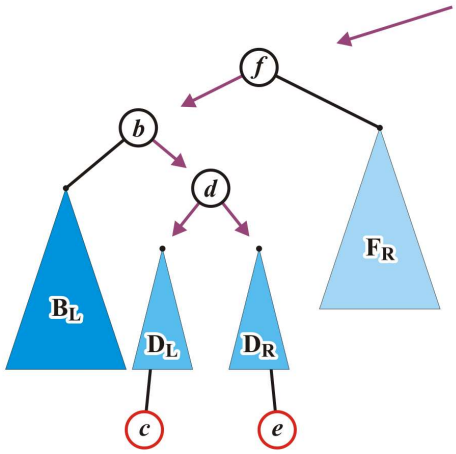


61

AVL trees

Maintaining Balance: Case 2
Double Rotation

We will reassign the following pointers

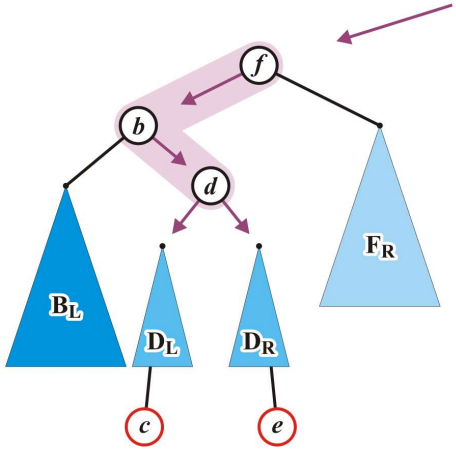


62

Maintaining Balance: Case 2
Double Rotation

Specifically, we will order these three nodes as a perfect tree

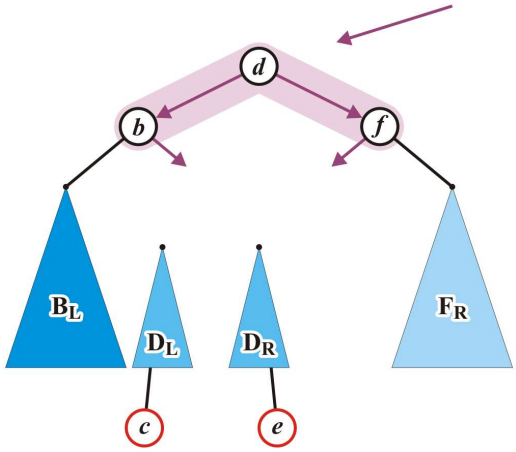
– Recall the second prototypical example



63

Maintaining Balance: Case 2
Double Rotation

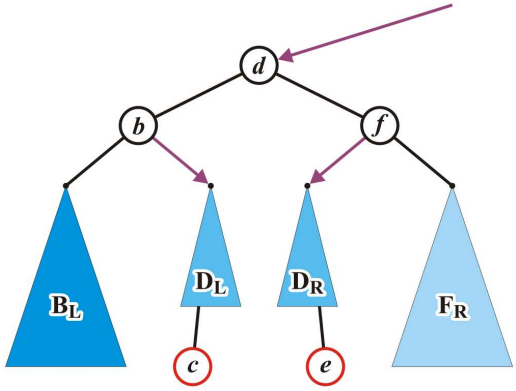
To achieve this, ***b*** and ***f*** will be assigned as children of the new root ***d***



64

Maintaining Balance: Case 2
Double Rotation

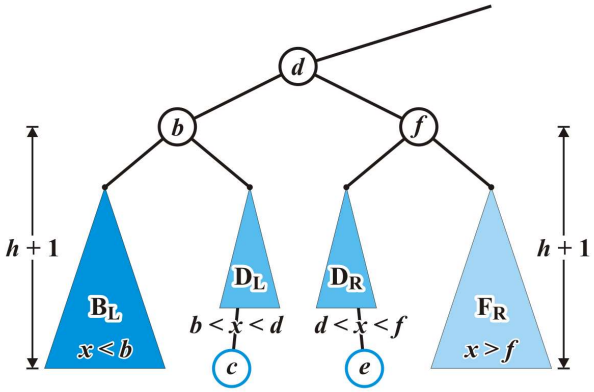
We also have to connect the two subtrees and original parent of f



Maintaining Balance: Case 2
Double Rotation

Now the tree rooted at d is balanced

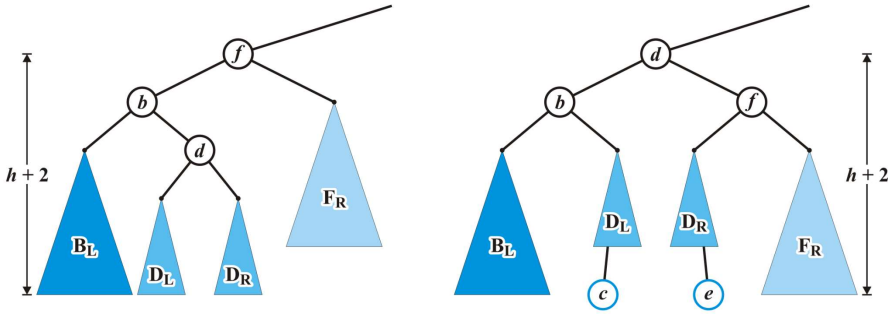
– After the correction, height of b and f become $h + 1$ and d is $h + 2$



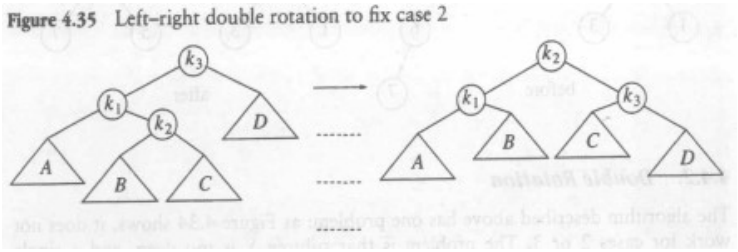
Maintaining Balance: Case 2

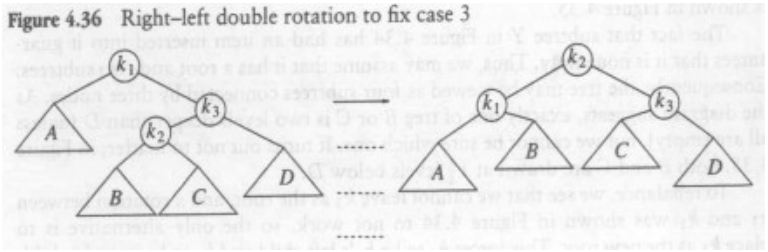
Again, the height of the root did not change

- The heights of all three nodes changed in this process



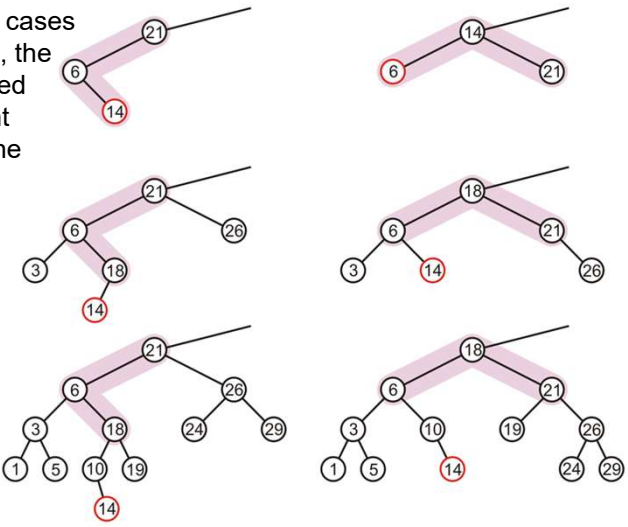
Summary: Case 2 and 3 Double Rotation





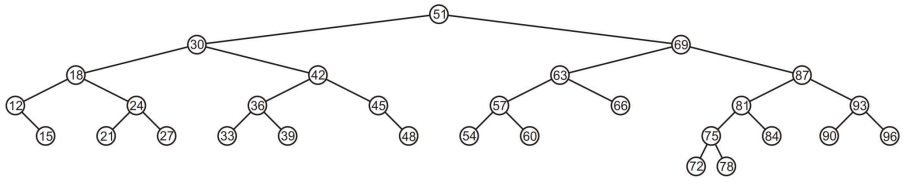
Maintaining Balance: Case 2

In our three sample cases with $h = -1, 0$, and 1, the node is now balanced and the same height as the tree before the insertion



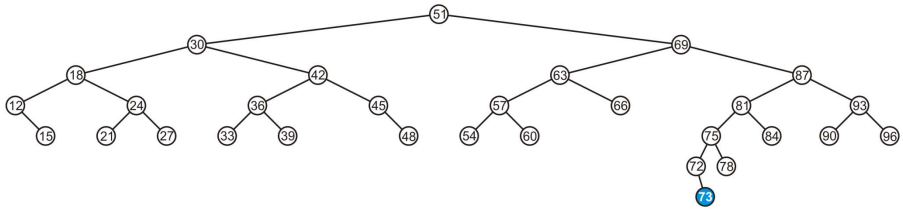
Insertion

Consider this AVL tree



Insertion

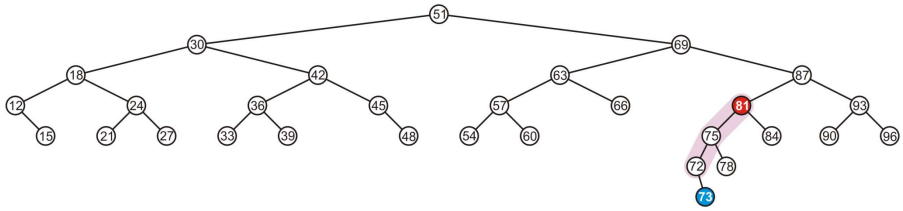
Insert 73



Insertion

The node 81 is unbalanced

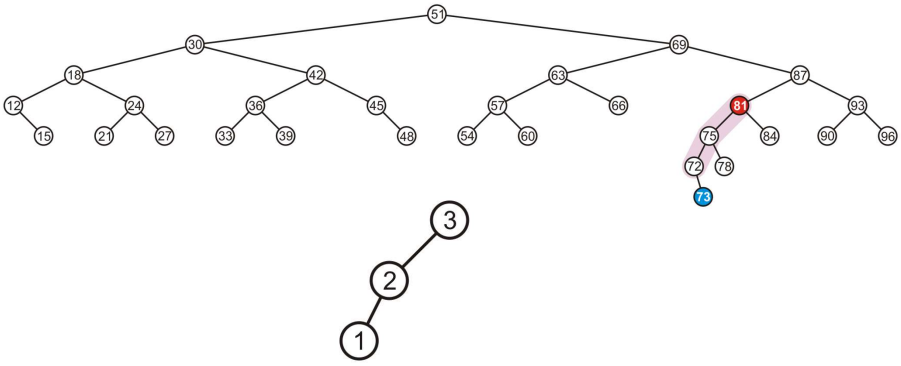
- A left-left imbalance



Insertion

The node 81 is unbalanced

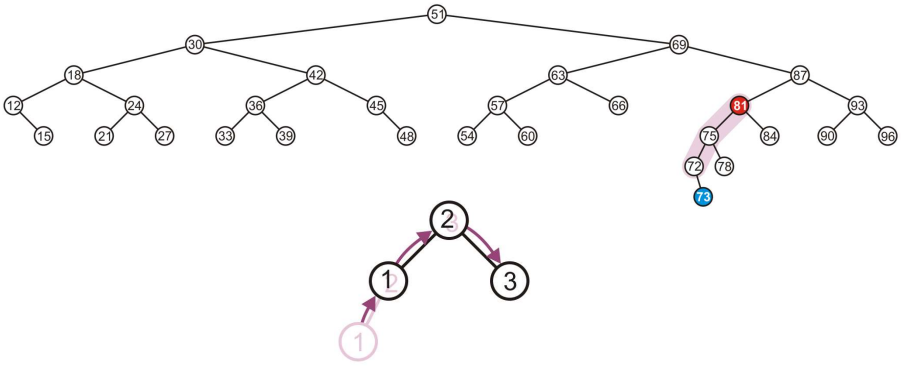
- A left-left imbalance



Insertion

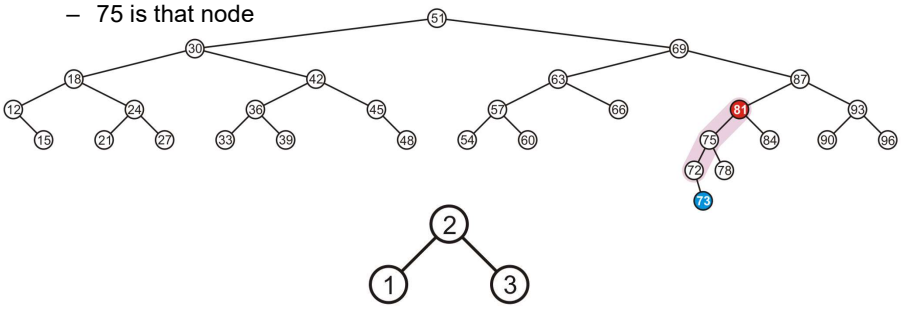
The node 81 is unbalanced

- A left-left imbalance
- Promote the intermediate node to the imbalanced node



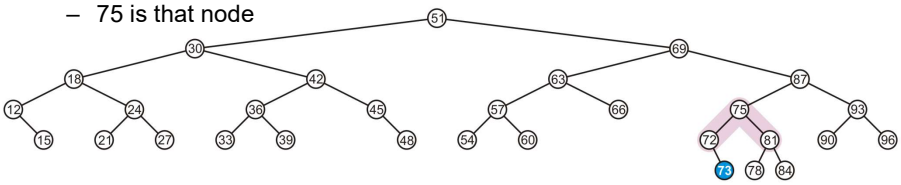
Insertion

The node 81 is unbalanced
- A left-left imbalance
- Promote the intermediate node to the imbalanced node
- 75 is that node



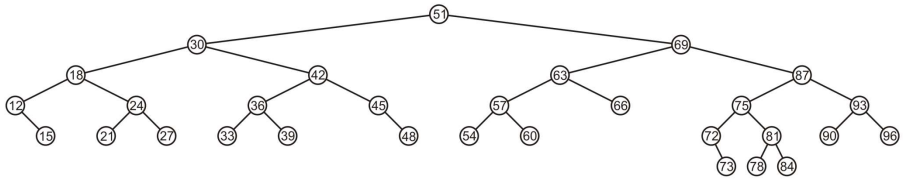
Insertion

The node 81 is unbalanced
- A left-left imbalance
- Promote the intermediate node to the imbalanced node
- 75 is that node



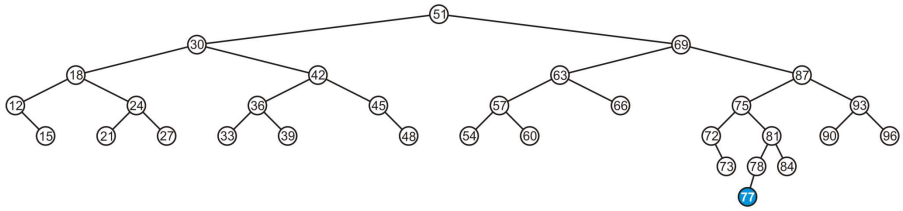
Insertion

The tree is AVL balanced



Insertion

Insert 77

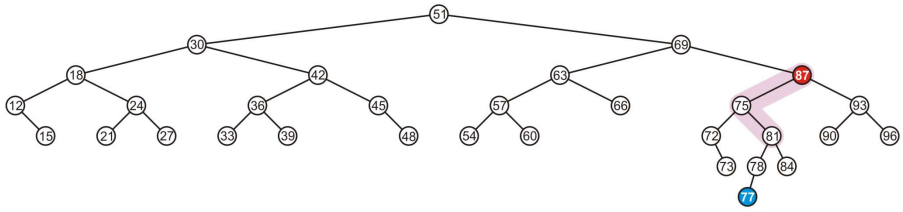


AVL trees

Insertion

The node 87 is unbalanced

- A left-right imbalance



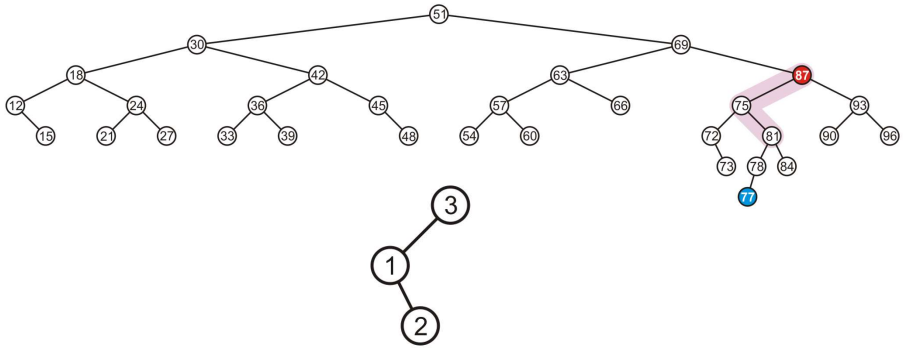
79

AVL trees

Insertion

The node 87 is unbalanced

- A left-right imbalance



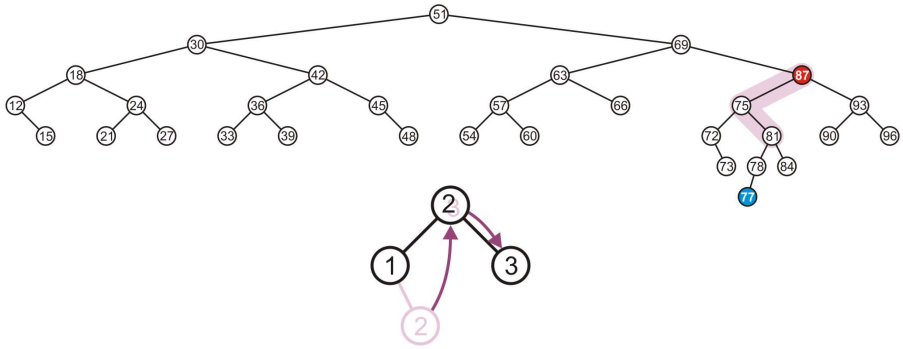
80

AVL trees

Insertion

The node 87 is unbalanced

- A left-right imbalance
- Promote the intermediate node to the imbalanced node



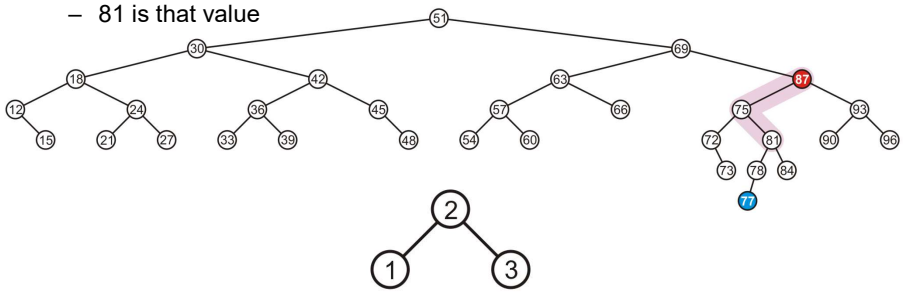
81

AVL trees

Insertion

The node 87 is unbalanced

- A left-right imbalance
- Promote the intermediate node to the imbalanced node
- 81 is that value



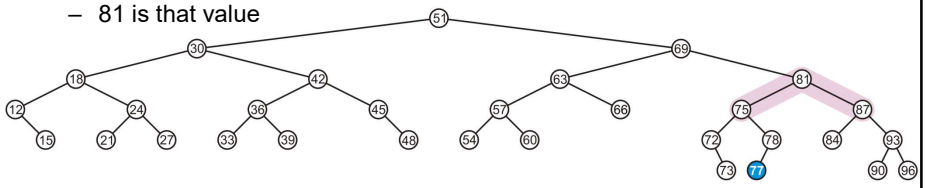
82

AVL trees

Insertion

The node 87 is unbalanced

- A left-right imbalance
- Promote the intermediate node to the imbalanced node
- 81 is that value

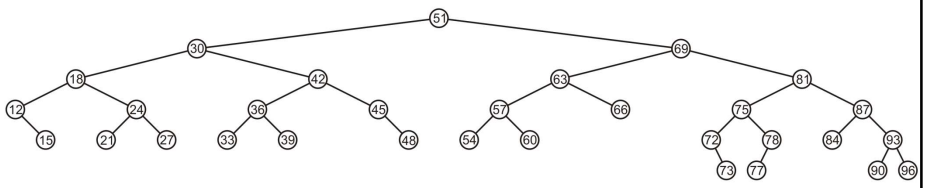


83

AVL trees

Insertion

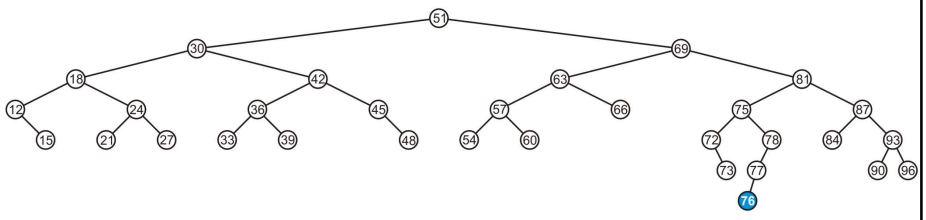
The tree is balanced



84

Insertion

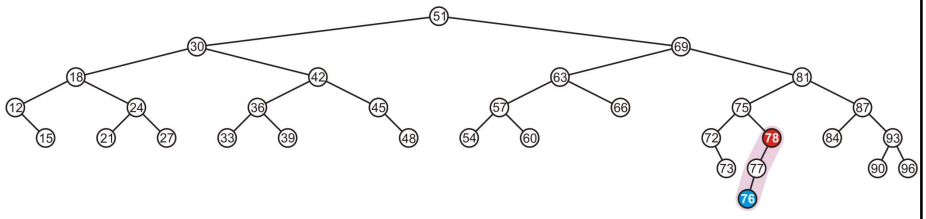
Insert 76



Insertion

The node 78 is unbalanced

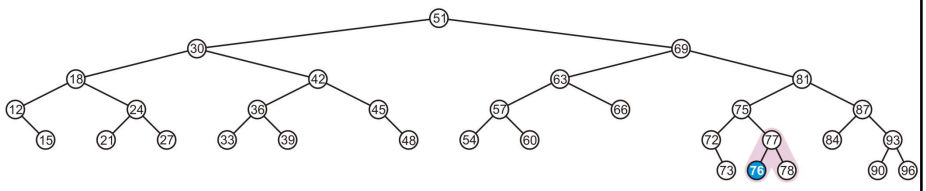
- A left-left imbalance



Insertion

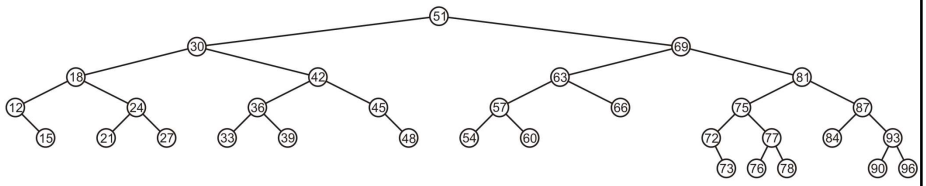
The node 78 is unbalanced

– Promote 77



Insertion

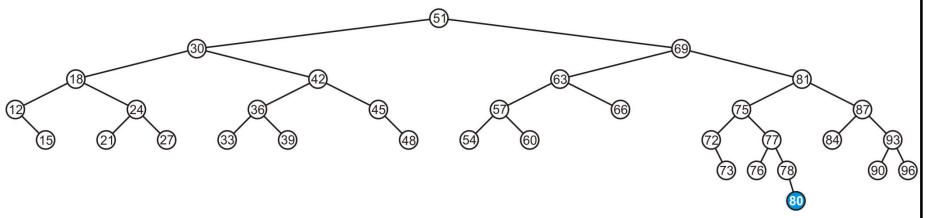
Again, balanced



AVL trees

Insertion

Insert 80



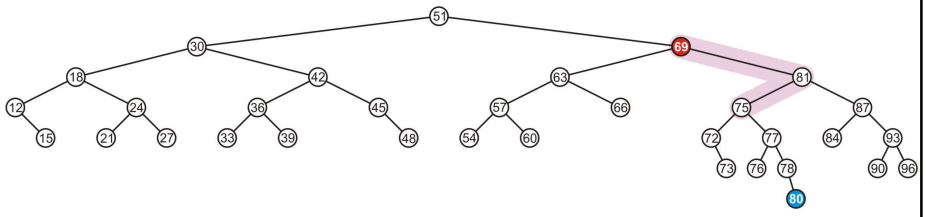
89

AVL trees

Insertion

The node 69 is unbalanced

- A right-left imbalance
- Promote the intermediate node to the imbalanced node

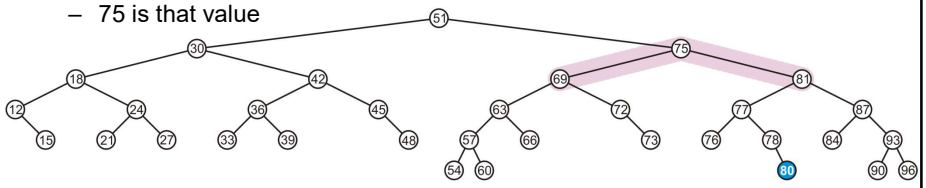


90

Insertion

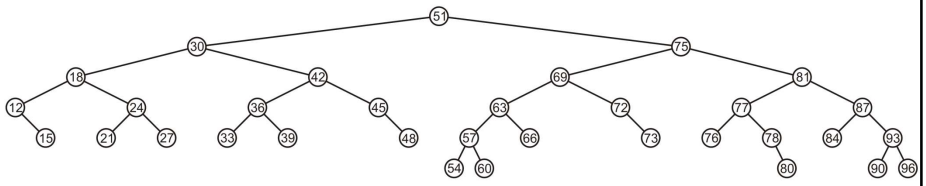
The node 69 is unbalanced

- A left-right imbalance
- Promote the intermediate node to the imbalanced node
- 75 is that value



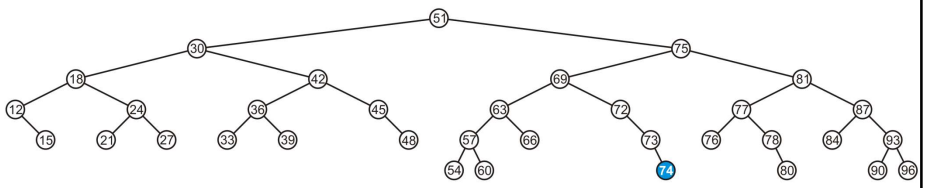
Insertion

Again, balanced



Insertion

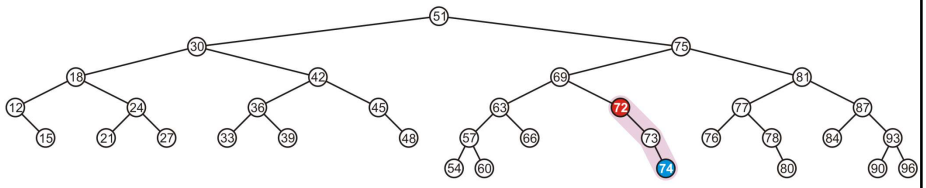
Insert 74



Insertion

The node 72 is unbalanced

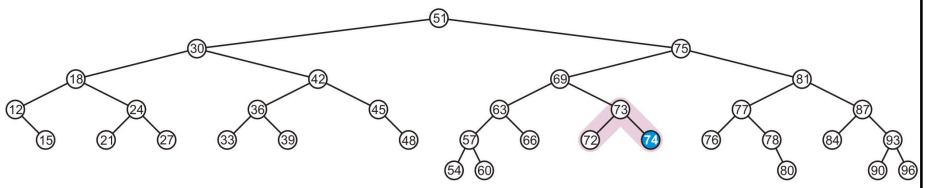
- A right-right imbalance
- Promote the intermediate node to the imbalanced node



Insertion

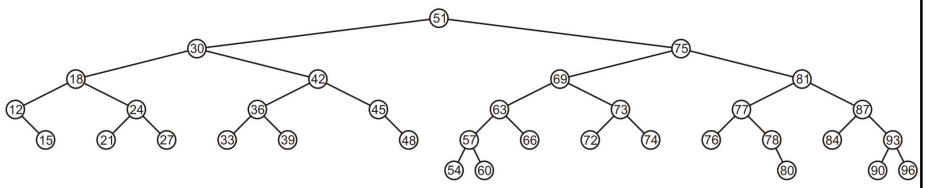
The node 72 is unbalanced

- A right-right imbalance
- Promote the intermediate node to the imbalanced node



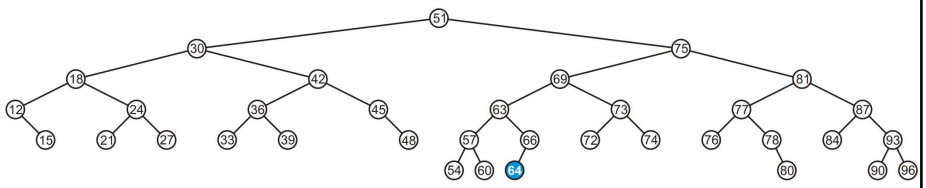
Insertion

Again, balanced



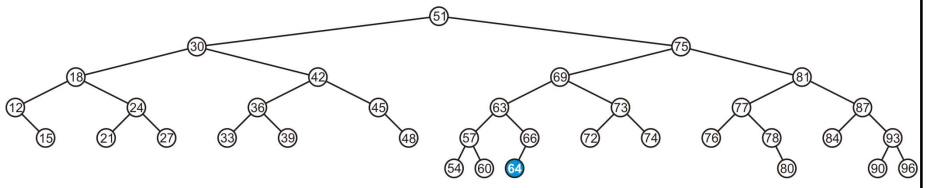
Insertion

Insert 64



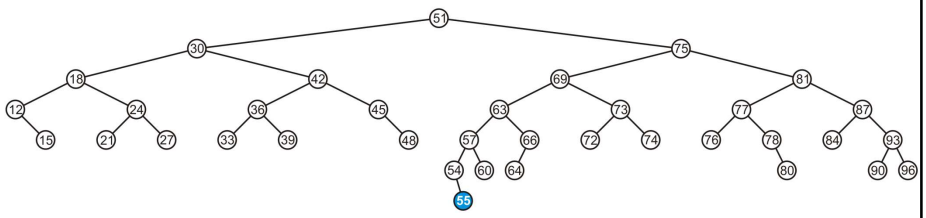
Insertion

This causes no imbalances



Insertion

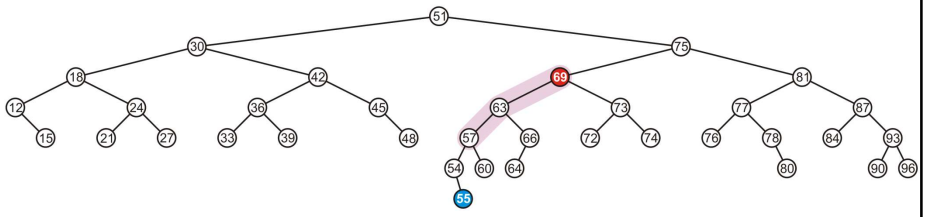
Insert 55



Insertion

The node 69 is imbalanced

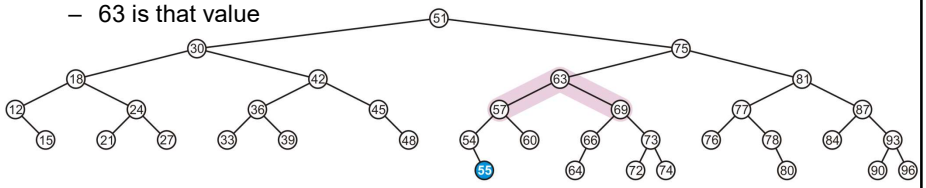
- A left-left imbalance
- Promote the intermediate node to the imbalanced node



Insertion

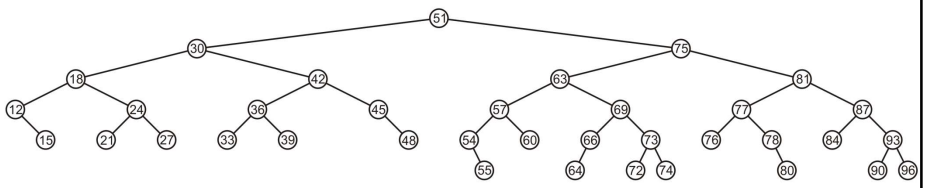
The node 69 is imbalanced

- A left-left imbalance
- Promote the intermediate node to the imbalanced node
- 63 is that value



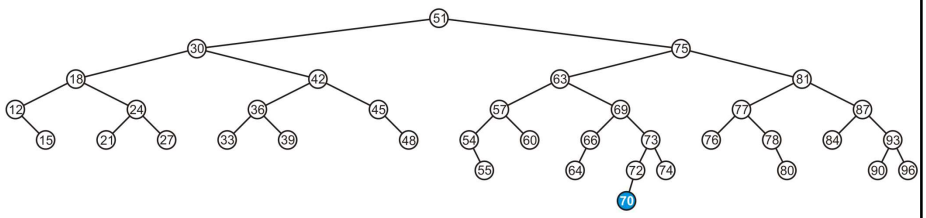
Insertion

The tree is now balanced



Insertion

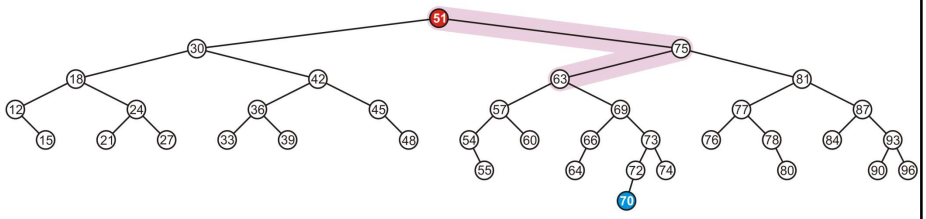
Insert 70



Insertion

The root node is now imbalanced

- A right-left imbalance
- Promote the intermediate node to the root

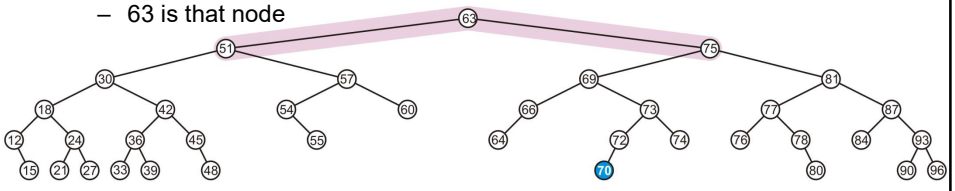


AVL trees

Insertion

The root node is imbalanced

- A right-left imbalance
- Promote the intermediate node to the root
- 63 is that node

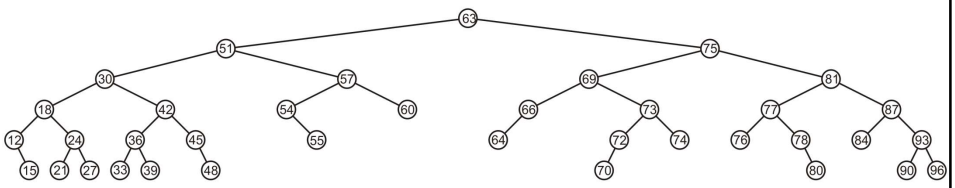


105

AVL trees

Insertion

The result is AVL balanced



106

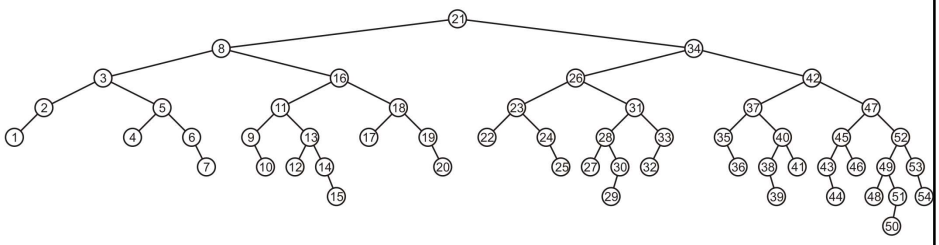
Erase

Removing a node from an AVL tree may cause more than one AVL imbalance

- Like insert, erase must check after it has been successfully called on a child to see if it caused an imbalance
- Unfortunately, it may cause many imbalances that must be corrected
  - Insertions will only cause one imbalance that must be fixed

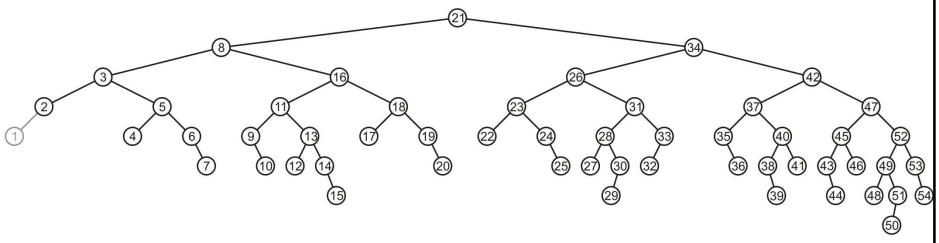
Erase

Consider the following AVL tree



Erase

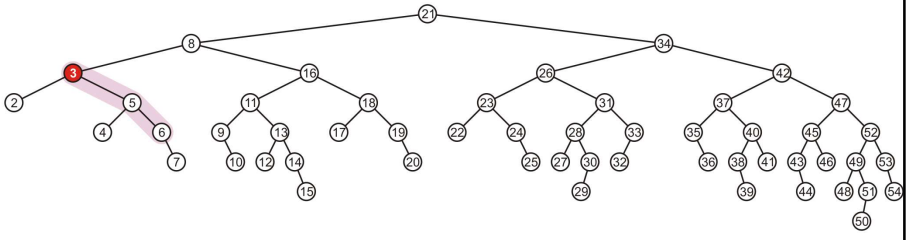
Suppose we erase the front node: 1



Erase

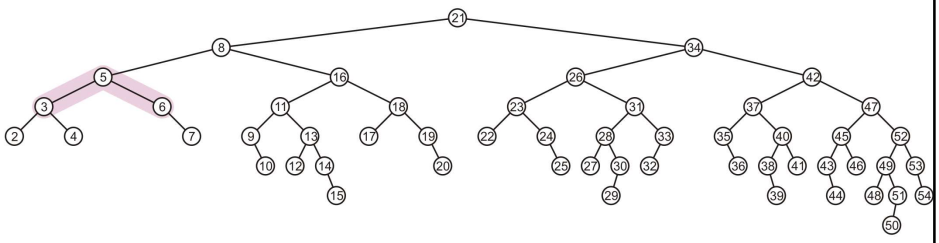
While its previous parent, 2, is not unbalanced, its grandparent 3 is

- The imbalance is in the right-right subtree



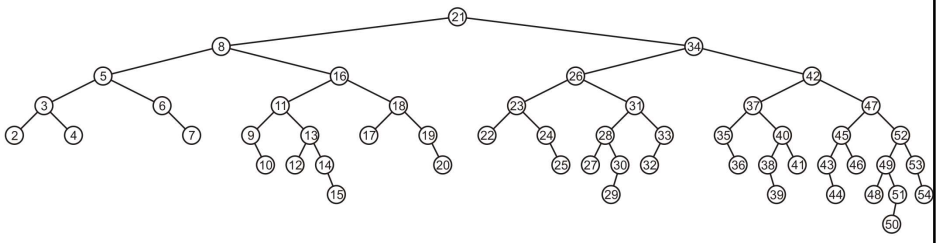
Erase

We can correct this with a simple balance



Erase

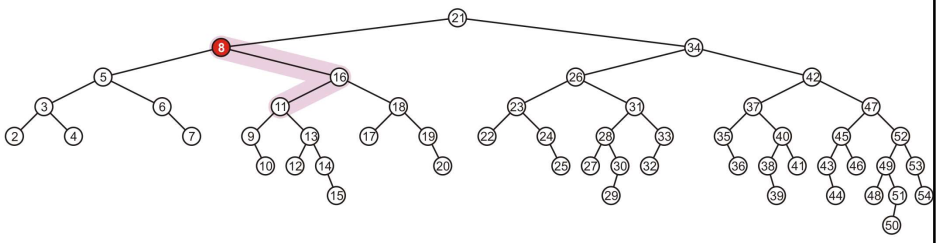
The node of that subtree, 5, is now balanced



Erase

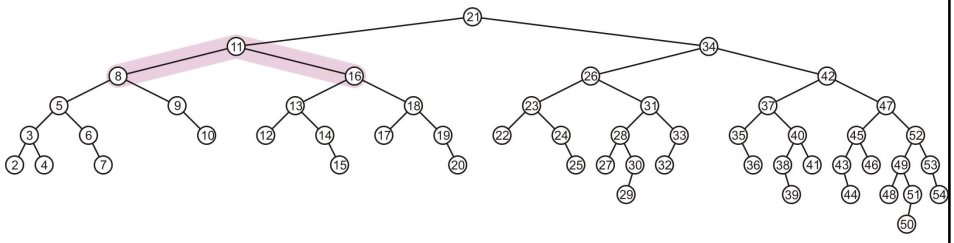
Recursing to the root, however, 8 is also unbalanced

- This is a right-left imbalance



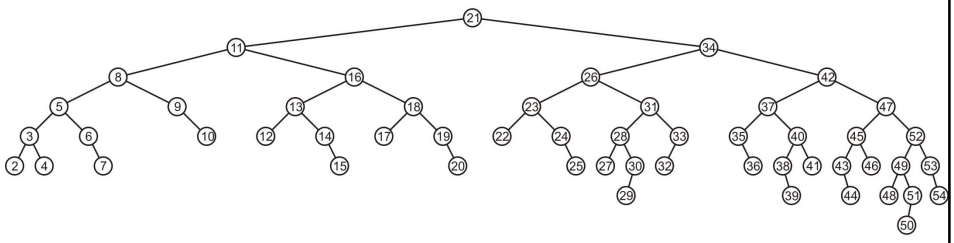
Erase

Promoting 11 to the root corrects the imbalance



Erase

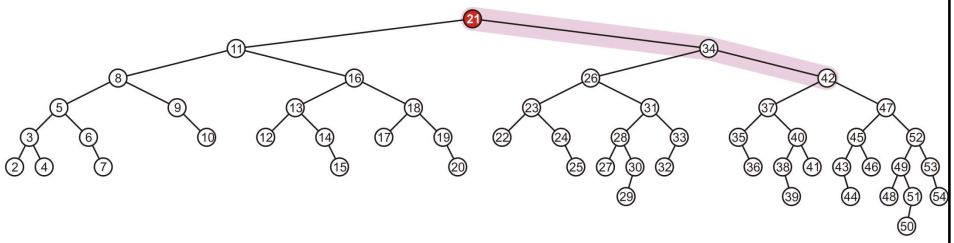
At this point, the node 11 is balanced



Erase

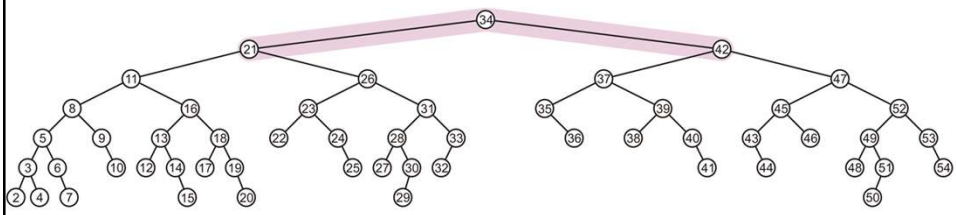
Still, the root node is unbalanced

- This is a right-right imbalance



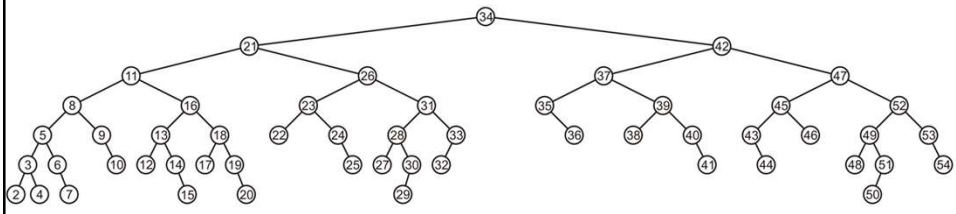
Erase

Again, a simple balance fixes the imbalance



Erase

The resulting tree is now AVL balanced



Implementation of AVL Tree

```
struct AvlNode
{
    ElementType Element;
    AvlTree  Left;
    AvlTree  Right;
    int      Height;
};
```

Because height of a AVL node is usually calculated, time complexity of height calculation should be O(1).

We add 1 field in struct AVL node: height of node

Summary

- AVL Tree is useful

	AVL Tree		Binary Search Tree	
	Average	Worst	Average	Worst
Insertion	O(logn)	O(logn)	O(logn)	O(n)
Deletion	O(logn)	O(logn)	O(logn)	O(n)
Searching	O(logn)	O(logn)	O(logn)	O(n)
Traversal	O(n)	O(n)	O(n)	O(n)

- If you have free time, please read Red-Black Tree, Slay Tree, and B-Tree
- Next week: Hashing and Priority Queues