DISCRETE MATH 37181 HOMEWORK 6

MURRAY ELDER

- 1. (a) Write down a precise definition for $f \in O(g)$ (g dominates f).
 - (b) Write down a precise definition for a function $f: A \to B$ to be one-to-one.
 - (c) Write down a precise definition for a function $f: A \to B$ to be onto.
 - (d) Write down a precise statement of the pigeonhole principle.
 - (e) Write down a precise statement of the generalised pigeonhole principle.
 - (f) Write down a precise definition for an equivalence relation on a set A.
 - (g) Write down a precise statement of a partition of a set A.
- 2. For each of the following functions, guess a Big-O form from one of those in Table 1 from the worksheet, then prove your guess.
 - (a) f(n) = 4n + 7
 - (b) $g(n) = 5n^2 + 3n \log_2 n$
 - (c) $h(n) = 1 + 2 + \dots + n$
- 3. Show that g dominates f in each of the following
 - (a) f(n) = 6n + 10, $g(n) = 0.05n^2$
 - (b) $f(n) = 3n^2$, $g(n) = 2^n + 2n$.
- 4. Prove that:

If $S \subseteq \mathbb{N}_+$ and |S| > 6 then S contains three distinct elements x, y, z such that $x + y + z \equiv 0 \mod 3$.

- 5. Let A,B be finite sets and $f:A\to B$ a function. Prove that if f is one-to-one then $|A|\leqslant |B|.$
- 6. Let A,B,C be sets and $f:A\to B,g:B\to C$. Recall that $g\circ f:A\to C$ is the function defined as $g\circ f(a)=g(f(a))$ for all $a\in A$.

Prove that:

If $f: A \to B, g: B \to C$ are one-to-one then $g \circ f$ is one-to-one.

Brief solutions:

- 1. (a) Let $f, g : \mathbb{N}_+ \to \mathbb{R}$. We say that g dominates f if there exist constants $m \in \mathbb{R}^+$ and $k \in \mathbb{Z}^+$ such that $|f(n)| \leq m|g(n)|$ for all $n \in \mathbb{N}, n \geq k$.
 - (b) $\forall x, y \in A \text{ if } f(x) = f(y) \text{ then } x = y.$
 - (c) $\forall y \in B \ \exists x \in A \ \text{such that} \ f(x) = y$.
 - (d) If m pigeons occupy n pigeonholes and m > n then some pigeonhole has at least two pigeons in it.
 - (e) If m pigeons occupy n pigeonholes and m > kn then some pigeonhole has more than k pigeons in it.
 - (f) A relation $\subseteq A \times A$ which is reflexive (includes (a, a) for every $a \in A$), symmetric (if $(a, b) \in \mathcal{R}$ then so is (b, a)) and transitive (if $(a, b), (b, c) \in \mathcal{R}$ then so is (a, c)).
 - (g) See the worksheet from last week: a set of sets A_i so that A is the union of all the A_i and $A_i \cap A_j = \emptyset$ for all $i \neq j$. So (imprecisely) a partition is a way of dividing up the set into disjoint smaller sets which cover the whole set.
- 2. (a) Guess O(n). There exists m = 5, k = 7 so that

$$f(n) = 4n + 7 \leqslant 4n + n \text{ (since } n \geqslant 7) = 5n$$

for all $n \ge 7$.

Guess $O(n^2)$ since the log is smaller than n. Proof: if $n \ge 1$ then $\log_2 n \le n$ since $n \le 2^n$ (prove this by induction to be very rigorous).

$$5n^2 + 3n\log_2 n \leqslant 5n^2 + 3n.n \text{ (since } n \geqslant \log_2 n) = 8n^2$$
 so $m = 8, k = 1.$

(c) Induction we know this equals $\frac{n(n+1)}{2} = \frac{1}{2}n^2 + \frac{1}{2}n$ so $O(n^2)$, quadratic:

$$\frac{1}{2}n^2 + \frac{1}{2}n \leqslant \frac{1}{2}n^2 + \frac{1}{2}n^2 = n^2$$

for any $n \ge 1$ so k = m = 1.

3. (a) $f(n) = 6n + 10 \le 6n + n$ (assuming $n \ge 10$) = $7n = \frac{7}{0.05}(0.05n) = 140(0.05n) \le 140(0.05n^2)$

if $n \ge 1$, so with k = 10 and m = 140 we have $f \in O(g)$.

Note there are many different ways to show this.

(b) I will show $3n^2 - 2n \le 2^n$ for all $n \ge ?$ by induction. So I will have m = 1 and k = ?.

The statement $P(n): 3n^2 - 2n \leq 2^n$ is true for n = 1, so we have P(1). Assume P(s). Then to show P(s+1) we have:

$$3(s+1)^{2} - 2(s+1) = 3(s^{2} + 2s + 1) - 2s - 2 = 3s^{2} + 6s + 3 - 2s - 2$$
$$= 3s^{2} - 2s + 6s + 1 \le 2^{s} + 6s + 1$$

by inductive assumption. To finish I need to show $6s + 1 \leq 2^s$, which I will prove separately below. However, this is only true for $s \geq ?$ so I will have to change the value of k = ? to k = 6 and start over.

Lemma 1. For all $n \ge 6$, $6n + 1 \le 2^n$

Proof. Induction: true for n=6 since $37 < 2^6 = 64$. (also true for smaller, but 6 is fine for Big O proofs.)

Assume true for $t \ge 6$, then

$$6(t+1) + 1 = 6t + 1 + 6 < 6t + 1 + 6t + 1$$

(since 6 < 6t + 1)

$$\leq 2^t + 2^t = 2^{t+1}$$
.

Then by PMI P(n) is true for all $n \ge 6$.

Now I start again, using the lemma. Let P(n) be the statement: $3n^2 - 2n \le 2^n$ P(6) since LHS = 3.36 - 12 = 96 and $RHS = 2^6 = 64$ Not true. So I change my value of k = ? again.

FINAL PROOF:

Let P(n) be the statement: $3n^2 - 2n \le 2^n$. P(10) is true since LHS = 3.100 - 20 = 280 and $RHS = 2^{10} = 1024$. Assume P(s). P(s+1):

$$LHS = 3(s+1)^2 - 2(s+1) = 3(s^2 + 2s + 1) - 2s - 2 = 3s^2 + 6s + 3 - 2s - 2$$
$$= 3s^2 - 2s + 6s + 1 \le 2^s + 6s + 1$$

by inductive assumption.

$$\leq 2^s + 2^s$$

by Lemma 1

$$=2.2^{s}=2^{s+1}$$

so by PMI P(n) is true for all $n \ge 10$.

4. Pigeons are the elements in S, and pigeonholes are the remainders of each number mod 3, so a box labeled 0, 1 and 2.

Since we have at least 7 pigeons going into 3 boxes, at least one box has 3 elements, call these x, y, z. If the three elements have the same remainder mod 3, then $x+y+z\equiv 0$ mod 3 (more detail: x=3p+i, y=3q+i, z=3r+i each with the same remainder i, then x+y+z=3(p+q+r+i).)

- 5. Suppose |A| > |B|. Let A be pigeons, B pigeonholes, and place each $a \in A$ into hole f(a). Then since |A| > |B| by PHP some hole has two elements of A in it, say $x, y \in A$, and f(x) = f(y) since they are placed in the same hole. This shows f is not one-to-one. The result is the contrapositive.
- 6. Suppose $g \circ f(x) = g \circ f(y)$ for some $x, y \in A$. Then g(f(x)) = g(f(y)) and since g is one-to-one, this means f(x) = f(y). But now since f is one-to-one this means x = y. This $g \circ f$ is one-to-one.